

Assignment 1

Group#33

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Dynamic arrays

1. The big O complexity of reverse() is $O(\log N)$
2. The big O complexity of remove() is $O(N)$
The big O complexity of remove2() is $O(1)$
3. The big O complexity of find() is $O(N)$
4. The big O complexity of maxPalindrome() is $O(N^2)$ n^2

Question 5.

a. $\sum_{i=1}^n i^2 = 1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(n+1)(2n+1)}{6}$
 $\rightarrow 2n^3 + n^2 + 2n^2 + n$

*removing constants

most dominant term : $O(n^3)$

b. $\sum_{i=1}^{\log n} 2^i = 2^1 + 2^2 + 2^3 + \dots + 2^{\log n}$
 $\rightarrow (2^{(\log n + 1)}) - 2$
 $\rightarrow 2 * (2^{(\log n)}) - 2$

*removing constants

As $2^{(\log n)} = n$, the most dominant term : $O(n)$

c. $2x^2 + 16x^3$

*remove constants

most dominant term : $O(n^3)$

d. $(x - \log x^3)(x - 2\sqrt{x}) + 4x \log x^2$
 $\rightarrow x^2 - 2x\sqrt{x} * (-x \log x^3) * (-2\sqrt{x}) + 4x \log x^2$

x^2 is the term with the highest complexity degree,

Hence, the most dominant term : $O(n^2)$

Question 6 : $4^n \notin 2^{n+1}$

$$4^n = (2^2)^n = 2^{2n}$$

$$T(n) = C * f(n)$$

We want to show that $4^n \leq C * 2^n$ is not true.

$$2^{2n} \leq C * 2^n$$

$$\rightarrow \frac{2^{2n}}{2^n} \leq C$$

$$\rightarrow 2^{2n-n} \leq C$$

$$\rightarrow 2^n \leq C$$

We have proven that $C \geq 2^n$, we know that C has to be a constant and $4^n \notin 2^n$ is therefore not true.

Question 7 :

```
int calculatePolynomialSum ( int x, int[] a) {  
    int sum = a[a.length -1] * x;  
    for (int i = a.length - 2, i > 0, i-- ) {  
        sum= x*( sum + a[i] );  
    }  
    sum += a[0];  
    return sum;  
}
```