

Anonymous Communication

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Anonymous Communication

Disclaimer:

The content to be presented aims only for helping students understand principles of anonymous communication. It should not be used for abusive Internet activities.

Communication



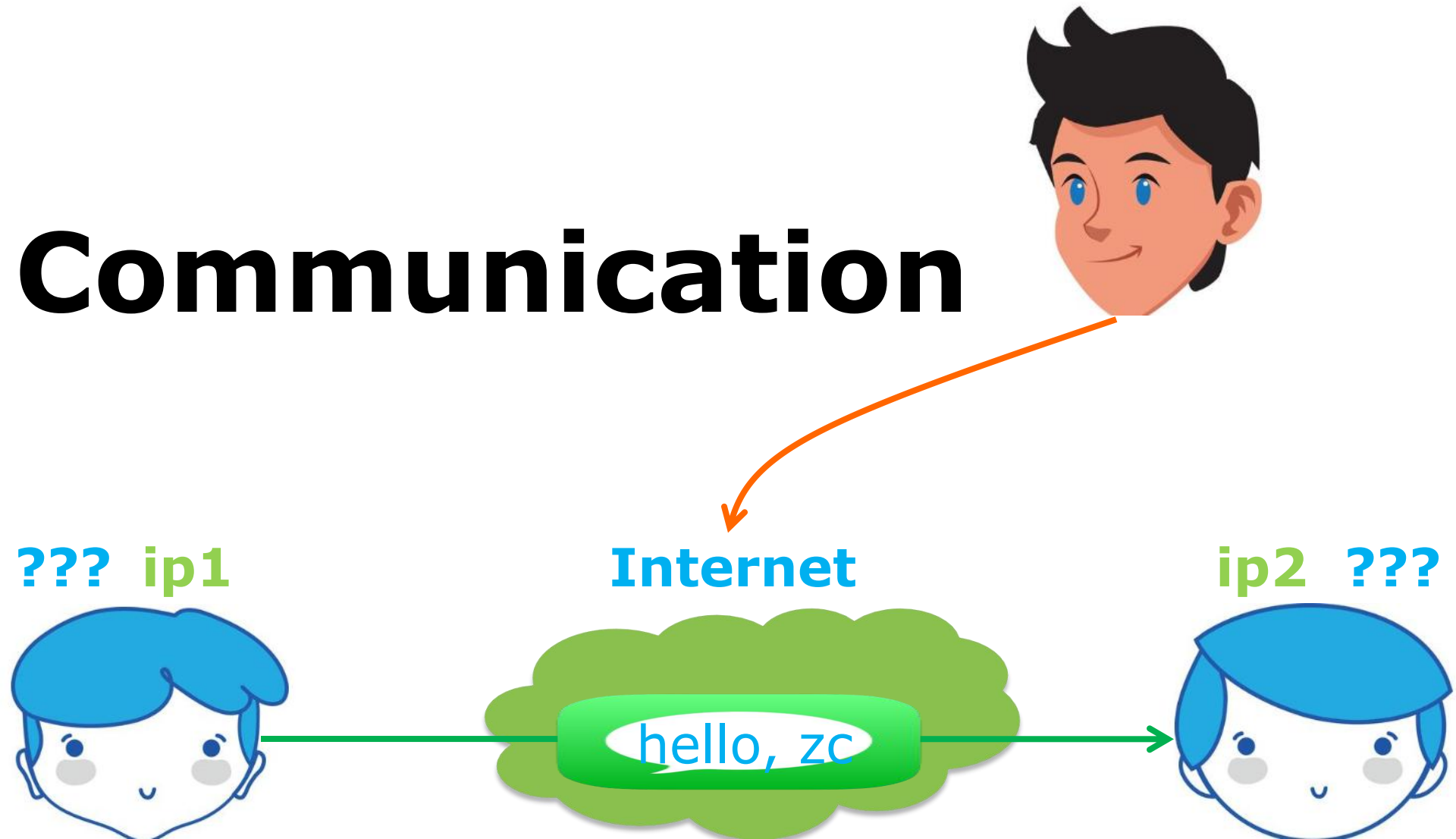
Communication



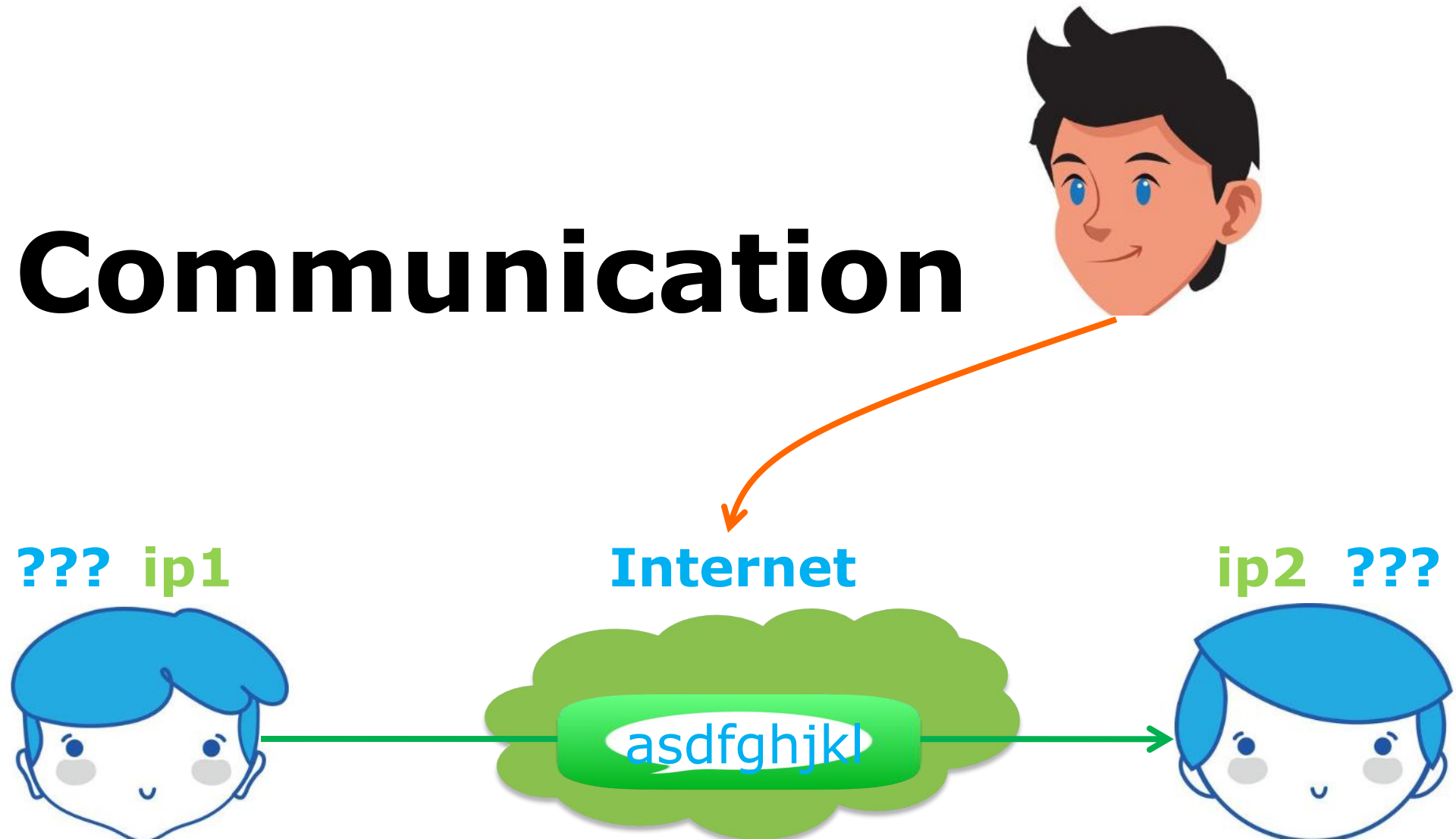
Communication



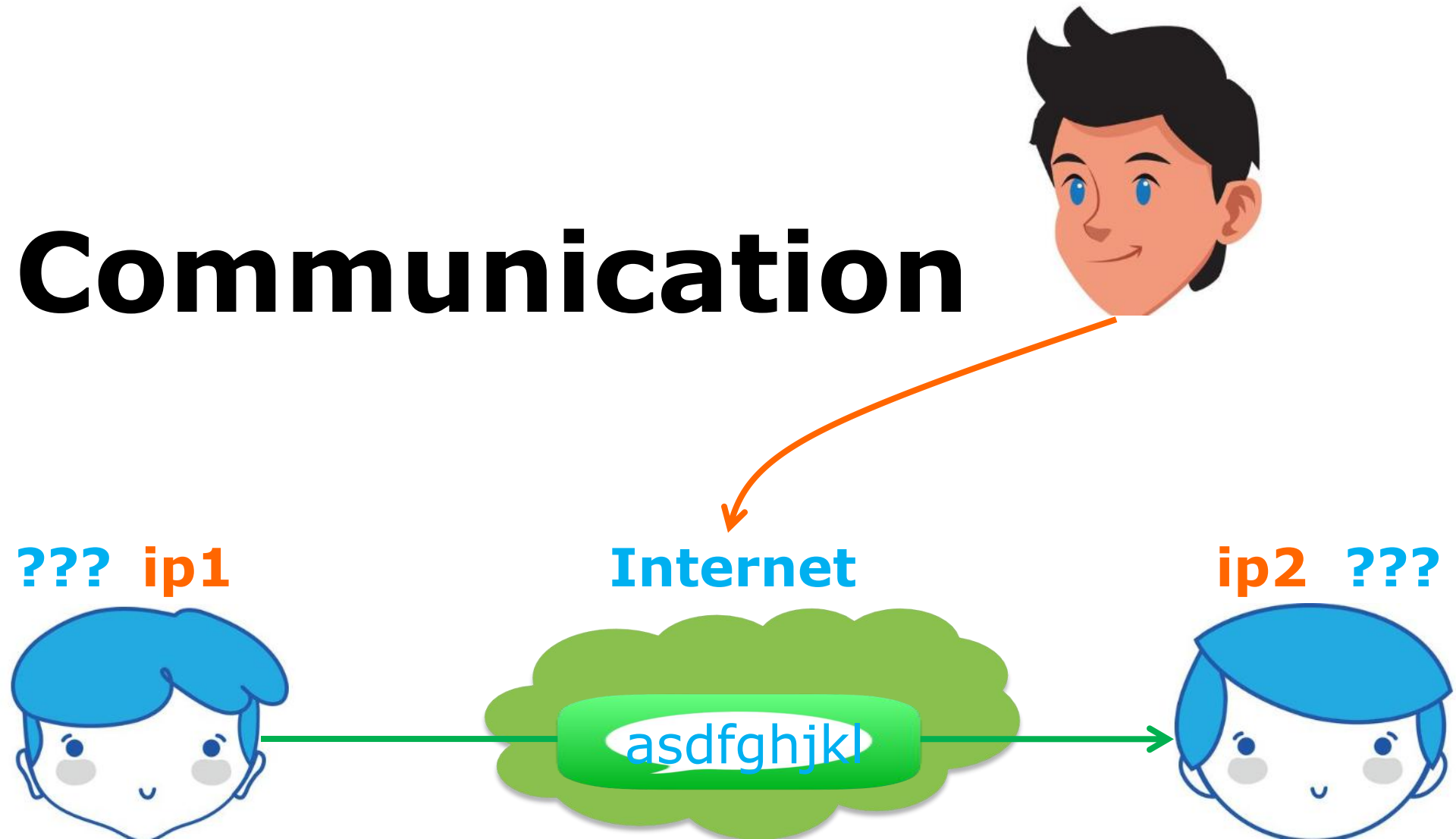
Communication



Communication



Communication



Communication



0	4	8	16	19	31
Version	IHL	Type of Service	Total Length		
Identification			Flags	Fragment Offset	
Time To Live		Protocol	Header Checksum		
Source IP Address					
Destination IP Address					
Options					Padding

Communication



anonymity & privacy:

who is communicating?

who are you talking to?

what type of activities?

what type of information?

Anonymous Communication

why wanted?



Anonymity for Mortals

- Unmonitored access to health and medical information
- Preservation of democracy:
anonymous election/jury
- Censorship circumvention:
anonymous access to otherwise restricted information
- ...

Anonymity for Attackers

Misbehaviors without getting caught:

- Terrorism
- Darknet
- Spam
- Pirate
- ...

Anonymous Communication

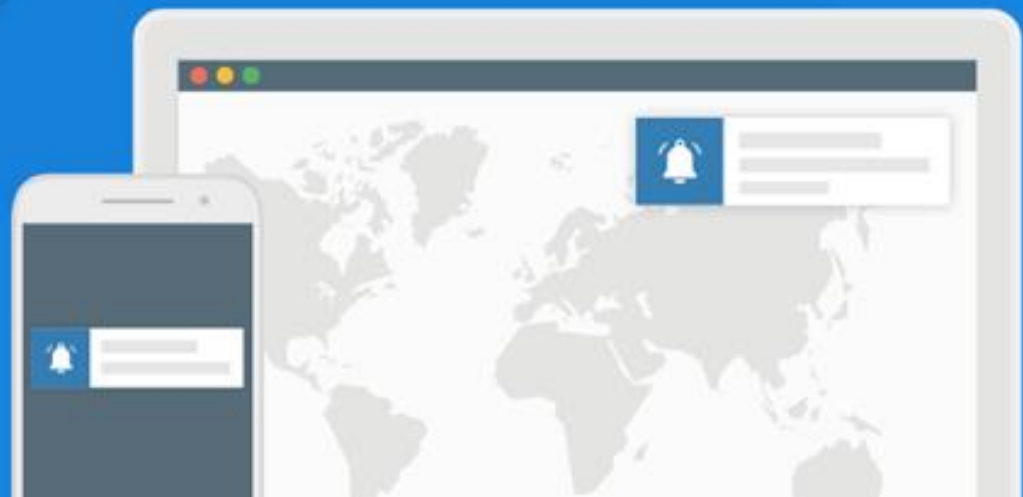
how to?



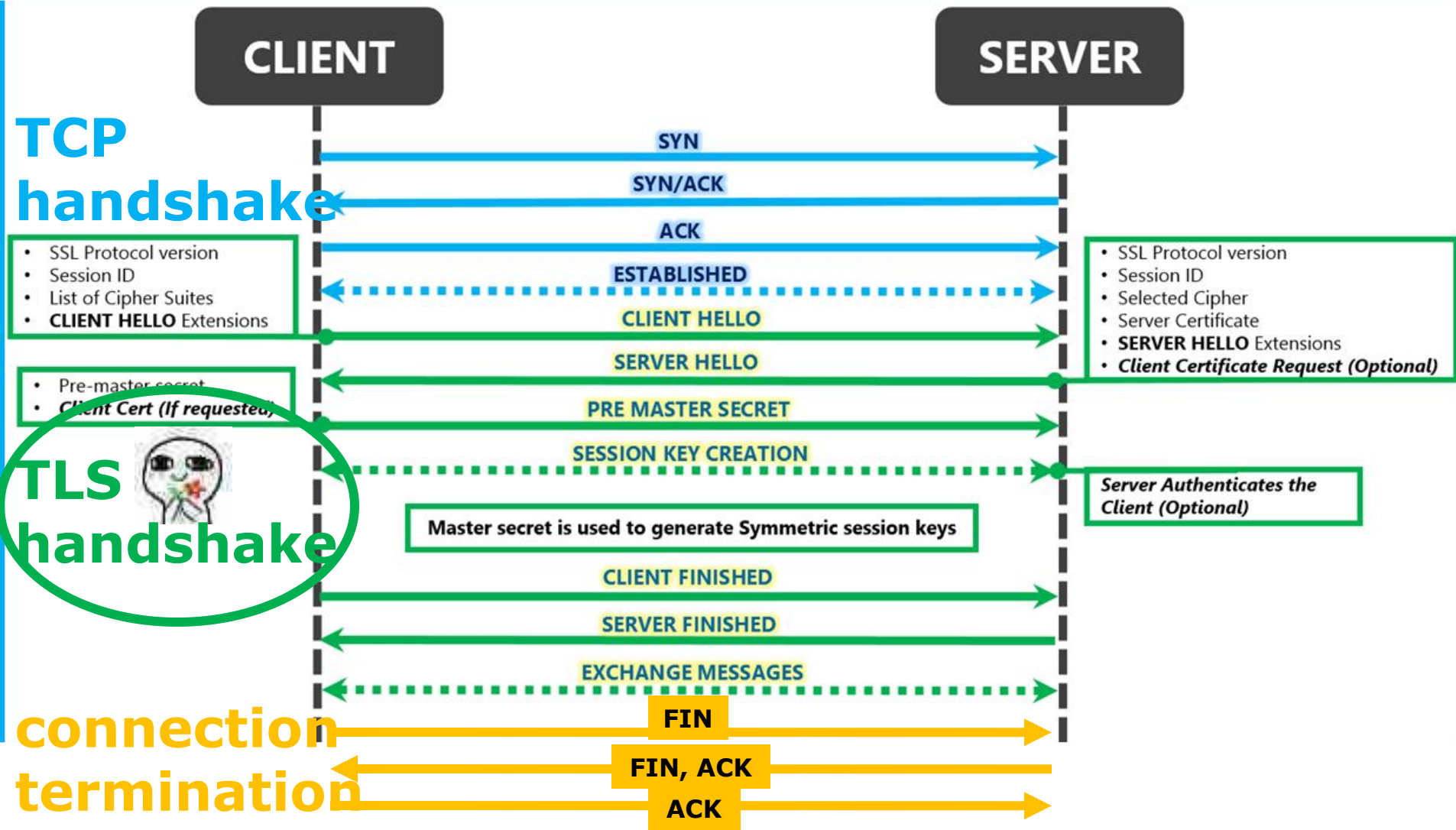
HTTPS?



