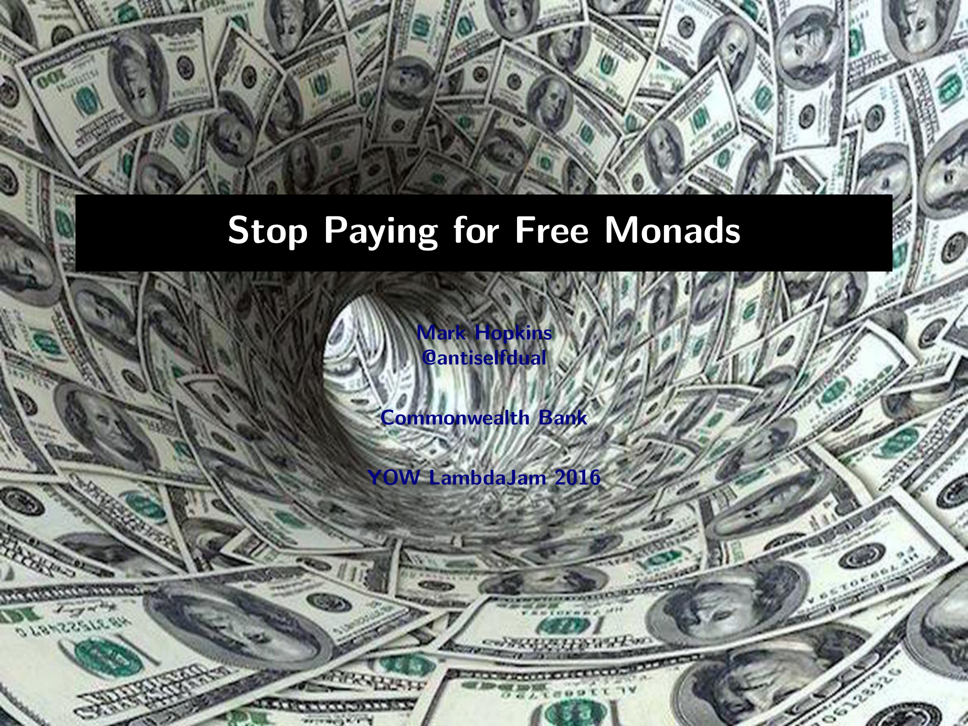


Stop Paying for Free Monads

Mark Hopkins
@antiselfdual

Commonwealth Bank

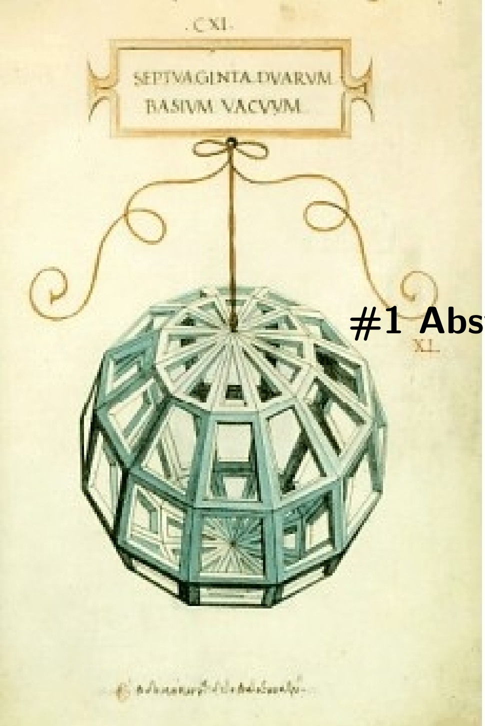
YOW LambdaJam 2016

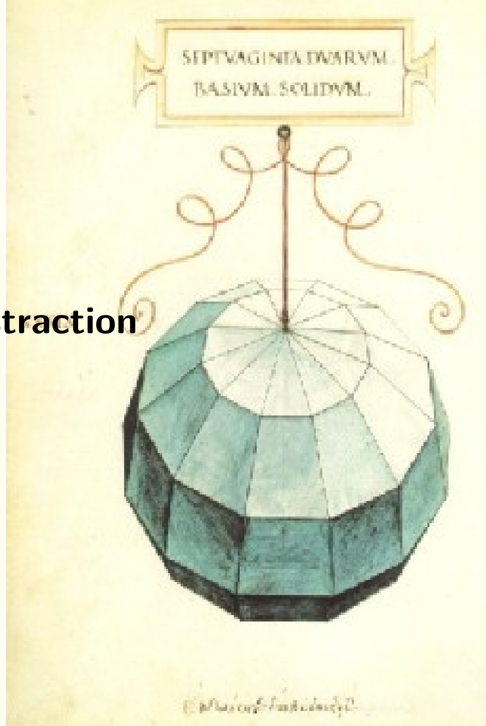
```haskell
prog tbl = do
  rows <- loadDb tbl
  dir <- mkTempDir
  report <- performAnalysis rows dir
  delete dir
  upload report s3Credentials $ "/reports/" ++ show tbl
  log $ "Wrote TPS report for " + show tbl
```

Abstraction and **modularity** for **effectful** programs.

Goals

#1 Abstraction





Interface/implementation separation

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- swap out implementations
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  - evaluation
  - pretty-printing
  - serialization

Some different ways of viewing this:

- interface vs implementation

