

New Techniques and Algorithms for Multiobjective and Lexicographic Goal-Based Shortest Path Problems

PhD Dissertation

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Outline

I. Introduction

II. State of the art

III. Algorithmic contributions

IV. Empirical analyses

V. Conclusions & future work

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- Motivation
- Research goals

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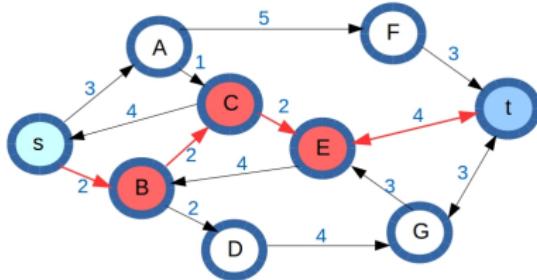
V. Conclusions & future work

The Shortest Path Problem (SPP)

Definition

The SPP consists in finding a minimal cost path between two nodes in a graph, where the graph is composed of arcs labeled with the cost of the transition between nodes.

- Solved by Dijkstra in 1959 and by Hart et al. (A^*) in 1968 exploiting specific problem knowledge.



The Shortest Path Problem (SPP)

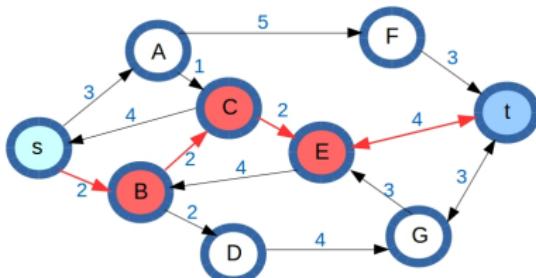
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Research fields

- Artificial Intelligence
- Operational Research



Applications to the SPP

- Path planning in game maps.
 - The game map is represented as a graph.
 - The characters movement through the map is solved as a SPP.



Applications to the SPP

- Motion planning in mobile robots.
 - The terrain is represented as a graph.
 - The robot must find a path avoiding obstacles.



Applications to the SPP

- Route planning in road maps.
 - Real road networks are graphs.
 - Arcs represent road junctions.
 - Arcs are typically labeled with the distance between junctions.



Multicriteria Search Problem (MSP)

Definition

MSP is the natural extension of SPP to the multicriteria case, where arcs are labeled with **vectors** instead of scalar values.

SPP	→	single-objective optimization problem
Decision problems	→	multiple conflicting criteria

Human decisions

Humans make choices according to their personal preferences, which involve simultaneously multiple, and often, conflicting criteria.



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New car?

- Price
- Speed
- Brand
- Color

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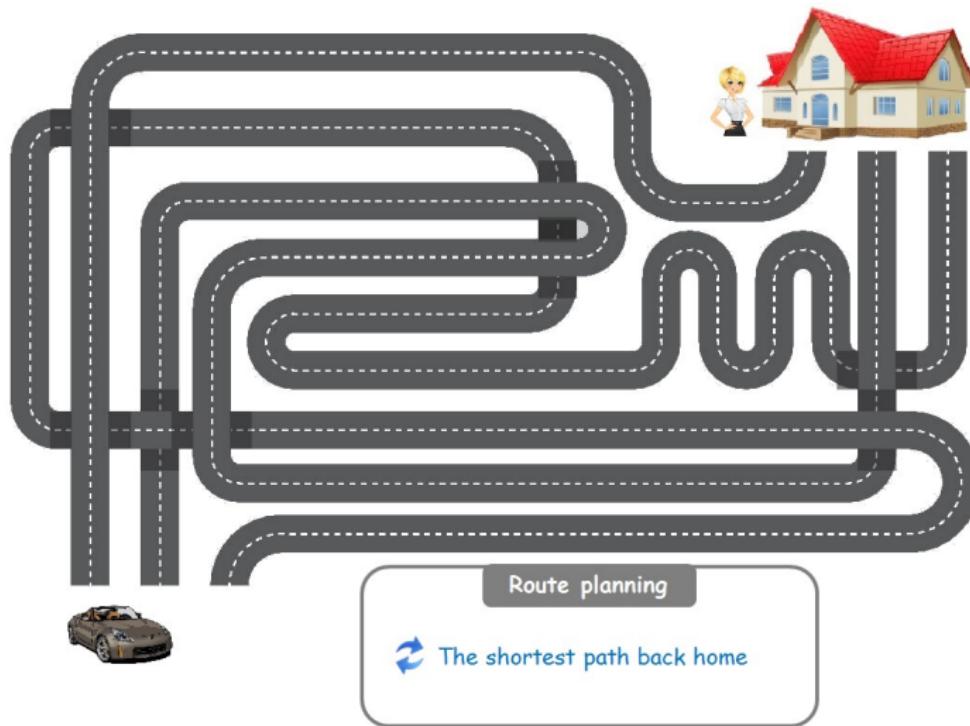
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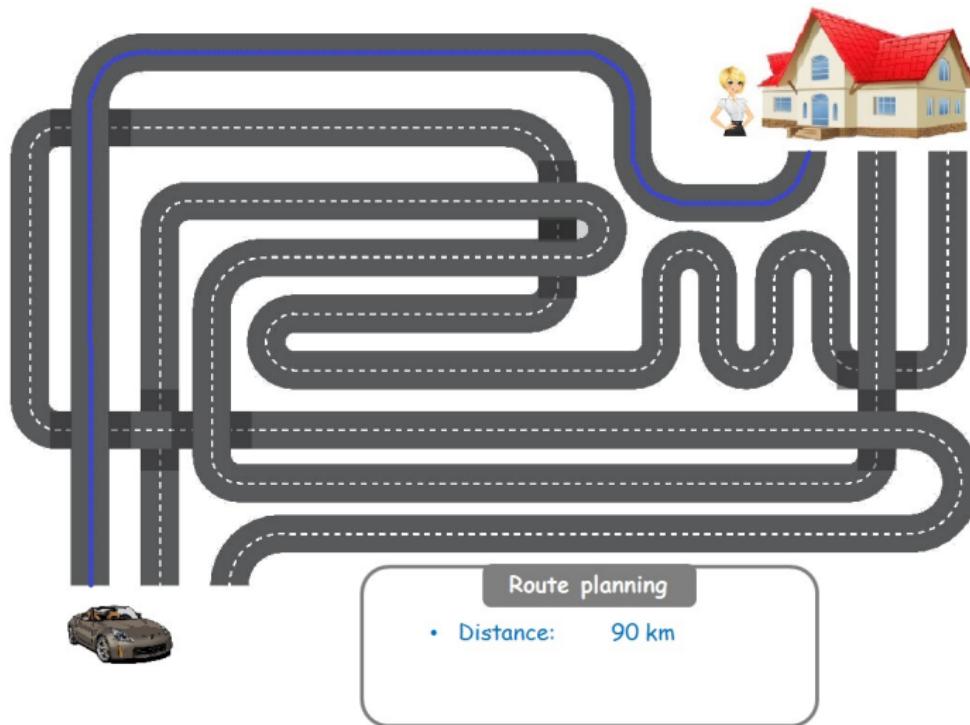
A route home?

- Distance?