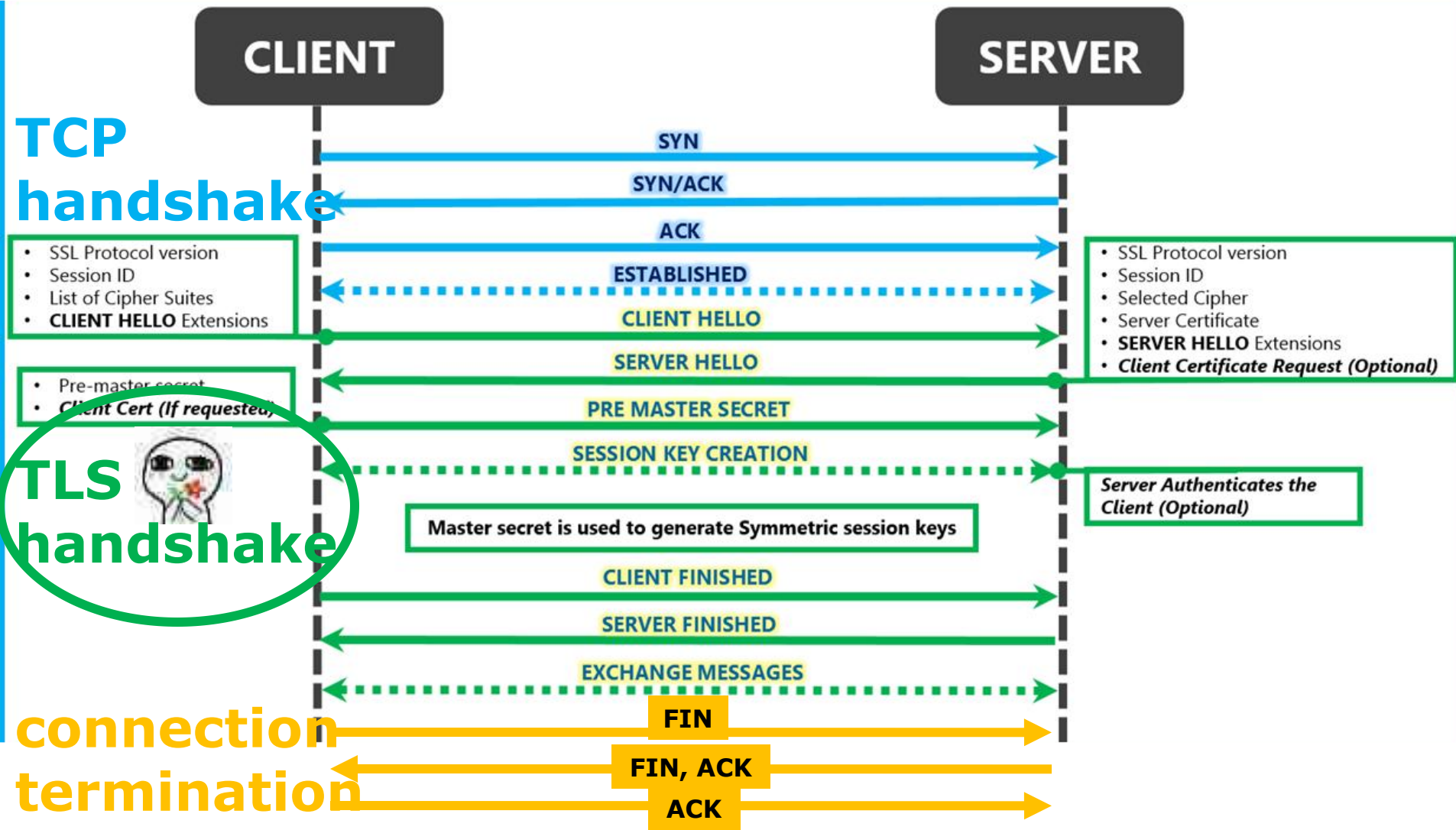
PushAlert Web Push Notification



HTTPS for Confidentiality



HTTPS Not for Anonymity

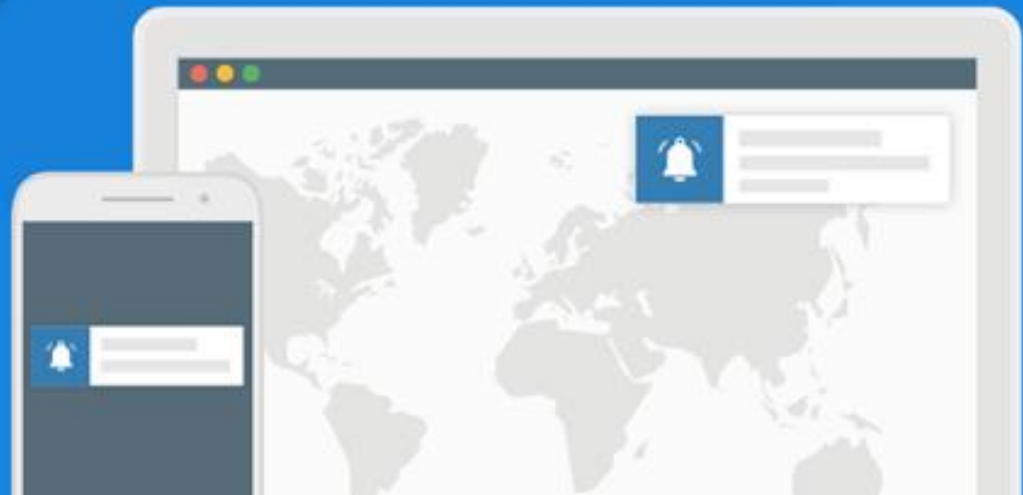


HTTPS?

<https://www.vote4or.com>



PushAlert Web Push Notification



Anonymous Communication



how to?

how to?

Anonymous Communication

hide destination address;



Anonymous Communication



hide destination address;
how to deliver packets to destination?

Anonymous Communication

RELAY!



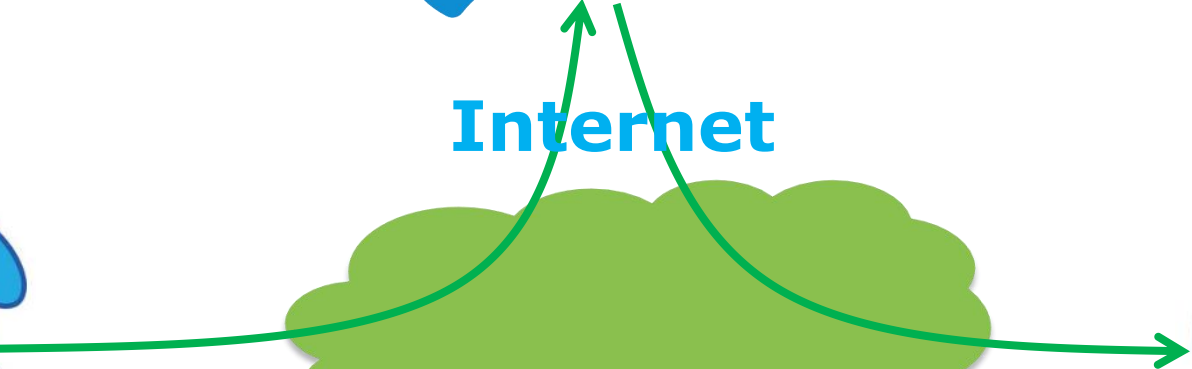
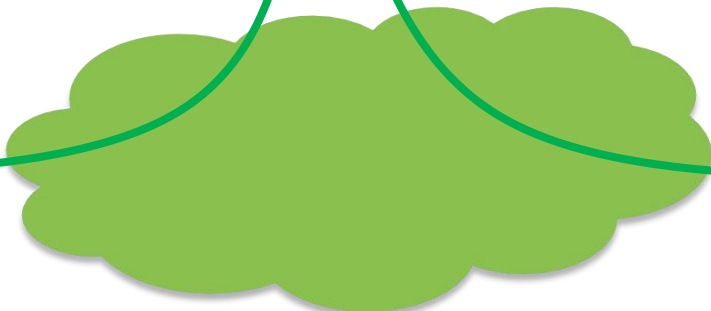
RELAY!

ip1

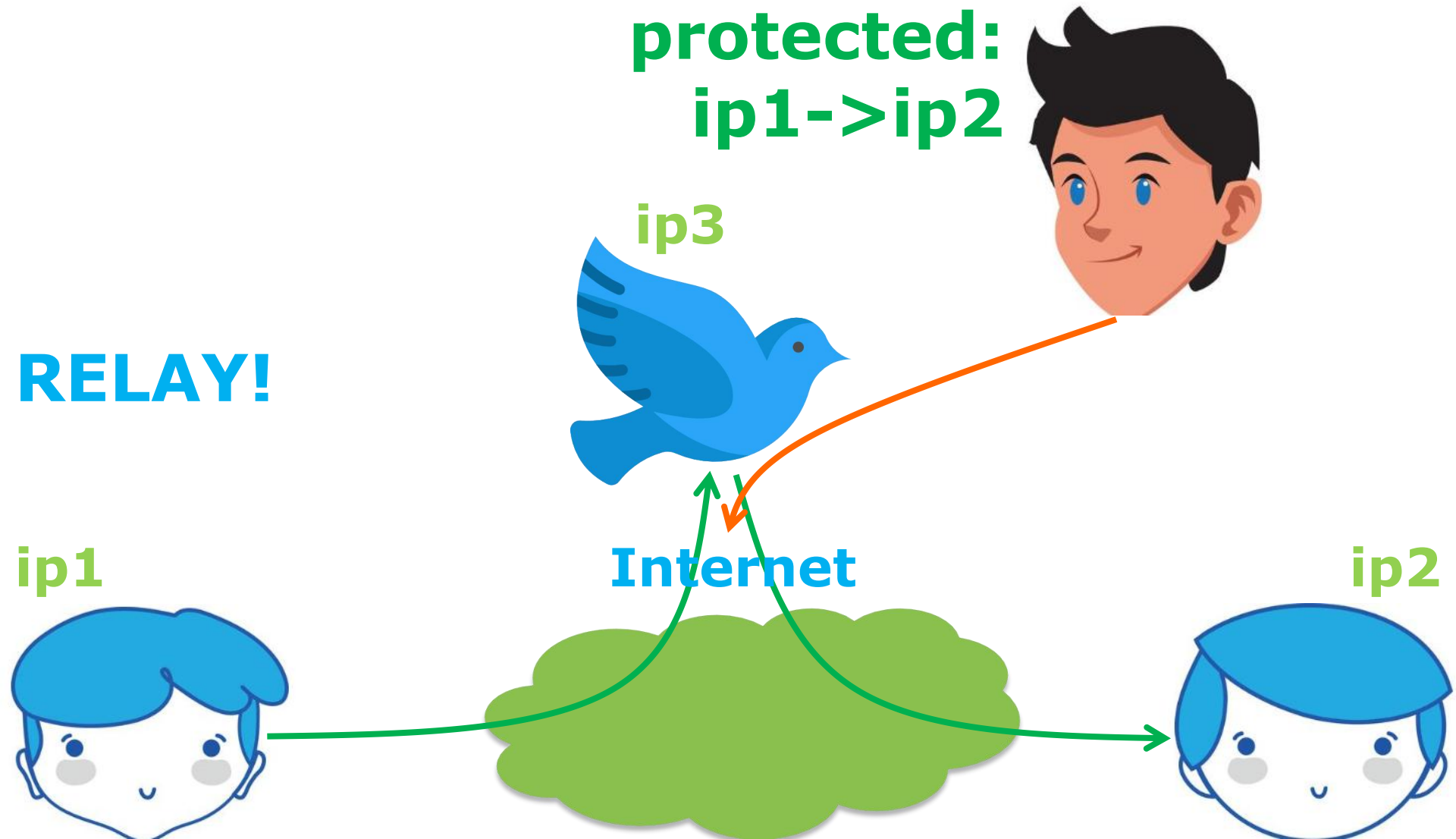
Internet

ip2

ip3



monitored:
ip1->ip3, ip3->ip2
protected:
ip1->ip2



overlay communication

monitored:
ip1->ip3, ip3->ip2

protected:
ip1->ip2

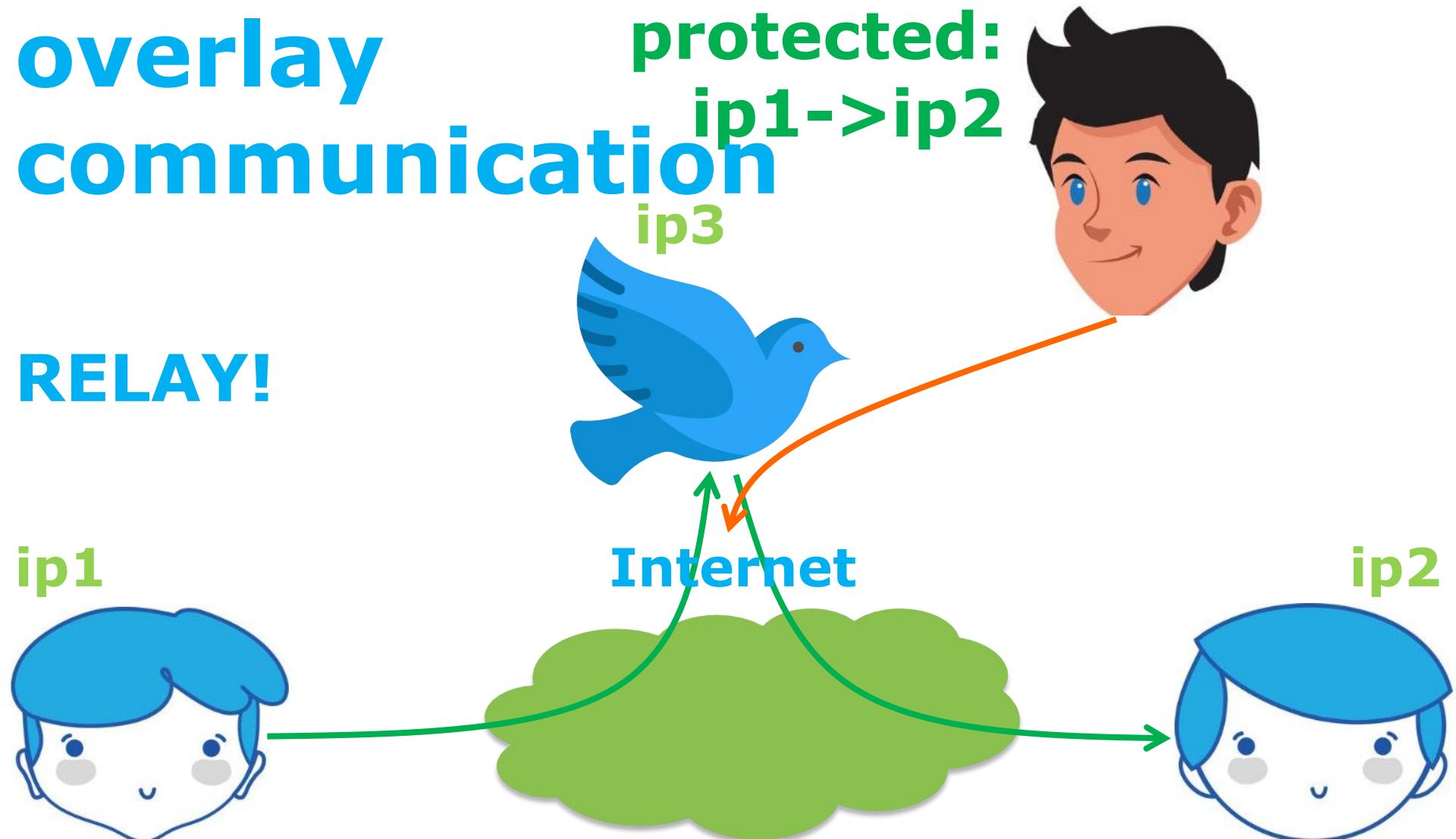
ip3

RELAY!