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- interface vs implementation
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- interface vs implementation
- syntax vs semantics
- a DSL with multiple interpreters
- compiling a source language to a target language

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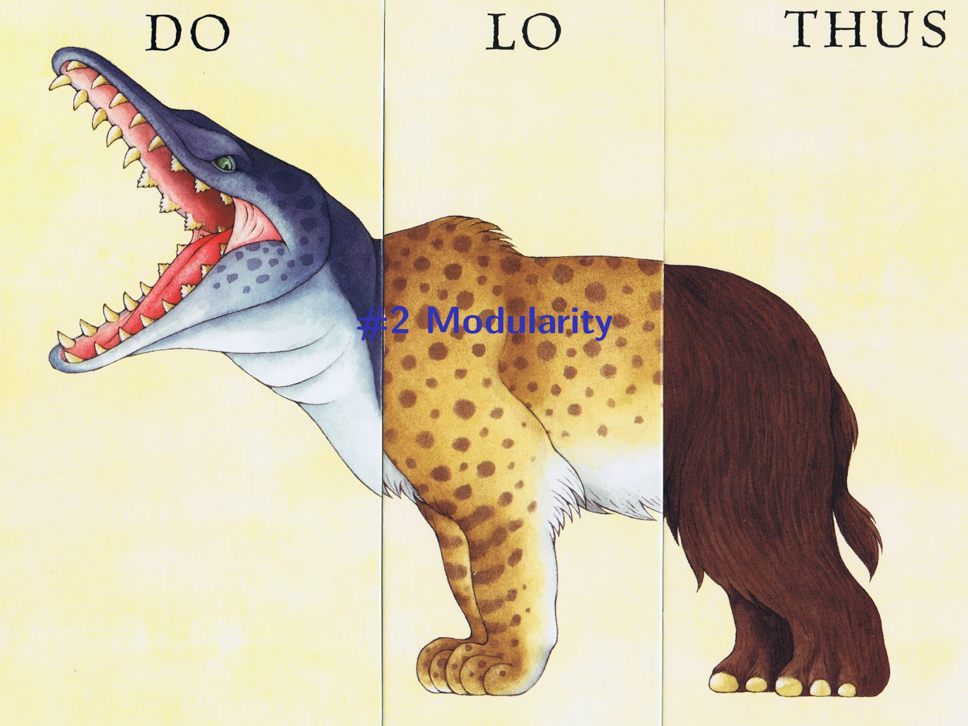
Additionally, we want **low-cost abstraction**.
We don't want to be penalised with

- poorly performing code
- code that's overly difficult to read or write.

#2 Modularity



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So we need some kind of “addition of DSLs”.

There should be a kind of **stability** too:
If I write a program, and then extend the language, my program should still be valid in the new language.

