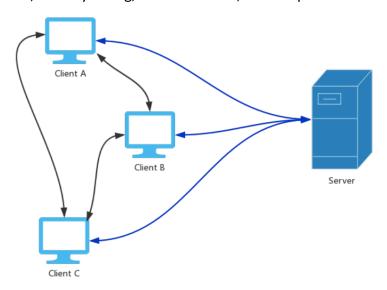
# Secure IM Design

## 1. Architecture:

The system is based on client-server architecture.

- Server listens to a port with a TCP server socket. It can manage online and offline users, and help them to authenticate each other.
- Client can build a TCP connection with the server. Based on the connection, the client can login, logout, get user list from the server and get authorization to talk to other clients.
- Client also listens to a port with a UDP server socket. Based on it, clients can exchange messages with each other.
- Services like PFS, identity hiding, DoS Protection, etc. are provided.



# 2. Assumptions:

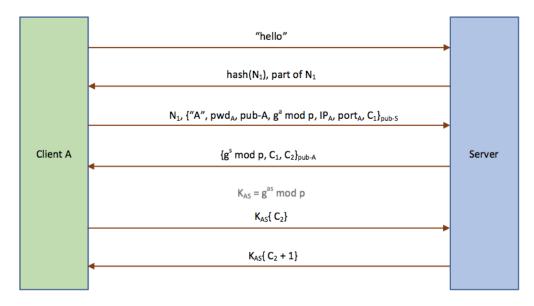
- Each user knows his/her password
- All clients know the server's public key
- Server knows all users' names, along with salted hashes of their passwords and salts
- Everyone knows g and p, which are used for Diffie-Hellman key exchange

# 3. Cryptographic Algorithms and Key Sizes

- Symmetric Encryption Algorithm: AES
- Symmetric Encryption Key Size: 256 bits
- Symmetric Encryption Block Cipher Mode: CTR
- Asymmetric Encryption Algorithm: RSA
- Asymmetric Encryption Key Size: 2048 bits
- Hashing Algorithm: SHA256
- Key Exchange Algorithm: Elliptic Curve Diffie-Hellman

#### 4. Protocols:

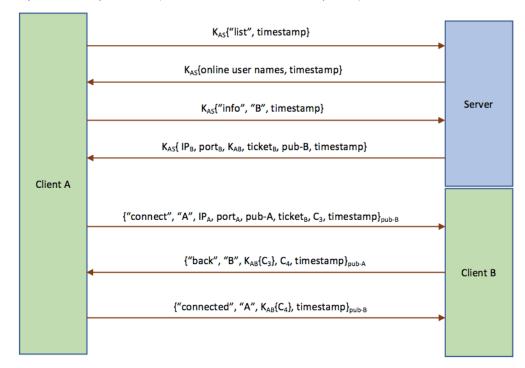
a. Authentication protocol (login)



When client A tries to login the server, it will start the following process:

- 1) Client A sends a "hello" message to the server.
- 2) Server generates a 128-bit nonce  $N_1$ , sends back part of it (the 21-128 bits) and its hash value.
- 3) Client A finds N<sub>1</sub> through computations and generates a DH component a. Then it generates a RSA key pair (pri-A and pub-A), sends N<sub>1</sub> and ("A", pwd<sub>A</sub>, pub-A, g<sup>a</sup> mod p, IP<sub>A</sub>, port<sub>A</sub>, C<sub>1</sub>) encrypted with the server's public key to the server.
- 4) The server first checks if the value of  $N_1$  is right. If it's right, it decrypts the data with its private key.
  - Then it uses the salt<sub>A</sub> and pwd<sub>A</sub> to compute the hash value, and compares with the value stored in the database to check if the pwd<sub>A</sub> is right.
  - If  $pwd_A$  is right, it generates a DH component s, and sends back ( $g^s \mod p$ ,  $C_1$ ,  $C_2$ ) encrypted with the public key of A.
- 5) Client A uses its private key to decrypt the message. Now client A and server shares the session key  $K_{AS} = g^{ab} \mod p$ . Then it sends  $C_2$  encrypted with  $K_{AS}$  to the server.
- 6) The server decrypts the message with  $K_{AS}$ , and confirms  $C_2$  is right. Then it sends  $C_2+1$  encrypted with  $K_{AS}$  back to the client.
- 7) Client A decrypts the message with  $K_{AS}$ , and confirms it's equal to  $C_2+1$ . By now, client A successfully logins to the server.

