Some Equivalence Laws of Propositional Logic

```
(P \wedge Q) \vee R \equiv (P \vee R) \wedge (Q \vee R)
                                                                distributivity law
                     \equiv P
         P \vee P
                                                                idempotency law for \vee
         P \lor Q \equiv Q \lor P
                                                                commutativity of \lor
 P \lor (Q \lor R) \equiv (P \lor Q) \lor R
                                                                associativity of \lor
      P \vee \text{true} \equiv \text{true}
                                                                true is right zero of \vee
      \mathsf{true} \vee P
                                                                 true is left zero of \vee
                          true
      P \vee \text{false} \equiv
                                                                 false is right one of \lor
      \text{false} \vee P
                                                                 false is left one of \vee
         P \wedge P \equiv
                                                                idempotency law for \wedge
         P \wedge Q \equiv
                           Q \wedge P
                                                                commutativity of \wedge
 P \wedge (Q \wedge R) \equiv (P \wedge Q) \wedge R
                                                                associativity of \wedge
       P \wedge \text{true} \equiv
                                                                 true is right one of \wedge
      \mathsf{true} \wedge P
                                                                 true is left one of \wedge
      P \wedge \text{false} \equiv \text{false}
                                                                 false is right zero of \land
      false \wedge P
                     \equiv
                           false
                                                                 false is left zero of \wedge
          \neg \neg P \equiv
                           P
                                                                double negation law
        P \Rightarrow Q \equiv
                           \neg P \lor Q
                                                                implication in terms of \vee
        P \Rightarrow Q \equiv
                           \neg Q \Rightarrow \neg P
                                                                contrapositive law
     true \Rightarrow P
                           P
                                                                 true absorbed in implication
                     \equiv
     false \Rightarrow P
                                                                 false implies anything
                     \equiv
                           true
     P \Rightarrow \text{true}
                     \equiv
                           true
                                                                 anything implies true
     P \Rightarrow \text{false} \equiv
                                                                implication and negation law
        P \Leftrightarrow Q \equiv (P \land Q) \lor (\neg P \land \neg Q)
                                                                bi-implication in terms of \vee and \wedge
        P \Leftrightarrow Q \equiv (P \Rightarrow Q) \land (Q \Rightarrow P)
                                                                bi-implication in terms of implication
         P \wedge Q \equiv \neg (\neg P \vee \neg Q)
                                                                De Morgan's law
         P \vee Q \equiv \neg (\neg P \wedge \neg Q)
                                                                De Morgan's law
 P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)
                                                                distributivity law
(P \lor Q) \land R \equiv (P \land R) \lor (Q \land R)
                                                                distributivity law
 P \lor (Q \land R) \equiv (P \lor Q) \land (P \lor R)
                                                                distributivity law
(P \wedge Q) \vee R \equiv (P \vee R) \wedge (Q \vee R)
                                                                distributivity law
(P \lor Q) \Rightarrow R \equiv (P \Rightarrow R) \land (Q \Rightarrow R)
                                                                distributivity law
P \Rightarrow (Q \land R) \equiv (P \Rightarrow Q) \land (P \Rightarrow R)
                                                                distributivity law
```

Some Equivalence Laws of Predicate Logic

Some Equivalence Laws of Set Operators

Some Equivalence Laws of Relation and Function Operators

```
(x,y) \in r^{-1} \quad \equiv \quad (y,x) \in r
                                                                                      from definition of relational inverse
              x \in \text{dom}(r) \equiv \exists y : T \cdot (x, y) \in r
                                                                                      from definition of domain
               x \in \operatorname{ran}(r) \equiv \exists y : T \cdot (y, x) \in r
                                                                                      from definition of range
            (x,z) \in r \stackrel{\circ}{\scriptscriptstyle 9} s \quad \equiv \quad \exists \ y: T \cdot (x,y) \in r \ \land \ (y,z) \in s \quad \text{from definition of relational composition}
                  (r^{-1})^{-1}
                                                                                      double inverse
                     x \in X \equiv (x, x) \in r
                                                                                      provided r is reflexive
(x,y) \in r \equiv (y,x) \in r

(x,y) \in r \land (y,z) \in r \Rightarrow (x,z) \in r
                                                                                      provided r is symmetric
                                                                                      provided r is transitive
               (x,y) \in r^*
                                 \equiv \exists n : \mathbb{N} \cdot (x, y) \in r^n
                                                                                      from definition of r^*
               (x,y) \in r^+ \equiv \exists n : \mathbb{N}^+ \cdot (x,y) \in r^n
                                                                                      from definition of r^+
                                                                                      composition of iterated relations
                     r^{st} \stackrel{	ext{o}}{	ext{9}} r^{st}
                                                                                      composition of transitive closure
                                                                                      composition of transitive closure
                                                                                      composition of transitive closure
                                                                                      composition of transitive closure
                                                                                      composition of transitive closure
```