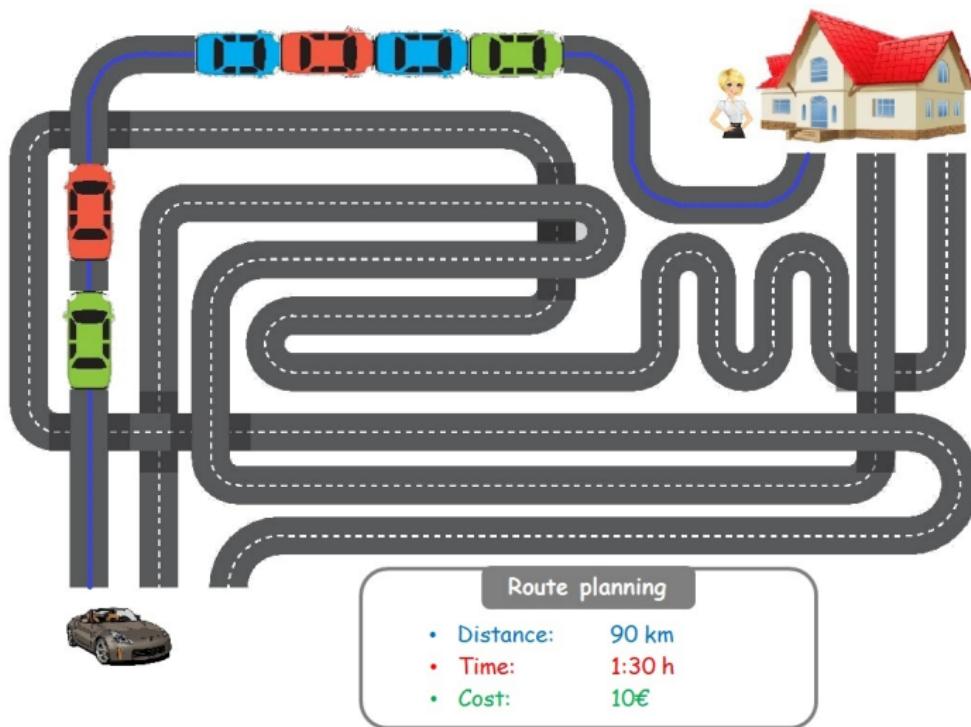
Route planning in road maps



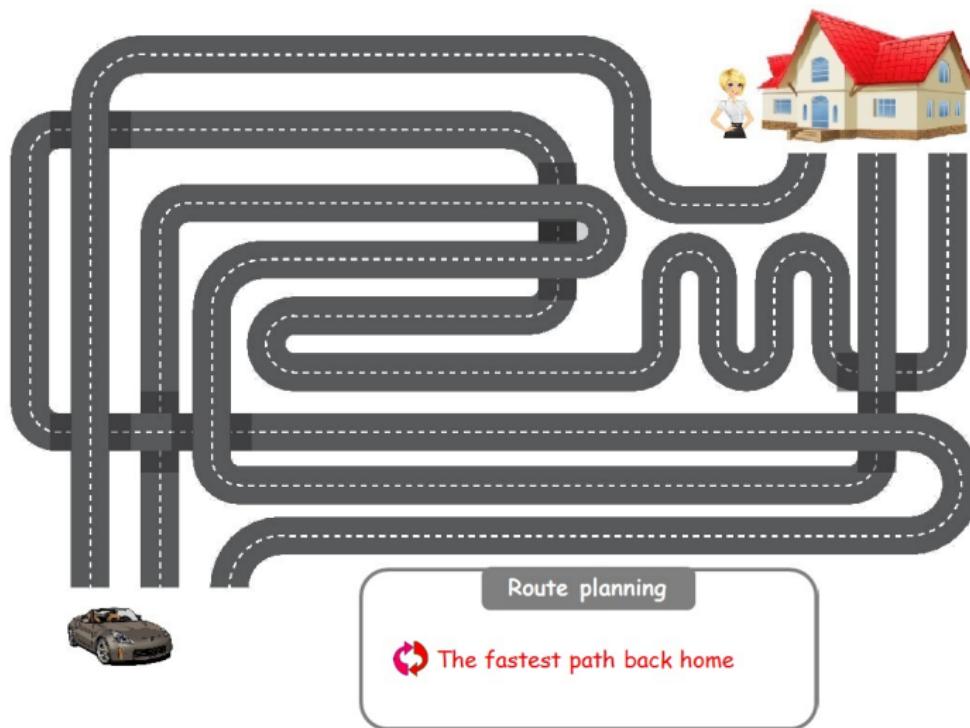
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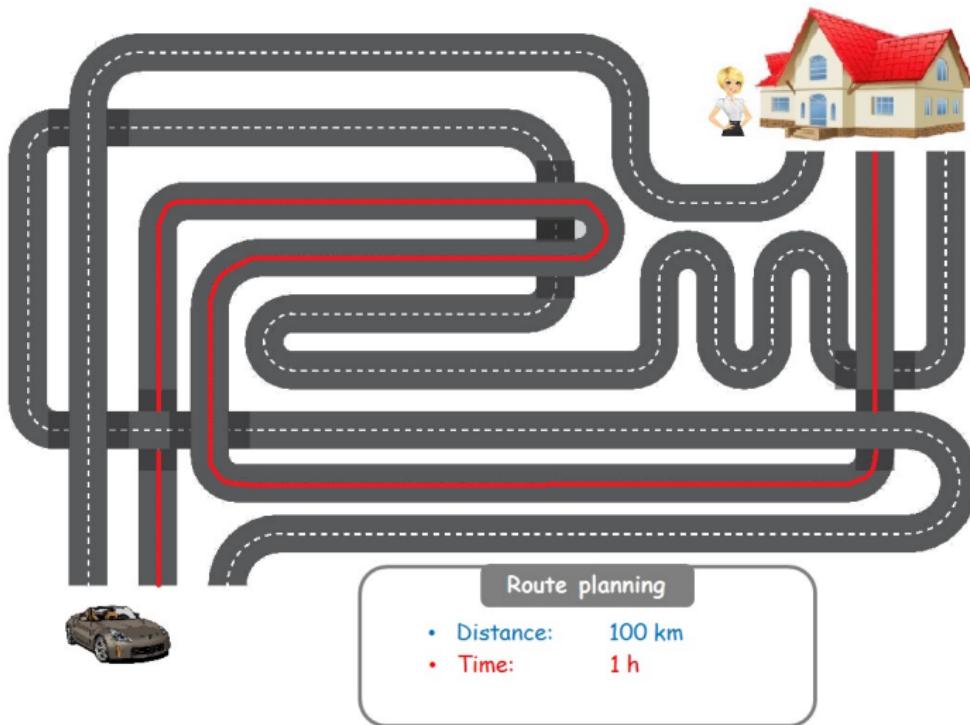
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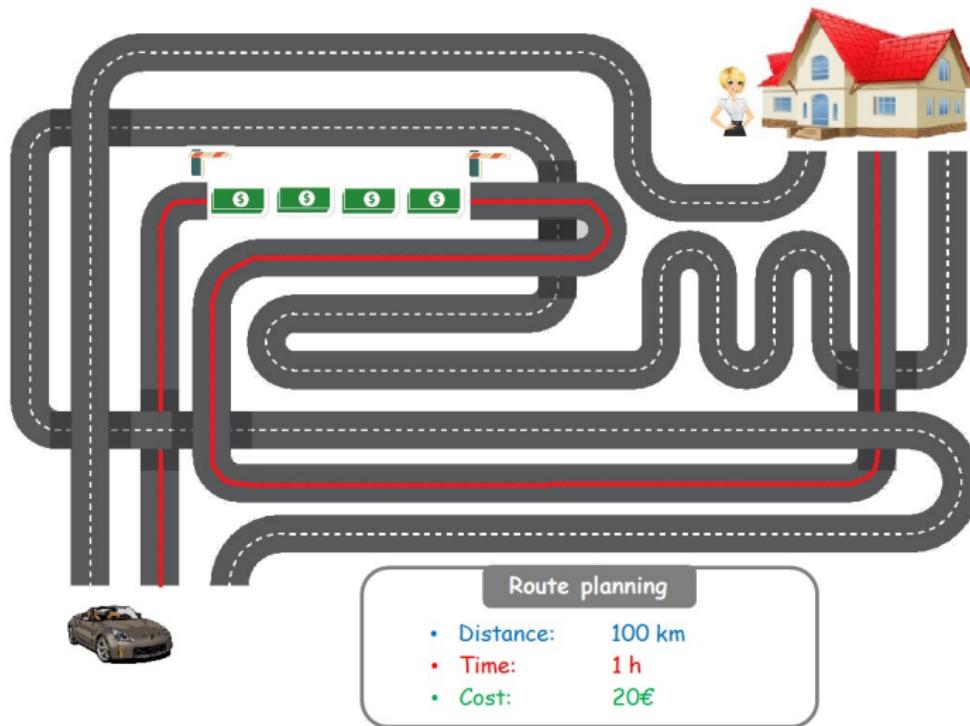
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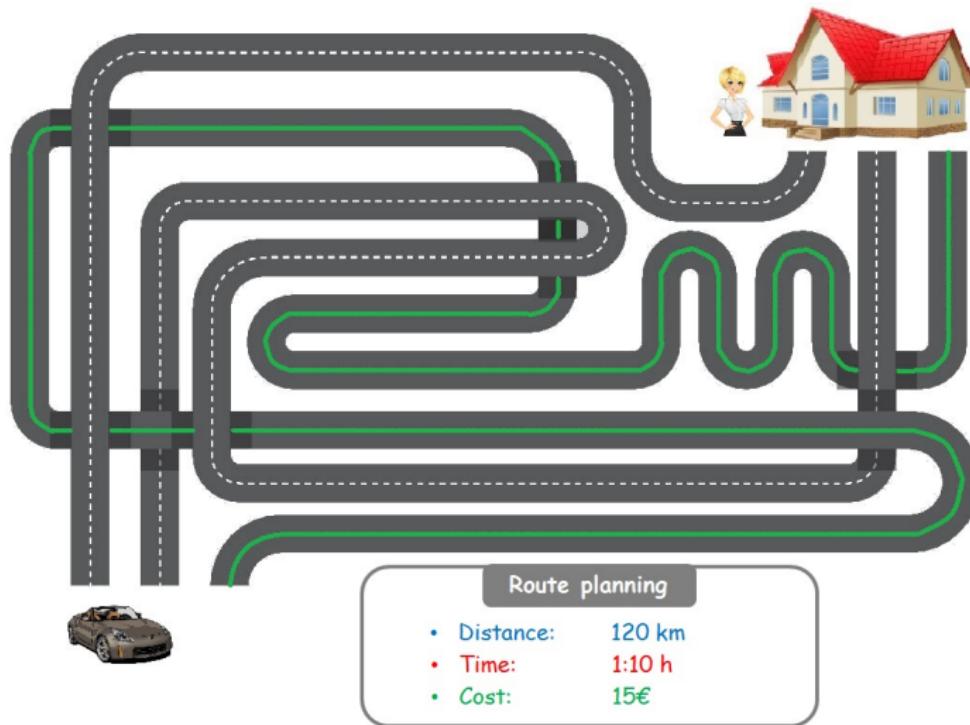
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Multicriteria Search Problem (MSP)

Definition

Dominance or Pareto preference \prec :

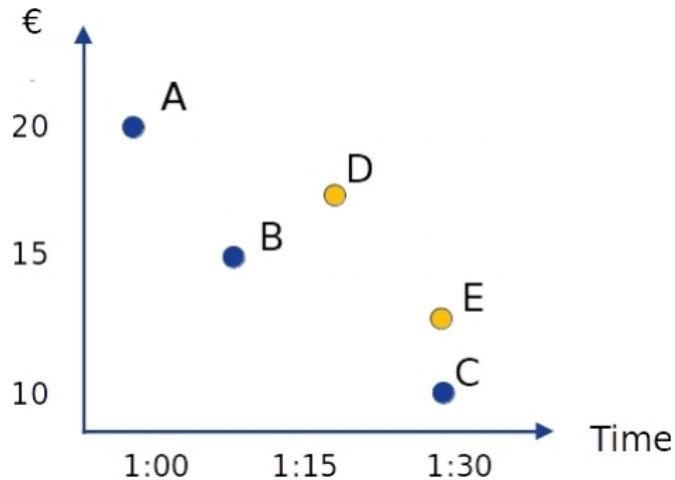
$$\vec{y} \prec \vec{y}' \Leftrightarrow \forall i (1 \leq i \leq q) \quad y_i \leq y'_i \wedge \vec{y} \neq \vec{y}'$$

where $\vec{y} = (y_1, y_2, \dots, y_q)$ represents a solution path in the cost space.

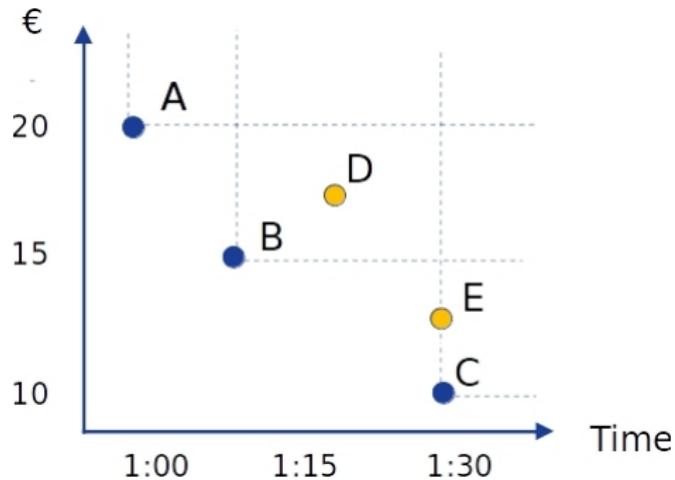
Example

Route cost representation: $\vec{y} = (y_1, y_2, y_3) \rightarrow \vec{y} = (120, 70, 15)$

Pareto optimality



Pareto optimality

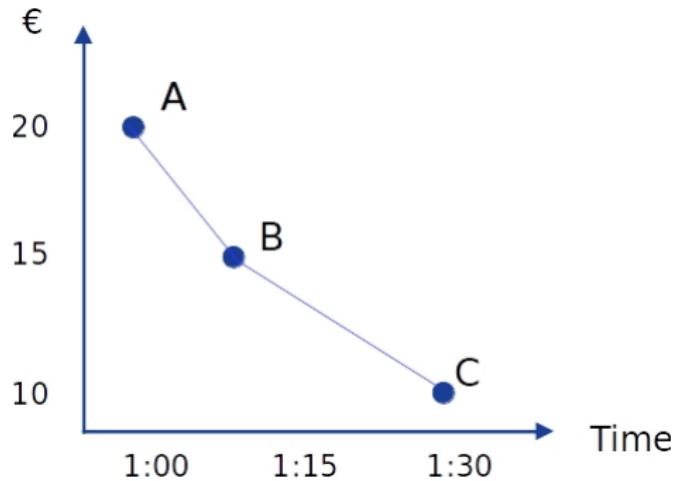


Multicriteria Search Problem (MSP)

Definition

The solution to a MSP, **without any knowledge of the user preferences**, is the set of **non-dominated** solution paths, or **Pareto frontier**.

