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Parallel computing

Space direction

- Domain decomposition methods (DDM): efficient approach to parallelization of the space direction.
 - · divide the whole domain in subdomains,
 - solve iteratively and independantly on each subdomain a smaller problem.
 - Corrections steps to propagate information from one subdomain to the others.
- Parallelization is not limited to space direction
 Time direction is also a candidate for parallelization

Parallel computing

Time Parallelization

- Multiple shooting Methods: Chartier, Philippe (1993), . . .
- Domain Decomposition type: Saha, Tradel, Tremaine (1996), Gander (1996), Gander, Halpern, Nataf (1999), ...
- Parareal method: Lions, Maday, Turinici (2001), . . .
- Direct solvers: Axelson, Verwer (1985), Worley (1991), Sleen, Sloan, Thomée (1999), Maday, Ronquist (2008), Güttel (2012), ...

M. Gander
50 Years of Time Parallel Time Integration,
Multiple Shooting and Time Domain Decomposition
Springer Verlag, 2015.

Parallel computing

Discretization of Evolution Equations

Parallel computation of Linear Systems

Applications Laplace's equation

Poisson's Equation

Limitations of the method

Conclusion

joint work with F. Hecht initiated with A. Loumi and P. Parnaudeau

Implicit discretisation

Consider the linear ODE

$$u'(t) = Au(t)$$

with initial condition $u(0) = u_0$.

FD Discretisation: implicit scheme (backward Euler)

$$(I - \delta t_n A) \mathbf{u}^{n+1} = \mathbf{u}^n.$$

For *m* time steps

$$A_{n+m}\cdots A_{n+1}A_n$$
 $\mathbf{u}^{n+m}=\mathbf{u}^n$.

with

$$A_k = I - \delta t_k A$$

Numerical discretisation

Let us recall what a sequential procedure would be.

- Computing $X = A_1 \cdots A_m$ before solving the linear system > very expensive
- A sensible sequential computation of the solution:
 - Compute successively \mathbf{u}^{n+j} (for $j=1,\cdots,m$) solution of

$$A_j\mathbf{u}^{n+j}=\mathbf{u}^{n+j-1}$$



Total $cost = m \times the cost of solving one linear system.$

Numerical discretisation

$$u'(t) = Au(t)$$
, with initial condition $u(0) = u_0$. Implicit scheme

$$X\mathbf{u}^{n+m}=\mathbf{u}^n$$
.

with

$$X = (I - \delta t_{n+m}A) \cdots (I - \delta t_{n+1}A)(I - \delta t_nA)$$

Key idea:

$$\mathbf{u}^{n+m} = \mathbf{X}^{-1}\mathbf{u}^n$$

Numerical discretisation

Take m=2.

$$X = (I - h_2 A)(I - h_1 A)$$

For $h_2 \neq h_1$, we have

$$X^{-1} = \alpha_1 (I - h_1 A)^{-1} + \alpha_2 (I - h_2 A)^{-1}$$

with

$$\alpha_1 = \frac{h_1}{h_1 - h_2}, \quad \alpha_2 = \frac{h_2}{h_2 - h_1}.$$

Numerical discretisation

Take m = 2. The solution of

$$X\mathbf{u}^{n+2} = \mathbf{u}^n$$

is broken down into

$$\mathbf{x} = \alpha_1 \mathbf{x}_1 + \alpha_2 \mathbf{x}_2$$

Each vector x_i is the unique solution of the linear system

$$A_i x_i = y$$

There are two independent linear systems to solve.



Computing *x* this way is as expensive as solving a single linear system (neglecting the $\alpha_i x_i$ summation).

Parallel computation of Linear Systems

 $(A_i)_{i=1}^m$ are $n \times n$ nonsingular real matrices

$$X = A_1 \cdots A_m$$
.

Problem: how to compute quickly the solution $x \in \mathbb{R}^n$ of the linear system

$$Xx = y$$

where $y \in \mathbb{R}^n$ is any given vector?

