```
// Hopcroft - Karp O(M*sqrt(N))
const int MAXV = 1001;
const int MAXV1 = 2*MAXV;
int N, M;
vector<int> ady[MAXV];
int D[MAXV1], Mx[MAXV], My[MAXV];
bool BFS(){
  int u, v, i, e;
  queue<int> cola;
  bool f = 0;
  for (i = 0; i < N+M; i++) D[i] = 0;</pre>
  for (i = 0; i < N; i++)
  if (Mx[i] == -1) cola.push(i);
  while (!cola.empty()) {
    u = cola.front(); cola.pop();
    for (e = ady[u].size()-1; e >= 0; e--) {
      v = ady[u][e];
      if (D[v + N]) continue;
      D[v + N] = D[u] + 1;
      if (My[v] != -1) {
        D[My[v]] = D[v + N] + 1;
        cola.push(My[v]);
      else f = 1;
    }
  return f;
}
int DFS(int u) {
  for (int v, e = ady[u].size()-1; e >=0; e--) {
    v = ady[u][e];
    if (D[v+N] != D[u]+1) continue;
    D[v+N] = 0;
    if (My[v] == -1 || DFS(My[v])) {
      Mx[u] = v; My[v] = u;
      return 1;
    }
  }
  return 0;
```

```
int Hopcroft_Karp(){
  int i, flow = 0;
  for (i = max(N,M); i >=0; i--)
    Mx[i] = My[i] = -1;
  while (BFS())
    for (i = 0; i < N; i++)
        if (Mx[i] == -1 && DFS(i))
        ++flow;
  return flow;
}</pre>
```

```
// Heavy Light Descomposition
int N, M;
vector<int> V[MN];
vector<int> G[MN];
vector<bool> L[MN];
/// cant- la cantidad de nodos
/// pos- la pos. donde aparece
/// nn- el nod en el cual aparece
/// pd- el link con el padre full superior
/// G-Dp
/// L-lazv
int cant[MN], pos[MN], nn[MN], pd[MN];
void Dfs(int nod, int pad){
 int t = V[nod].size(), newn;
 if (t == 1 && nod != 1){
    pos[nod] = 0;
    nn[nod] = nod;
    cant[nod] = 1;
    pd[nod] = pad;
    return;
  }
  int mei = nod:
 for (int i = 0; i < t; i ++){
    newn = V[nod][i];
    if (newn == pad) continue;
    Dfs (newn, nod);
    if (cant[mej] < cant[nn[newn]])</pre>
      mei = nn[newn];
  pos[nod] = cant[mei];
  cant[mei] ++;
 nn[nod] = mej;
  pd[mei] = pad;
typedef pair<int, int> par;
typedef pair<int, par> tri;
typedef vector<tri> vt;
typedef vector<par> vp;
/// me da el recorrido desde a hasta b en vector<tri>
/// f posicion s.f in, s.f fin
```

```
vt rec(int a, int b) {
  vp A1, B1;
  A1.clear(), B1.clear();
  for (int i = a; i != -1; i = pd[nn[i]])
    A1.push back(par(nn[i], pos[i]));
  for (int i = b; i != -1; i = pd[nn[i]])
    B1.push back(par(nn[i], pos[i]));
  vt C1;
  C1.clear();
  reverse(A1.begin(), A1.end());
  reverse(B1.begin(), B1.end());
  int t = 0:
  while (t < A1.size() && t < B1.size() && A1[t] == B1[t])</pre>
  if (t >= A1.size() || t >= B1.size() || (t < B1.size() &&</pre>
  t < A1.size() && A1[t].first != B1[t].first))
    t --;
  if ((t < A1.size() && t < B1.size()) && A1[t].first ==</pre>
    B1[t].first){
    C1.push back(tri(A1[t].first, par(min(A1[t].second,
    B1[t].second), max(A1[t].second, B1[t].second))));
    t ++;
  for (int i = t; i < A1.size(); i ++)</pre>
    C1.push back (tri(A1[i].first, par(A1[i].second,
    cant[A1[i].first] - 1)));
  for( int i = t; i < B1.size(); i ++)</pre>
    C1.push back (tri(B1[i].first, par(B1[i].second,
    cant[B1[i].first] - 1)));
  return C1;
}
void havy light() {
  Dfs (1, -1); // root
  for (int i = 1; i \le N; i ++) /// rellenar con 4*cant
    if(cant[i]){
      G[i] = vector<int> (cant[i]*4, 0);
      L[i] = vector<bool> (cant[i]*4, false);
      G[i][1] = cant[i], L[i][1] = true;
}
```

```
// Segment Tree Persistente
const int N = 100000 + 100, LOGN = 20;
const int TOT = 4*N + N*LOGN;
int sum[TOT], L[TOT], R[TOT];
int sz = 1;
int newNode(int s = 0){
  sum[sz] = s;
 return sz++;
int build(int b, int e){
 if (b==e) return newNode();
 int mid = (b + e) \gg 1;
 int cur = newNode();
 L[cur] = build(b, mid);
 R[cur] = build(mid+1 , e);
 return cur;
}
int update(int node, int b, int e, int p){
 if(b == e) return newNode(sum[node] + 1);