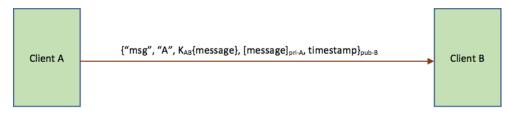
## b. Key establish protocol (and authentication with peers)



When client A wants to authenticate with client B, it will start the following process:

- 1) Client A sends a "list" message to the server, encrypted with KAS
- 2) Server responds with all online user names, encrypted with KAS
- 3) Client A sends a "info" message specifying "B" to the server, encrypted with KAS
- 4) Server generates ticket<sub>B</sub> for A, and send it back together with (IP<sub>B</sub>, port<sub>B</sub>,  $K_{AB}$ , pub-B), encrypt with  $K_{AS}$ 
  - ticketB = ["A", K<sub>AB</sub>, timestamp to expire]<sub>pri-S</sub>
- 5) Client A creates a "connect" message including ("A", IP<sub>A</sub>, port<sub>A</sub>, pub-A, ticket<sub>B</sub>, C<sub>3</sub>), encrypts it with pub-B, and sends it to B
- 6) Client B decrypts the message with its private key, generates  $C_4$  and encrypts with  $K_{AB}$ , and sends it back along with  $C_3$ .
- 7) Client A decrypts the message with its private key, confirms  $C_3$  is right, and adds B into the authenticated client list. Then it sends back a "connected" message including its information and  $C_4$  back to client B.
- 8) Client B decrypts the message with its private key, confirms C<sub>4</sub> is right. Till now, A and B have built an authenticated connection and can send messages to each other.

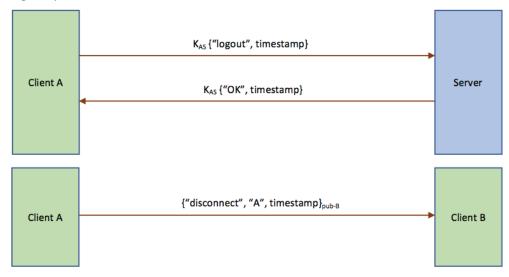
#### c. Messaging protocol



The messaging protocol is relatively easy:

Client A specifies the type as "msg", add "A" and the message encrypted with K<sub>AB</sub>, encrypts the whole message with pub-B, and sends it to client B.

#### d. Logout protocol



When client A tries to logout from the server, it will start the following process:

- 1) Client A sends a "logout" message to the server, encrypted with K<sub>AS</sub>
- 2) Then the server removes A from the online user list, deletes all A's information including K<sub>AS</sub>, pub-A, etc., and sends backs "OK"
- 3) Client A notifies all connected client (in this case only client B) a "disconnect" message and then exits. Client B receives this message and remove A from the authenticated client list.

# 5. Security Services:

a. Protection from DoS attacks:

Before a client authenticates to the server, it must solve a challenge from the server, which takes about 2 seconds and is helpful to prevent the DoS attacks.

PS: we don't provide DoS attack protection for the clients. But since single client is not as important as the server, we think this service can be omitted.

b. Protection from Replay attacks:

Since every message in this protocol contains either nonce or timestamp, which are used to check if the given message is legal. So, attackers cannot simply send any recorded message to perform replay attacks.

c. Identities Hiding:

Since all identities are encrypted and cannot by seen be attackers in clear, identities hiding service is provided.

d. Perfect Forward Secrecy:

We use Elliptic Curve Diffie-Hellman Key Exchange algorithm to generate the session keys between clients and server, so PFS is provided.

e. Protection from Password Guessing attacks:

In the server side, users' information is saved in a .csv file. For each password, it is hashed along with a random salt value, which is helpful to prevent the off-line password guessing attack.

In the client side, each user can only try 3 times for incorrect password input, which is helpful to prevent the on-line password guessing attack.