In addition, I also want **modular interpreters**.

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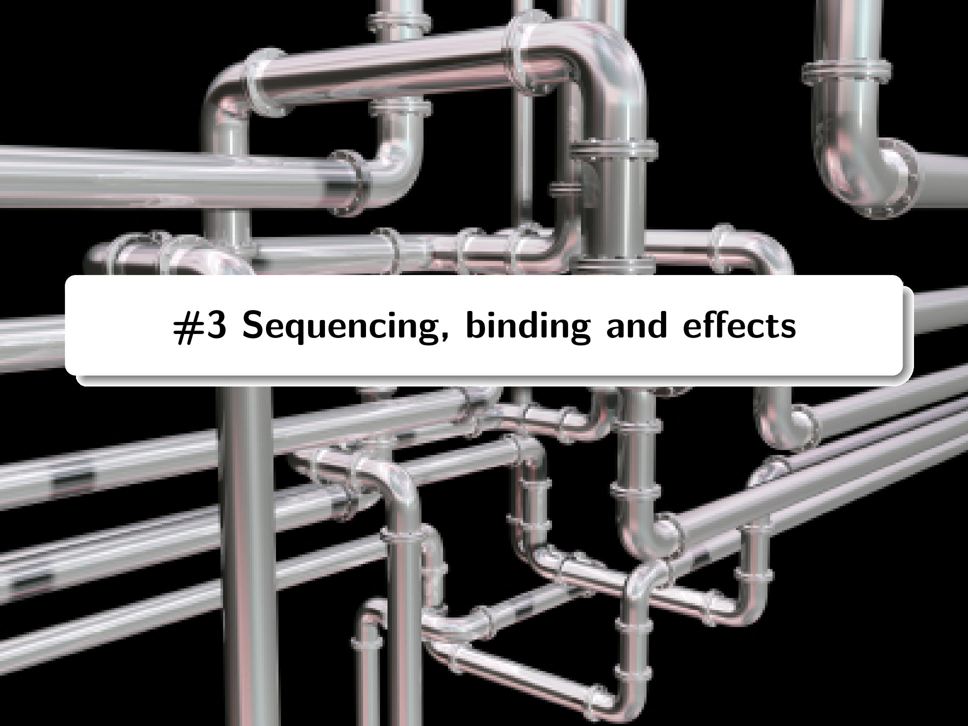
If I write interpreters for a number of different languages. . .

In addition, I also want **modular interpreters**.
If I write interpreters for a number of different languages...

...I want to reuse them to interpret a program written using them together.

In other words, we need a solution of the *expression problem*.

#3 Sequencing, binding and effects



The (pure) functional programming vision

Better software through programs we can reason about.

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Better software through programs we can reason about.

No matter how complex our programs their behaviour remains predictable.

The lambda calculus

Functions are values: their meaning does not depend on their context.

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Functions are values: their meaning does not depend on their context.
The laws that function composition satisfies

```
f . id = f = id . f

f . (g . h) = (f . g) . h
```

allow us to reason about our code and to safety refactor.

i.e. **equational reasoning**.

How can we add side effects but retain predictability?

Moggi's insight was to *use the type system*.[1]

[1] Eugenio Moggi, Notions of computation as monads, 1991

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Types will stand in for effects.

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This gives us three basic type classes for structuring computations

- comonads `w a ->   b`
- monads `  a -> m b`
- categories (and arrows) `  a >>>  b`

