ACM Templates

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ACM Templates by XTS

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1 图论

```
dijstra
1.1
struct edge {
    int v, d;
};
struct node {
    int d, u;
    bool operator<(const node &x) const {</pre>
        return d > x.d;
    }
};
vector <edge> e[maxn];
void add(int u, int v, int d) {
    e[u].push_back(edge{v, d});
}
int dij(int s, int nd) {
    priority_queue <node> q;
    for (int i = 0; i <= N; i++)d[i] = inf;</pre>
    d[s] = 0;
    pre[s] = -1;
    q.push(node{0, s});
    while (!q.empty()) {
        node x = q.top();
        q.pop();
        int u = x.u;
        if (vis[u])continue;
        vis[u] = 1;
        for (auto E:e[u]) {
            if (d[E.v] > d[u] + E.d) {
                d[E.v] = d[u] + E.d;
                pre[E.v] = u;
```

```
q.push(node{d[E.v], E.v});
            }
        }
    }
}
1.2
    倍增 LCA
const int maxn = 5e5 + 7;
const int maxp = 21; //最大高度 log 1e5->18 1e4->15
int n, m, s;
int fa[maxn], dep[maxn], pa[maxn][maxp];
vector<int> G[maxn];
void dfs(int u, int fa)
{
    pa[u][0] = fa;
    dep[u] = dep[fa] + 1;
    for (int i = 1; i < maxp; i++)</pre>
        pa[u][i] = pa[pa[u][i - 1]][i - 1];
    for (int &v : G[u])
    {
        if (v == fa)
            continue;
        dfs(v, u);
    }
}
int lca(int u, int v)
{
    if (dep[u] < dep[v])
        swap(u, v);
    int t = dep[u] - dep[v];
    for (int i = 0; i < maxp; i++)</pre>
        if (t & (1 << i))
            u = pa[u][i];
```

```
for (int i = maxp - 1; i >= 0; i--)
    {
       int uu = pa[u][i], vv = pa[v][i];
       if (uu != vv)
       {
           u = uu;
           v = vv;
       }
    }
    return u == v ? u : pa[u][0];
}
    树上路径交
1.3
//x-y 和 xx-yy 的路径交即为 lca(x,xx), lca(x,yy), lca(y,xx), lca(y,yy) 中最深两点的路径
//注意判断是否有交
int intersection(int x, int y, int xx, int yy) {
    int t[4] = {lca(x, xx), lca(x, yy), lca(y, xx), lca(y, yy)};
    sort(t, t + 4);
    int r = lca(x, y), rr = lca(xx, yy);
    if (dep[t[0]] < min(dep[r], dep[rr]) || dep[t[2]] < max(dep[r], dep[rr]))</pre>
       return 0;
    int tt = lca(t[2], t[3]);
    int ret = 1 + dep[t[2]] + dep[t[3]] - dep[tt] * 2;
    return ret;
}
1.4
    树链剖分
//调用 predfs(root,1) dfs(root,root) 多组数据把 clk 赋 0
const int N = 1e5 + 7;
vector<int> G[N];
int n, m, s, mod;
int fa[N], dep[N], idx[N], out[N], ridx[N], sz[N], son[N], top[N], v[N], clk;
void predfs(int u, int d)
{
    dep[u] = d;
```

```
sz[u] = 1;
    int maxs = son[u] = -1;
    for (int &v : G[u])
        if (v == fa[u])
            continue;
        fa[v] = u;
        predfs(v, d + 1);
        sz[u] += sz[v];
        if (maxs == -1 \mid \mid sz[v] > sz[maxs])
            maxs = v;
    }
}
void dfs(int u, int tp)
{
    top[u] = tp;
    idx[u] = ++clk;
    ridx[clk] = u;
    if (son[u] != -1)
        dfs(son[u], tp);
    for (int &v : G[u])
        if (v != fa[u] \&\& v != son[u])
            dfs(v, v);
    out[u] = clk;
}
void go(int u, int v)
{
    int uu = top[u], vv = top[v];
    while (uu != vv)
        if (dep[uu] < dep[vv])</pre>
        {
            swap(uu, vv);
            swap(u, v);
        //do sth. in [idx[uu], idx[u]]
        u = fa[uu];
```

```
uu = top[u];
    }
    if (dep[u] < dep[v])</pre>
        swap(u, v);
    //do sth. in [idx[v], idx[u]]
}
int up(int u, int d)
{
    while (d)
    {
        if (dep[u] - dep[top[u]] < d)</pre>
        {
            d -= dep[u] - dep[top[u]];
            u = top[u];
        }
        else
            return ridx[idx[u] - d];
        u = fa[u];
