## Hidden Markov Model

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#### Viterbi: Problem Definition and Transformation

- $P(S_{1:k} \mid O_{1:k}) \stackrel{\text{def}}{=} P(S_1, \ldots, S_k \mid O_1, \ldots, O_k)$
- What do we need?
  - $arg \max P(S_{1:k} | O_{1:k})$
  - Can we derive even  $P(S_1 \mid O_1)$ ?

k=1的时候

# s是state O是observation

### Viterbi: Problem Definition and Transformation

- $P(S_{1:k} \mid O_{1:k}) \stackrel{\text{def}}{=} P(S_1, \ldots, S_k \mid O_1, \ldots, O_k)$
- What do we need? 在给定观测样本序列前提下,得到最大概率的状态序列排序
  - $arg max P(S_{1:k} | O_{1:k})$
  - Can we derive even  $P(S_1 \mid O_1)$ ?
  - Need to use the Bayes rule:

$$P(S_{1:k} \mid O_{1:k}) = \frac{P(O_{1:k} \mid S_{1:k}) \cdot P(S_{1:k})}{P(O_{1:k})}$$

· Because of the "arg max", we have

$$\begin{aligned} \arg\max_{S_{1:k}} P(S_{1:k} \mid O_{1:k}) &= \arg\max_{S_{1:k}} P(O_{1:k} \mid S_{1:k}) \cdot P(S_{1:k}) \\ &= \arg\max_{S_{1:k}} P(O_{1:k}, S_{1:k}) \end{aligned}$$

• Since we can compute the probabilities, the naive method is to consider all  $O(N^k)$  state sequences, where N is the number of states.

## Viterbi: Recursive Computation

 A common trick (for Markov Chains) is to check if we can use recursion (thanks for the abundant conditional independences given by the Markov assumption).

Note:

$$P(O_{1:k+1} \mid S_{1:k+1}) = P(O_{k+1} \mid O_{1:k}, S_{1:k+1}) \cdot P(O_{1:k} \mid S_{1:k+1})$$

$$= P(O_{k+1} \mid S_{k+1}) \cdot P(O_{1:k} \mid S_{1:k})$$

$$P(S_{1:k+1}) = P(S_{1:k}) \cdot P(S_{k+1} \mid S_k)$$

Then

$$P(O_{1:k+1}, S_{1:k+1})$$
= $P(O_{k+1} | S_{k+1}) \cdot P(O_{1:k} | S_{1:k}) \cdot P(S_{1:k}) \cdot P(S_{k+1} | S_k)$ 
= $P(O_{k+1} | S_{k+1}) \cdot P(S_{k+1} | S_k) \cdot P(O_{1:k}, S_{1:k})$ 

## Viterbi: Recursive Computation /2

$$P(O_{1:k+1}, S_{1:k+1}) = P(O_{1:k}, S_{1:k}) \cdot \left( P(O_{k+1} \mid S_{k+1}) \cdot P(S_{k+1} \mid S_k) \right)$$
  
=  $P(O_{1:k}, S_{1:k}) \cdot f(S_k, S_{k+1})$ 

- Note that the value of  $S_k$  is important when considering applying "max".
- Define  $\delta(k, v) \stackrel{\text{def}}{=} \max_{S_{1:k-1}} P(O_{1:k}, S_{1:k-1}, v)$ , we have the recurrence:

$$\delta(k+1,v) = \max_{u} (\delta(k,u) \cdot f(u,v))$$

with the boundary condition:

$$\delta(1,v)=P(O_1,v)$$

Then

$$\max_{S_{1:k+1}} P(O_{1:k+1}, S_{1:k+1}) = \max_{v} \max_{S_{1:k}} P(O_{1:k+1}, S_{1:k}, v) = \max_{v} \delta(k+1, v)$$

 The arg max solution can be obtained via backtracking or some additional data structure.