Equational reasoning

Reinstating the rules that we had for functions gives us the laws these have to obey:

```
f =>= extract = f = extract =>= f
f >=> return  = f =  return >=> f
f >>> id      = f =      id >>> f

(f =>= g) =>= h = f =>= (g =>= h)
(f >=> g) >=> h = f >=> (g >=> h)
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```

We can use these to refactor safely, i.e. we have equational reasoning in the presence of effects.

This generalises beyond side effects.
We use **effect** (or **context**) as a way of talking about the enhancement that `m a` brings over the raw type `a`.

We want to reuse these patterns that FP has evolved to deal with effects, sequencing and binding.

#4 Effects in the specification

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And this needs to obey appropriate laws.


Solution 1: Reify

Datatypes à la carte, Wouter Swierstra, 2008

Motto:

> *Turn operations into constructors*

Let's start with a simple typed language of console interactions.

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```
data Console a where
  Ask  :: String -> Console String
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```haskell
data Console a where
  Ask  :: String -> Console String
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```

Each operation becomes a constructor.

(co-)Yoneda embedding

CPS transform to get a functor:

```
data Console a =
    Ask String (String -> a)
  | Tell String a
    deriving Functor
```

Adding functors

So far, so good. What about modularity?

Idea: combining languages can literally be a **sum** of functors.

Adding functors

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Idea: combining languages can literally be a **sum** of functors.

```
data (f :+: g) a = Inl f a | Inr g a

instance (Functor f, Functor g) => Functor (f :+: g) where
  fmap k (Inl fa) = fmap k fa
  fmap k (Inr ga) = fmap k ga
```

For instance we can break `Console` up

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data Ask a = Ask String (a -> String) deriving Functor
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type Console = Ask :+: Tell
```

Free monads

We can make a monad out of any functor, using the “free monad” construction.[2]

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Free monads

We can make a monad out of any functor, using the “free monad” construction.[2]
What is it, and where does it come from?

It turns out that being “free” is closely connected to our idea of operations as data types.

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```
class Monad m where
  return :: a -> m a
  join :: m (m a) -> m a
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```
class Monad m where
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data Free f a where
  Return :: a -> Free f a
  Join :: f (Free f a) -> Free f a
```

```haskell
data Free f a = Return a
              | Join (f (Free f a))
```

`Free f a` is a monad whenever `f` is a functor.

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What exactly are we claiming when we say it's the "free monad on `f`"?

A small digression into abstract nonsense

Categories have **objects**, and **morphisms** between objects.

Haskell types

Haskell types form a category: a morphism between two types is just a function.

```
f :: a -> b
```

Functors

Functors between Haskell types forms a category, where a morphisms between two functors is a **natural transformation**. A natural transformation is just a polymorphic function

```
n :: f a -> g a
```

Monads

Monads on Haskell types also form a category.

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A monad is a functor with some additional structure (>>= and `return`).
So a morphism between monads will be a natural transformation preserving that structure:

```
n . (f >=> g) = (n . f) >=> (n . g)
n . return    = return
```

where

```
(f >=> g) a = f a >>= g
```

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In fact this is important for what will follow.

It means if we write a program in the free monad, then translate, we can be sure we continue to obey the monad laws.

Otherwise, we'll lose predictability i.e. the ability to reason about our code.

Adjunctions

Usually in Haskell, functors go from `Hask` to `Hask`.
But in general, functors go from one category $\mathcal{C}$ to another $\mathcal{D}$.

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But in general, functors go from one category $\mathcal{C}$ to another $\mathcal{D}$.

An **adjunction** is the name for the situation where we have functors $L :: \mathcal{C} \to \mathcal{D}$, $R :: \mathcal{D} \to \mathcal{C}$. satisfying

$$\mathcal{D}(Lc, d) \cong \mathcal{C}(c, Rd)$$

for all objects c in $\mathcal{C}$ and d in $\mathcal{D}$.

$\mathcal{C}(c_1, c_2) =$ the collection of $\mathcal{C}$ morphisms from c_1 to c_2.
$\mathcal{D}(d_1, d_2) =$ the collection of $\mathcal{D}$ morphisms from d_1 to d_2.

When R is an inclusion of $\mathcal{C}$ into $\mathcal{D}$ we write $F = L$ and say

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they just correspond to the $\mathcal{C}$ morphisms out of c.

In our case, our adjunction looks like

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$$\mathsf{Mnd}(\mathsf{Free}\, F, M) \cong \mathsf{Func}(F, M)$$

If F breaks up into a sum of functors ...

$$F = F_1 + \cdots + F_k,$$

then

$$\begin{aligned} & \mathsf{Mnd}(\mathsf{Free}\,(F_1 + \cdots + F_k), M) \\ \cong \;& \mathsf{Func}(F_1 + \cdots + F_k, M) \\ \cong \;& \mathsf{Func}(F_1, M) \times \cdots \times \mathsf{Func}(F_k, M) \end{aligned}$$

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So this is the sense in which we have modular interpreters.

Example