        --d;
    }
    return u;
}
int finds(int u, int rt)
\{ //  找 u 在 rt 的哪个儿子的子树中
    while (top[u] != top[rt])
    {
        u = top[u];
        if (fa[u] == rt)
            return u;
        u = fa[u];
    }
    return ridx[idx[rt] + 1];
}
```

1.5 点分治

//求树上长度小于等于 k 的路径数

```
int n, m, k, cnt;
long long ans;
struct edge
{
    int to, d;
};
bool vis[maxn];
int h[maxn];
vector<edge> G[maxn];
int q[maxn], fa[maxn], sz[maxn], mx[maxn], dep[maxn];
void get_dep(int u, int fa, int d)
{
    dep[u] = d;
    h[++cnt] = d;
    for (int i = 0; i < G[u].size(); i++)</pre>
    {
        edge e = G[u][i];
        int v = e.to;
        if (vis[v] \mid \mid v == fa)
            continue;
        get_dep(v, u, d + e.d);
    }
}
int get_rt(int u)
{
    int p = 0, cur = -1;
    q[p++] = u;
    fa[u] = -1;
    while (++cur < p)</pre>
    {
```

```
u = q[cur];
        mx[u] = 0;
        sz[u] = 1;
        for (int i = 0; i < G[u].size(); i++)</pre>
        {
            edge e = G[u][i];
            if (!vis[e.to] && e.to != fa[u])
                fa[q[p++] = e.to] = u;
        }
    }
    for (int i = p - 1; i >= 0; i--)
    {
        u = q[i];
        mx[u] = max(mx[u], p - sz[u]);
        if (mx[u] * 2 <= p)
            return u;
        sz[fa[u]] += sz[u];
        mx[fa[u]] = max(mx[fa[u]], sz[u]);
    }
}
long long cal(int u, int x)
{
    long long res = 0;
    cnt = 0;
    get_dep(u, -1, 0);
    sort(h + 1, h + 1 + cnt);
    int r = cnt;
    for (int l = 1; l < r; l++)
    {
        while (h[1] + h[r] > x && r > 1)
            r--;
        res += r - 1;
    }
    return res;
}
void dfs(int u)
```

```
{
    u = get_rt(u);
    vis[u] = true;
    ans += cal(u, k);
    for (int i = 0; i < G[u].size(); i++)</pre>
    {
        edge e = G[u][i];
        int v = e.to;
        if (vis[v])
            continue;
        ans -= cal(v, k - e.d * 2);
        dfs(v);
    }
}
1.6
     强连通分量
const int maxn = 2e4 + 7;
int n, m, top;
vector<int> G[maxn];
int st[maxn], pre[maxn], low[maxn], sccno[maxn], clk, cnt;
void dfs(int u)
{
    pre[u] = low[u] = ++clk;
    st[++top] = u;
    for (int &v : G[u])
    {
        if (!pre[v])
            dfs(v);
            low[u] = min(low[u], low[v]);
        }
        else if (!sccno[v])
            low[u] = min(low[u], pre[v]);
    }
```

```
if (low[u] == pre[u])
        cnt++;
        for (;;)
        {
            int x = st[top--];
            sccno[x] = cnt;
            if (u == x)
                break;
        }
    }
}
void find_scc(int n)
{
    clk = cnt = top = 0;
    memset(sccno, 0, sizeof(sccno));
    memset(pre, 0, sizeof(pre));
    for (int i = 1; i <= n; i++)
    {
        if (!pre[i])
            dfs(i);
    }
}
1.7
     双连通分量
//
const int maxn = 1e3 + 7;
struct edge {
    int u, v;
};
int n, m, cnt;
stack <edge> st;
int pre[maxn], iscut[maxn], bccno[maxn], clk, bcc_cnt;
```

```
vector<int> G[maxn], bcc[maxn];
void init() {
    for (int i = 1; i <= n; i++)G[i].clear();</pre>
    n = 0;
}
int dfs(int u, int fa) {
    int lowu = pre[u] = ++clk;
    int child = 0;
    for (int i = 0; i < G[u].size(); i++) {</pre>
        int v = G[u][i];
        edge e = edge{u, v};
        if (!pre[v]) {
            st.push(e);
            child++;
            int lowv = dfs(v, u);
            lowu = min(lowu, lowv);
            if (lowv >= pre[u]) {
                 iscut[u] = 1;
                 bcc[++bcc_cnt].clear();
                 for (;;) {
                     edge x = st.top();
                     st.pop();
                     if (bccno[x.u] != bcc_cnt) {