Consider m distinct real numbers h_i and

$$X = \prod_{i=1}^{m} (I + h_i A)$$

Proposition

The inverse matrix of X (with nonsingular matrices A_i) is

$$X^{-1} = \sum_{i=1}^{m} \alpha_i A_i^{-1}$$

with

$$\alpha_i = \prod_{k \neq i} (1 - h_k/h_i)^{-1}.$$

(partial fraction decomposition of rational functions!!)

Accordingly, the solution Xx = y is broken down into

$$x = \sum_{i=1}^{m} \alpha_i x_i$$

Each vector x_i is the unique solution of the linear system

$$A_i x_i = y$$

There are *m* independent linear systems to solve.



Cost

Computing x this way is as expensive as solving a single linear system (neglecting the $\alpha_i x_i$ summation).

In practice, the interconnections between processors are far from being neglectible.

$$M\frac{u^{n+1}-u^n}{h_1}=Bu^{n+1}$$

with

M: mass matrix

B: stiffness matrtix

For m=2

$$(I - h_1 A)(I - h_2 A)u^{n+2} = u^n$$

with

$$A = M^{-1}B$$

Exactly the same framework!

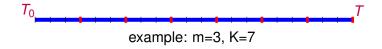
Consider the homogeneous equation

$$u'(t) = Au(t)$$

with initial condition $u(0) = u_0$. Implicit scheme using m distinct time steps h_i . Notations:

- intial and final times: T₀, T
- $t_n = T_0 + n$ time steps : $h_1 + h_2 + \cdots + h_m + h_1 + h_2 + \cdots$

Remark: solution computed only at times t_{km} , $k = 1, \dots, K$



FEM: Laplace's equation (2D)

Approximation from t_{km} to $t_{(k+1)m}$

$$A_1 u_{km+1} = u_{km}$$

$$A_2 u_{km+2} = u_{km+1}$$

$$\vdots \qquad \vdots$$

$$A_m u_{km+m} = u_{km+m-1}$$

•
$$Xu_{(k+1)m} = u_{km}$$

with $X = A_1 \cdots A_m$

FEM: Laplace's equation (2D)

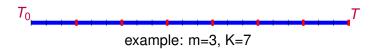
For
$$k = 1, \dots K$$

• Compute $\mathbf{u}_{(k+1)m}$ solution of

$$Xu_{(k+1)m}=u_{km}+g_{km}$$

$$u_{(k+1)m} = \sum_{i=1}^{m} \alpha_i X_i$$
 (Reduce step)
 x_i computed in parallel

$$x_i$$
 computed in parallel

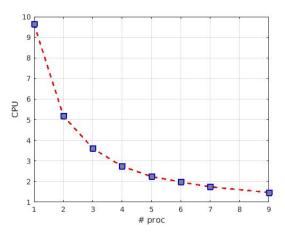


FreeFem Domain: unit square

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nb vertices =3721 , nb triangles =7200 , nb boundary edges 240
real T = .1; int Ndt=3000;
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Applications

FEM: Laplace's equation (2D)



Let us introduce some definitions related to the computation time.

• $f(1, T, N) = f(1, mK\delta t, N)$ is the sequential time to compute the solution at the final time $T = mK\delta t$, for a problem of size N:

$$f(1, T, N) = mK f(1, \delta t, N).$$

 f(m, T, N) is the parallel time to solve a problem of size N using *m* processors

$$f(m, T, N) = Kf(1, \delta t, N) + KC(m, N).$$

with C(m, N) the cost of the communications between m processors sharing a data of size N.

