

# 6.896: Probability and Computation

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lecture 2


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# Recall: the MCMC Paradigm

**Input:** a. very large, but finite, set  $\Omega$  ;  
b. a positive weight function  $w : \Omega \rightarrow \mathbb{R}^+$ .

**Goal:** Sample  $x \in \Omega$ , with probability  $\pi(x) \propto w(x)$ .

in other words:  $\pi(x) = \frac{w(x)}{Z}$   the “partition function”

$$Z = \sum_{x \in \Omega} w(x)$$

**MCMC approach:**

construct a Markov Chain (think sequence of r.v.'s)  $(X_t)_t$   
converging to  $\pi$ , i.e.

$$\Pr[X_t = y \mid X_0 = x] \rightarrow \pi(y) \text{ as } t \rightarrow +\infty \text{ (independent of } x)$$

# Markov Chains

**Def:** A *Markov Chain* on  $\Omega$  is a stochastic process  $(X_0, X_1, \dots, X_t, \dots)$  such that

a.  $X_t \in \Omega, \forall t$

b.  $\Pr[X_{t+1} = y \mid X_t = x, X_{t-1} = x_{t-1}, \dots, X_0 = x_0] \equiv \Pr[X_{t+1} = y \mid X_t = x]$

the *transition probability* from state  $x$  to state  $y$   $\stackrel{\text{def}}{=} P(x, y)$

Properties of the matrix  $P$ :

Non-negativity:  $\forall x, y \in \Omega, P(x, y) \geq 0$ ;

Stochasticity:  $\sum_{y \in \Omega} P(x, y) = 1, \forall x \in \Omega$ .

such a matrix is  
called  
*stochastic*

# Card Shuffling

Sample a random permutation of a deck of cards

$\Omega = \{\text{all possible permutations}\}$

$w(x) = 1$ , for all permutations  $x$

Markov Chain:



and repeat forever

$X_t$ : state of the deck after the  $t$ -th riffle;  $X_0$  is initial configuration of the deck;

$X_{t+1}$  is independent of  $X_{t-1}, \dots, X_0$  conditioning on  $X_t$ .

# Evolution of the Chain

$p_x^{(t)} \in \mathbb{R}_+^{1 \times |\Omega|}$  : distribution of  $X_t$  conditioning on  $X_0 = x$ .

then

$$p_x^{(t+1)} = p_x^{(t)} P$$

$$p_x^{(t)} = p_x^{(0)} P^t$$

# Graphical Representation

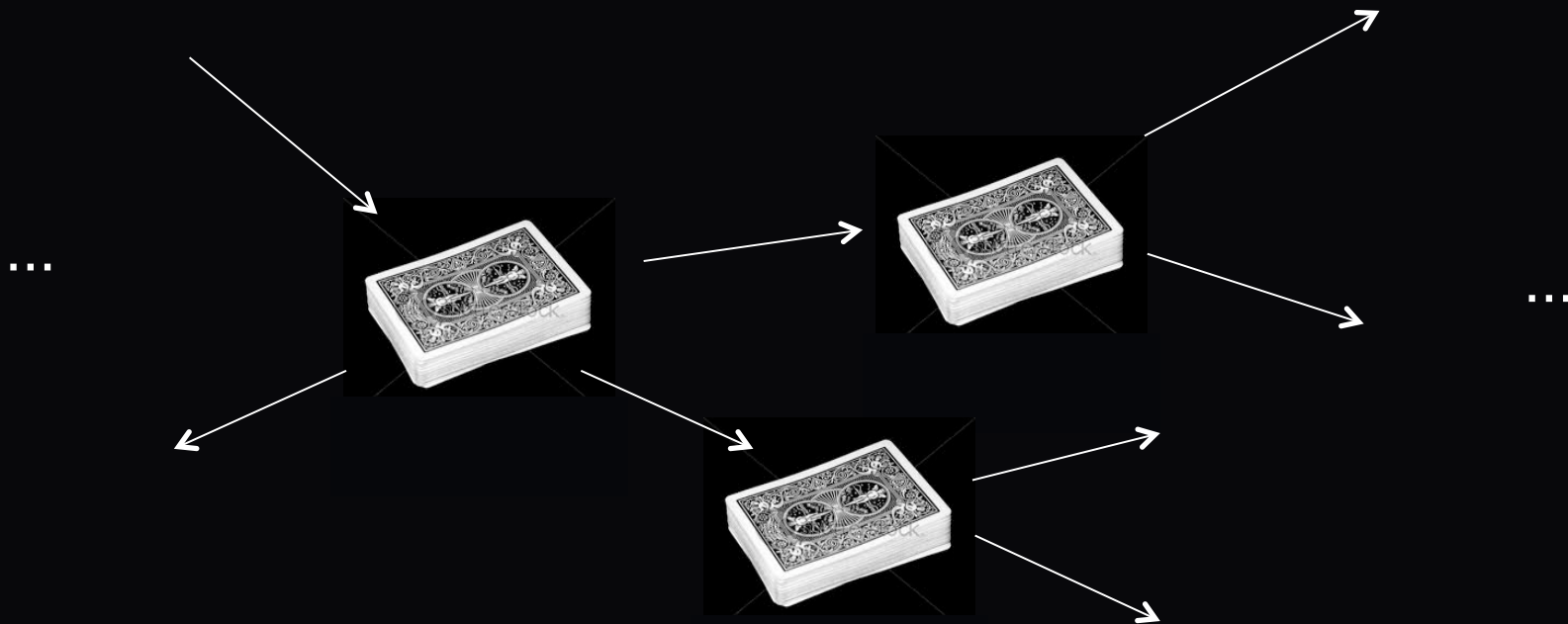
Represent Markov chain by a graph  $G(P)$ :

