

Study the Introspective Sort, which is a default sorting algorithm implemented in the Array.Sort method. You must address the following questions:

1. What is special about the Introspective Sort?

A : Introspective Sort combines QuickSort, HeapSort, and Insertion Sort. It starts with QuickSort, switches to HeapSort if recursion depth exceeds $O(\log n)$, and uses Insertion Sort for small subarrays for efficiency.

2. Does this method result in an unstable sorting; that is, if two elements are seen equal, might their order be not preserved after the data collection is sorted? Or does it perform a stable sorting, which preserves the order of elements that are seen equal?

A : QuickSort and HeapSort are both unstable, so Introsort does not preserve the order of equal elements.

3. What is the time complexity of this sorting algorithm?

A :

Best: $O(n \log n)$

Average: $O(n \log n)$

Worst-case: $O(n \log n)$

In Task 1.1P on the `Vector<T>` class, you have completed/been provided with a number of methods (and properties), such as `Count`, `Capacity`, `Add`, `IndexOf`, `Insert`, `Clear`, `Contains`, `Remove`, and `RemoveAt`. Answer the following questions:

1. What is the time complexity of each of these operations?

A:

`Count/Capacity`: $O(1)$

`Add`: $O(1)$

`IndexOf/Contains`: $O(n)$.

Insert/Remove/RemoveAt: $O(n)$

Clear: $O(1)$

2. Does your implementation match the complexity of the equivalent operations provided by the Microsoft .NET Framework for its `List<T>` collection?

A: Yes they are matches, because both use dynamic arrays with identical amortized complexities.

Answer and explain whether the following statements are right or wrong.

- 3.a. A $\theta(n^3)$ algorithm always takes longer to run than a $\theta(\log n)$ algorithm.
- 3.b. The best-case time complexity of the [Bubble Sort](#) algorithm is always $O(n)$.
- 3.c. The worst-case time complexity of the [Insertion Sort](#) algorithm is always $O(n^2)$.
- 3.d. The [Selection Sort](#) is an in-place sorting algorithm.

3.a: It's wrong, it only true for large n , not always.

3.b: It's right sorted input requires $O(n)$ comparisons.

3.c: It's right, reversed input causes $O(n^2)$ swaps.

3.d: It's right because it uses $O(1)$ extra space.

- 3.e. The worst-case space complexity of the [Insertion Sort](#) algorithm is $O(1)$.
- 3.f. $2 + 2 = O(1)$
- 3.g. $n^3 + 10^6n = O(n^3)$
- 3.h. $n \log n = O(n)$
- 3.i. $\log n = o(n)$
- 3.j. $\log n + 2^{100} = \theta(\log n)$
- 3.k. $n \log n = \Omega(n)$
- 3.l. $n + 2^n = O(n)$
- 3.m. $n + \frac{100n}{\log n} = o(n)$
- 3.n. $n! = O(n)$

3.e: It's right because the worst case will be Sorts in place.

3.f: It's right because the time will be constant, and there is no other condition.

3.g: It's wrong because 10^6n grows much faster, with n than n does and cannot be reduced to $O(n)$.

3. h. It's wrong because $n \log(n)$ grows faster.

3.i. It's right because $\log n$ is getting smaller and smaller.

3.j. It's right because $\log n$ will eventually exceed any constant term as long as n is large enough.

3.k. It's right because $(n \log n)$ grows at least as fast as n .

3.l. It's wrong because it will be more rely on 2^n .

3.m. It's wrong because both n have the same increase rate.

3.n. It's wrong because $(n!)$ grows faster than any polynomial.

Give your best- and worst-case asymptotic runtime analysis to the following code snippets.

4.1. Let flag be a random Boolean variable, whose value is either true or false.

```
int count = 0;
for (int i = 0; i < n; i++)
{
    if (flag)
    {
        for (j = i+1; i < n; j++)
        {
            count = count + i + j;
        }
    }
}
```

A: Best-case: $O(n)$ (if flag is false).

Worst-case: $O(n^2)$ (if flag is true).

4.2. Let $\text{random}()$ be a function producing a real random number in the range $[0,1)$.

```

int count = 0;
for (int i = 0; i < n; i++)
{
    int num = random();
    if( num < 0.01 )
    {
        count = count + 1;
    }
}
int num = count;
for (int j = 0; j < num; j++)
{
    count = count + j;
}

```

A: Best-case: $O(n)$.

Worst-case: $O(n)$.

4.3. Let Compare(x, y) be a method with the time complexity of $\theta(1)$.

```

for (int i = 0; i < n; i++)
{
    for (int j = 0; j < n-i-1; j++)
    {
        if (a[j] > a[j+1])
        {
            Compare(a[j], a[j + 1]);
        }
    }
}

```

A: Best-case: $O(n^2)$.

Worst-case: $O(n^2)$.

Similar to bubble sort