Pareto optimality



Goal Programming

Definition

Goal Programming (GP) is one of the **most successful techniques** to solve multicriteria problems. It is also a **natural way** to express user preferences as targets or aspiration levels: $y_i \leq t_i$.

Definition

Lexicographic Goal Programming (LGP) **ranks all goals in order of importance**, dividing them into lexicographic levels of importance.

Definition

The deviation from a goal in a minimization problem represents a **positive distance** between the i -th goal and its aspiration level:
 $d_i = \max(0, y_i - t_i)$.

Goal Programming

Example

Three criteria (economic cost (euros), travel time (minutes), and distance (km)) grouped in two priority levels where level 1 is infinitely more important than level 2. The importance of the criteria within a level is defined by weights as:

$$\begin{aligned} \text{Level 1: } & \text{cost} \leq 15, & w_1 = 0.5 \\ & \text{time} \leq 75, & w_2 = 0.5 \\ \text{Level 2: } & \text{distance} \leq 120, & w_3 = 1. \end{aligned}$$

Example

Let's assume a solution path with cost $\vec{y} = (10, 60, 150)$, its deviation from those goals is $\vec{d} = (0, 30) \rightarrow (\max(0, (10 - 15) \times 0.5) + \max(0, (60 - 75) \times 0.5), (150 - 120) \times 1)$

Goal Programming

Example

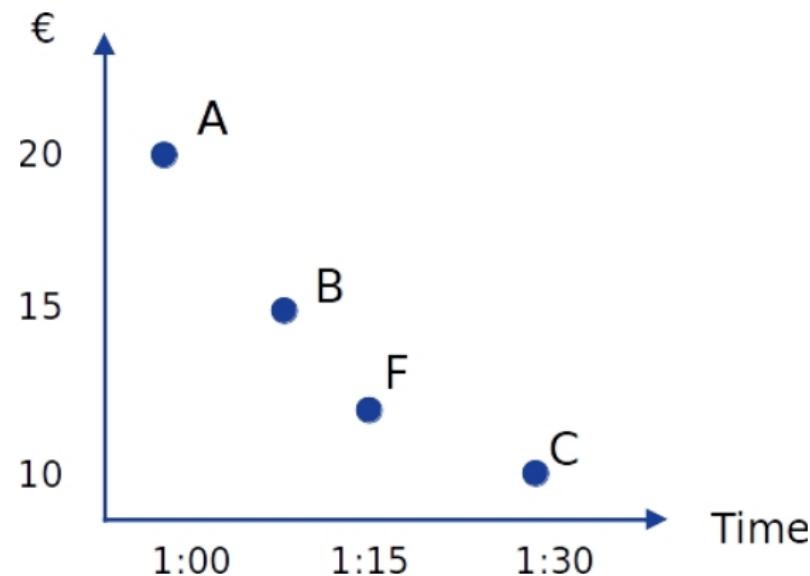
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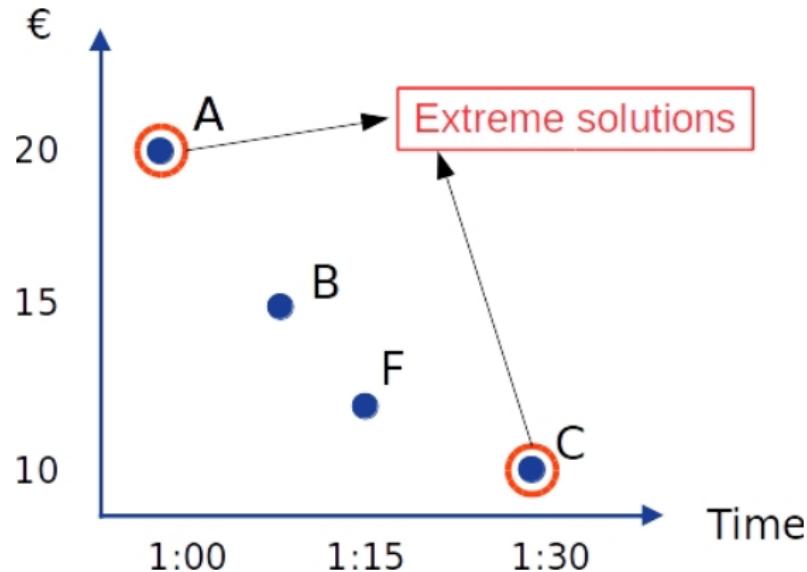
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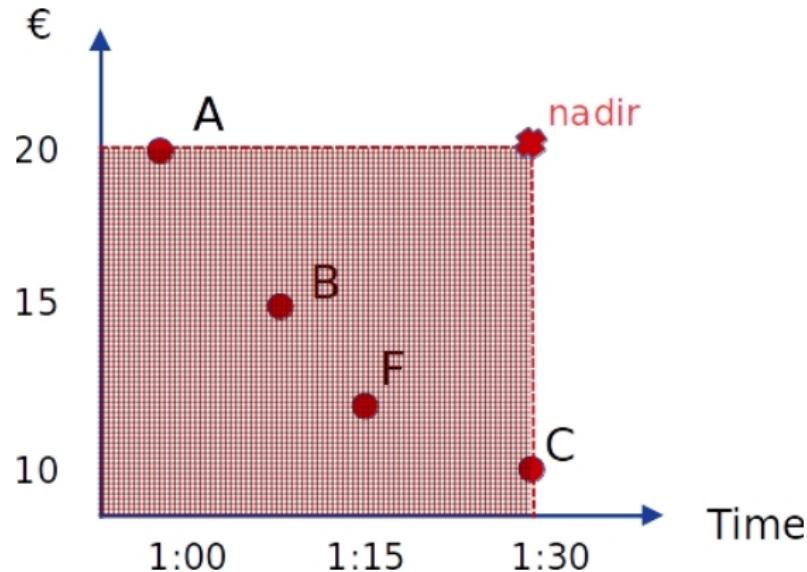
Lexicographic goals



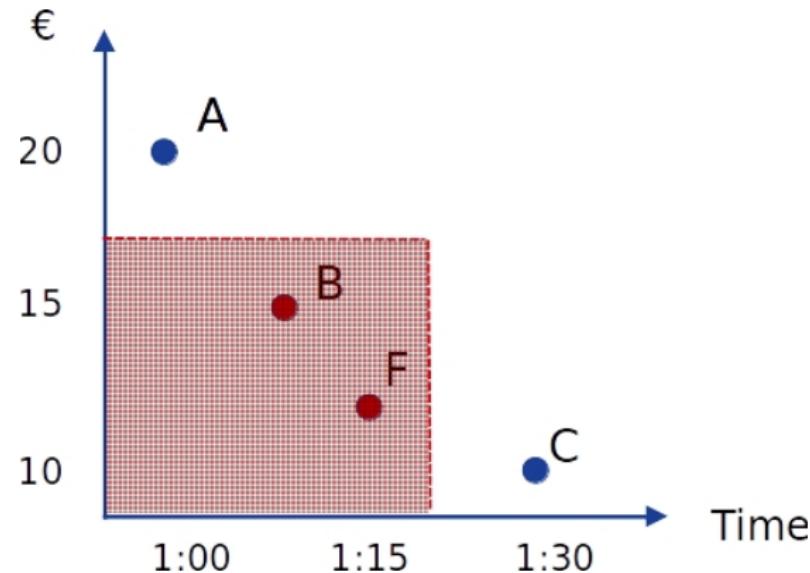
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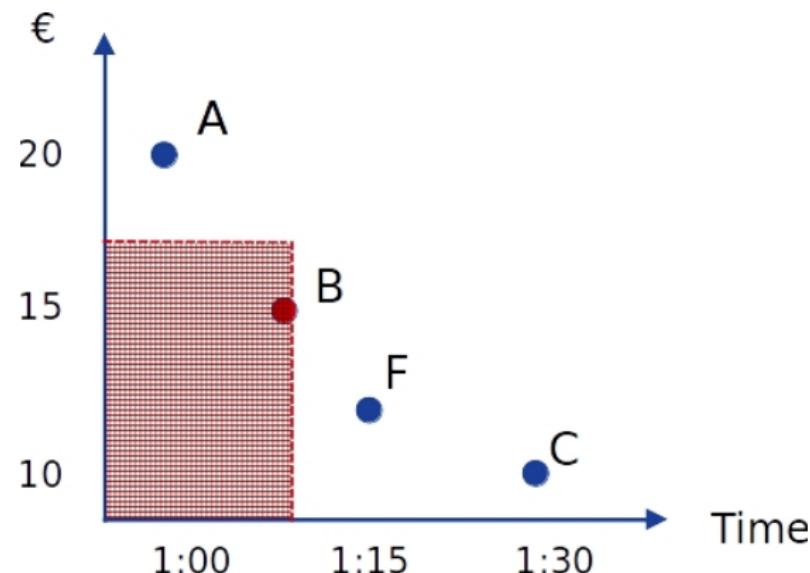
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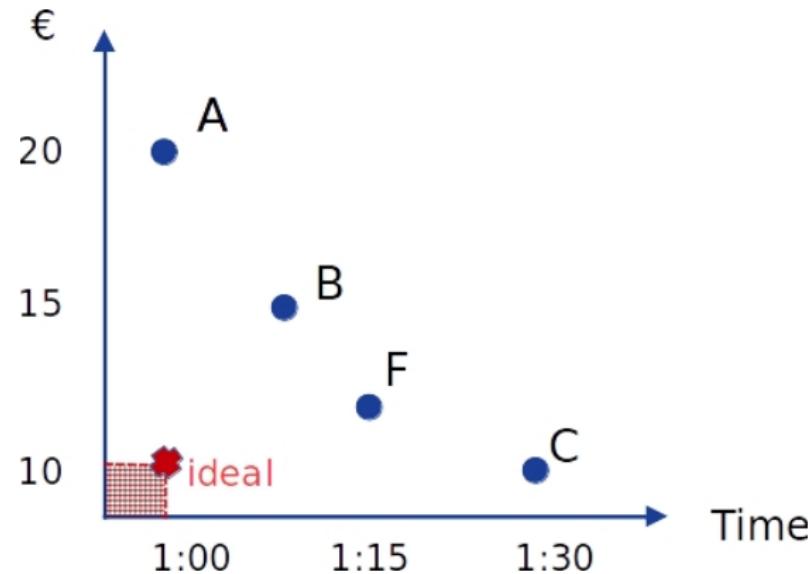
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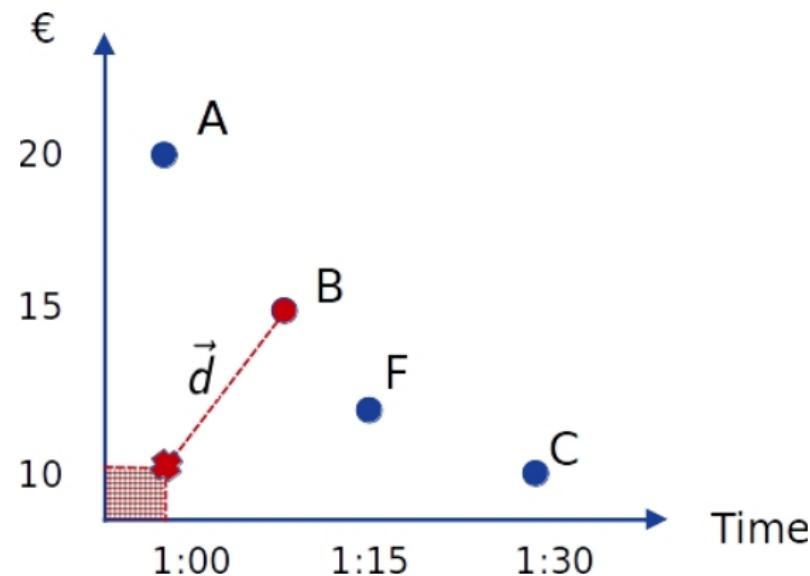
Lexicographic goals