ip1

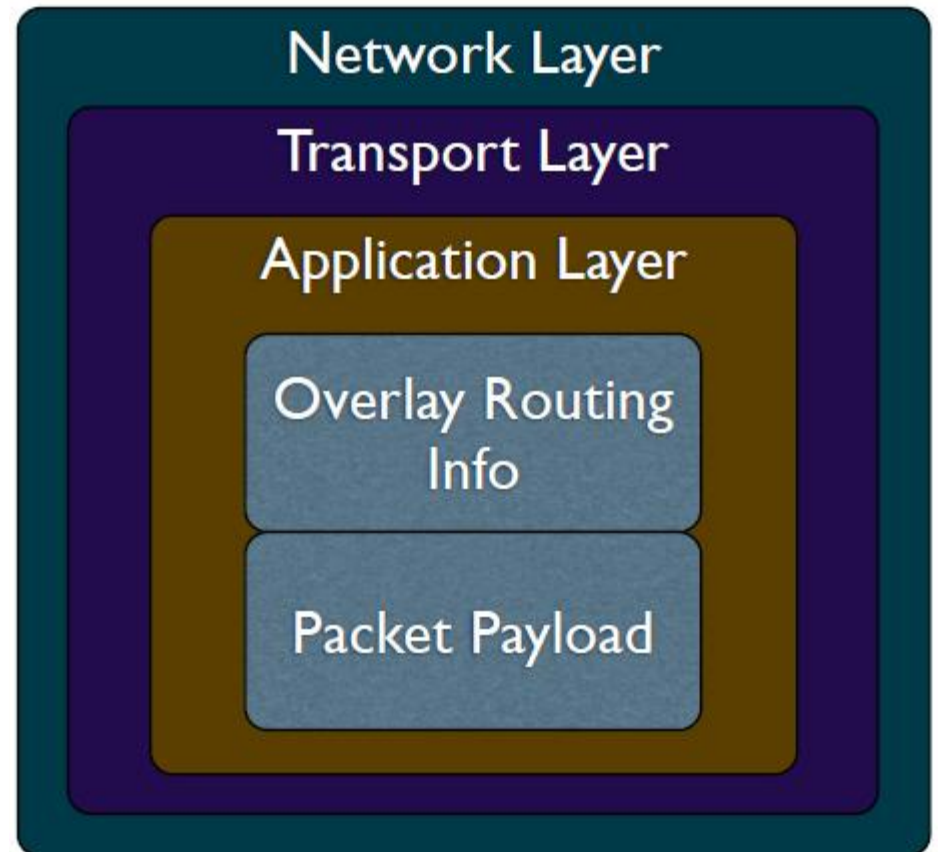
Internet

ip2

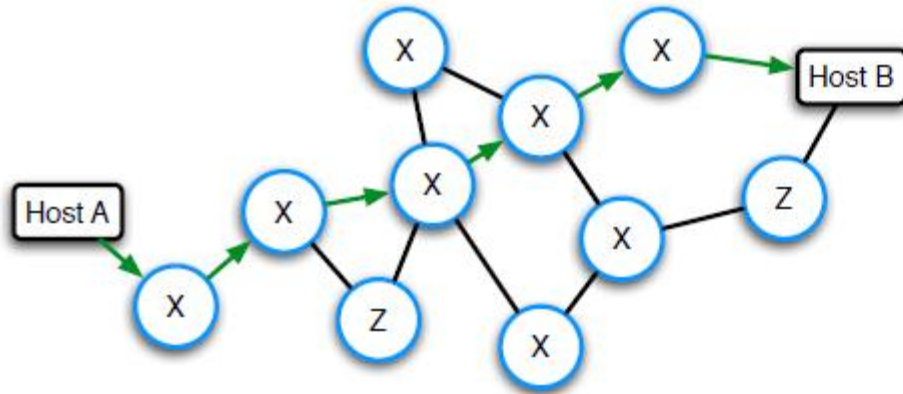


Overlay Network

- Handle routing at the application layer
- Tunnel messages inside other messages



Threat Model



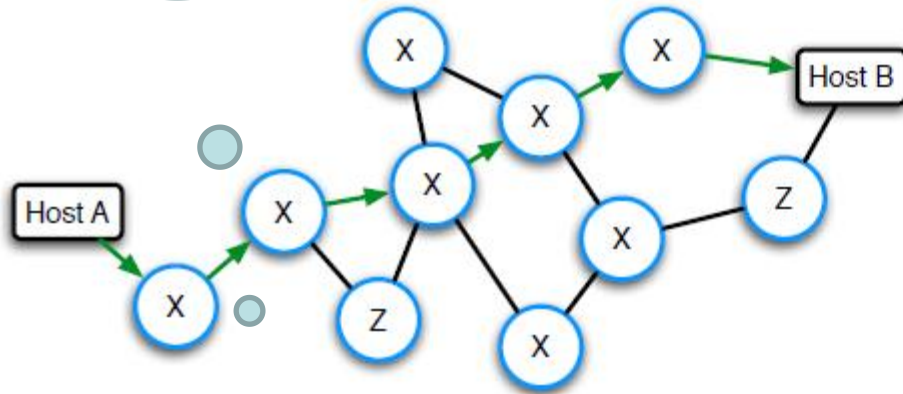
vs



- Insider Byzantine attacker with limited view of network:
an attacker might have tight control over a network (e.g., Z ASes), yet unlikely to observe entire Internet

Attacker is
part of the
network

Threat Model



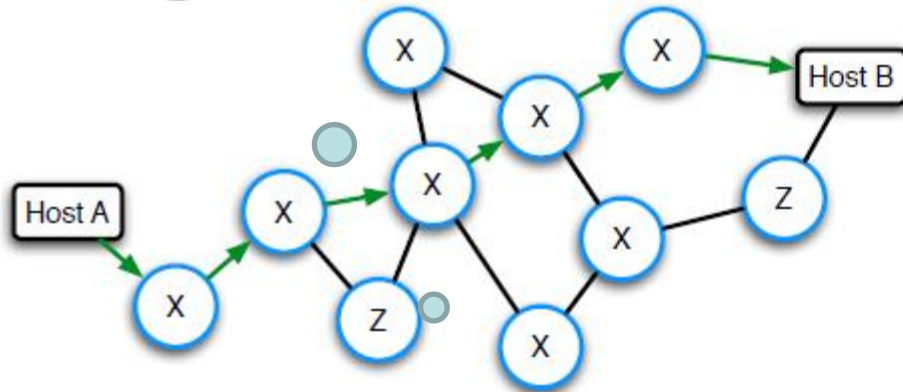
vs



- **Insider** Byzantine attacker with limited view of network: an attacker might have tight control over a network (e.g., Z ASes), yet unlikely to observe entire Internet

Attacker may
or may not
attack

Threat Model



vs

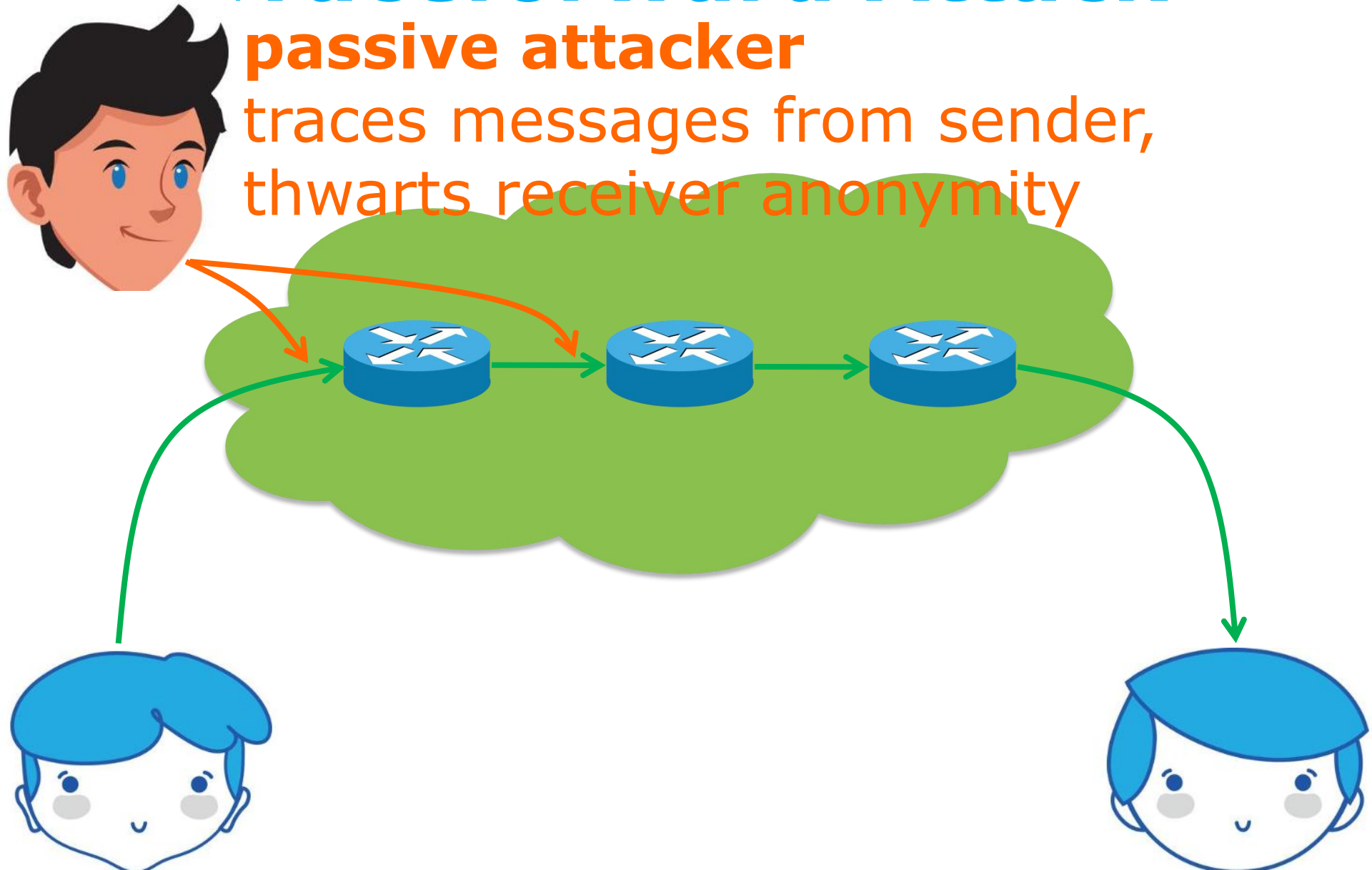


- Insider **Byzantine** attacker with limited view of network:
an attacker might have tight control over a network (e.g., Z ASes), yet unlikely to observe entire Internet

Traceforward Attack

passive attacker

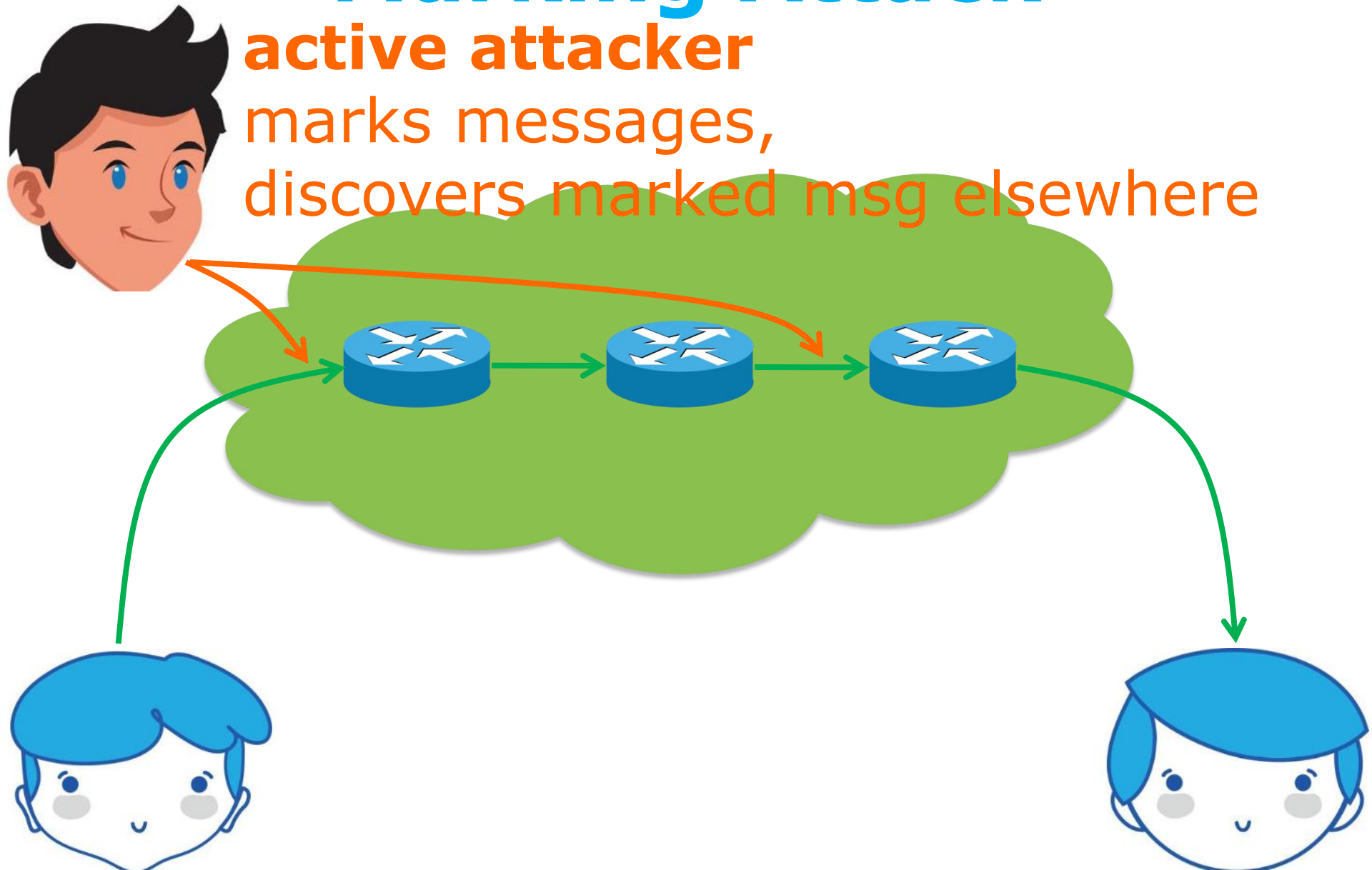
traces messages from sender,
thwarts receiver anonymity



Marking Attack

active attacker

marks messages,
discovers marked msg elsewhere



Anonymous Communication

anonymizing proxy

Anonymizing Proxy

PROXIFY

Anonymouse.org

hidemyass!
FREE WEB PROXY



 **FREEPROXY.CA**
Anonymous, free and proudly Canadian.

Zend2.com
Freedom of Speech

Anonymizing Proxy

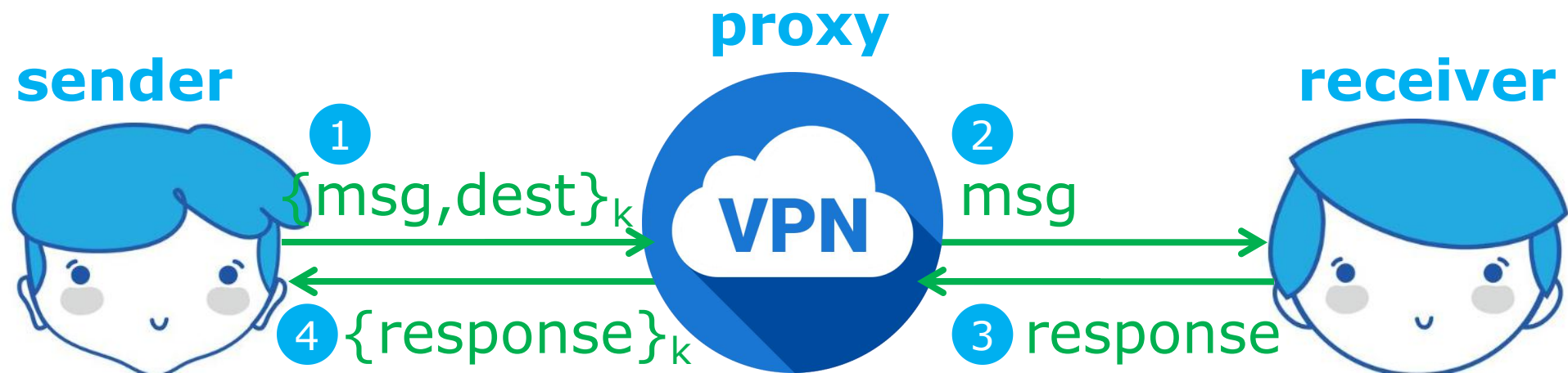
- intermediary between sender & receiver
- Sender relays all traffic through proxy
- Encrypt destination and payload