- nodes are identified with elements of the state-space  $\Omega$
- there is a directed edge between states  $x$  and  $y$  if  $P(x, y) > 0$ , with edge-weight  $P(x, y)$ ;
- no edge if  $P(x, y) = 0$ ;
- self loops are allowed (when  $P(x, x) > 0$ )

Much of the theory of Markov chains only depends on the topology of  $G(P)$ , rather than its edge-weights.

Many natural Markov Chains have the property that  $P(x, y) > 0$  iff  $P(y, x) > 0$ . In this case, we'll call  $G(P)$  *undirected* (ignoring the potential difference in the weights on an edge).

e.g. card Shuffling



“  $\rightarrow$  ” : reachable via a cut and riffle

e.g. of non-edge: no way to go from permutation 1234 to 4132

e.g. of directed edge: Can go from 123456 to 142536, but not vice versa

# Ir-reducibility and A-periodicity

**Def:** A Markov chain  $P$  is *irreducible* if for all  $x, y$ , there exists some  $t$  such that  $P^t(x, y) > 0$ .

[Equivalently,  $G(P)$  is strongly *connected*. In case the graphical representation is an undirected graph, then  $\Leftrightarrow G(P)$  being connected.]

**Def:** A Markov chain  $P$  is *aperiodic* if for all  $x, y$  we have

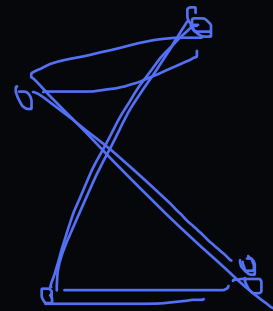
$$\gcd\{t : P^t(x, y) > 0\} = 1.$$



# True or False

For an irreducible Markov chain  $P$ , if  $G(P)$  is undirected then aperiodicity is equivalent to  $G(P)$  being **non-bipartite**.

A: true, look at lecture notes



## True or False (ii)

Define the period of  $x$  as  $\gcd\{t : P^t(x, x) > 0\}$ . For an irreducible Markov chain, the period of every  $x \in \Omega$  is the same.

A: true, 1 point exercise

[Hence, if  $G(P)$  is undirected, the period is either 1 or 2.]

## True or False (iii)

Suppose  $P$  is irreducible. Then  $P$  is aperiodic iff there exists  $t$  such that  $P_t(x,y) > 0$  for all  $x, y \in \Omega$ .

A: true, 1 point exercise to fill in the details of the sketch we discussed in class. For the forward direction, you may want to use the concept of the *Frobenius number* (aka the *Coin Problem*).

## True or False (iv)

Suppose  $P$  is irreducible and contains at least one self-loop (i.e.,  $P(x, x) > 0$  for some  $x$ ). Then  $P$  is aperiodic.

A: true, easy to see.