Lexicographic goals



Lexicographic goals



Definition

We define **lexicographic goal preferences** (\prec_G) as a partial order relation,

$$\vec{y} \prec_G \vec{y}' \Leftrightarrow \vec{d}(\vec{y}) \prec_L \vec{d}(\vec{y}') \vee (\vec{d}(\vec{y}) = \vec{d}(\vec{y}') \wedge \vec{y} \prec \vec{y}')$$

MSP with lexicographic goals

Definition

The solution to a MSP with lexicographic goals is the **set of non-dominated solution paths** that satisfy the goals, or the subset that **minimizes deviation** if these cannot be fully satisfied.

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Research goals

Tackle the MSP with lexicographic goals from two algorithmic perspectives:

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A posteriori Preferences of the decision maker are provided **after** the execution of the algorithm.

A priori Preferences of the decision maker are provided **before** the execution of the algorithm.

Research goals

Tackle the MSP with lexicographic goals from two algorithmic perspectives:

A posteriori Preferences of the decision maker are provided **after** the execution of the algorithm.

A priori Preferences of the decision maker are provided **before** the execution of the algorithm.

- **A posteriori** algorithms obtain the **whole Pareto frontier** and then, extract the goal-optimal solutions.
- **A priori** algorithms will return **only goal-optimal** solutions.

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Properties of multicriteria algorithms

Definition

An algorithm is **admissible** if it guarantees to find the optimal solution to the problem.

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Definition

We will measure the **efficiency** of algorithms according to the **number of explored labels** to find the solution.

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A posteriori algorithms

Extensions of A^* to the multiobjective case to calculate the Pareto frontier:

MOA* (Stewart & White, 1991)

TC (Tung & Chew, 1992)

NAMOA* (Mandow & Pérez de la Cruz, 2005)

All can be used with heuristic functions as lower bounds

A posteriori algorithms

Extensions of A^* to the multiobjective case to calculate the Pareto frontier:

MOA* (Stewart & White, 1991) → pathological behaviour

TC (Tung & Chew, 1992) → less efficient than NAMOA*

NAMOA* (Mandow & Pérez de la Cruz, 2005)

All can be used with heuristic functions as lower bounds

NAMOA* properties

Analogously to A^* :

- Admissible when provided with a consistent heuristic function.
- It explores an optimal number of labels in its class.
- It improves its efficiency with more informed lower bounds.

Relevant features of NAMOA*:

- Label selection policy
- Two sets of labels for each node n : $G_{op}(n)$ and $G_{cl}(n)$.
- A set COSTS of solutions.
- Discarding rules of dominated paths

Label selection policy

NAMOA* can be used with any path selection policy that assures the **best label** according to that policy **is a non-dominated label**.

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- Linear aggregation order

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$$\vec{y} \prec_L \vec{y}' \Leftrightarrow \exists j \ (1 \leq i \leq q) \ y_j < y'_j \wedge \forall i < j \ y_i = y'_i.$$

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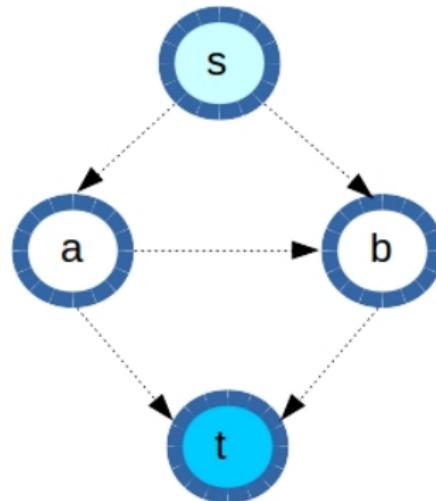
- Linear aggregation order

$$\vec{y} \prec_{lin} \vec{y}' \Leftrightarrow \sum_i y_i < \sum_i y'_i, \quad 1 \leq i \leq q$$

Discarding dominated paths on NAMOA*

NAMOA* discards early in the search those paths that will not lead to non-dominated solutions.

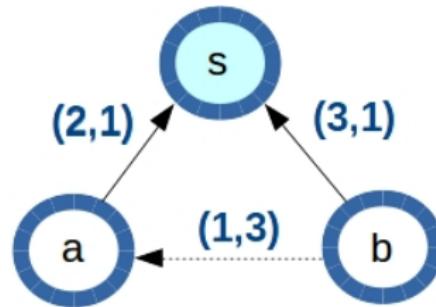
- Op-pruning
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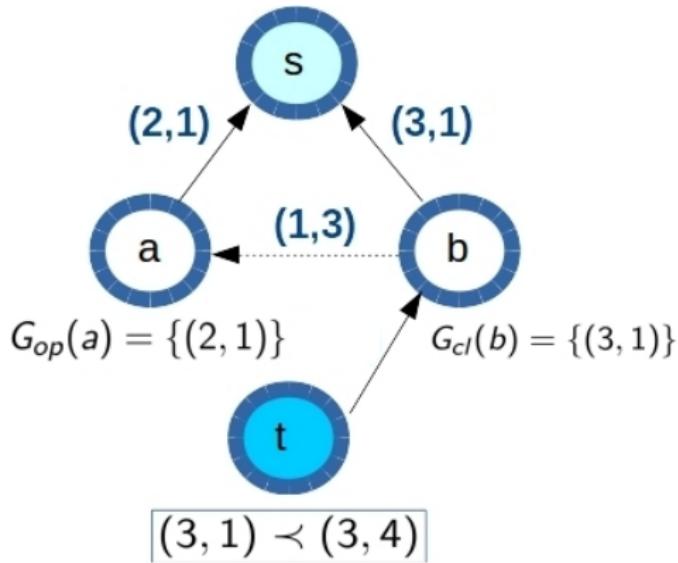
$$G_{op}(a) = \{(2, 1)\} \quad G_{op}(b) = \{(3, 1)\}$$

$$(3, 1) \prec (3, 4)$$

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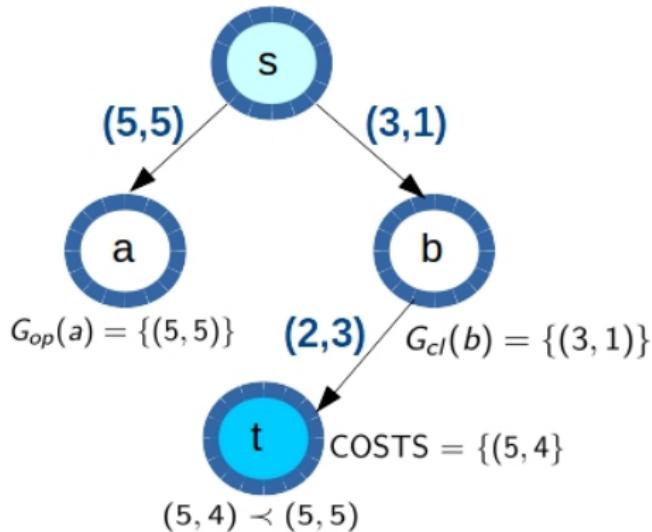
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A priori algorithms

Is there any specifically devised algorithm for lexicographic goals?

METAL-A* (Mandow & Pérez de la Cruz, 2001)

Based on MOA* → pathological behaviour

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→ Then, let's propose a new approach based on NAMOA*

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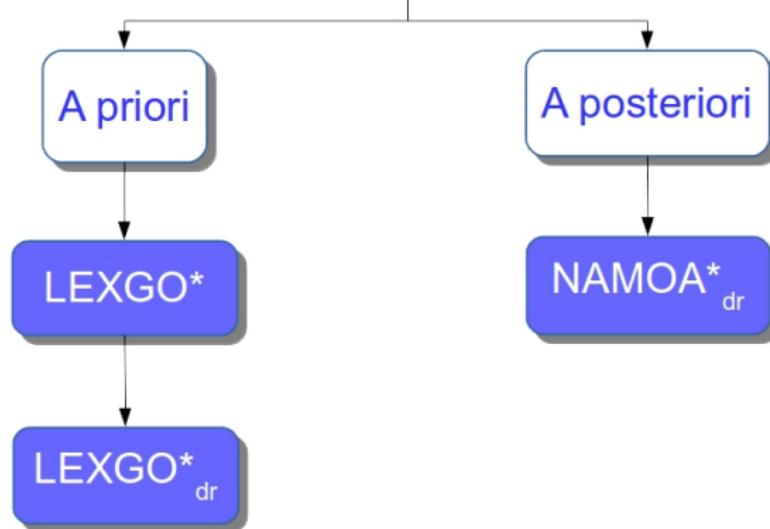
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- LEXGO*
- T-discardng
- NAMOA_{dr}*
- LEXGO_{dr}*

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Our contributions to the MSP with lexicographic goals



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Lexicographic goal preferences

Important!