                         bcc[bcc_cnt].push_back(x.u);
                         bccno[x.u] = bcc_cnt;
                     }
                     if (bccno[x.v] != bcc_cnt) {
                         bcc[bcc_cnt].push_back(x.v);
                         bccno[x.v] = bcc_cnt;
                     }
                     if (x.u == u \&\& x.v == v)break;
                 }
             }
        } else if (pre[v] < pre[u] && v != fa) {</pre>
            st.push(e);
```

```
lowu = min(lowu, pre[v]);
}

if (fa < 0 && child == 1)iscut[u] = 0;
return lowu;
}

void find_bcc(int n) {
   memset(pre, 0, sizeof(pre));
   memset(iscut, 0, sizeof(iscut));
   memset(bccno, 0, sizeof(bccno));
   clk = bcc_cnt = 0;
   for (int i = 1; i <= n; i++) {
        if (!pre[i])dfs(i, -1);
    }
}</pre>
```

2 字符串

2.1 回文树

```
//num[i] 表示以节点 i 表示的最长回文串的最右端点为回文串结尾的回文串个数
#include <bits/stdc++.h>
using namespace std;
const int maxa=2e5+20,cha=30,MAXN=1e6+20;//maxa 字符串最长长度, cha 字符集大小
typedef long long 11;
struct PalindromicTree
{
   int next[maxa][cha],fail[maxa],cnt[maxa],len[maxa],num[maxa];
   int tot,s[maxa],p,last;
   int newnode(int 1)
   {
       for(int i=0;i<cha;i++) next[p][i]=0;</pre>
       len[p]=1;
       cnt[p] = num[p] = 0;
       return p++;
   }
   void init()
   {
       tot=p=last=0;
       s[0]=-1,fail[0]=1;//0 为存储 偶数回文串 树根节点, 1 为存储 奇数回文串 树根节点
       newnode(0); /*p==0, 偶数回文串树根节点编号为 0, len 值为 0*/
       newnode(-1);/*p==1, 奇数回文串树根节点编号为 1,len 值为-1*/
   }
   int getfail(int x)
   {
       while(s[tot-len[x]-1]!=s[tot]) x=fail[x];
       return x;
   void insert(char c)
   {
       c-='a';
       s[++tot]=c;
       int cur=getfail(last);
       if(!next[cur][c])
       {
```

```
int now=newnode(len[cur]+2);
           fail[now] = next[getfail(fail[cur])][c];
           next[cur][c]=now;
           num[now] = num[fail[now]]+1;
       }
       last=next[cur][c];
       cnt[last]++;
   }
   void makecnt() //统计本质相同的回文串的出现次数
   {
       for(int i=p-1;i>=2;i--) cnt[fail[i]]+=cnt[i];//根节点 cnt 无意义, i>=2 即可
   }
}pt;
char str[maxa];
ll ans=0;
int main()
   scanf("%s",str);
   int len=strlen(str);
   pt.init();
   for(int i=0;i<len;i++) pt.insert(str[i]);</pre>
   pt.makecnt();
   cout << (pt.p-1)/2 << endl;
   return 0;
}
2.2 广义后缀自动机
//cnt[i] 表示节点 i 表示的本质不同的串的个数 (建树时求出的不是完全的, 最后 count() 函数
→ 跑一遍以后才是正确的)
#include <bits/stdc++.h>
using namespace std;
const int maxa=2e5+20, cha=30, MAXN=1e6+20; //maxa 字符串最长长度, cha 字符集大小
typedef long long 11;
struct SAM{
   int cnt,las,root;
   int ch[MAXN][30],fa[MAXN],len[MAXN];
   inline void init(){root = las = ++cnt;}
```

```
inline void insert(int x){
        int p = las;
        int np = las = ++cnt;
        len[np] = len[p] + 1;
        for(;p && !ch[p][x];p = fa[p]) ch[p][x] = np;
        if(!p) fa[np] = root;
        else{
            int q = ch[p][x];
            if(len[q] == len[p] + 1) fa[np] = q;
            else {
                int nq = ++cnt;
                fa[nq] = fa[q]; fa[q] = fa[np] = nq;
                memcpy(ch[nq],ch[q],sizeof(ch[nq]));
                len[nq] = len[p] + 1;
                for(;p \&\& ch[p][x] == q;p = fa[p]) ch[p][x] = nq;
            }
        }
    }
    inline long long sum(int x){
        return len[x] - len[fa[x]];
    }
}T;
char str[maxa];
ll ans=0;
int main()
{
    scanf("%s",str);
    int len=strlen(str);
    ll ans=0;
    T.init();
    for(int i = 0;i < len;i++) T.insert(str[i] - 'a');</pre>
    T.las = T.root;
    for(int i = len-1;i >=0;i--) T.insert(str[i] - 'a');
    for(int i = 0; i \le T.cnt; i++) ans = ans + T.sum(i);