 int mid = (b + e) \gg 1;
 int cur = newNode();
 if(p <= mid) {
   L[cur] = update(L[node], b, mid, p);
    R[cur] = R[node];
 }
  else {
    R[cur] = update(R[node], mid+1 , e, p);
    L[cur] = L[node];
  sum[cur] = sum[L[cur]] + sum[R[cur]];
  return cur:
int query(int node1, int node2, int b, int e, int k){
 if(b == e) return b;
 int s = sum[L[node2]] - sum[L[node1]];
 int mid = (b + e) \gg 1;
 if(s >= k) return query(L[node1], L[node2], b, mid, k);
  else return query(R[node1], R[node2], mid+1 , e, k-s);
}
```

```
int root[N];
int main()
  int n, m;
  cin >> n >> m;
  root[0] = build(1, n);
  vector<int> v(n), tmp(n);
  for(int i = 0; i < n; ++i){
    cin >> v[i]; tmp[i] = v[i];
  sort(tmp.begin(), tmp.end());
  tmp.resize(unique(tmp.begin(), tmp.end()) - tmp.begin());
  for(int i = 0; i < n; ++i)
    root[i+1] = update(root[i], 1 , n, lower bound (
    tmp.begin(), tmp.end(), v[i]) - tmp.begin() + 1);
  while(m--){
    int i, j, k;
    cin >> i >> j >> k;
    cout << tmp[query(root[i-1], root[j], 1 , n, k)-1];</pre>
}
```

```
// Pollard's Rho Integer Factoring Algorithm
typedef long long 11;
11 mulmod(ll a, ll b, ll c) { // (a*b)%c, minimizing overflow
 11 x = 0, y = a % c;
 while (b > 0) {
   if (b % 2 == 1)
     x = (x + y) \% c;
   y = (y * 2) % c;
    b /= 2;
 }
  return x % c;
}
11 pollard rho(ll n) {
 int i = 0, k = 2;
 11 \times = 3, v = 3; // random seed = 3, other values possible
 while (1) {
   i++;
   x = (mulmod(x, x, n) + n - 1) \% n;
    11 d = gcd(abs(y - x), n);
    if (d != 1 && d != n) return d;
    if (i == k) y = x, k *= 2;
  }
}
int main() {
 11 n = 2063512844981574047LL; // n is not a large prime
 11 ans = pollard_rho(n);
 if (ans > n / ans) ans = n / ans;
  cout << ans << ' ' << n / ans;</pre>
 // should be: 1112041493 1855607779
 return 0;
}
```

```
// Floyd's Cycle-Finding Algorithm O(\mu + \lambda) \mu->mu, \lambda->lambda
typedef pair<int,int> ii;
int f(int x);
ii floydCycleFinding(int x0) {
  // finding k*mu, hare's speed is 2x tortoise's
  int tortoise = f(x0);
  int hare = f(f(x0)); // f(x0) is the node next to x0
  while (tortoise != hare) {
    tortoise = f(tortoise);
    hare = f(f(hare));
  }
  // finding mu, hare and tortoise move at the same speed
  int mu = 0;
  hare = x0;
  while (tortoise != hare) {
    tortoise = f(tortoise);
    hare = f(hare);
    mu++;
  // finding lambda, hare moves, tortoise stays
  int lambda = 1;
  hare = f(tortoise);
  while (tortoise != hare) {
    hare = f(hare);
    lambda++;
  return ii(mu, lambda);
}
```

```
return factors; } // if N does not fit in 32-bit integer and is a prime number
// then 'factors' will have to be changed to vector<11>
// inside int main(), assuming sieve(1000000) has been called before
vi res = primeFactors(2147483647);
                                                    // slowest, 2147483647 is a prime
res = primeFactors(136117223861LL); // slow, 2 large pfactors 104729*1299709
res = primeFactors(142391208960LL);
                                                      // faster, 2^10*3^4*5*7^4*11*13
// functions involving prime factors
// numPF(N):Count the number of prime factors of N
11 numPF(11 N) {
  ll PF idx = 0, PF = primes[PF idx], ans = 0;
  while (N != 1 && (PF * PF <= N)) {
   while (N % PF == 0) { N /= PF; ans++; }
   PF = primes[++PF idx];
  if (N != 1) ans++;
  return ans;
// numDiffPF(N): count the number of different prime factors of N
ll numDiffPF(ll N) {
  ll PF idx = 0, PF = primes[PF idx], ans = 0;
  while (PF * PF <= N) {</pre>
    if (N % PF == 0) ans ++;
                                                            // count this pf only once
   while (N % PF == 0) N /= PF;
   PF = primes[++PF idx];
 if (N != 1) ans++;