Bellman's Principle of Optimality, **an optimal path is made up of suboptimal paths**, holds for multiobjective search problems, but **it does not hold for lexicographic goal preferences!**

Lexicographic goal preferences

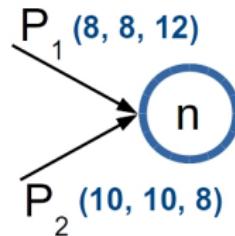
Example

Two paths P_1 and P_2 reach some node n with costs $\vec{g}(P_1) = (8, 8, 12)$ and $\vec{g}(P_2) = (10, 10, 8)$, respectively. The user goals are defined as follows:

$$\text{Level 1: } g_1 \leq 10, \quad w_1 = 0.5$$

$$g_2 \leq 10, \quad w_2 = 0.5$$

$$\text{Level 2: } g_3 \leq 10, \quad w_3 = 1.$$



Example

The deviation from goals for P_1 and P_2 is $\vec{d}(P_1) = (0, 2)$ and $\vec{d}(P_2) = (0, 0)$, respectively. Thus, $P_2 \prec_G P_1$.

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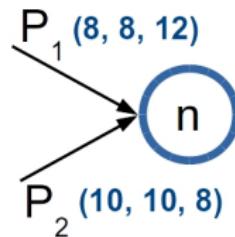
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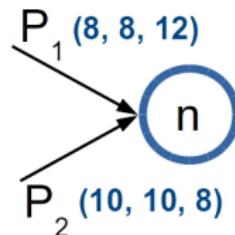
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Level 1: $g_1 \leq 10, w_1 = 0.5$

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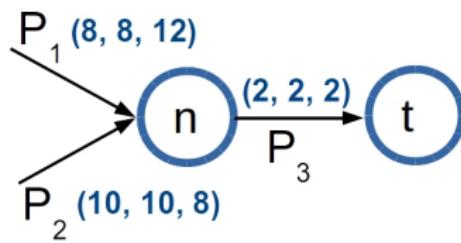
Example

Let's now consider there is only one additional path P_3 to the destination node t , such as $\vec{g}(P_3) = (2, 2, 2)$.

Therefore,

- $\vec{g}(P_1 P_3) = (10, 10, 14)$ and $\vec{g}(P_2 P_3) = (12, 12, 10)$.
- $\vec{d}(P_1 P_3) = (0, 4)$ and $\vec{d}(P_2 P_3) = (2, 0)$.
- Now, $P_1 P_3 \prec_G P_2 P_3$

We have pruned an optimal path!



Lexicographic goal preferences

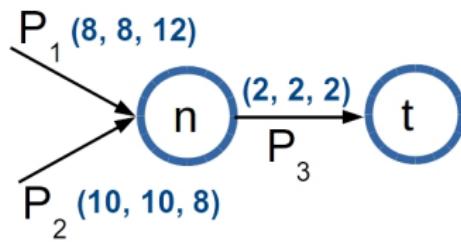
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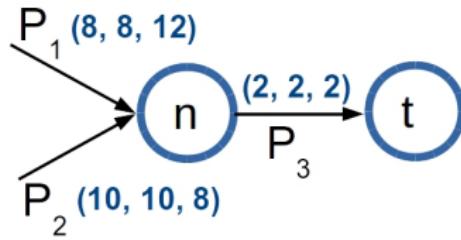
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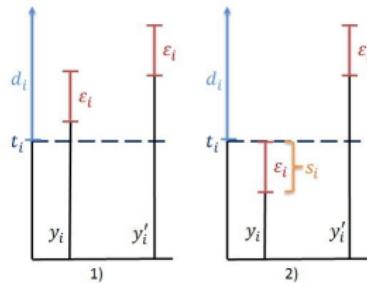
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Let's now assume paths P_1 and P_2 have costs $\vec{g}(P_1) = (12, 8, 16)$ and $\vec{g}(P_2) = (12, 12, 6)$, respectively. The user goals are now *different* and defined as follows:

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Example

The deviation from goals for P_1 and P_2 is now $\vec{d}(P_1) = (2, 3)$ and $\vec{d}(P_2) = (2, 1)$, respectively. $P_2 \prec_G P_1$, but, can P_1 be **safely** pruned?

Lexicographic goal preferences

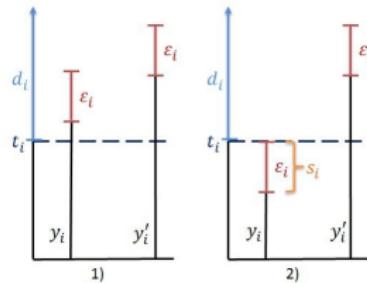
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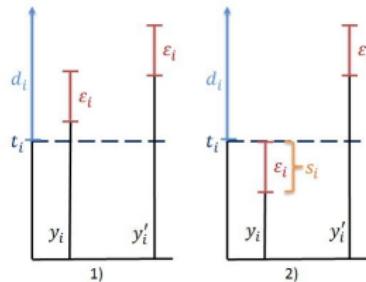
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$$g_3 \leq 10, \quad w_3 = 0.5.$$



Example

The deviation from goals for P_1 and P_2 is now $\vec{d}(P_1) = (2, 3)$ and $\vec{d}(P_2) = (2, 1)$, respectively. $P_2 \prec_G P_1$, but, can P_1 be **safely** pruned?

Discarding labels by deviation

Deviation-based pruning (\prec_P)

Slack variable

$$s_k = \max(0, t_k - y_k)$$

Cross slacks

$$\delta_j(\vec{y}, \vec{y}') = \sum_{k \in I_j} w_k \times \max(0, s'_k - s_k)$$

Example

Given $\vec{g}(P_1) = (12, 8, 16)$ and $\vec{g}(P_2) = (12, 12, 6)$,
where $\vec{d}(P_1) = (2, 3)$ and $\vec{d}(P_2) = (2, 1)$:

- $s_2(P_1) = (10 - 8) \times 0.5 = 1$.
- $s_3(P_2) = (10 - 6) \times 0.5 = 2$.
- $\delta_2(P_1, P_2) = 2 - 1 = 1$,

Finally, $d_2(P_2) - d_2(P_1) = 3 - 1 = 2 > 1$.

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Discarding labels by deviation

Deviation-based pruning (\prec_P)

$$\vec{y} \prec_P \vec{y}' \Leftrightarrow \forall i < j (d_i(\vec{y}) = d_i(\vec{y}') \wedge \delta_i(\vec{y}, \vec{y}') = 0) \wedge \exists j (d_j(\vec{y}) < d_j(\vec{y}') \wedge \delta_j(\vec{y}, \vec{y}') < d_j(\vec{y}') - d_j(\vec{y}))$$

Deviation-based filtering

$$\vec{d}_B \prec_L \vec{d}$$

Discarding labels by deviation

Deviation-based pruning (\prec_P)

$$\vec{y} \prec_P \vec{y}' \Leftrightarrow \forall i < j (d_i(\vec{y}) = d_i(\vec{y}') \wedge \delta_i(\vec{y}, \vec{y}') = 0) \wedge \exists j (d_j(\vec{y}) < d_j(\vec{y}') \wedge \delta_j(\vec{y}, \vec{y}') < d_j(\vec{y}') - d_j(\vec{y}))$$

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Deviation-based filtering

$$\vec{d}_B \prec_L \vec{d}$$

In addition to the pruning and filtering rules of NAMOA*, LEXGO* also uses its own rules to prune and filter paths by deviation.

LEXGO* shares NAMOA* formal properties

- The **new** deviation-based pruning and filtering rules **do not discard** goal-optimal solutions.
- LEXGO* **always** returns the set of all goal-optimal solutions.
- LEXGO* expands a **subset** of the labels expanded by NAMOA*.

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T-discardng

Lmiting factor in MSP

Runtime, rather than memory, is the limiting factor in multicriteria search algorithms performance, as well as in the size of problems that can be practically solved.

Lmiting factor in MSP (II)

The majority of the algorithm runtime can be attributed to dominance checks against G_{op} , G_{cl} and COSTS to discard dominated paths.

T-discardng

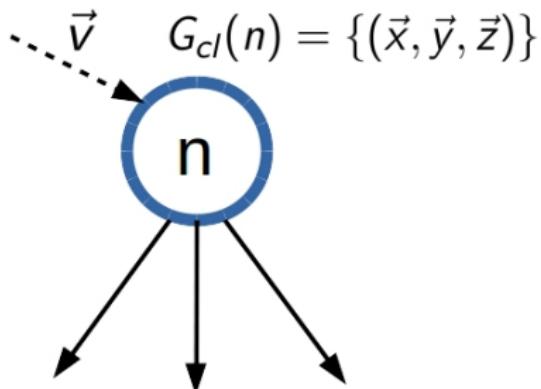
Let's suppose a typical scenario where a new path with cost \vec{v} reaches some node n . This node has already three closed paths with costs $(\vec{x}, \vec{y}, \vec{z})$. Let's now check if \vec{v} is dominated by any vector in $G_{cl}(n)$:

Example

- $\vec{x} \prec \vec{v} ? \rightarrow x_1 \leq v_1, x_2 \leq v_2 \dots$
- $\vec{y} \prec \vec{v} ? \rightarrow y_1 \leq v_1, y_2 \leq v_2 \dots$
- $\vec{z} \prec \vec{v} ? \rightarrow z_1 \leq v_1, z_2 \leq v_2 \dots$

Worst-case scenario:

- 3 vector comparisons
- 9 scalar comparisons



T-discardng

Definition

T-discardng (*truncated discarding*) is a dimensionality reduction technique to speed up dominance checks.

Assumptions

To apply t-discard an algorithm must have:

- Lexicographic order as label selection policy.
- One consistent heuristic function as lower bound.

T-discard

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To apply t-discard an algorithm must have:

- Lexicographic order as label selection policy.
- One consistent heuristic function as lower bound.

Example

When **lexicographic order** is used to select open paths:

- $\rightarrow (3, 5, 7)$
- $\rightarrow (4, 6, 4)$
- $\rightarrow (4, 6, 7)$
- $\rightarrow (6, 2, 4)$

T-discard

Assumptions

To apply t-discard an algorithm must have:

- Lexicographic order as label selection policy.
- One consistent heuristic function as lower bound.

Example

When **lexicographic order** is used to select open paths:

- $\rightarrow (3, -, -)$
- $\rightarrow (4, -, -)$
- $\rightarrow (4, -, -)$
- $\rightarrow (6, -, -)$

T-discard

Assumptions

To apply t-discard an algorithm must have:

- Lexicographic order as label selection policy.
- One consistent heuristic function as lower bound.

Property

Assuming the last selected open path had cost $\vec{c} = c_1, c_2, \dots, c_q$, the next path selected to be open with cost $\vec{c}' = c'_1, c'_2, \dots, c'_q$ **must have** $c'_1 \geq c_1$!