```haskell
data Ask a        = Ask String (String -> a)
data Tell a       = Tell String a
data Lookup k v a = Lookup k ((Maybe v) -> a)
```

```haskell
checkQuota :: Free (Ask :+: (Lookup String Int) :+: Tell) ()
checkQuota = do
  name <- Join . Inl $ Ask "What's your name" Return
  quota <- Join . Inr . Inl $ Lookup name Return
  Join . Inr . Inr $ Tell
    (   "Hi "
     ++ name ++ ", your quota is "
     ++ show (fromMaybe 0 quota)
    )
    (Return ())
```

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The injections Inl and Inr break our stability goal!

We'll have to modify the embeddings *everywhere* if we want to add another language.

Solution: polymorphism!

Introduce a type class to manage the injections for us

```
class f :<: g where
  inj :: f a -> g a
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Add instances to capture composition of injections.

```
instance f :<: f where ...
instance f :<: (f :+: g) where ...
instance (f :<: h) => f :<: (g :+: h) where ...
```

A polymorphic injection function

```
inject :: (f :<: g) => f (Free g a) -> Free g a
inject = Join . inj
```

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```
inject :: (f :<: g) => f (Free g a) -> Free g a
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```

Rewrite our program in terms of smart constructors

```
ask :: Ask :<: f => String -> Free f String
ask s = inject $ Ask s Return
```

Invisible abstraction

```haskell
checkQuota :: Free (Ask :+: Lookup String Int :+: Tell) ()
checkQuota = do
  name <- ask "What is your name?"
  quota <- lookup name
  tell $ "Hi " ++ name ++ ", your quota is " ++
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Invisible abstraction

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checkQuota ::
(Ask :<: f, Lookup String Int :<: f, Tell :<: f) =>
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    (show (fromMaybe 0 quota))
```

Interpretation

```haskell
class (Functor f, Monad m) => Interpret f where
  intp :: f a -> m a

instance (Interpret f m, Interpret g m) =>
  Interpret (f :+: g) m where
    intp (Inl fa) = intp fa
    intp (Inr ga) = intp ga

interpret :: Interpret f m => Free f a => m a
interpret (Return a) = return a
interpret (Join fa) = join $ intp (interpret <$> fa)
```

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interpret is a monad morphism.

```haskell
type MyMonad = ReaderT (Map String Int) IO

instance Interpret Ask MyMonad where ...
instance Interpret Tell MyMonad where ...
instance Interpret (Lookup String Int) MyMonad where ...
```

Let's run our program in `ReaderT (Map String Int) IO`.

```
> runReaderT (interpret checkQuota) $
 [ ("Stu", 33), ("Ann", 55) ]
What's your name?
Ann
Hi Ann, your quota is 55
```

Specification-side effects

It's not clear how we could easily add effects like early termination to our program: we'd need to create a `FreeError f`.

i.e. an analogy of Free f but for `MonadFree`, obeying the appropriate laws and satisfying a universal property.

Other binding constructs

- If we wanted to use comonads instead, we could use cofree comonads.[3]

  We'd have to invent "codata types à la carte".

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- If we wanted arrows, it's not clear what to do.

  Although free arrows ought to exist, we don't yet have a Haskell implementation.