  return ans;
// sumPF(N): sum the prime factors of N
11 sumPF(11 N) {
 ll PF idx = 0, PF = primes[PF idx], ans = 0;
  while (PF * PF <= N) {</pre>
   while (N % PF == 0) { N \neq PF; ans += PF; }
   PF = primes[++PF idx];
  if (N != 1) ans += N;
  return ans;
}
// numDiv(N): count the number of divisors of N
ll numDiv(ll N) {
  11 PF idx = 0, PF = primes[PF idx], ans = 1;
                                                               // start from ans = 1
  while (N != 1 && (PF * PF <= N)) {
   11 power = 0;
                                                                    // count the power
   while (N % PF == 0) { N \neq PF; power++; }
   ans *= (power + 1);
                                                           // according to the formula
   PF = primes[++PF_idx];
  }
  if (N != 1) ans *= 2;
                                        // (last factor has pow = 1, we add 1 to it)
  return ans;
}
```

```
// sumDiv(N): sum the divisors of N
ll sumDiv(ll N) {
  ll PF idx = 0, PF = primes[PF idx], ans = 1;
                                                                // start from ans = 1
  while (N != 1 && (PF * PF <= N)) {
    11 power = 0;
    while (N % PF == 0) { N /= PF; power++; }
    ans *= ((11) pow((double) PF, power + 1.0) - 1)/(PF - 1);
                                                                          // formula
   PF = primes[++PF idx];
  if (N != 1) ans *= ((11)pow((double)N, 2.0) - 1) / (N - 1);
                                                                         // last one
  return ans;
}
// EulerPhi(N): count the number of positive integers < N that are
// relatively prime to N.
ll EulerPhi(ll N) {
  11 PF idx = 0, PF = primes[PF idx], ans = N;
                                                                 // start from ans = N
  while (N != 1 \&\& (PF * PF <= N)) {
    if (N % PF == 0) ans -= ans / PF;
                                                         // only count unique factor
    while (N % PF == 0) N /= PF;
   PF = primes[++PF idx];
  if (N != 1) ans -= ans / N;
                                                                        // last factor
  return ans;
}
// square matrix exponentiation
#define MAX N 105
                                           // increase/decrease this value as needed
struct Matrix {
  int mat[MAX N][MAX N];
                                                        // we will return a 2D array
Matrix matMul(Matrix a, Matrix b) {
                                                                            // O(n^3)
  Matrix ans; int i, j, k;
  for (i = 0; i < MAX N; i++)</pre>
    for (j = 0; j < MAX N; j++)
      for (ans.mat[i][j] = k = 0; k < MAX N; k++)
                                                                // if necessary, use
        ans.mat[i][j] += a.mat[i][k] * b.mat[k][j];
                                                                // modulo arithmetic
  return ans;
Matrix matPow(Matrix base, int p) {
                                                                       // O(n^3 log p)
  Matrix ans; int i, j;
  for (i = 0; i < MAX N; i++)</pre>
    for (j = 0; j < MAX N; j++)
      ans.mat[i][j] = (i == j);
                                                           // prepare identity matrix
  while (p) {
                             // iterative version of Divide & Conquer exponentiation
                                                     // if p is odd (last bit is on)
    if (p & 1) ans = matMul(ans, base);
    base = matMul(base, base);
                                                                    // square the base
    p >>= 1;
                                                                      // divide p by 2
  }
  return ans;
```