T-discardng

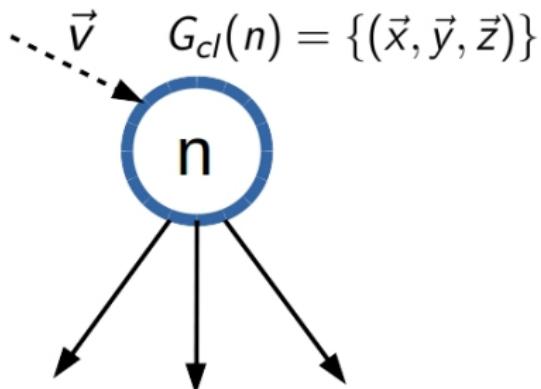
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Example

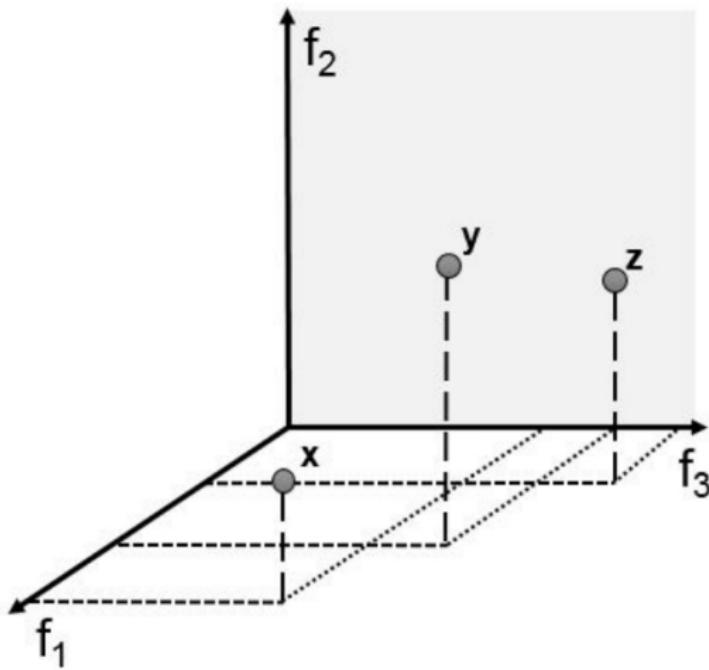
- $\vec{x} \prec \vec{v} ? \rightarrow x_1 \leq v_1, x_2 \leq v_2 \dots$
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Worst-case scenario:

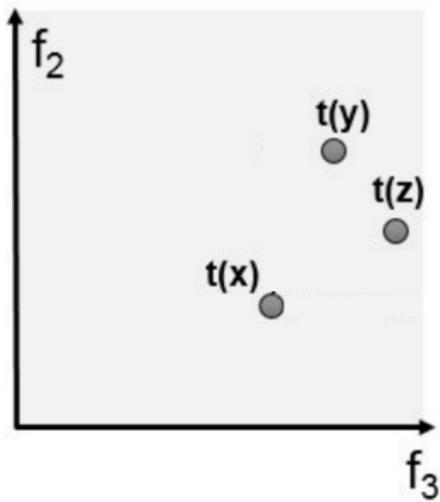
- 3 vector comparisons
- 9 scalar comparisons



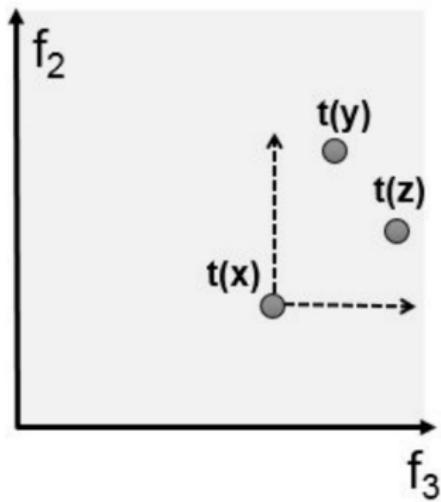
T-discriminating



T-discriminating



T-discriminating



T-discardng

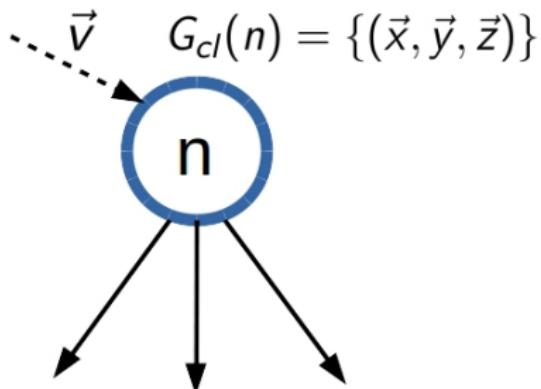
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Worst-case scenario:

- 3 vector comparisons
- 9 scalar comparisons



T-discardng

Is a vector \vec{v} dominated by any vector in the previous set $(\vec{x}, \vec{y}, \vec{z})$?

Standard discarding

\vec{v} must be compared to $(\vec{x}, \vec{y}, \vec{z})$ to check if it is dominated.

T-discardng

The truncated vector $t(\vec{v})$ must be compared only to $t(\vec{x})$ to check if it is dominated.

T-discard

Benefits of t-discard

- Dominance checks can be done with vectors reduced in one dimension.
- The size of sets $T(G_{cl}(n))$ and $T(COSTS)$ is also reduced.

Important!

T-discard only implies checking dominance against different sets:

- Very simple implementation.
- Any label discarded by standard pruning \iff is discarded by t-discard.
- Does not affect any formal property of the algorithm.

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The applicability of t-discarding to NAMOA* is **full** and straightforward:

CI-pruning and filtering are now performed against:

- $T(G_{cl}(n))$
- T(COSTS)

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The applicability of t-discarding to LEXGO* is **partial**:

CI-pruning and filtering are now performed against:

- $T(G_{cl}(n))$
- T(COSTS)

but **only when $\vec{d} = 0$.**

Summary

	NAMOA*	LEXGO*	NAMOA _{dr} *	LEXGO _{dr} *
Op-pruning	$G_{op}(n)$	$G_{op}(n)$	$G_{op}(n)$	$G_{op}(n)$
Cl-pruning	$G_{cl}(n)$	$G_{cl}(n)$	$T(G_{cl}(n))$	$T(G_{cl}(n))^*$
Filtering	COSTS	COSTS	T(COSTS)	T(COSTS)*

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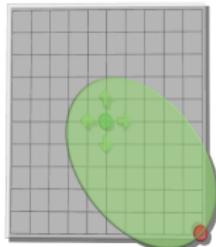
Empirical analyses

- Are these new algorithms effective?
- Will LEXGO* become faster when goals get more restrictive?
- Have we achieved improvements over NAMOA*?
To what extent? Why?

Empirical analyses

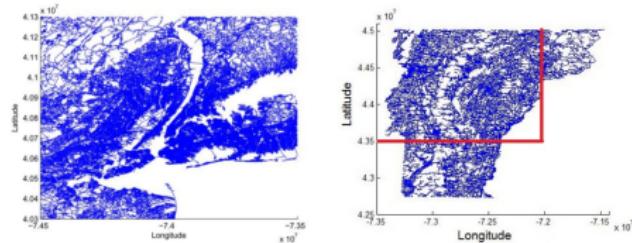
Random grids

- 100×100
- Solution depth from 20 to 100
- Random costs in range [1,10]



Realistic Road maps

- 20 random problems over New York City & Vermont State from 9th DIMACS Challenge
- Three integer costs
 - Distance
 - Travel time
 - Economic cost



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A priori

LEXGO* vs NAMOA*

LEXGO* vs NAMOA*

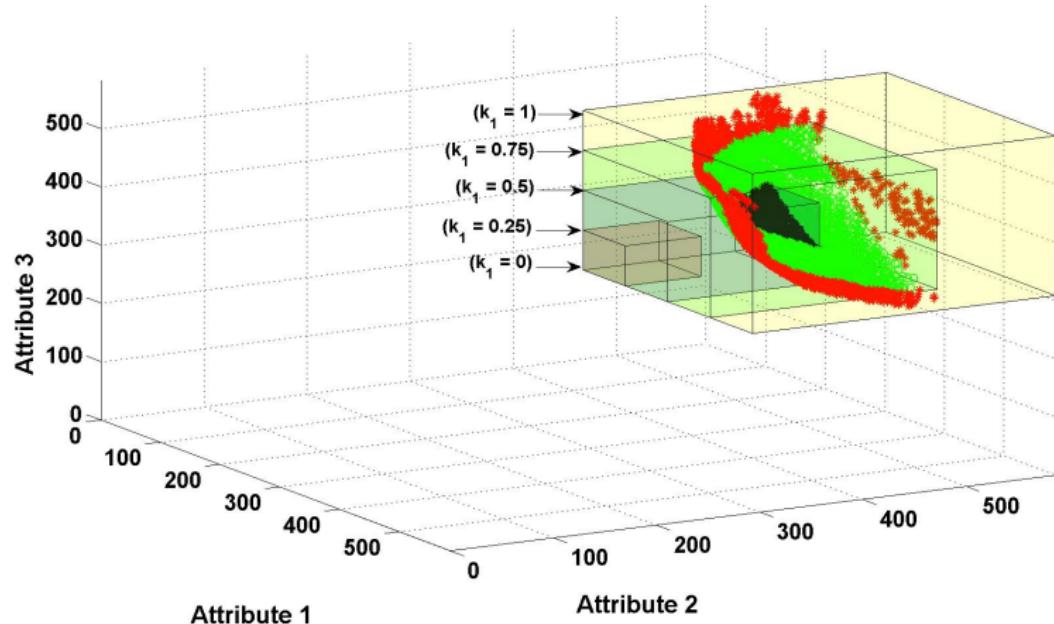


Figure: 3D Pareto frontier according to goal satisfiability.