Anonymizing Proxy

- intermediary between sender & receiver
- Sender relays all traffic through proxy
- Encrypt destination and payload
- **Asymmetric technique:**
receiver not involved (or informed of)
anonymity

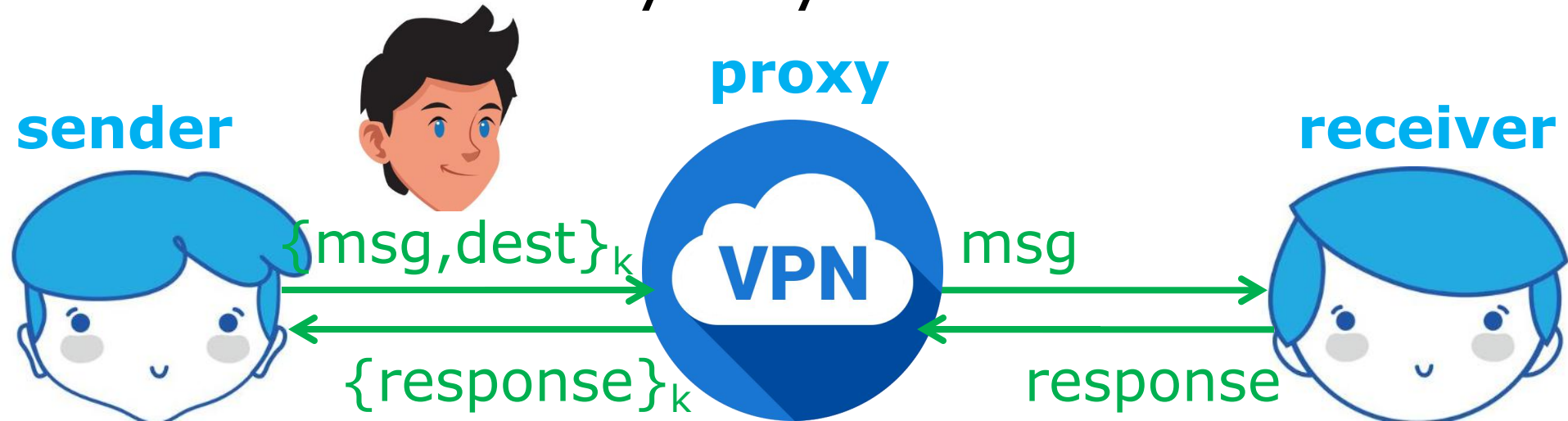
Anonymizing Proxy

- k : shared key of sender and proxy



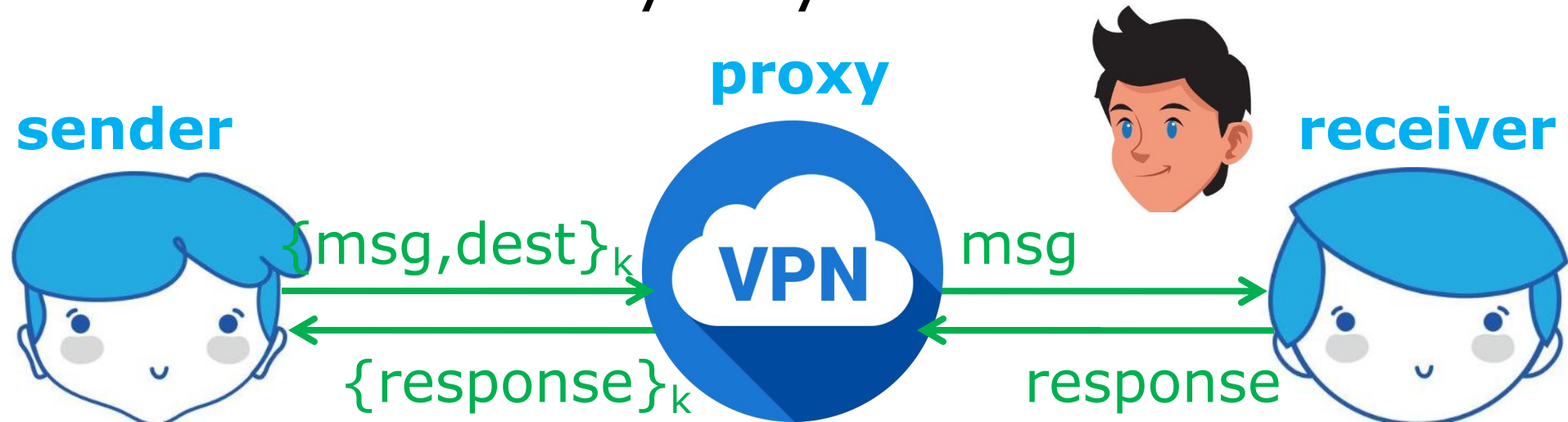
Anonymizing Proxy

- k : shared key of sender and proxy
- if attacker is located between sender and proxy:
sender anonymity: 0
receiver anonymity: 1



Anonymizing Proxy

- k : shared key of sender and proxy
- if attacker is located between proxy and receiver:
sender anonymity: 1
receiver anonymity: 0

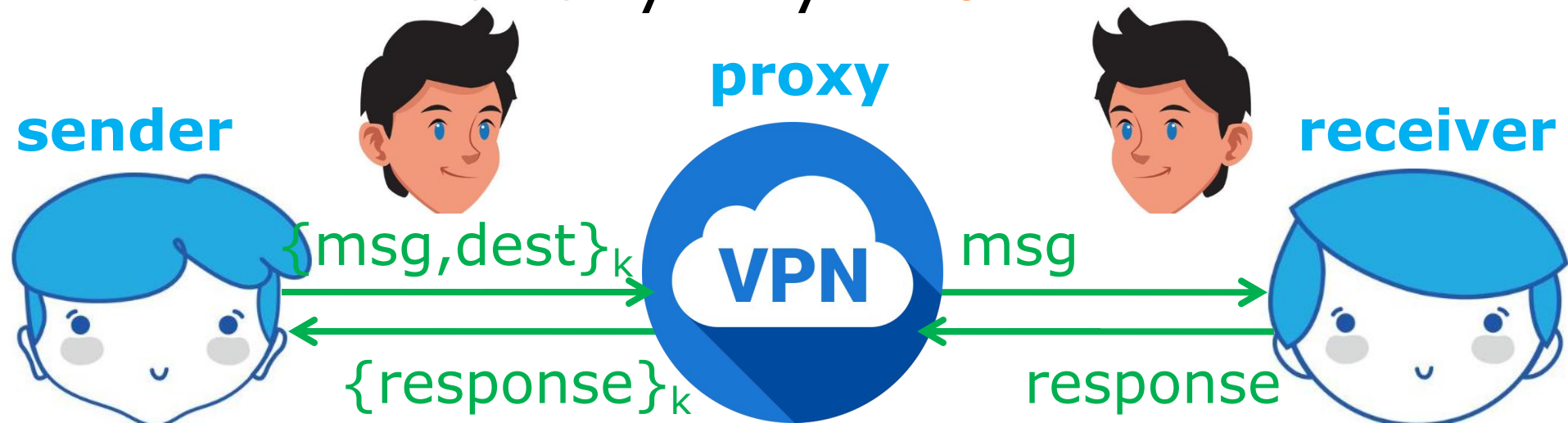


Anonymizing Proxy

- k : shared key of sender and proxy
- if two attackers collude to correlate ingress and egress proxy traffic:

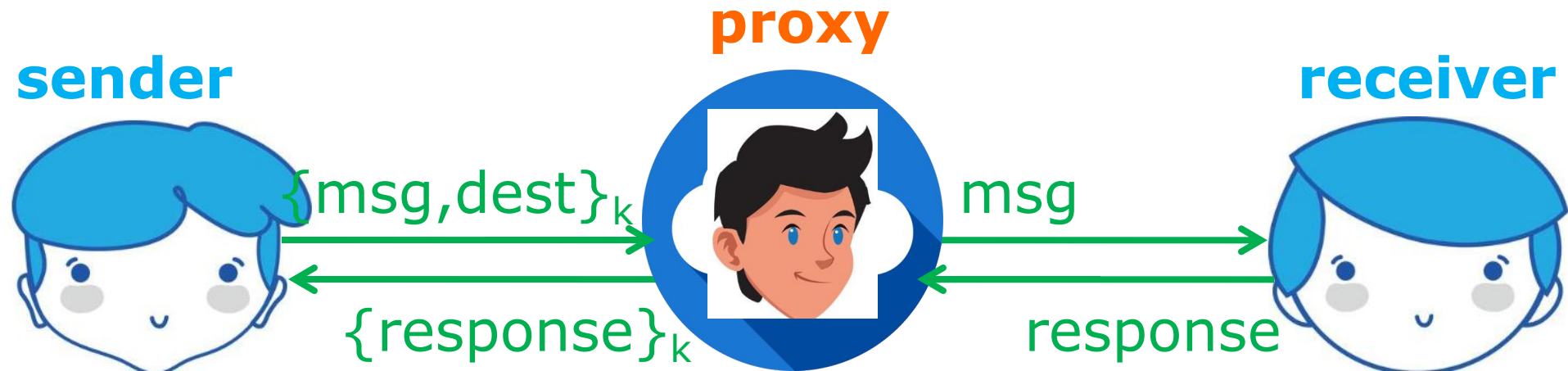
sender anonymity: 0

receiver anonymity: 0



Anonymizing Proxy

- If **attacker is the proxy** per se:
decrypt all messages;
crack both confidentiality & anonymity;
sender anonymity: 0
receiver anonymity: 0

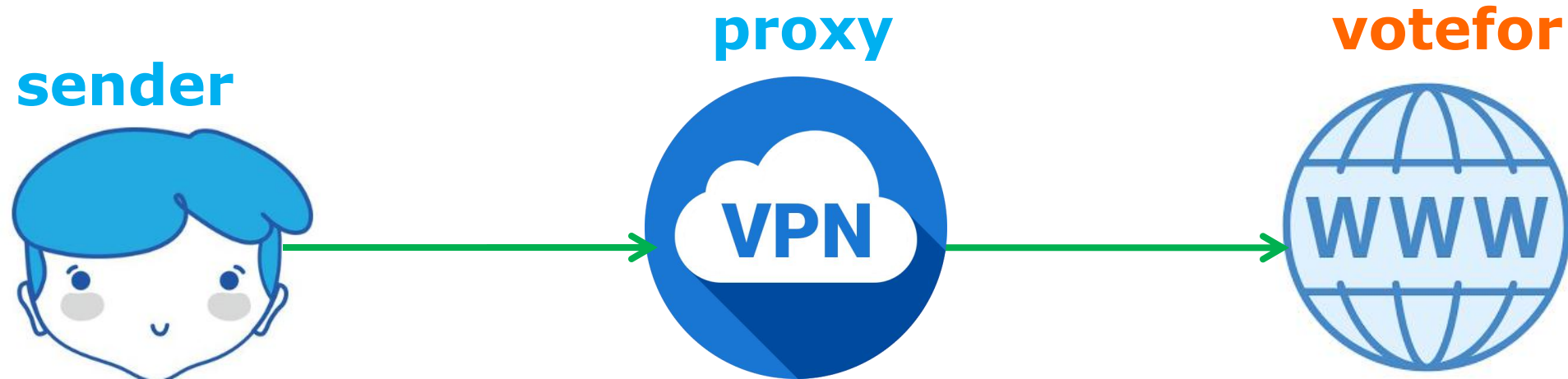


what if receiver is attacker?

what if receiver is attacker?
protect sender anonymity from receiver

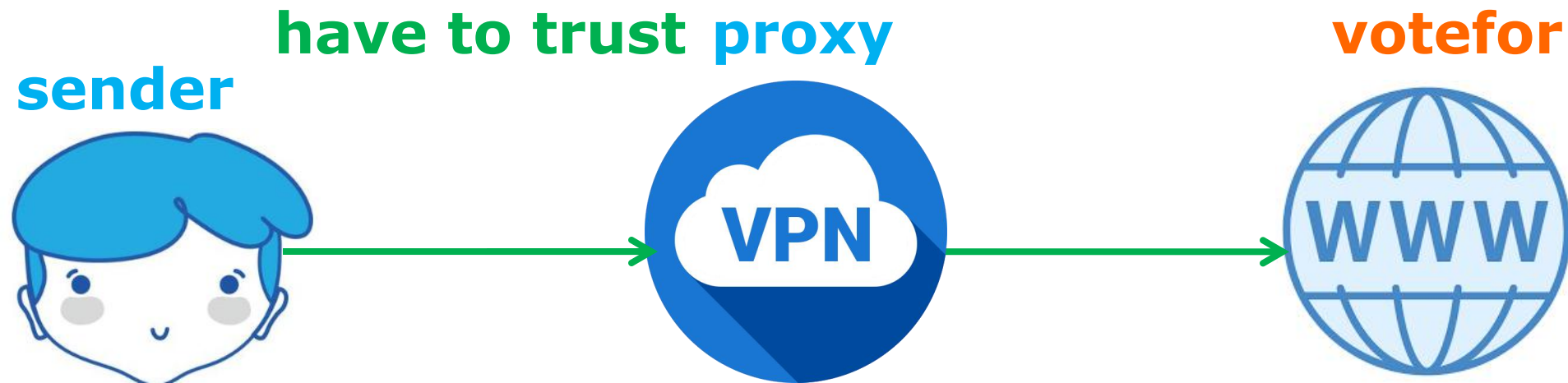
Receiver as Attacker

- Known-location attacker
- Use proxy to protect sender anonymity:
do not vote for letting votefor website
know that I accessed votefor /wink



Receiver as Attacker

- Known-location attacker
- Use proxy to protect sender anonymity:
do not vote for letting votefor website
know that I accessed votefor /wink



Anonymizing Proxy

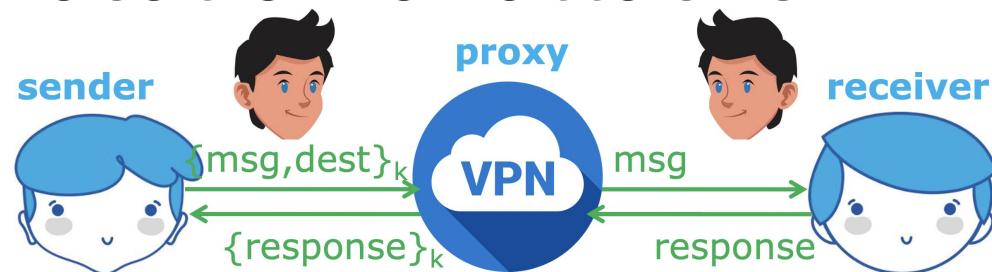
Advantages

- Easy to configure
- Require no active participation of receiver, which need not be aware of anonymity service
- Have been widely deployed on Internet

Anonymizing Proxy

Disadvantages

- Require trusted third party
proxy may release logs,
or sell them,
or blackmail sender
- Anonymity largely depends on the
(likely unknown) location of attacker



how to evade attacker?

how to evade attacker?
dynamize proxy location

Crowds Algorithm

- Basic idea:
get lost in a crowd
- Jump from one crowd to another
- Members of a crowd called Jondos