    return 0;
```

}

3 数据结构

4 数学

4.1 带预处理 BGSG

```
#include <bits/stdc++.h>
#define LL long long
using namespace std;
unordered_map<LL, LL> mp;
int loop, up;
LL n, x0, a, b, p, v;
LL _pow(LL a, LL b, LL Mod)
    LL ret = 1;
    LL mul = a % Mod;
    while (b > 0)
    {
        if (b & 1)
            ret = ret * mul % Mod;
        mul = mul * mul % Mod;
        b = b >> 1;
    }
    return ret;
}
LL inv(LL x, LL P)
{
    return _pow(x, P - 2, P);
}
void pre_BSGS(int p, int a)
{ //预处理一部分 a ~y
    mp.clear();
    up = ceil(pow(p, 2.0 / 3));
    int t = 1;
    for (int i = 0; i <= up; i++)</pre>
    {
        if (i == up)
            loop = t;
        mp[t] = i;
        t = 1LL * t * a % p;
```

```
}
}
LL BSGS(LL B, LL N, LL P)
{ //B^ans=N(%p)
    int m = ceil(pow(p, 1.0 / 3));
    int obj = inv(N, P);
    for (int i = 1; i <= m; i++)
    {
        obj = 1LL * obj * loop % P;
        if (mp.count(obj))
        {
            return 1LL * i * up - mp[obj];
        }
    }
    return -1;
}
int main()
{
    int T;
    scanf("%d", &T);
    while (T--)
    {
        scanf("%lld%lld%lld%lld", &n, &x0, &a, &b, &p);
        int Q;
        scanf("%d", &Q);
        pre_BSGS(p, a);
        while (Q--)
        {
            scanf("%lld", &v);
            LL bi = b + ((a - 1) * v \% p);
            bi = bi % p;
            bi = bi * inv((b + x0 * (a - 1)) % p, p) % p;
            LL ans = BSGS(a, bi, p); //a ans=bi(%p)
            if (ans >= n)
                ans = -1;
            printf("%lld\n", ans);
        }
```

```
}
    return 0;
}
     十进制矩阵快速幂
4.2
#include <bits/stdc++.h>
using namespace std;
long long Mod;
struct Matrix
    long long a[2][2];
    Matrix operator*(const Matrix &b) const
    {
        Matrix tmp;
        for (int i = 0; i < 2; i++)
            for (int j = 0; j < 2; j++)
                tmp.a[i][j] = 0;
        for (int i = 0; i < 2; i++)
            for (int j = 0; j < 2; j++)
                for (int k = 0; k < 2; k++)
                {
                    tmp.a[i][j] += a[i][k] * b.a[k][j];
                    tmp.a[i][j] %= Mod;
                }
        return tmp;
    }
};
Matrix Matrixpow(Matrix x, int b)
{
    Matrix ans1;
    ans1.a[0][0] = 1;
    ans1.a[1][1] = 1;
    ans1.a[1][0] = 0;
    ans1.a[0][1] = 0;
    Matrix ret = x;
    while (b)
    {
        if (b & 1)
```

```
ans1 = ans1 * ret;
        ret = ret * ret;
        b = b / 2;
    }
    return ans1;
}
int main()
{
    ios::sync_with_stdio(false);
    long long x0, x1, a, b;
    cin >> x0 >> x1 >> a >> b;
    string s;
    cin >> s;
    cin >> Mod;
    int len = s.length();
    Matrix ans;
    ans.a[0][0] = 1;
    ans.a[1][1] = 1;
    ans.a[1][0] = 0;
    ans.a[0][1] = 0;
    Matrix mult;
    mult.a[0][0] = a;
    mult.a[0][1] = b;
    mult.a[1][0] = 1;
    mult.a[1][1] = 0;
    for (int i = len - 1; i >= 0; i--)
    {
        ans = ans * Matrixpow(mult, s[i] - '0');
        mult = Matrixpow(mult, 10);
    }
    long long tot = ans.a[1][1] * x0 % Mod + ans.a[1][0] * x1 % Mod;
    tot = tot % Mod;
    cout << tot << endl;</pre>
    return 0;
}
```

4.3 exlucas

```
#include <bits/stdc++.h>
#define rep(ii, a, b) for (int ii = a; ii \le b; ii++)
#define ll long long
using namespace std;
ll qpow(ll a, ll b, ll nmod)
{
    11 res = 1;
    while (b)
        if (b & 1)
            res = res * a % nmod;
        a = a * a \% nmod;
        b >>= 1;
    }
    return res;
}
void exgcd(ll a, ll b, ll &x, ll &y, ll &d)
{
    b ? (exgcd(b, a \% b, y, x, d), y -= x * (a / b)) : (x = 1, y = 0, d = a);