LEXGO* vs NAMOA*

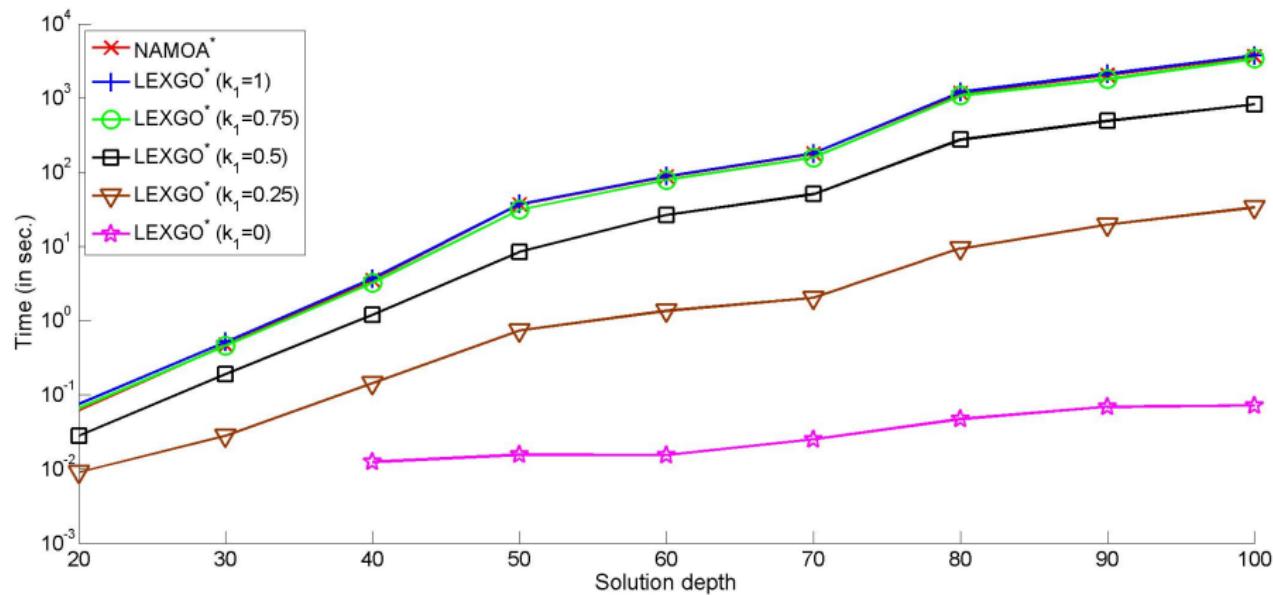


Figure: Runtimes of NAMOA* and LEXGO* in seconds (logarithmic scale).

LEXGO* vs NAMOA*

Table: Runtimes for $d = 100$ experiments.

		LEXGO* _{lex}				
NAMOA* _{lex}		1	0.75	0.5	0.25	0
Runtime (s)	%	%	%	%	%	%
3,662.9	102.7	92.3	22.3	0.9	0.001	

A priori

LEXGO* vs NAMOA*

LEXGO_{dr}* vs LEXGO*

$\text{LEXGO}_{\text{dr}}^*$ vs LEXGO^*

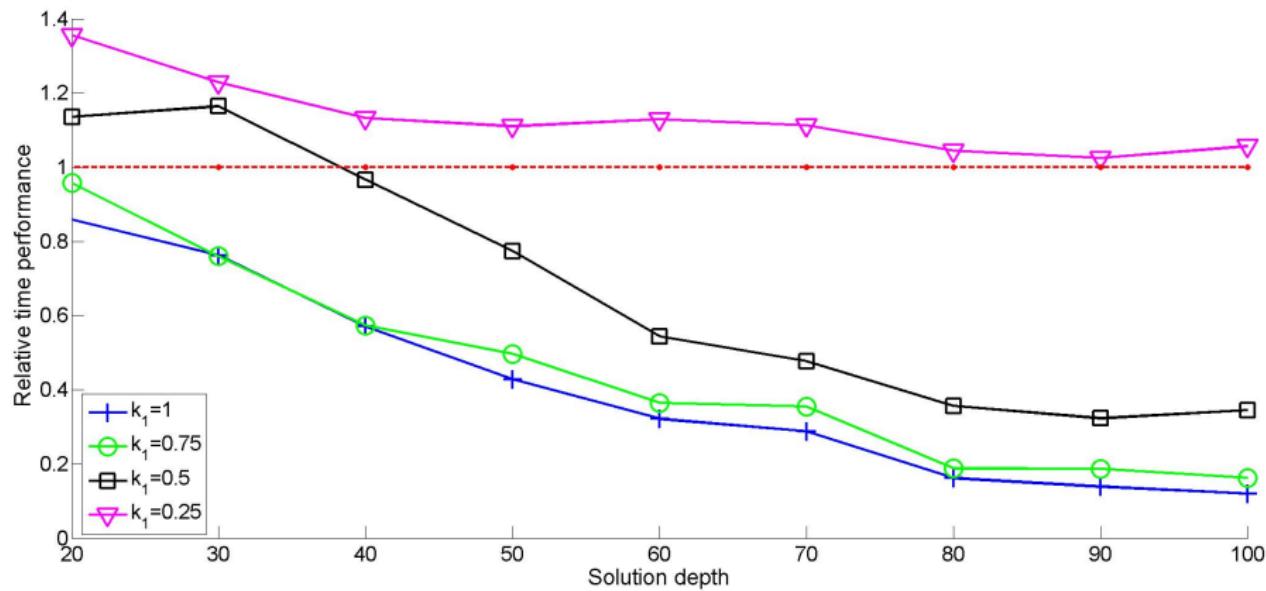


Figure: Relative runtime performance of $\text{LEXGO}_{\text{dr}}^*$ over LEXGO^* .

$\text{LEXGO}_{\text{dr}}^*$ vs LEXGO^*

Table: Runtimes in seconds for $d = 100$ experiments.

k_1	$\text{LEXGO}_{\text{lex}}^*$	$\text{LEXGO}_{\text{dr}}^*$	Speedup
1	3,763.78	254.40	14.81
0.75	3,381.33	300.86	11.27
0.5	819.28	282.49	2.90
0.25	33.47	35.37	0.94
0	0.07	0.10	0.7

A posteriori

NAMOA^{*}_{dr} vs NAMOA^{*}

NAMOA_{dr}^{*} vs NAMOA^{*}

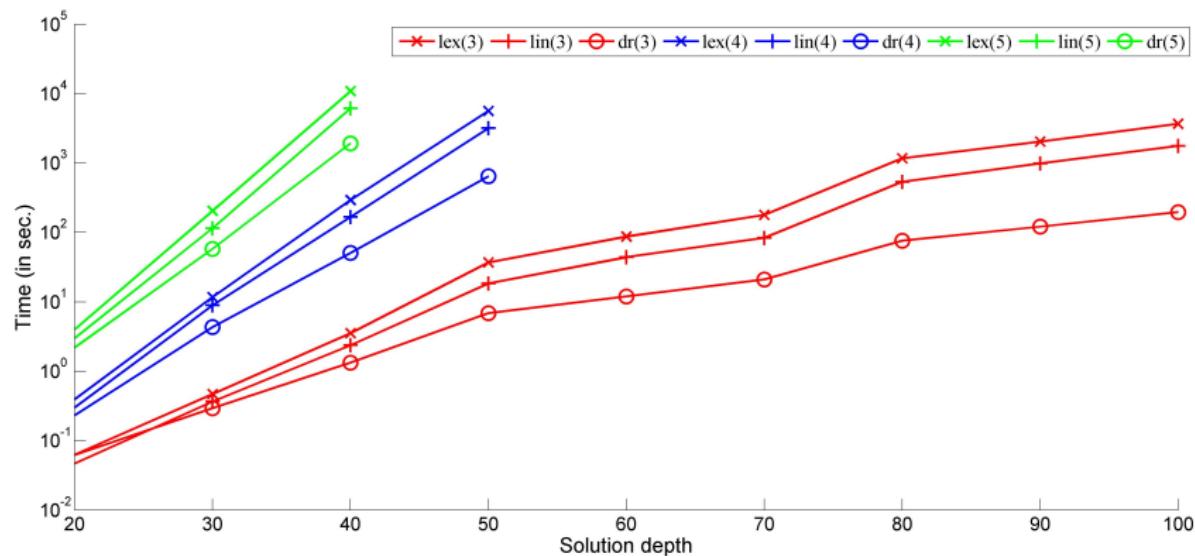


Figure: Runtimes for NAMOA_{lex}^{*}, NAMOA_{lin}^{*}, and NAMOA_{dr}^{*} with $q = \{3, 4, 5\}$

A posteriori vs A priori

NAMOA^{*}_{dr} vs LEXGO^{*}_{dr}

$\text{LEXGO}_{\text{dr}}^*$ vs $\text{NAMOA}_{\text{dr}}^*$

Table: Runtimes of $\text{LEXGO}_{\text{dr}}^*$ and $\text{NAMOA}_{\text{dr}}^*$.

	$\text{LEXGO}_{\text{dr}}^*$					
$\text{NAMOA}_{\text{dr}}^*$	$(k_1 = 1)$	$(k_1 = 0.75)$	$(k_1 = 0.5)$	$(k_1 = 0.25)$	$(k_1 = 0)$	
Runtime (s)	100%	129.7%	153.4%	144%	18%	0.04%

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A priori

LEXGO* vs NAMOA*

LEXGO* vs NAMOA*

Table: Experiments for Vermont State map.

		LEXGO* _{lex}				
NAMOA* _{lex}		1	0.75	0.5	0.25	0
Avg. runtime (s)	%	%	%	%	%	%
5,048.47	100.39	74.67	29.06	3.51	0.01	

A priori

LEXGO* vs NAMOA*

LEXGO_{dr}* vs LEXGO*

$\text{LEXGO}_{\text{dr}}^*$ vs LEXGO^*

Table: Runtimes in seconds for Vermont State map.

	1	0.75	0.5	0.25	0	k_1
$\text{LEXGO}_{\text{lex}}^*$	5,068.00	3,769.51	1,467.00	177.33	0.03	
$\text{LEXGO}_{\text{dr}}^*$	94.90	83.59	57.38	157.02	0.03	

A posteriori

NAMOA^{*}_{dr} vs NAMOA^{*}

NAMOA_{dr}^{*} vs NAMOA^{*}

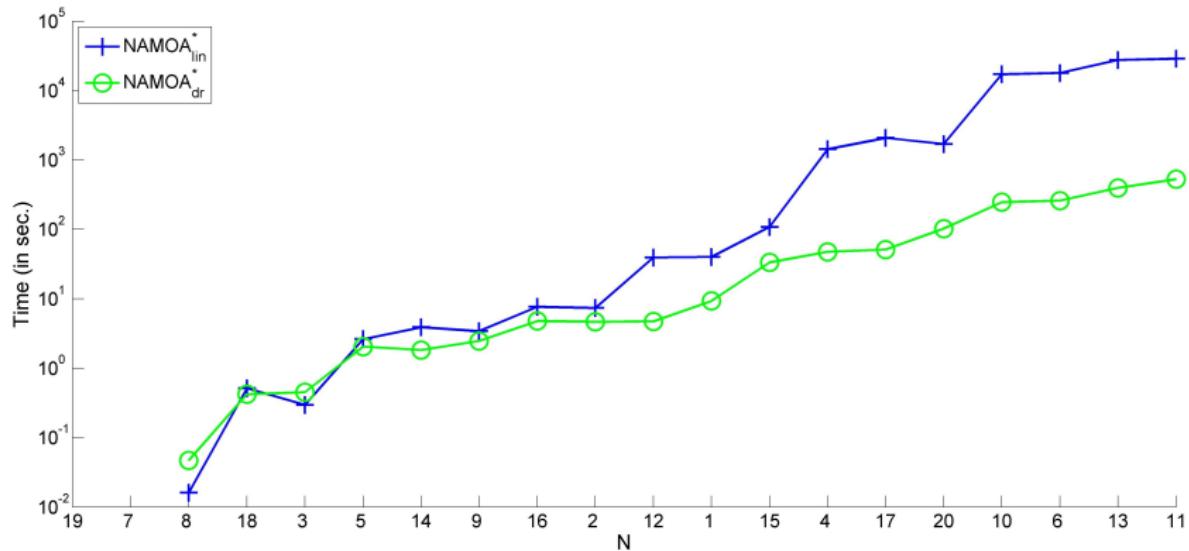


Figure: Runtimes of NAMOA_{lin}^{*} and NAMOA_{dr}^{*} for Vermont map problems.