Crowds Algorithm

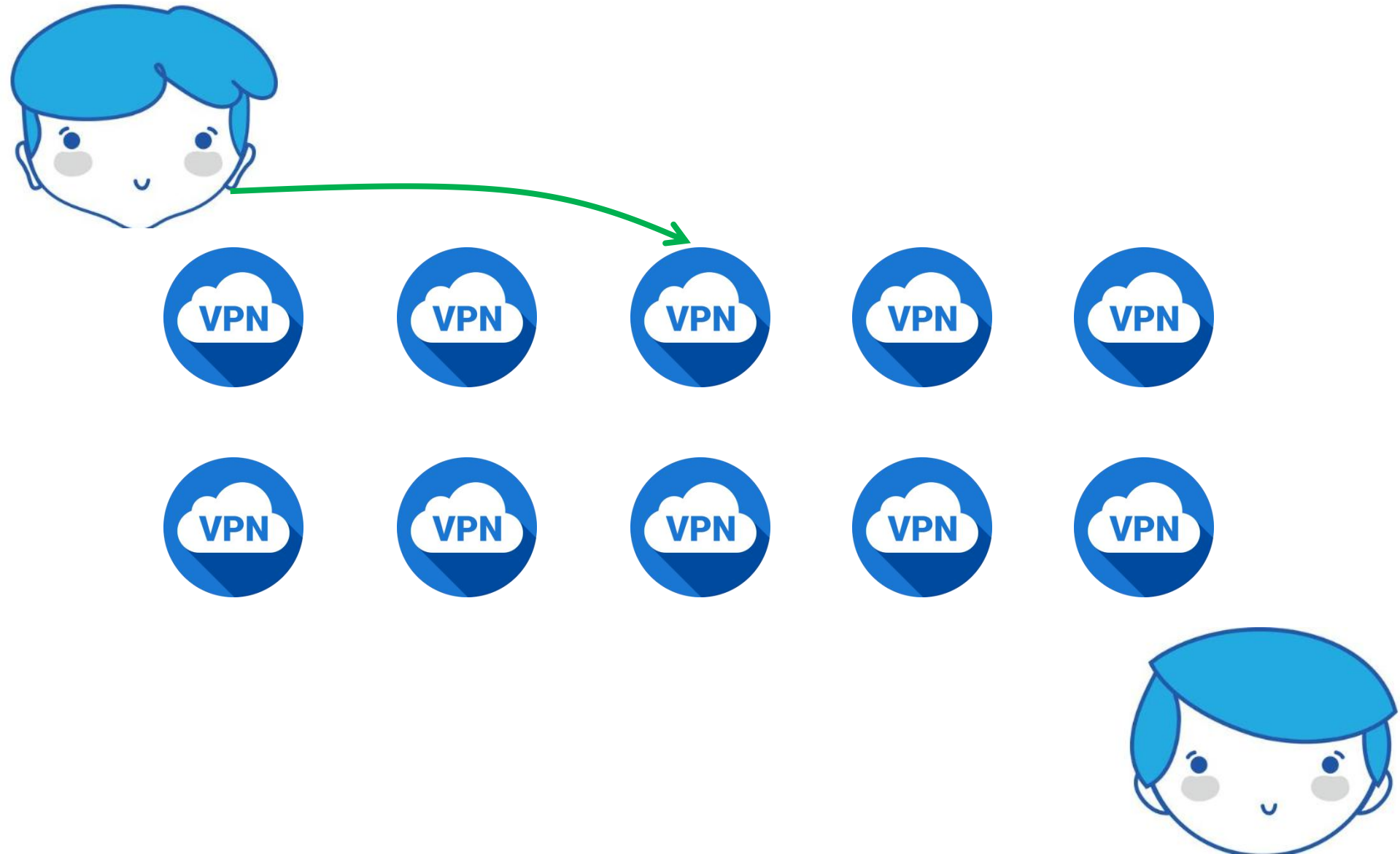
Algorithm:

- Relay message to random jondo
- With probability p , jondo forwards message to another jondo
- With probability $1-p$, jondo delivers message to its intended destination

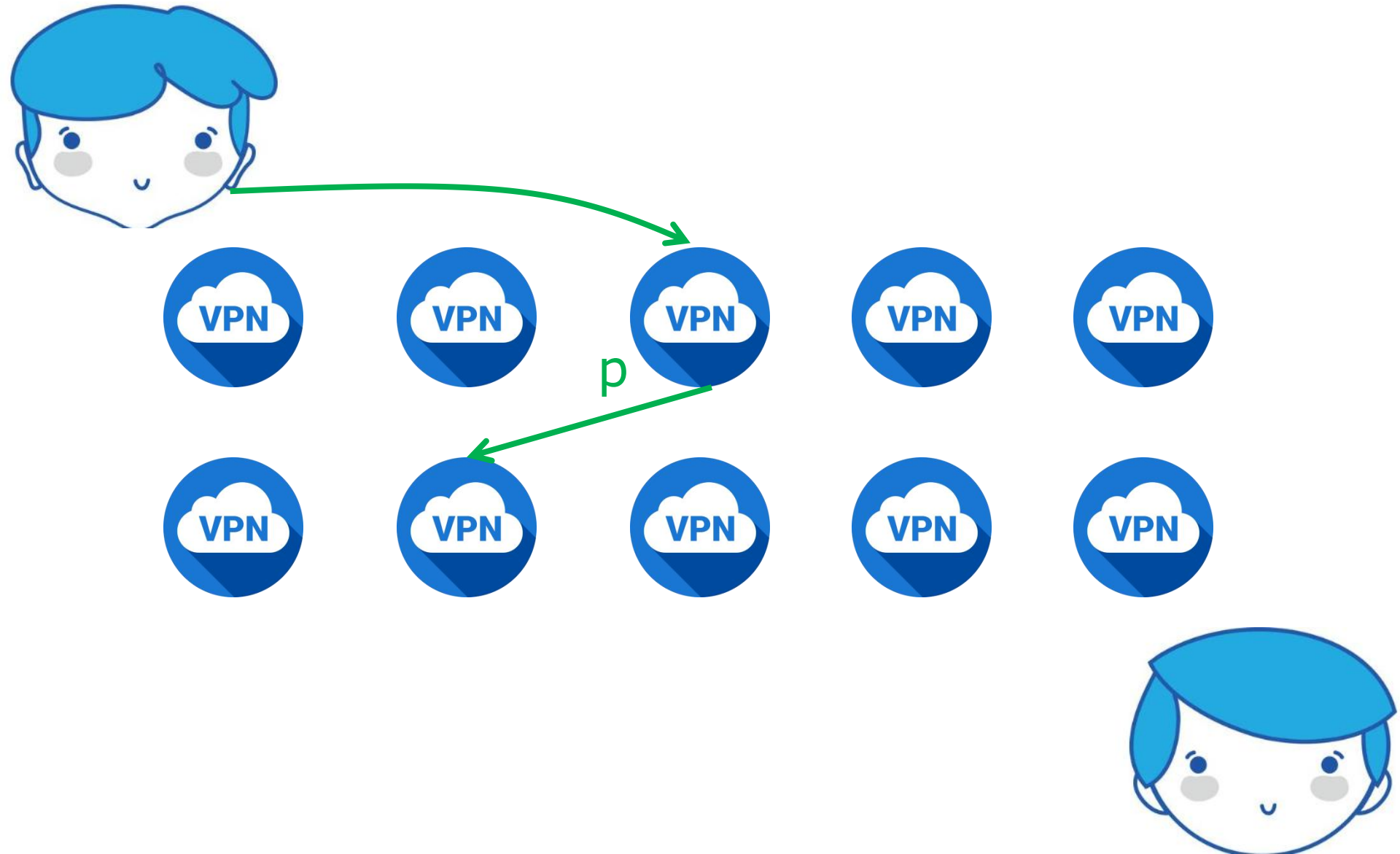
Crowds Algorithm



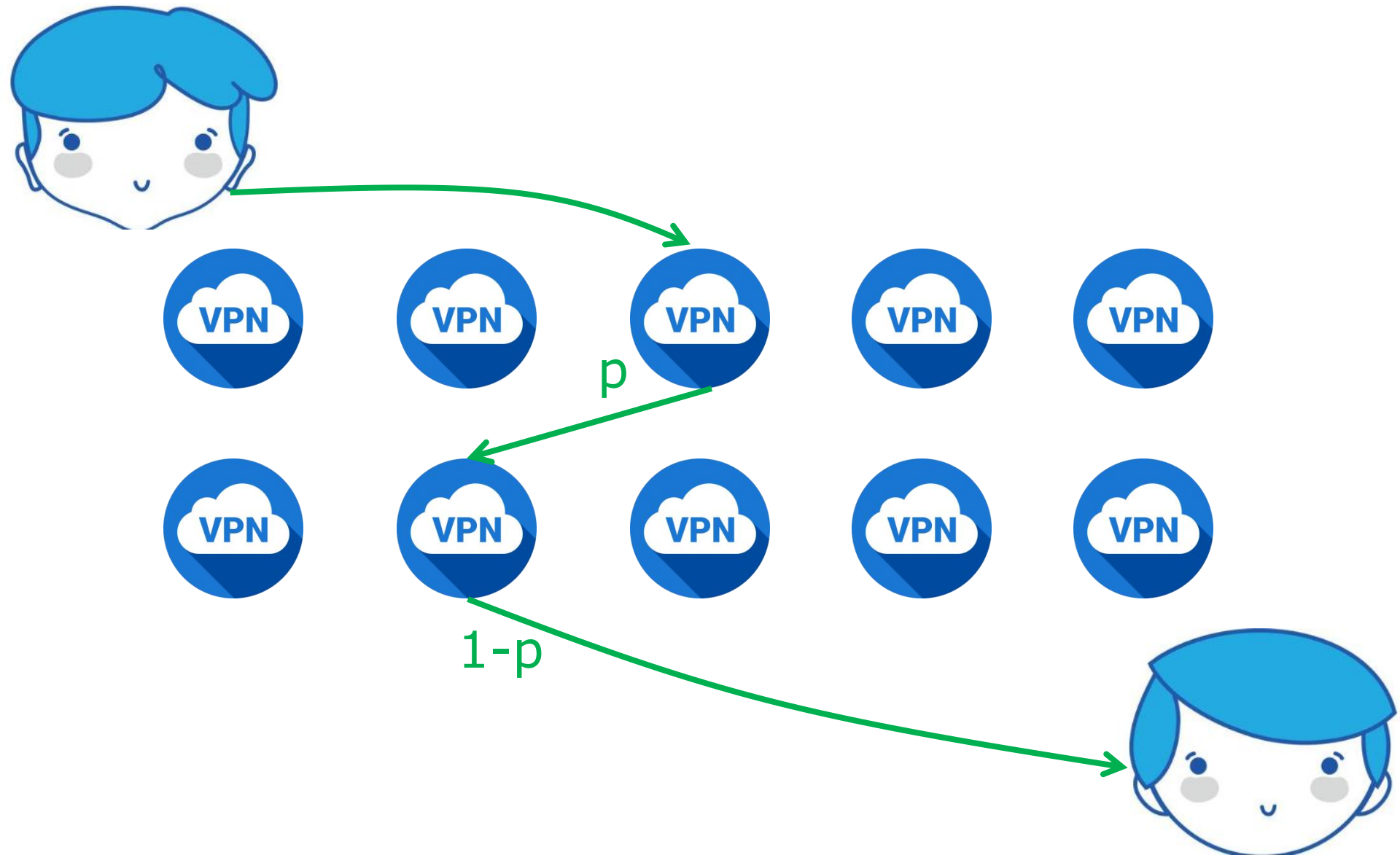
Crowds Algorithm



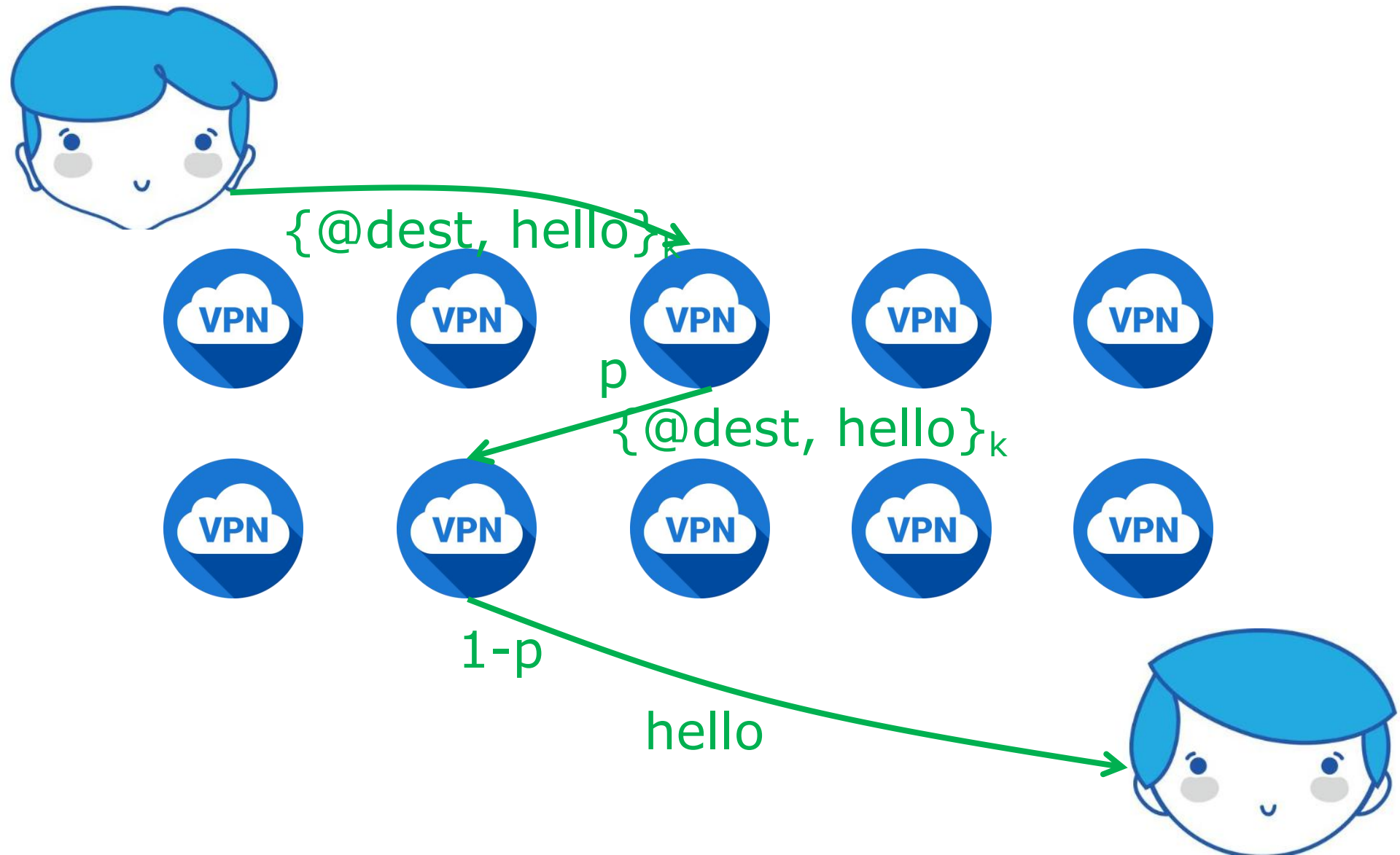
Crowds Algorithm



Crowds Algorithm

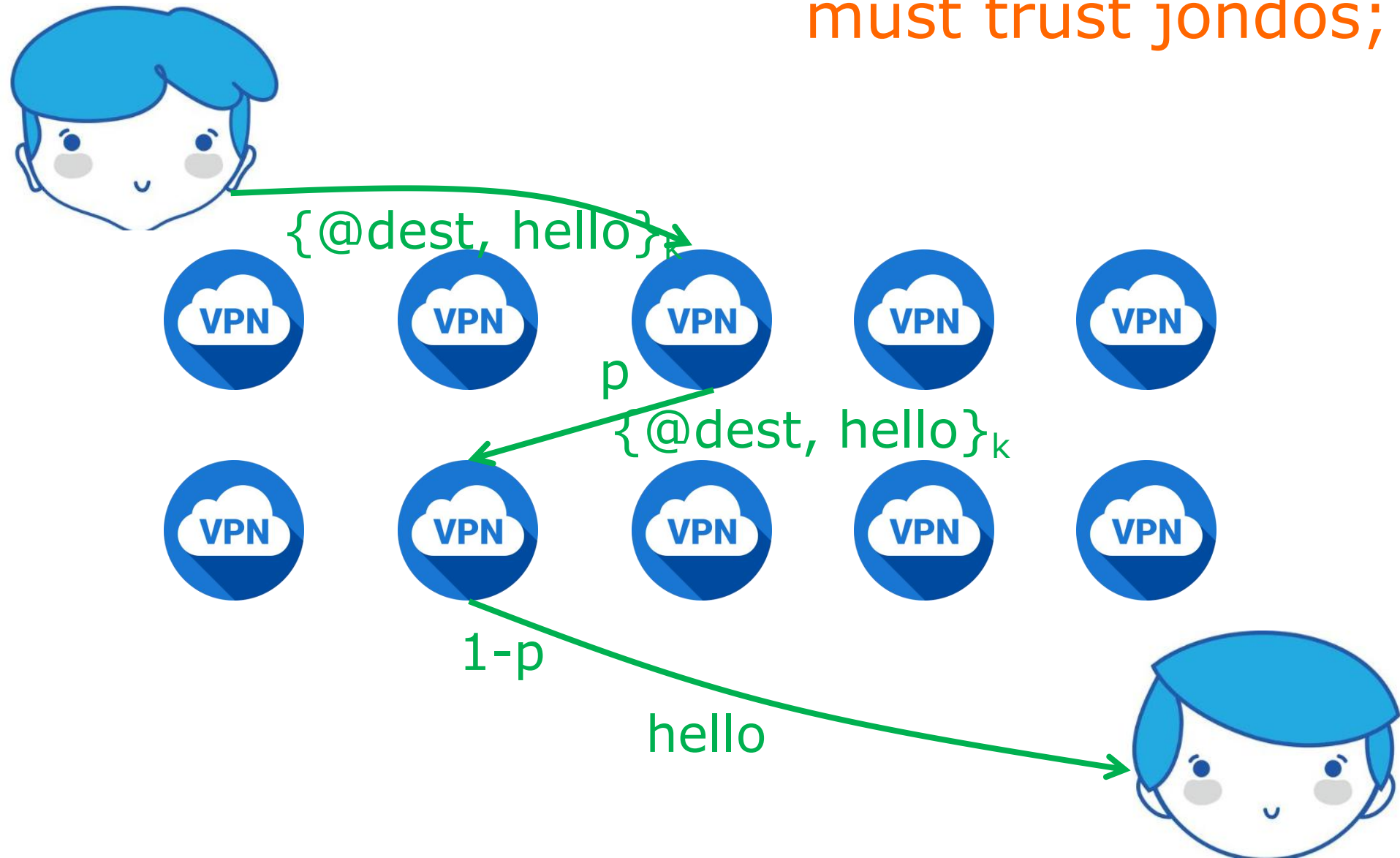


Crowds Algorithm



Crowds Algorithm

must trust jondos;