}
ll inv(ll a, ll m)
    11 x, y, d;
    exgcd(a, m, x, y, d);
    return x >= 0 ? x % m : x % m + m;
}
11 crt(ll a, ll m, ll m0)
{ //m0 | m; (m0, m/m0)=1
    return m / m0 * inv(m / m0, m0) % m * a % m;
}
ll fact(ll n, ll p, ll pk)
 \{ //(n!/p^x)\%(p^k) 
    if (n <= 1)
        return 1;
    ll ans = 1, tmp = n \% pk;
```

```
rep(i, 1, pk)
    {
        if (i % p)
            ans = ans * i \% pk;
    }
    ans = qpow(ans, n / pk, pk);
    rep(i, 1, tmp)
    {
        if (i % p)
            ans = ans * i \% pk;
    }
    return ans * fact(n / p, p, pk) % pk;
}
ll c(ll n, ll m, ll p, ll pk)
11 sum = 0;
    for (ll i = n; i; i \neq p)
        sum += i / p;
    for (ll i = m; i; i \neq p)
        sum -= i / p;
    for (ll i = n - m; i; i /= p)
        sum -= i / p;
    return qpow(p, sum, pk) * fact(n, p, pk) % pk * inv(fact(m, p, pk), pk) % pk *
    → inv(fact(n - m, p, pk), pk) % pk;
}
ll fac[40][2], pf; //0 p; 1 pk
void getfac(ll n)
{
    11 tmp = sqrt(n);
    rep(i, 2, tmp)
    {
        if (n % i == 0)
        {
            fac[++pf][0] = i, fac[pf][1] = 1;
            while (n \% i == 0)
                n /= i, fac[pf][1] *= i;
```

```
}
    }
    if (n > 1)
        fac[++pf][0] = n, fac[pf][1] = n;
}
ll exlucas(ll n, ll m, ll p)
{
    11 \text{ ans} = 0;
    getfac(p);
    rep(i, 1, pf)
        ans = (ans + crt(c(n, m, fac[i][0], fac[i][1]), p, fac[i][1])) % p;
    }
    return ans;
}
int main()
{
    ll n, m, p;
    cin >> n >> m >> p;
    cout << exlucas(n, m, p) << endl;</pre>
    return 0;
}
4.4 BM
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
using VI = vector<int64_t>;
class Linear_Seq
{
public:
    static const int N = 50020; ///多项式系数最大值
    int64_t res[N], c[N], md[N], COEF[N] /**COEF 是多项式系数 */, Mod;
    vector<int> Md;
    inline static int64_t gcdEx(int64_t a, int64_t b, int64_t &x, int64_t &y)
    {
        if (!b)
        {
```

```
x = 1;
        y = 0;
        return a;
    }
    int64_t d = gcdEx(b, a \% b, y, x);
    y = (a / b) * x;
    return d;
}
static int64_t Inv(int64_t a, int64_t Mod)
{
    int64_t x, y;
    return gcdEx(a, Mod, x, y) == 1 ? (x % Mod + Mod) % Mod : -1;
};
inline void mul(int64_t *a, int64_t *b, int k)
{ ///下边的线性齐次递推用的.
    fill(c, c + 2 * k, 0);
    for (int i(0); i < k; ++i)</pre>
        if (a[i])
            for (int j(0); j < k; ++j)
                c[i + j] = (c[i + j] + a[i] * b[j]) % Mod;
    for (int i(2 * k - 1); i >= k; --i)
        if (c[i])
            for (size_t j(0); j < Md.size(); ++j)</pre>
                c[i - k + Md[j]] = (c[i - k + Md[j]] - c[i] * md[Md[j]]) % Mod;
    copy(c, c + k, a);
}
VI BM(VI s)
{ ///BM 算法求模数是质数的递推式子的通项公式, 可以单独用
   VI C(1, 1), B(1, 1);
    int L(0), m(1), b(1);
    for (size_t n(0); n < s.size(); ++n)</pre>
    {
        int64_t d(0);
        for (int i(0); i <= L; ++i)
            d = (d + (int64_t)C[i] * s[n - i]) \% Mod;
        if (!d)
            ++m;
```

```
else
        {
            VI T(C);
            int64_t c(Mod - d * Inv(b, Mod) % Mod);
            while (C.size() < B.size() + m)</pre>
                C.push_back(0);
            for (size_t i(0); i < B.size(); ++i)</pre>
                C[i + m] = (C[i + m] + c * B[i]) \% Mod;
            if (2 * L <= (int)n)</pre>
            {
                L = n + 1 - L;
                B = T;
                b = d;
                m = 1;
            }
            else
                ++m;
        }
    }
    return C;
}
int solve(int64_t n, VI A, VI B)
{ //线性齐次递推:A 系数,B 初值 B[n]=A[0]*B[n-1]+...