A posteriori vs A priori

NAMOA^{*}_{dr} vs LEXGO^{*}_{dr}

$\text{LEXGO}_{\text{dr}}^*$ vs $\text{NAMOA}_{\text{dr}}^*$

Table: Runtimes in seconds of $\text{LEXGO}_{\text{dr}}^*$ and $\text{NAMOA}_{\text{dr}}^*$ in Vermont State map problems.

$\text{NAMOA}_{\text{dr}}^*$	1	0.75	0.5	0.25	0	k_1
84.66	94.89	83.59	57.38	157.02	0.03	

Why do the new algorithms have such a runtime performance?

LEXGO* vs NAMOA*

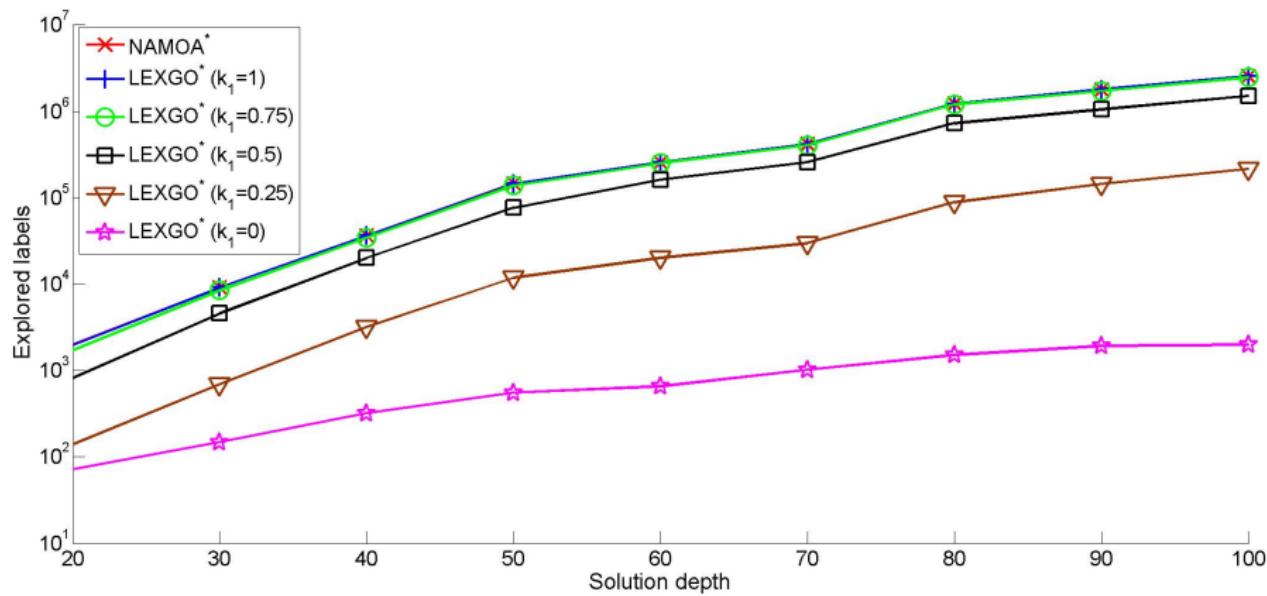


Figure: Explored labels by NAMOA* and LEXGO* in grids.

Deviation-based pruning

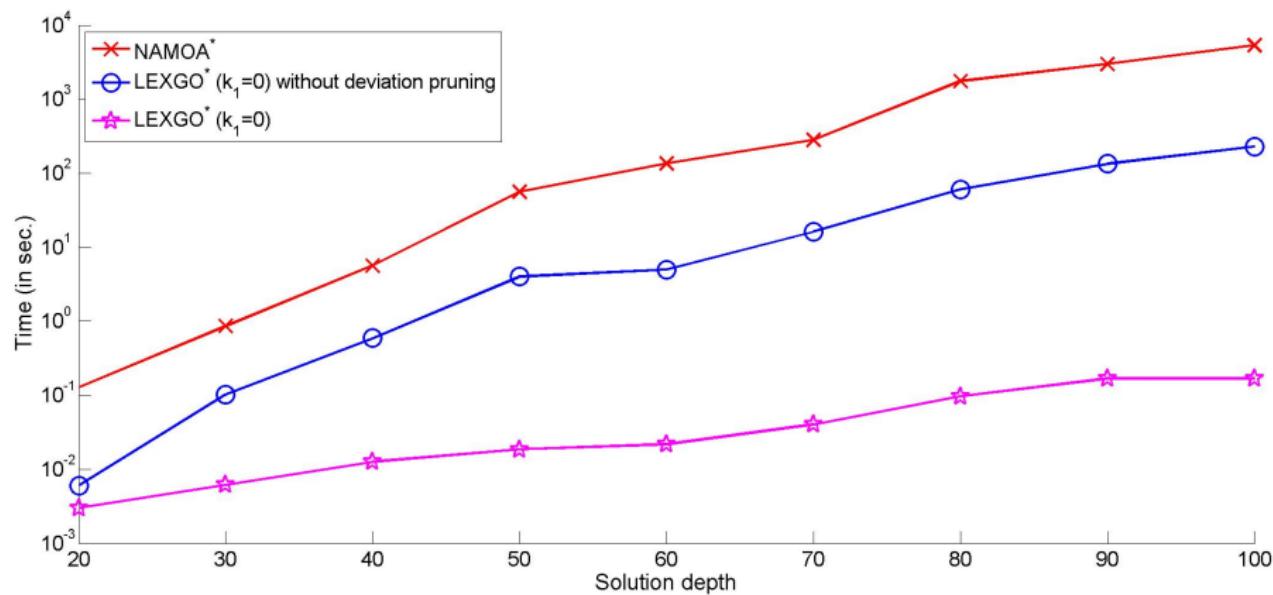


Figure: Impact of deviation pruning in $k_1 = 0$ experiments.

NAMOA_{dr}^{*} vs NAMOA^{*}

Table: Set sizes for Vermont State map.

Map	$(\frac{\sum T(G_{cl})}{\sum G_{cl}}) \%$	$(\frac{\sum T(C^*)}{\sum C^*}) \%$
VT _{cut}	3.43	2.45

NAMOA_{dr}^{*} vs NAMOA^{*}

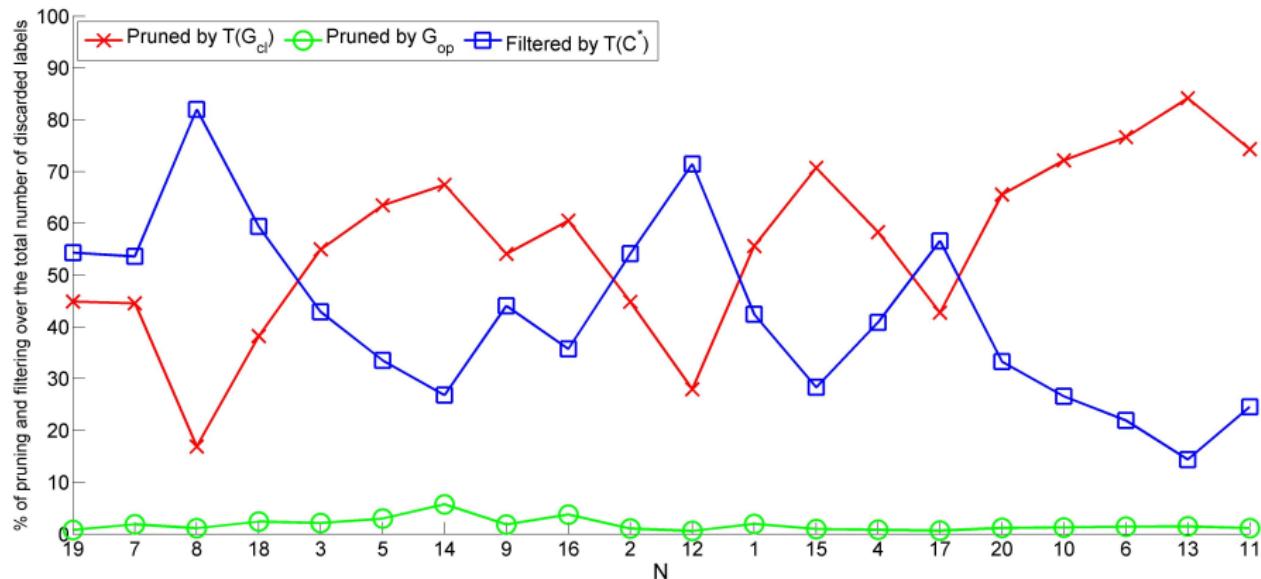


Figure: Discarded labels in Vermont problems.

New York experiments

Now, we can even solve **more** problems!

Problem #	NAMOA _{dr} [*]	NAMOA _{lin} [*]
#11	30 min.	31 hours
#6	3 hours	25 days

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Conclusions

- ① We have tackled the MSP with lexicographic preferences from two approaches: *a priori* and *a posteriori*.
- ② On the *a priori* approach, we have contributed two new algorithms: LEXGO* and LEXGO_{dr}*
- ③ On the *a posteriori* approach: NAMOA_{dr}* based on NAMOA*.
- ④ All algorithmic approaches have been formally proved to be admissible.
- ⑤ The contributions of this thesis outperform the state of the art in an order of magnitude for medium size problems and two orders of magnitude for the hardest problems.
- ⑥ NAMOA_{dr}* LEXGO_{dr}* represent the new state of the art in MSP with lexicographic goals.

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Future work

- ① Investigate more formal proofs about LEXGO*, concerning the optimality in its class of algorithms and the expansion of labels when using more informed lower bounds.
- ② Extend LEXGO* to be used with other GP models or other formulas to measure the deviation from goals.
- ③ Research the applicability of t-discard to other multiobjective or multicriteria search algorithms with lower bounds.
- ④ Bring multicriteria search algorithms to real route planning applications.

Thank you for your attention !!!

Gracias por su atención !!!

New Techniques and Algorithms for Multiobjective and Lexicographic Goal-Based Shortest Path Problems

PhD Dissertation

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Supervised by: Dr. Lawrence Madow

Dept. Lenguajes y Ciencias de la Computación
E.T.S. Ingeniería Informática
Universidad de Málaga

July 7, 2015