Crowds Algorithm

must trust jondos;
if any message is intercepted,
receiver is trivially exposed.



how to evade



untrusted proxies?



how to evade



untrusted proxies?



proxy++

to evade



untrusted proxies?





hard for an attacker
to simultaneously control
too many proxies



proxy++

to evade

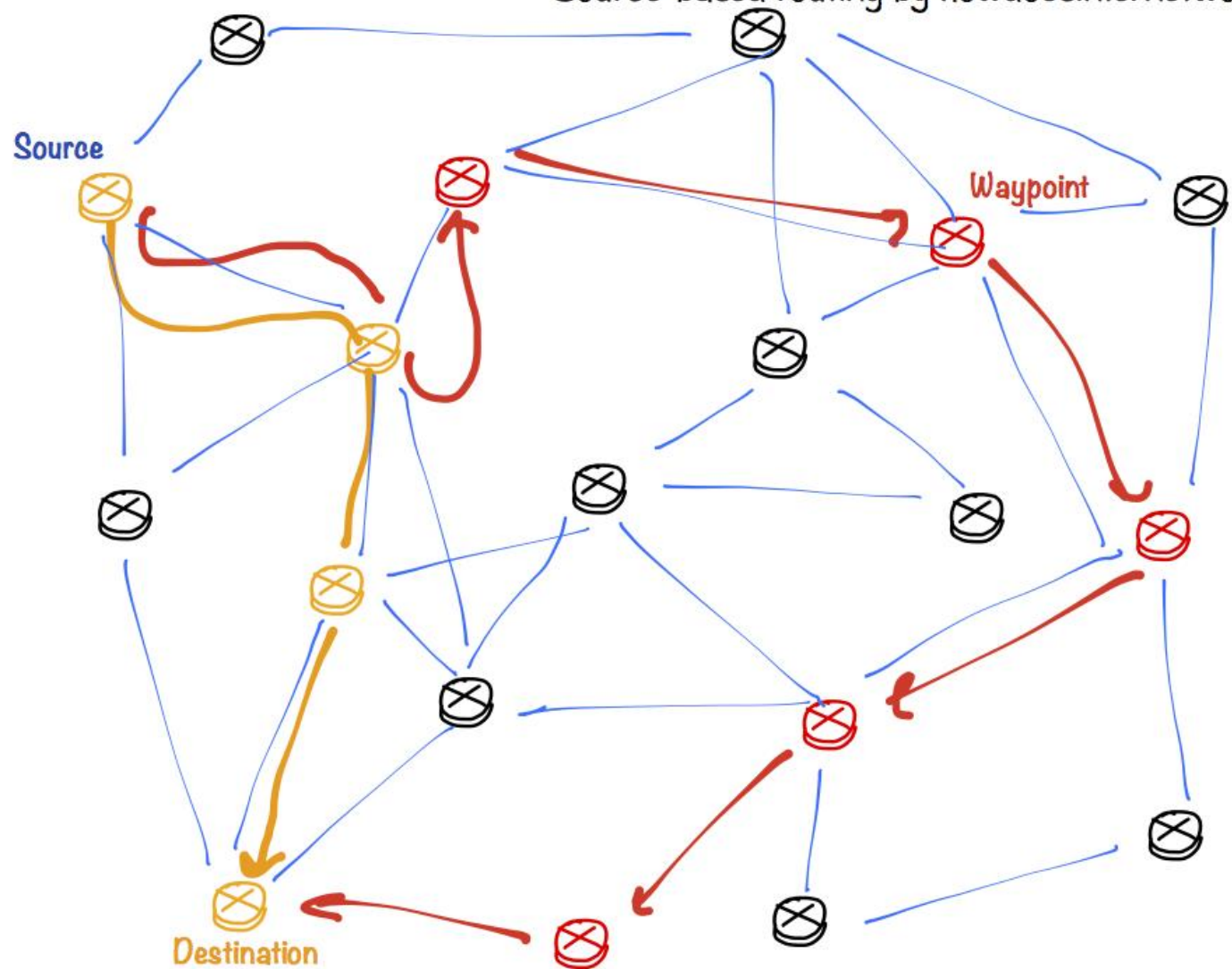


untrusted proxies?



source routing

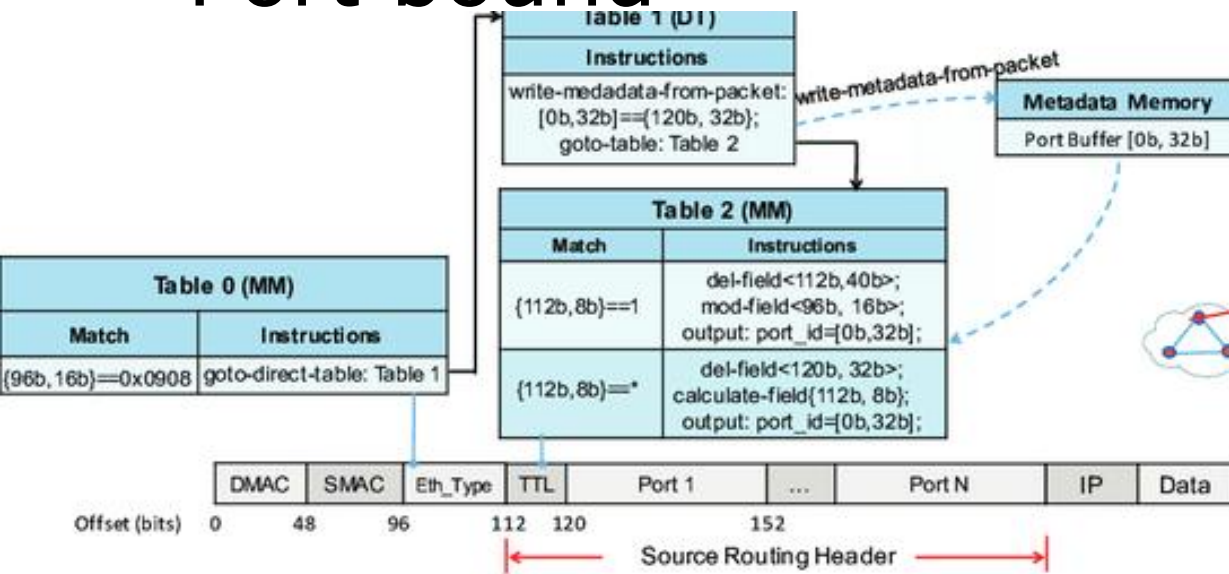
specify on-path routers by source



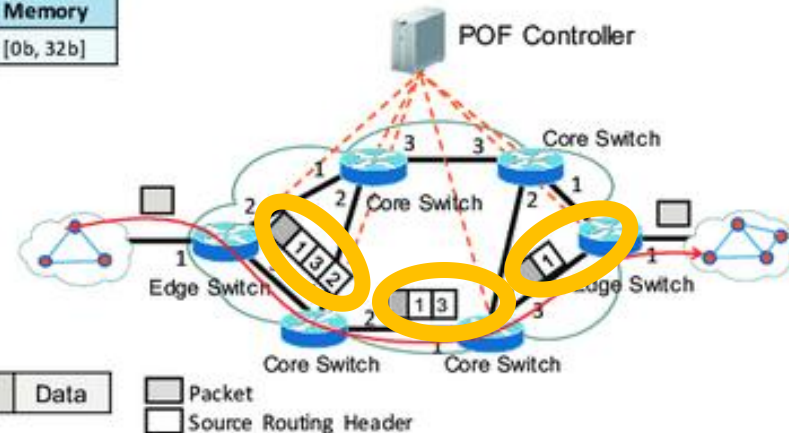
- All possible paths
- Best path
- Source routing forced path

POF-based Source Routing

- POF: Protocol Oblivious Forwarding
- Port bound



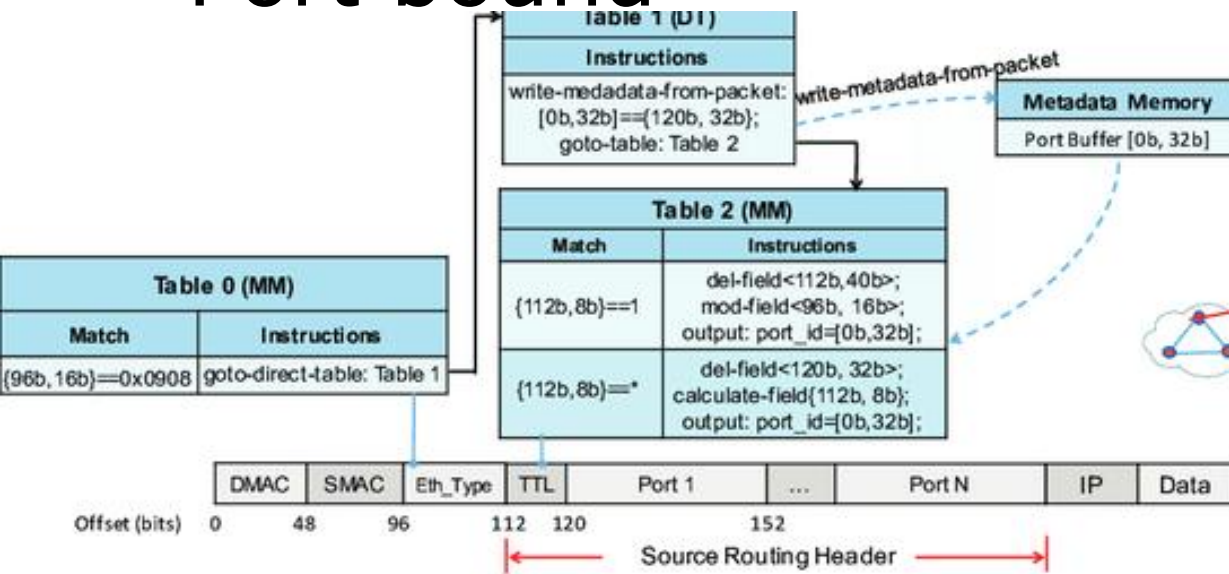
forwarding
table



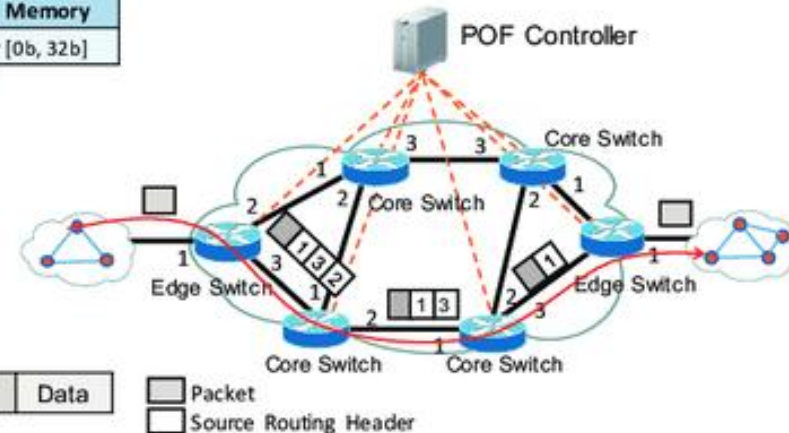
forwarding
process

POF-based Source Routing

- POF: Protocol Oblivious Forwarding
- Port bound



forwarding
table

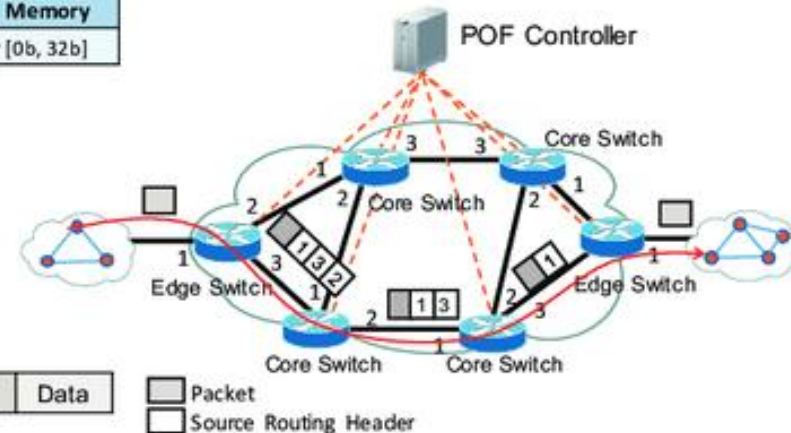
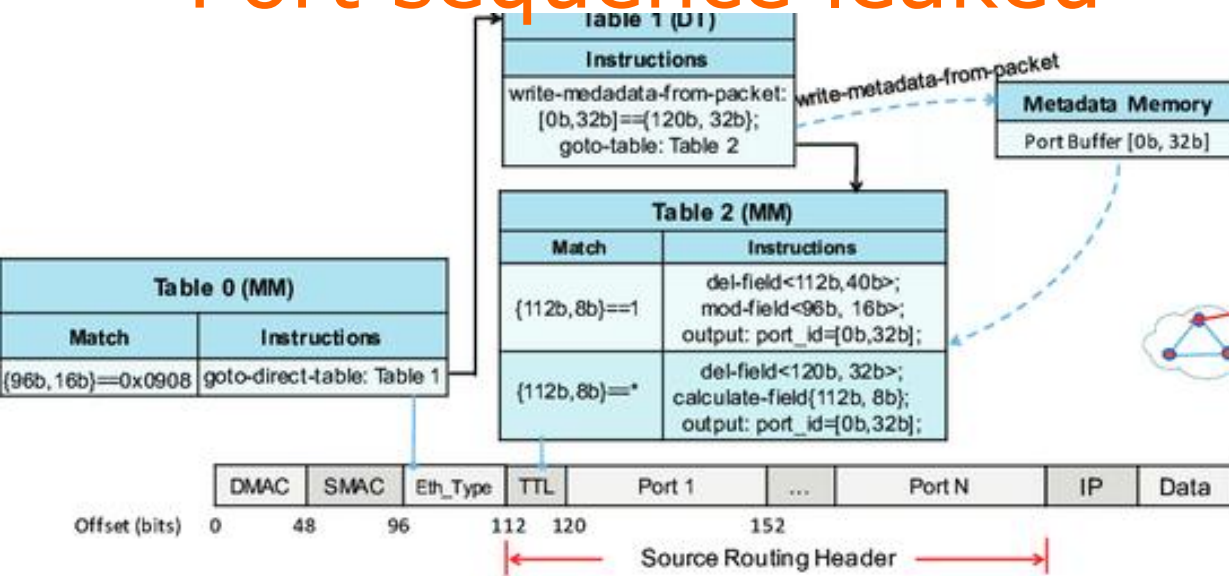


forwarding
process

anonymity protected?