    ///这里可以可以单独用,给出递推系数和前几项代替矩阵快速幂求递推式第 n 项.
    int64_t ans(0), cnt(0);
    int k(A.size());
    for (int i(0); i < k; ++i)</pre>
        md[k - i - 1] = -A[i];
    md[k] = 1;
    Md.clear();
    for (int i(0); i < k; ++i)</pre>
        res[i] = 0;
        if (md[i])
            Md.push_back(i);
    }
    res[0] = 1;
    while ((1LL << cnt) <= n)
```

```
++cnt;
    for (int p(cnt); ~p; --p)
    {
        mul(res, res, k);
        if ((n >> p) & 1)
        {
             copy(res, res + k, res + 1);
             res[0] = 0;
             for (size_t j(0); j < Md.size(); ++j)</pre>
                 \verb"res[Md[j]] = (\verb"res[Md[j]]] - \verb"res[k]" * \verb"md[Md[j]]") \% \ \verb"Mod";
        }
    }
    for (int i(0); i < k; ++i)</pre>
        ans = (ans + res[i] * B[i]) \% Mod;
    return ans + (ans < 0 ? Mod : 0);
}
inline static void extand(VI &a, size_t d, int64_t value = 0)
    if (d <= a.size())</pre>
        return;
    a.resize(d, value);
}
static int64_t CRT(const VI &c, const VI &m)
{ ///中国剩余定理合并
    int n(c.size());
    int64_t M(1), ans(0);
    for (int i = 0; i < n; ++i)
        M = m[i];
    for (int i = 0; i < n; ++i)
    {
        int64_t x, y, tM(M / m[i]);
        gcdEx(tM, m[i], x, y);
        ans = (ans + tM * x * c[i] % M) % M;
    }
    return (ans + M) % M;
}
static VI ReedsSloane(const VI &s, int64_t Mod)
```

{ ///求模数不是质数的递推式系数

```
auto L = [](const VI &a, const VI &b) {
    int da = (a.size() > 1 || (a.size() == 1 && a[0])) ? a.size() - 1 :

→ -1000;

    int db = (b.size() > 1 || (b.size() == 1 && b[0])) ? b.size() - 1 :
    \rightarrow -1000;
    return max(da, db + 1);
};
auto prime_power = [&](const VI &s, int64_t Mod, int64_t p, int64_t e) {
    vector<VI> a(e), b(e), an(e), bn(e), ao(e), bo(e);
    VI t(e), u(e), r(e), to(e, 1), uo(e), pw(e + 1);
    pw[0] = 1;
    for (int i(pw[0] = 1); i <= e; ++i)
        pw[i] = pw[i - 1] * p;
    for (int64_t i(0); i < e; ++i)
        a[i] = {pw[i]};
        an[i] = {pw[i]};
        b[i] = \{0\};
        bn[i] = {s[0] * pw[i] % Mod};
        t[i] = s[0] * pw[i] % Mod;
        if (!t[i])
        {
            t[i] = 1;
            u[i] = e;
        }
        else
            for (u[i] = 0; t[i] % p == 0; t[i] /= p, ++u[i])
                ;
    for (size_t k(1); k < s.size(); ++k)</pre>
    {
        for (int g(0); g < e; ++g)
        {
            if (L(an[g], bn[g]) > L(a[g], b[g]))
            {
                int id(e - 1 - u[g]);
```

```
ao[g] = a[id];
        bo[g] = b[id];
        to[g] = t[id];
        uo[g] = u[id];
        r[g] = k - 1;
    }
}
a = an;
b = bn;
for (int o(0); o < e; ++o)
{
    int64_t d(0);
    for (size_t i(0); i < a[o].size() && i <= k; ++i)
        d = (d + a[o][i] * s[k - i]) % Mod;
    if (d == 0)
    {
        t[o] = 1;
        u[o] = e;
    }
    else
    {
        for (u[o] = 0, t[o] = d; !(t[o] % p); t[o] /= p, ++u[o])
        int g(e - 1 - u[o]);
        if (!L(a[g], b[g]))
        {
            extand(bn[o], k + 1);
            bn[o][k] = (bn[o][k] + d) \% Mod;
        }
        else
        {
            int64_t coef = t[o] * Inv(to[g], Mod) % Mod * pw[u[o] -
             \rightarrow uo[g]] % Mod;
            int m(k - r[g]);
            extand(an[o], ao[g].size() + m);
            extand(bn[o], bo[g].size() + m);
            auto fun = [&](vector<VI> &vn, vector<VI> &vo, bool f) {
                 for (size_t i(0); i < vo[g].size(); ++i)</pre>
```

```
{