# Stationary Distribution

**Def:** A probability distribution  $\pi$  over  $\Omega$  is a *stationary distribution* for  $P$  if  $\pi = \pi P$ .

**Theorem (Fundamental Theorem of Markov Chains) :**

If a Markov chain  $P$  is *irreducible* and *aperiodic* then it has a unique stationary distribution  $\pi$ .

In particular,  $\pi$  is the unique (normalized such that the entries sum to 1) left eigenvector of  $P$  corresponding to eigenvalue 1.

Finally,  $P^t(x, y) \rightarrow \pi(y)$  as  $t \rightarrow \infty$  for all  $x, y \in \Omega$ .

In light of this theorem, we shall sometimes refer to an irreducible, aperiodic Markov chain as **ergodic**.

# Reversible Markov Chains

**Def:** Let  $\pi > 0$  be a probability distribution over  $\Omega$ . A Markov chain  $P$  is said to be *reversible wrt  $\pi$*  if

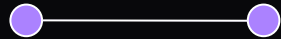
$$\forall x, y \in \Omega: \pi(x) P(x, y) = \pi(y) P(y, x).$$

Note that any symmetric matrix  $P$  is trivially reversible (w.r.t. the uniform distribution  $\pi$ ).

**Lemma:** If a Markov chain  $P$  is reversible w.r.t.  $\pi$ , then  $\pi$  is a stationary distribution for  $P$ .

# Reversible Markov Chains

Representation by *ergodic flows*:



*detailed balanced condition*

$$Q(x, y) := \boxed{\pi(x) \cdot P(x, y) \equiv \pi(y) P(y, x)}$$

the amount of probability mass flowing from  $x$  to  $y$  under  $\pi$

From flows to transition probabilities:

$$P(x, y) = \frac{Q(x, y)}{\sum_x Q(x, y)} \quad (\text{verify})$$

From flows to stationary distribution:

$$\frac{\pi(x)}{\pi(y)} = \frac{P(y, x)}{P(x, y)} \quad (\text{verify})$$

# Mixing of Reversible Markov Chains

**Theorem (Fundamental Theorem of Markov Chains) :**

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In particular,  $\pi$  is the unique (normalized such that the entries sum to 1) left eigenvector of  $P$  corresponding to eigenvalue 1.

Finally,  $P^t(x, y) \rightarrow \pi(y)$  as  $t \rightarrow \infty$  for all  $x, y \in \Omega$ .

Proof of FTMC: For reversible Markov Chains (today on the board-see lecture notes); full proof next time (probabilistic proof).



# Mixing in non-ergodic chains

When  $P$  is irreducible (but not necessarily aperiodic), then  $\pi$  still exists and is unique, but the Markov chain does not necessarily converge to  $\pi$  from every starting state.

For example, consider the two-state Markov chain with  $P = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$ .

This has the unique stationary distribution  $\pi = (1/2, 1/2)$ , but does not converge from either of the two initial states.

Notice that in this example  $\lambda_0 = 1$  and  $\lambda_1 = -1$ , so there is another eigenvalue of magnitude 1.

# Lazy Markov Chains

**Observation:** Let  $P$  be an irreducible (but not necessarily aperiodic) stochastic matrix. For any  $0 < \alpha < 1$ , the matrix  $P' = \alpha P + (1 - \alpha) I$  is stochastic, irreducible and aperiodic, and has the same stationary distribution as  $P$ .

This operation going from  $P$  to  $P'$  corresponds to introducing a self-loop at all vertices of  $G(P)$  with probability  $1 - \alpha$ .

Such a chain  $P'$  is usually called a *lazy version of  $P$* .

## e.g. Card Shuffling

Argue that the following shuffling methods converge to the uniform distribution:

- Random Transpositions

Pick two cards  $i$  and  $j$  uniformly at random with replacement, and switch cards  $i$  and  $j$ ; repeat.

- Top-in-at-Random:

Take the top card and insert it at one of the  $n$  positions in the deck chosen uniformly at random; repeat.

- Riffle Shuffle:

a. Split the deck into two parts according to the binomial distribution  $\text{Bin}(n, 1/2)$ .

b. Drop cards in sequence, where the next card comes from the left hand  $L$  (resp. right hand  $R$ ) with probability  $\frac{|L|}{|L|+|R|}$  (resp.  $\frac{|R|}{|L|+|R|}$ ).

c. Repeat.