POF-based Source Routing

- POF: Protocol Oblivious Forwarding
- Port sequence leaked



forwarding
table

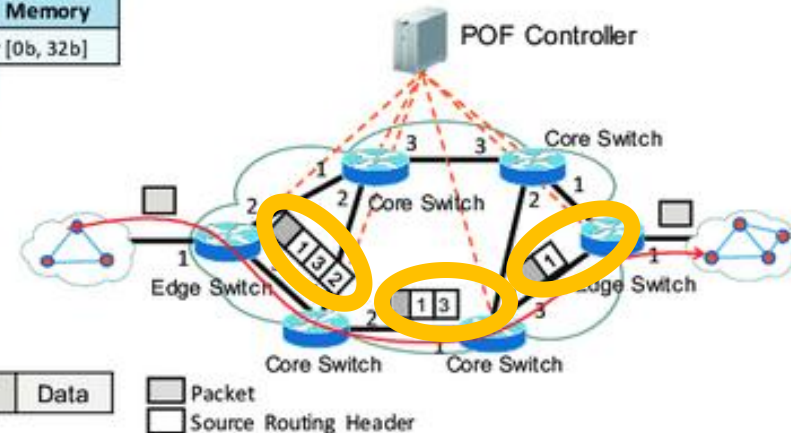
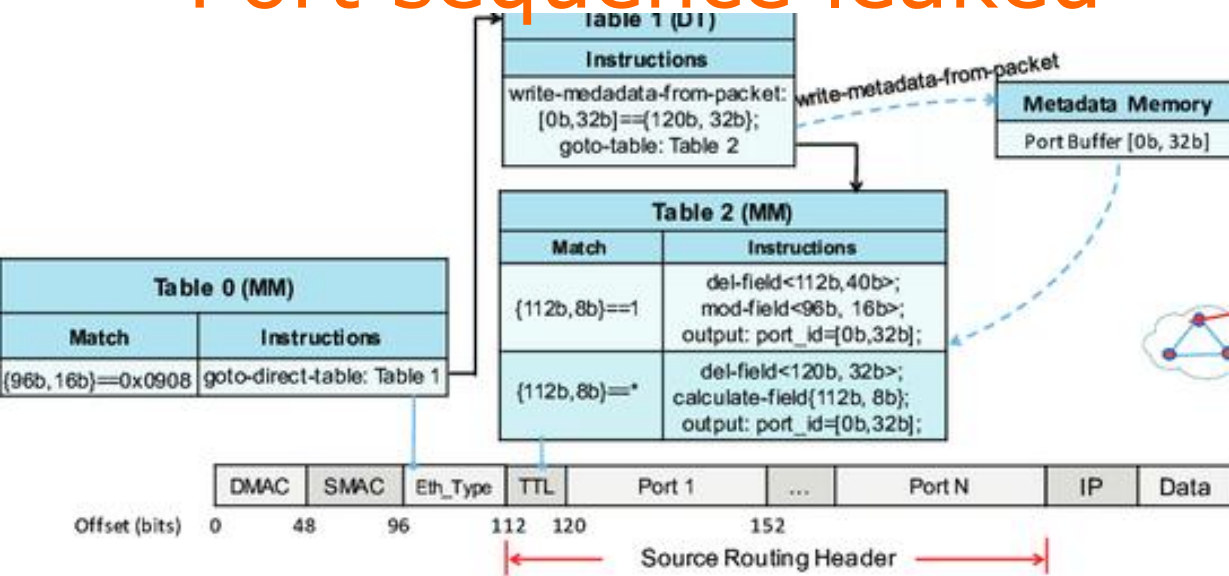
forwarding
process

anonymity protected? nah!

source routing
how to anonymize?

POF-based Source Routing

- POF: Protocol Oblivious Forwarding
- Port sequence leaked



forwarding
table

forwarding
process

should hide ports from non-neighbors

onion routing
source-routing based
anonymous overlay communication

onion routing

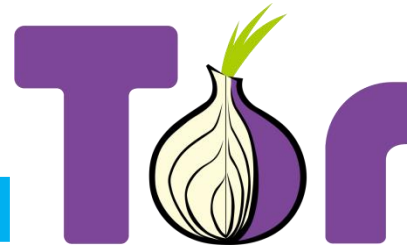
source-routing based

anonymous overlay communication

onion routing

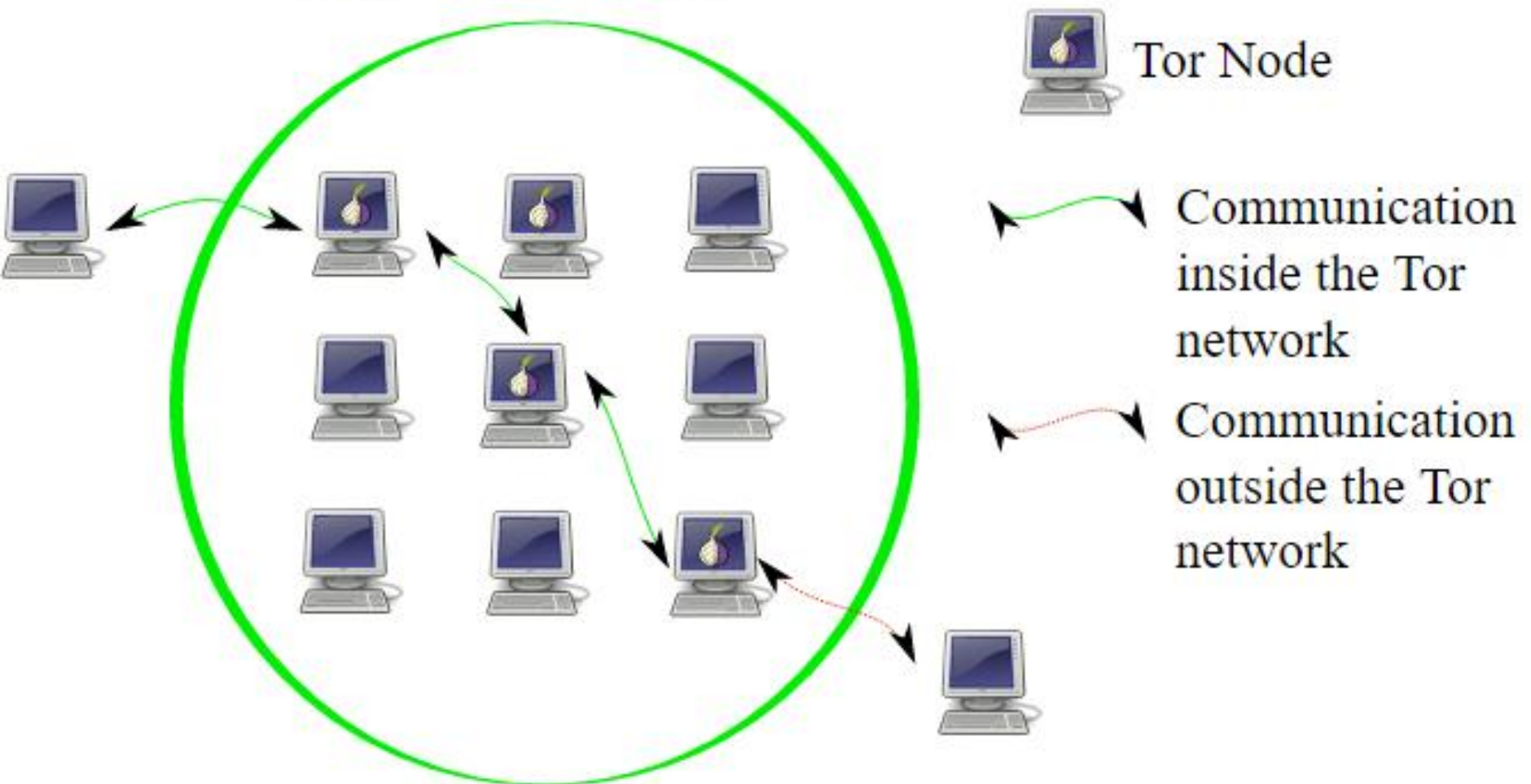
source-routing based

anonymous overlay communication

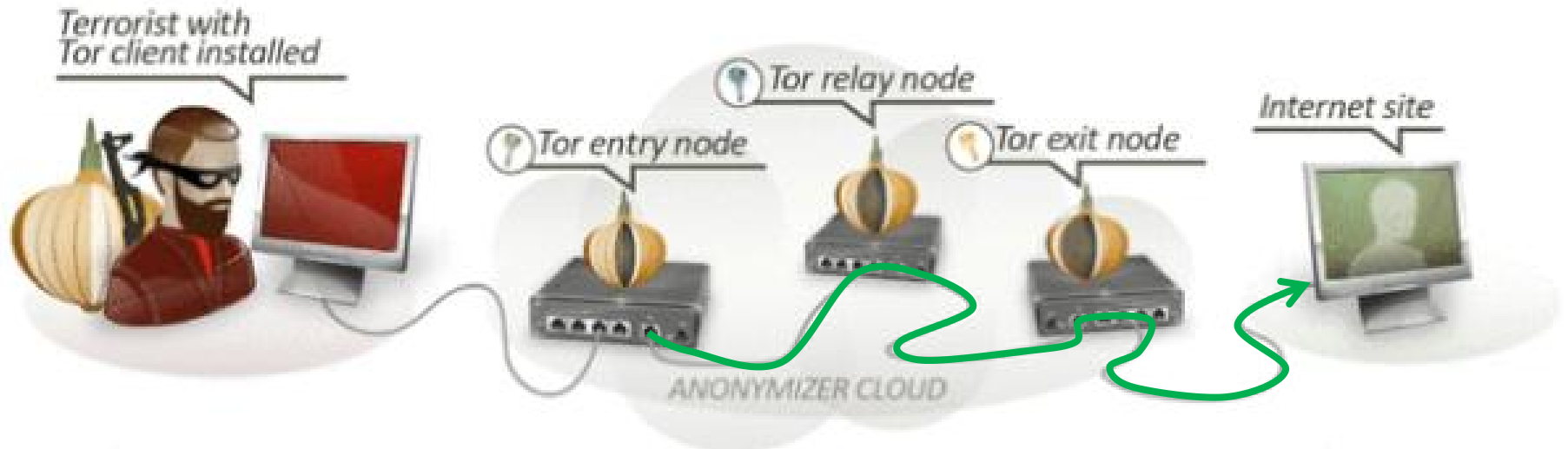


Onion Routing

The Internet

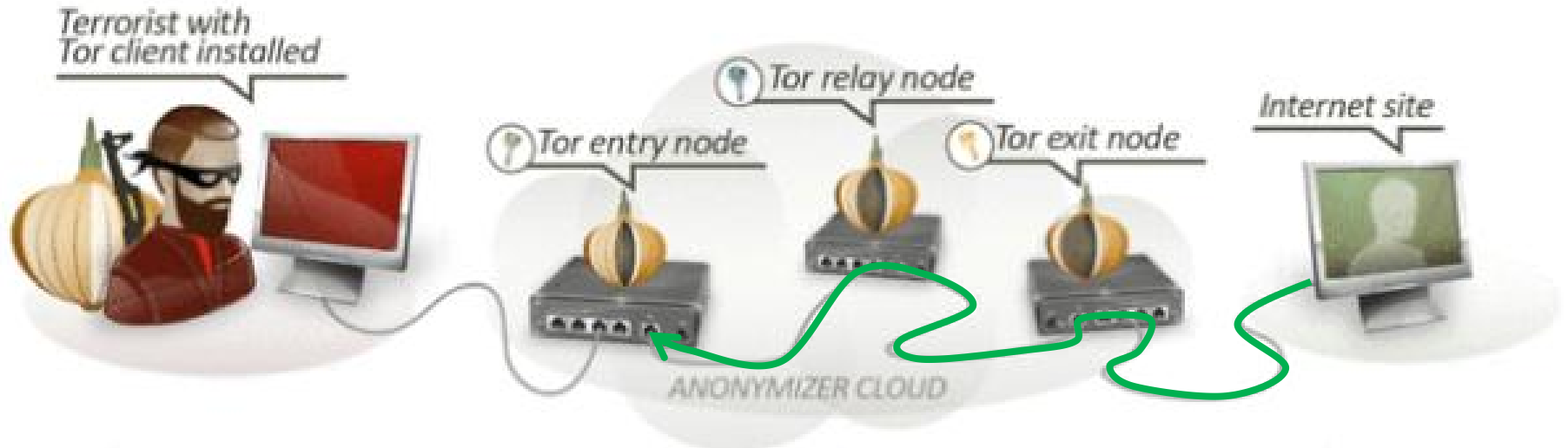


Onion Routing



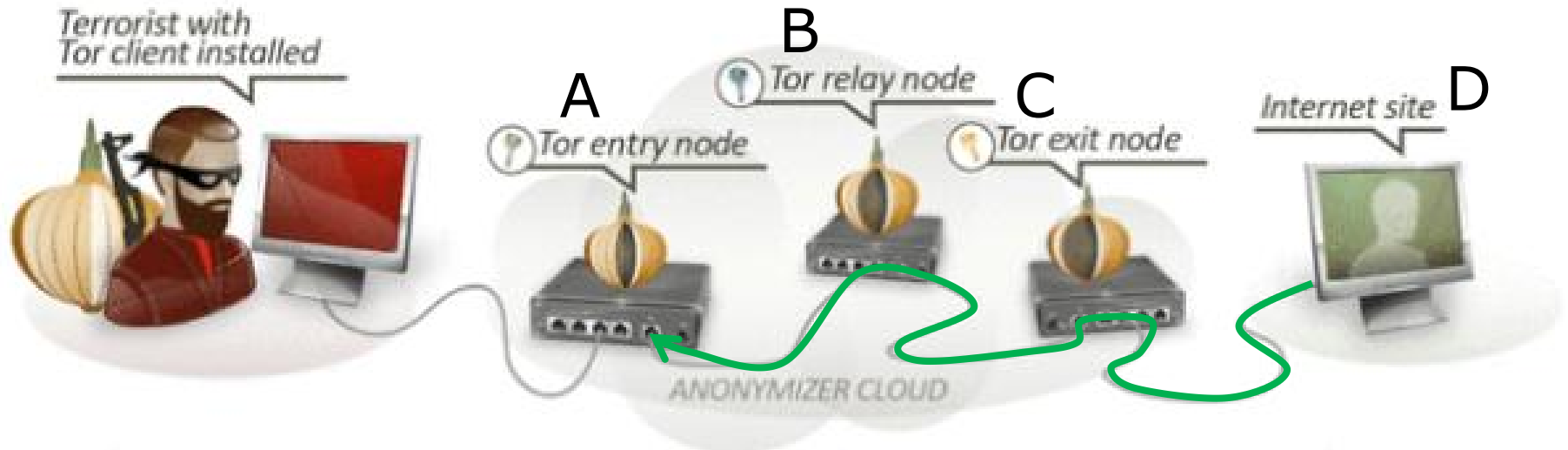
- Connect to Tor entry
- Randomly select a series of Tors
- Relay messages across them
- Tor exit relays messages to destination

Onion Routing



- Reply traffic from destination traverses the reverse path
- Maintains a bidirectional persistent multi-hop path between source and destination

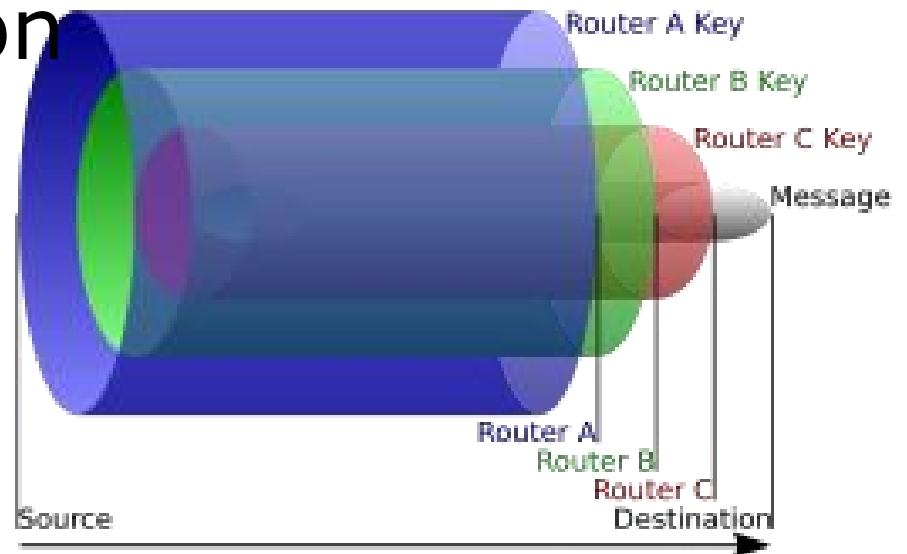
Onion Routing



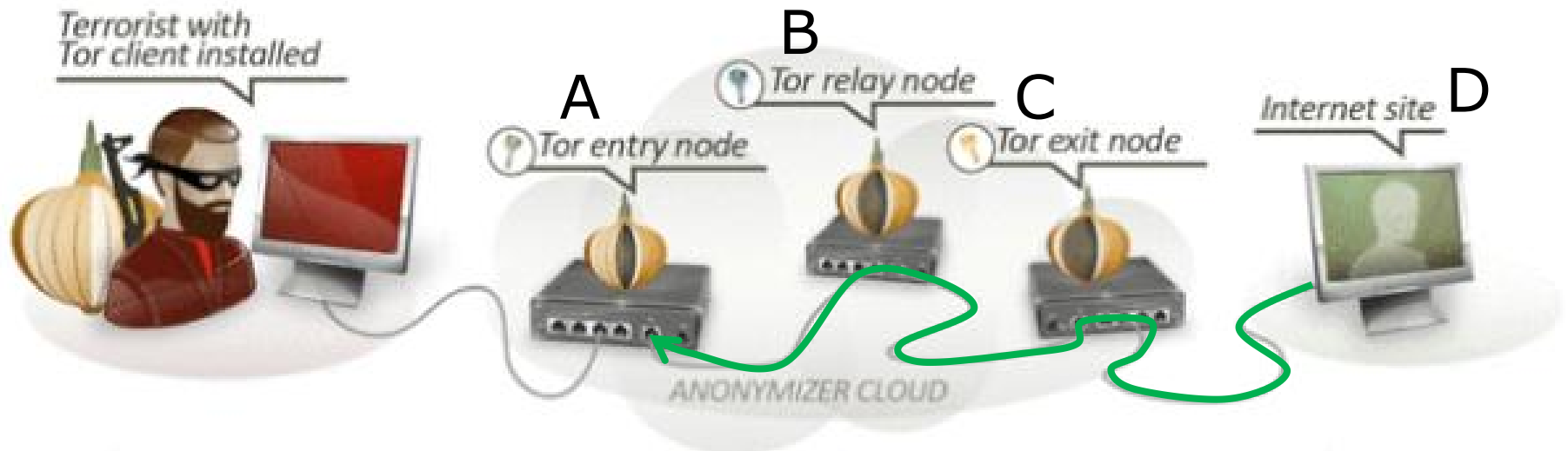
- Layered Encryption

$\{\{\{\{msg\}_D\}_C\}_B\}_A$

sufficient?