Efficient parallel solvers

Two quantities are of interest in parallel computing:

Speedup

$$S(m,T,N) = \frac{f(1,T,N)}{f(m,T,N)} = m \frac{f(1,\delta t,N)}{f(1,\delta t,N) + C(m,N)}.$$

Efficiency

$$E(m,T,N)=\frac{1}{m}S(m,T,N)=\frac{f(1,\delta t,N)}{f(1,\delta t,N)+C(m,N)}.$$

In a perfect world, the communications have no cost \Longrightarrow ideal speedup and ideal efficiency

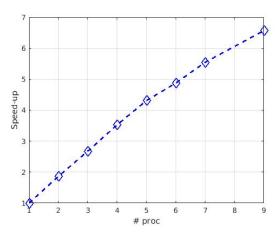
$$S^*(m, T, N) = m, \quad E^*(m, T, N) = 1.$$

Unfortunately our world is not perfect and the communication time C(m, N) is not at all neglectible.

Applications

FEM : Laplace's equation (2D)

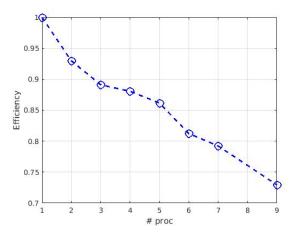
Speedup



Applications

FEM : Laplace's equation (2D)

Efficiency



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Consider the non-homogeneous equation

$$u'(t) = Au(t) + f$$

with initial condition $u(0) = u_0$. Implicit scheme using m distinct time steps h_i

Notations:

- intial and final times: T₀, T
- $t_n = T_0 + n$ time steps : $h_1 + h_2 + \cdots + h_m + h_1 + h_2 + \cdots$

Remark: solution computed only at times t_{km} , $k = 1, \dots, K$



Approximation from t_{km} to $t_{(k+1)m}$

$$A_1 u_{km+1} = u_{km} + h_1 f_{km+1}$$
 $A_2 u_{km+2} = u_{km+1} + h_2 f_{km+2}$
 \vdots
 \vdots
 $A_m u_{km+m} = u_{km+m-1} + h_m f_{km+m}$

$$Xu_{(k+1)m}=u_{km}+g_{km}$$
 with $X=A_1\cdots A_m$ and

$$g_{km} = h_1 f_{km+1} + A_1 \left\{ h_2 f_{km+2} + A_2 \left\{ h_3 f_{km+3} + \dots + A_{m-1} \left\{ h_m f_{km+m} \right\} \right\} \right\}$$

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- $h = \frac{1}{m} \sum_{i} h_{i}, T = Km^{2}h$
- Outer loop $k = 1, \cdots$
 - input: u_{km^2}
 - For $j = 1, \dots, m$ compute u_{km^2+im} by solving sequentially

$$Xu_{km^2+jm} = u_{km^2+(j-1)m} + \underbrace{g_{km^2+(j-1)m}}_{11}$$

• output: *u*_{(k+1)m²}

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- For $k=1,\cdots$
 - Compute in parallel $g_{km^2+(i-1)m}$ for $j=1,\cdots,m$
 - For $j = 1, \dots, m$
 - compute sequentially u_{km²+im} by solving

$$Xu_{km^2+jm} = u_{km^2+(j-1)m} + g_{km^2+(j-1)m}$$

each u_{km^2+im} being computed in parallel



See the computer sessions of the workshop

Limitation(s) of the method

Accuracy. In practical situations, the solution x_i is known only through an approximation \tilde{x}_i with an error ε_i :

$$\|\tilde{\mathbf{x}}_i - \mathbf{x}_i\| \leq \varepsilon_i$$
.

¿From which we deduce an upper bound on the error:

$$\|\tilde{\mathbf{x}} - \mathbf{x}\| \le (\max_{1 \le i \le m} \varepsilon_i) \sum_{i=1}^m |\alpha_i|$$

this bound may be very large as *m* grows.

Conclusion

- New method to solve linear time-dependent systems using a fractional decomposition of the matrix resulting from the discretisation of the time evolution operator.
- This method has been applied to solve the bidimensional heat equation. We obtain good results.
- However, for precision and stability reasons, the use of our method should be limited to moderate number of processors.
- We are currently investigating several developpements of the method.
 - combine low-order methods (computed in parallele) to get a high-order one
 - combine with space Domain Decomposition.
 - ...
- In summary, the method is efficient, but should be limited to a small number of processors

References

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