                             vn[o][i + m] -= coef * vo[g][i] % Mod;
                             if (vn[o][i + m] < 0)
                                 vn[o][i + m] += Mod * (f ? 1 : -1);
                         }
                         while (vn[o].size() && !vn[o].back())
                             vn[o].pop_back();
                    };
                    fun(an, ao, 1);
                    fun(bn, bo, -1);
                }
            }
        }
    }
    return make_pair(an[0], bn[0]);
};
vector<tuple<int64_t, int64_t, int>> fac;
for (int64_t i(2); i * i <= Mod; ++i)</pre>
    if (!(Mod % i))
    {
        int64_t cnt(0), pw(1);
        while (!(Mod % i))
        {
            Mod /= i;
            ++cnt;
            pw *= i;
        }
        fac.emplace_back(pw, i, cnt);
    }
if (Mod > 1)
    fac.emplace_back(Mod, Mod, 1);
vector<VI> as;
size_t n = 0;
for (auto &&x : fac)
{
    int64_t Mod, p, e;
    VI a, b;
    std::tie(Mod, p, e) = x;
```

```
auto ss = s;
            for (auto \&\&x : ss)
                x \%= Mod;
            std::tie(a, b) = prime_power(ss, Mod, p, e);
            as.emplace_back(a);
            n = max(n, a.size());
        }
        VI a(n), c(as.size()), m(as.size());
        for (size_t i(0); i < n; ++i)</pre>
        {
            for (size_t j(0); j < as.size(); ++j)</pre>
            {
                m[j] = std::get<0>(fac[j]);
                c[j] = i < as[j].size() ? as[j][i] : 0;
            }
            a[i] = CRT(c, m);
        }
        return a;
    }
    int64_t solve(VI a, int64_t n, int64_t Mod, bool prime = true)
    {
        VI c;
        this->Mod = Mod;
        if (prime)
            c = BM(a); ///如果已经知道系数了,直接输入到 c 就行了,不用调用 BM().
        else
            c = ReedsSloane(a, Mod);
        c.erase(c.begin());
        for (size_t i(0); i < c.size(); ++i)</pre>
            c[i] = (Mod - c[i]) \% Mod;
        return solve(n, c, VI(a.begin(), a.begin() + c.size()));
    }
};
11 _pow(ll a, ll b, ll Mod)
{
    ll ret = 1LL % Mod;
    11 mul = a % Mod;
    while (b)
```

```
{
        if (b & 1)
            ret = ret * mul % Mod;
        mul = mul * mul % Mod;
        b >>= 1;
    }
    return ret;
};
ll f[2020], sum[2020];
int main()
{
    VI vec;
    vec.clear();
    for (int i = 0; i <= 2000; i++)
        vec.push_back(sum[i]);
    Linear_Seq bmi;
    printf("%lld\n", bmi.solve(vec, n, P, false));
}
4.5
    杜教筛
#include <bits/stdc++.h>
#define LL long long
const int P = 1e9 + 7;
using namespace std;
LL _pow(LL a, LL b)
{
    LL ret = 1;
    LL mul = a \% P;
    while (b)
    {
        if (b & 1)
            ret = ret * mul % P;
        mul = mul * mul % P;
        b = b >> 1;
    }
    return ret;
}
```

```
namespace dujiao
unordered_map<LL, LL> mapi;
const int M = 1e6 + 7;
LL sum[M] = \{0\};
int cnt = 0;
LL prim[M], phi[M];
bool vis[M];
void init(int maxn)
{
    mapi.clear();
    memset(vis, 0, sizeof(vis));
    phi[1] = 1;
    for (int i = 2; i <= maxn; i++)</pre>
    {
        if (!vis[i])
            prim[++cnt] = i;
            phi[i] = i - 1;
        }
        for (int j = 1; j <= cnt && prim[j] * i <= maxn; j++)</pre>
        {
            vis[i * prim[j]] = 1;
            if (i % prim[j] == 0)
            {
                 phi[i * prim[j]] = phi[i] * prim[j];
                 break;
            }
            else
                 phi[i * prim[j]] = phi[i] * (prim[j] - 1);
        }
    }
    for (int i = 1; i <= maxn; i++)</pre>