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  {-# LANGUAGE StandaloneDeriving    #-}
  {-# LANGUAGE MultiParamTypeClasses #-}
  {-# LANGUAGE FlexibleContexts      #-}
  {-# LANGUAGE FlexibleInstances     #-}
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  {-# LANGUAGE OverlappingInstances  #-}
  ```
  ```

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  {-# LANGUAGE FlexibleInstances     #-}
  {-# LANGUAGE UndecidableInstances  #-}
  {-# LANGUAGE OverlappingInstances  #-}
  ```
- No backtracking in Haskell's type solver means for `:<:` we can only have nesting on one side of `:+:`

Solution 2: Operations remain functions

Oleg Kiselyov, Typed Tagless Final Interpreters, 2012

Motto

Let the language do the work

A language is a type class parametrized by a data constructor.

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i.e. the parameter has kind `* -> *`.

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```
class Console r where
  ask  :: String -> r String
  tell :: String -> r ()
```

A language is a type class parametrized by a data constructor. i.e. the parameter has kind `* -> *`.

```
class Console r where
  ask  :: String -> r String
  tell :: String -> r ()

class KeyVal k v r where
  lookup :: k -> r (Maybe v)
  set    :: k -> v -> r ()
```

An interpreter is a type with an instance.

```haskell
instance MonadIO m => Console m where
  ask s = putStrLn s >> getLine
  tell = putStrLn
```

An interpreter is a type with an instance.

```haskell
instance MonadIO m => Console m where
  ask s = putStrLn s >> getLine
  tell = putStrLn

instance (Ord k, MonadState (Map k v) m) =>
  KeyVal k v m where
    lookup = gets . DM.lookup
    set k v = modify $ DM.insert k v
```

Adding languages simply means adding constraints.

It's easy to split languages.

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```
class Lookup k v where
  lookup :: k -> r (Maybe v)

class Set k v r where
  set :: k -> v -> r ()
```

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```haskell
class Lookup k v where
  lookup :: k -> r (Maybe v)

class Set k v r where
  set :: k -> v -> r ()
```

... and to combine them.

```haskell
type KeyVal k v r = (Lookup k v r, Set k v r)
```

Requires `-XConstraintKinds`.

So our languages are modular.

We have modular interpreters just by adding new instances for our type.

For a monadic computation, just add a `Monad` constraint.

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```haskell
getQuota ::
( Console r
, Lookup String Int r
, Monad r
) => r ()
getQuota = do
  name <- ask "What's your name?"
  quota <- fromMaybe 0 <$> lookup name
  tell $ "Hi " + name + ", your quota is " ++ show quota
```

Adding effects to the specification is easy.
Just add a different constraint in place of `Monad`.

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Just add a different constraint in place of Monad.

```
changePwd ::
( Console r
, KeyVal String String r
, MonadThrow r
) => r ()
changePwd = do
  n <- ask "What's your name?"
  p <- ask "What's your password?"
  matches <- (== Just p) <$> lookup n
  unless matches $ throwM WrongPassword
  np <- ask "Enter new password"
  np2 <- ask "Re-enter new password"
  unless (np == np2) $ throwM PasswordsDidNotMatch
  put n np
  tell "Password successfully updated"
```

To have a comonadic or arrowized computation, just add the relevant constraint.

Interpretation is the identity function.

Interpretation is the identity function.
i.e. just select a type (that has an instance for all constraints).

Let's run our program in `StateT (Map String String) IO`.

```
> runStateT changePwd $
  fromList [
    ("anne", "pwd123")
  , ("mark", "p@ssw0rd")
  ]

What's your name?
mark
What's your password?
p@ssw0rd
Enter new password
1337
Re-enter new password
1337
((),fromList [("sue","pwd123"),("mark","1337")])
```

Comparison

À la carte

Free monads over sums of functors.

Interpretation is a monad morphism, assembled in a modular fashion from interpretations (natural transformations) for each component language.

Typed tagless

Languages are (higher-kinded) type classes.
We combine languages by adding constraints.
Interpretations are types with the necessary instances.

Compared to à la carte, the tagless approach

- has minimal runtime overhead

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On the other hand, some things are easier with the free monad approach

e.g.

- free monads allow stepping through instruction-by-instruction

Related work

- Stacked `FreeT`.

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  see the `compdata` package and "Typed tagless final interpreters".

References

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