        sum[i] = (sum[i - 1] + (i * phi[i]) % P) % P;
inline LL s_fg(LL n)
{
    LL ans = n * (n + 1) \% P;
```

```
ans = ans * (2 * n + 1) \% P;
    ans = ans * pow(6, P - 2) \% P;
    return ans;
}
LL go(LL n)
{
    if (n < M)
        return sum[n];
    if (mapi[n])
        return mapi[n];
    LL ans = s_fg(n);
    for (int i = 2; i <= n;)
        int j = n / (n / i);
        LL t1 = 1LL * j * (j + 1) / 2 - 1LL * i * (i - 1) / 2;
        t1 = t1 \% P;
        ans = (ans - t1 * go(n / i) % P + P) % P;
        i = j + 1;
    }
    return mapi[n] = ans;
}
LL solve(LL n)
    return go(n);
}
} // namespace dujiao
int main()
{
    int T;
    scanf("%d", &T);
    int M = 1e6 + 7;
    dujiao::init(M);
    while (T--)
    {
        LL n;
        int a, b;
        scanf("%lld%d%d", &n, &a, &b);
        LL ans = dujiao::go(n);
```

```
ans = (ans - 1 + P) \% P;
        ans = ans * pow(2, P - 2) \% P;
        printf("%lld\n", ans);
    }
    return 0;
}
4.6
    线性基
#include <bits/stdc++.h>
#define ll long long
const int Maxn = 5e4 + 20;
using namespace std;
ll a[Maxn][50];
int cnt[Maxn];
struct L_B
{
    long long d[35]; //若为 LL, 开 63++;
    L_B()
    {
        memset(d, 0, sizeof(d));
    }
    void init()
    {
        memset(d, 0, sizeof(d));
    void insert(ll x)
    {
        for (int i = 32; i >= 0; i--)
        { //or 64
            if (x & (111 << i))
            {
                if (!d[i])
                {
                    d[i] = x;
                    break;
                x = d[i];
            }
```

```
}
}
bool check(ll x)
{
    for (int i = 32; i >= 0; i--)
    {
        if (x & (111 << i))
            x = d[i];
    }
    return x == 011;
}
L_B merge(L_B k)
    L_B res, tmp = k;
    res.init();
    for (int i = 0; i <= 32; i++)
        11 x = d[i], y = 0;
        bool flag = false;
        for (int j = 32; j >= 0; j--)
        {
            if ((x >> j) & 1)
            {
                if (k.d[j])
                    x ^= k.d[j], y ^= tmp.d[j];
                else
                {
                    k.d[j] = x;
                    tmp.d[j] = y;
                    flag = true;
                    break;
                }
            }
        }
        if (!flag)
            res.d[i] = y;
    }
    return res;
```

```
}
} tree[(Maxn) << 2];</pre>
void build(int o, int 1, int r)
{
    tree[o].init();
    if (1 == r)
    {
        for (int j = 1; j <= cnt[1]; j++)</pre>
             tree[o].insert(a[l][j]);
    }
    else
    {
        int mid = (1 + r) / 2;
        build(o << 1, 1, mid);</pre>
        build(o << 1 | 1, mid + 1, r);
        tree[o] = tree[o << 1].merge(tree[o << 1 | 1]);</pre>
    }
}
bool query(int o, int l, int r, int ql, int qr, ll x)
{
    if (ql <= l && qr >= r)
        return tree[o].check(x);
    int mid = (1 + r) / 2;
    bool ans = true;
    if (ql <= mid)</pre>
        ans = ans & query(o << 1, 1, mid, q1, qr, x);
    if (qr >= mid + 1)
        ans = ans & query(o << 1 | 1, mid + 1, r, ql, qr, x);
    return ans;
}
int main()
{
    int n, m;
    cin >> n >> m;
    for (int i = 1; i <= n; i++)
    {
        cin >> cnt[i];
```

```
for (int j = 1; j <= cnt[i]; j++)</pre>
            cin >> a[i][j];
    build(1, 1, n);
    while (m--)
    {
        int 1, r;
        11 x;
        cin >> 1 >> r >> x;
        if (query(1, 1, n, 1, r, x))
        {
           puts("YES");
        }
        else
            puts("NO");
    }
    return 0;
}
```

5 杂项

6 计算几何