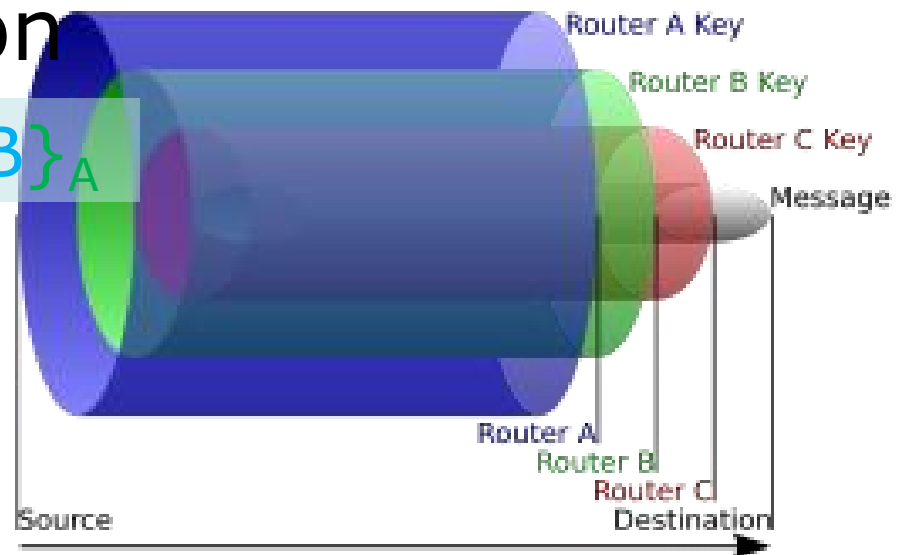
Onion Routing



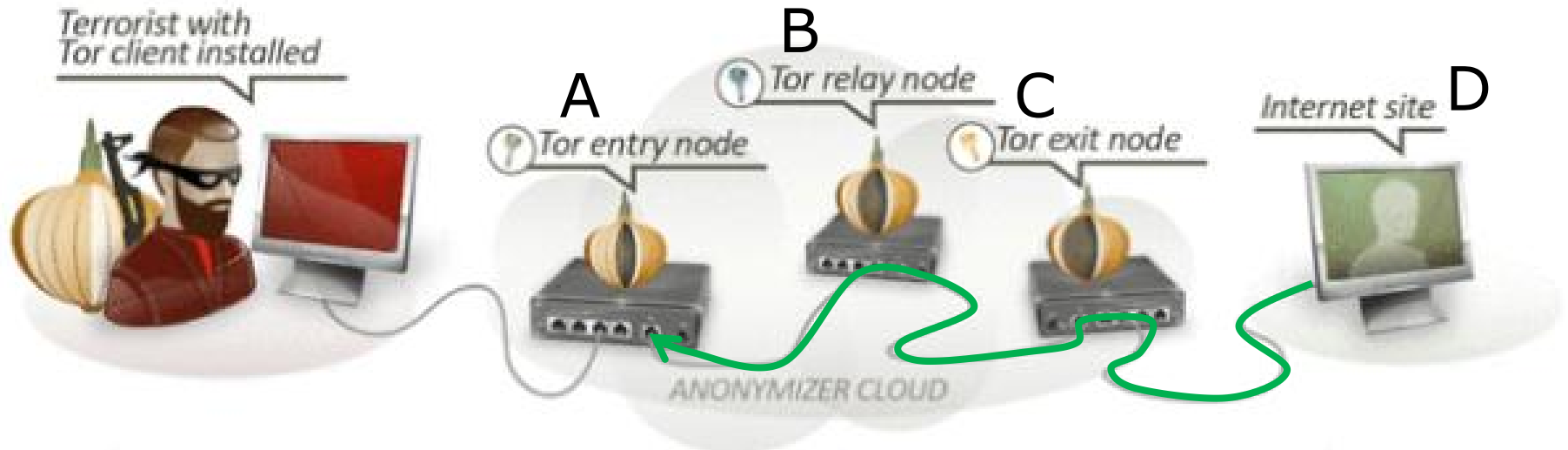
- Layered Encryption

$\{\{\{\{msg\}_D, D\}_C, C\}_B, B\}_A$

sufficient



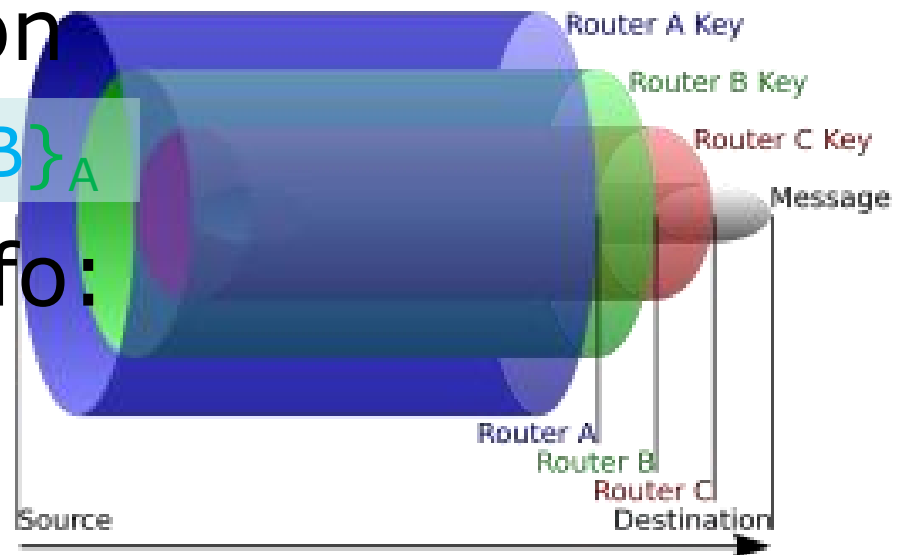
Onion Routing



- Layered Encryption

$\{\{\{\{msg\}_D, D\}_C, C\}_B, B\}_A$

- Leaked routing info:
neighborship only

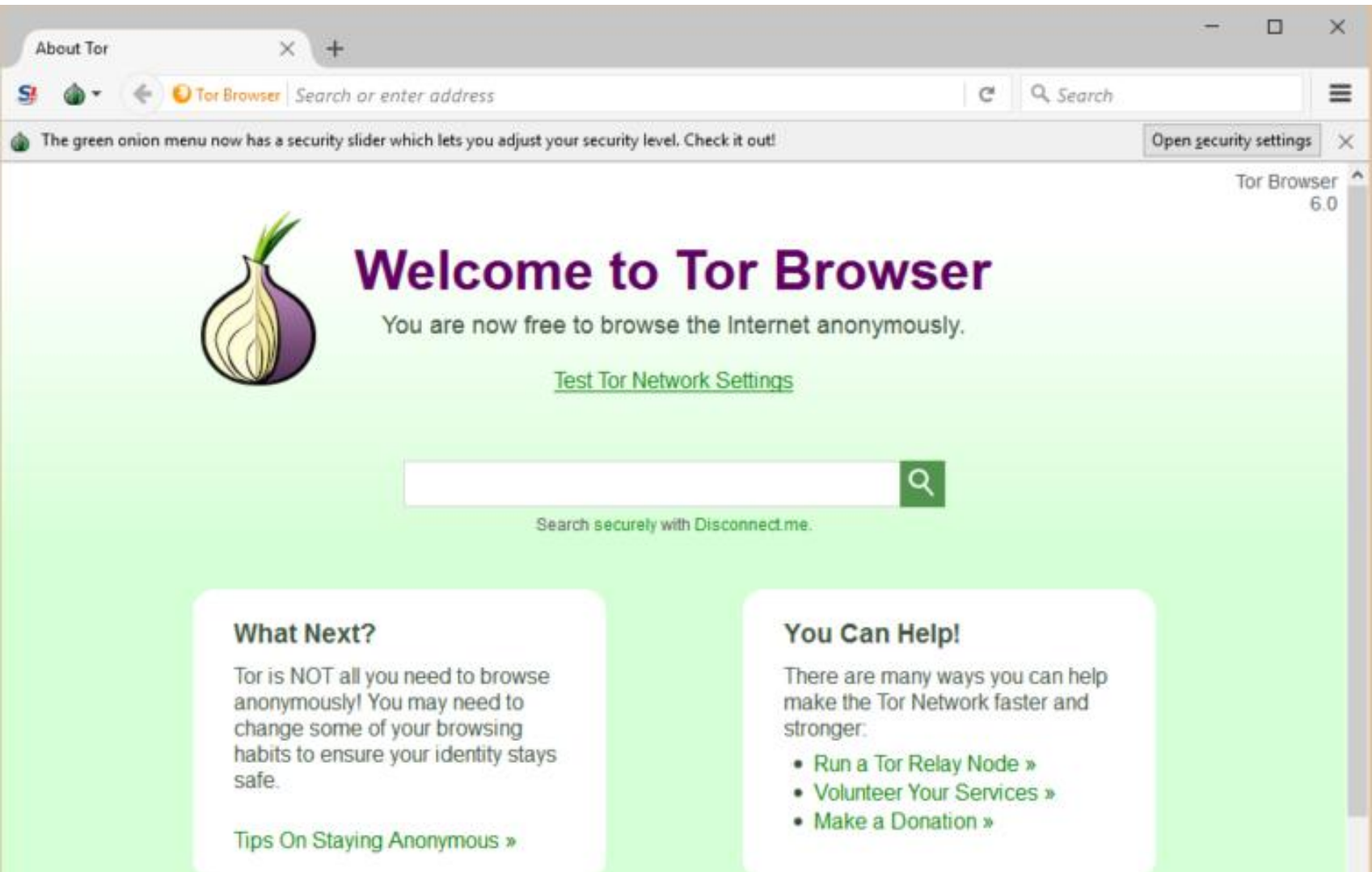


onion routing
applications?

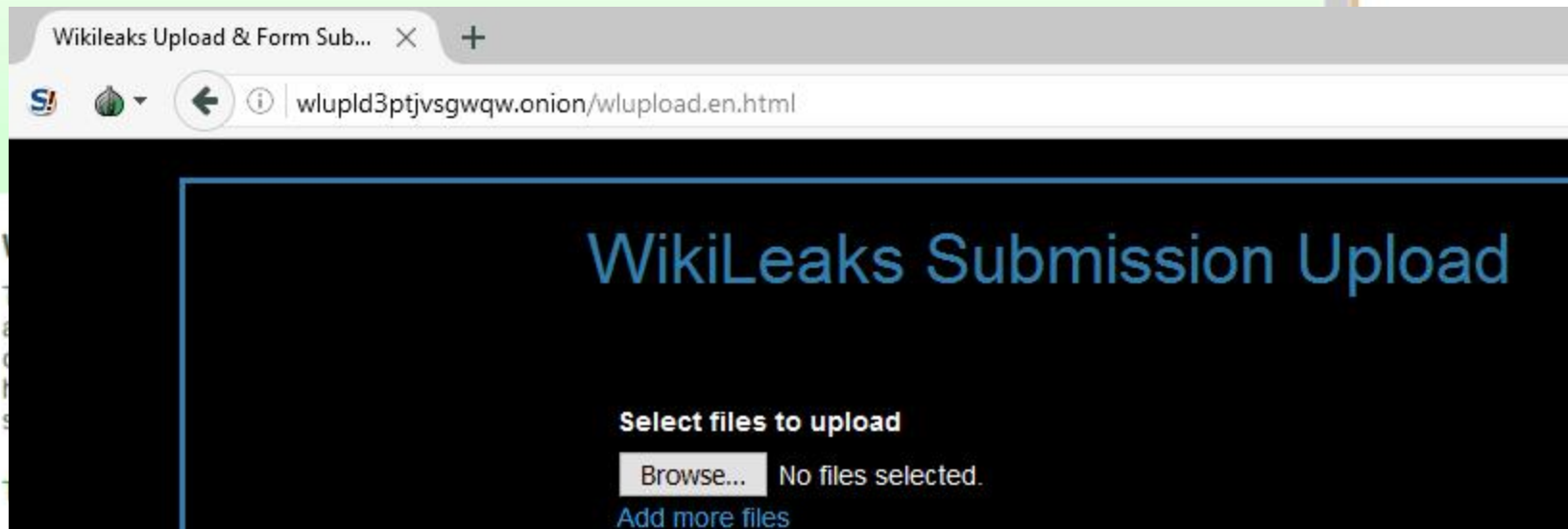
Darknet

- Portions of the Internet purposefully not open to public view or hidden networks whose architecture is superimposed on that of Internet.
- Install Tor
- Access `darknet.onion` through it

Darknet



Darknet



how to de-anonymization?

in it to win it!

Tor Traffic Correlation

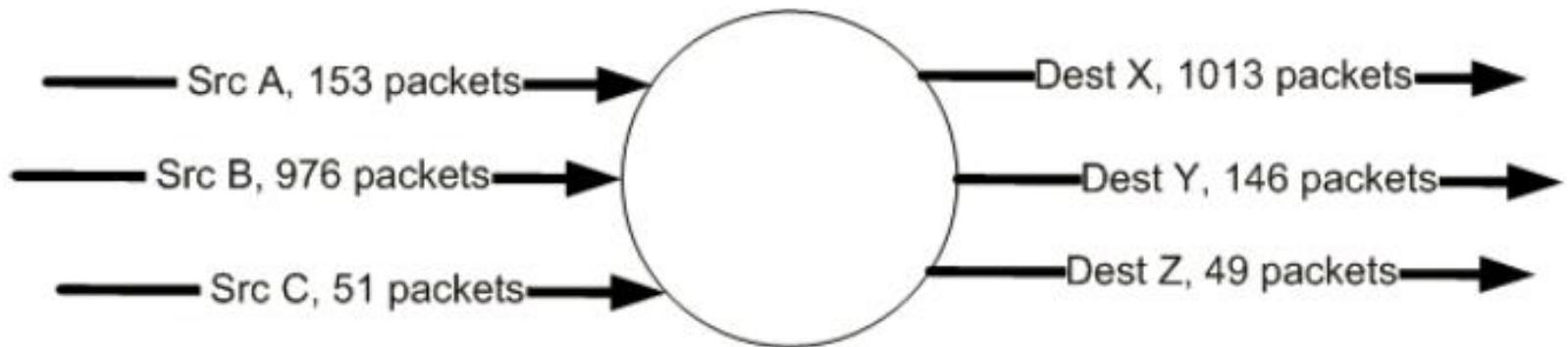
- Passive monitoring
- Active attraction:
deploy a Tor router;
attract Tor traffic;
perform traffic analysis and correlation;

Path Selection Attack

- Tor path selection algorithm:
weight nodes by selfreported bandwidth
select each node using weighted
probability distribution;
- Attack:
malicious relay reports very high bw to
increase selection probability;
if it controls the first hop, de- sender;
if it controls the last hop, de- receiver;

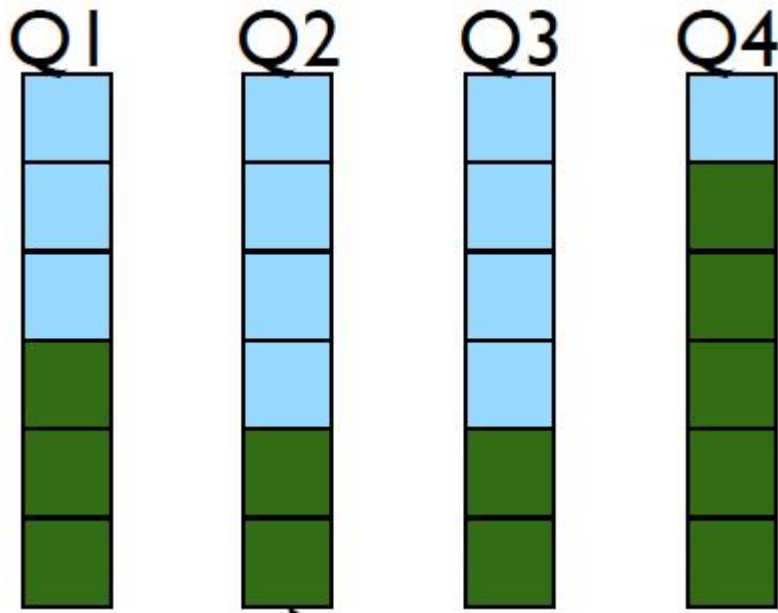
Counting Attack

- Correlate incoming and outgoing flows by counting the number of packets



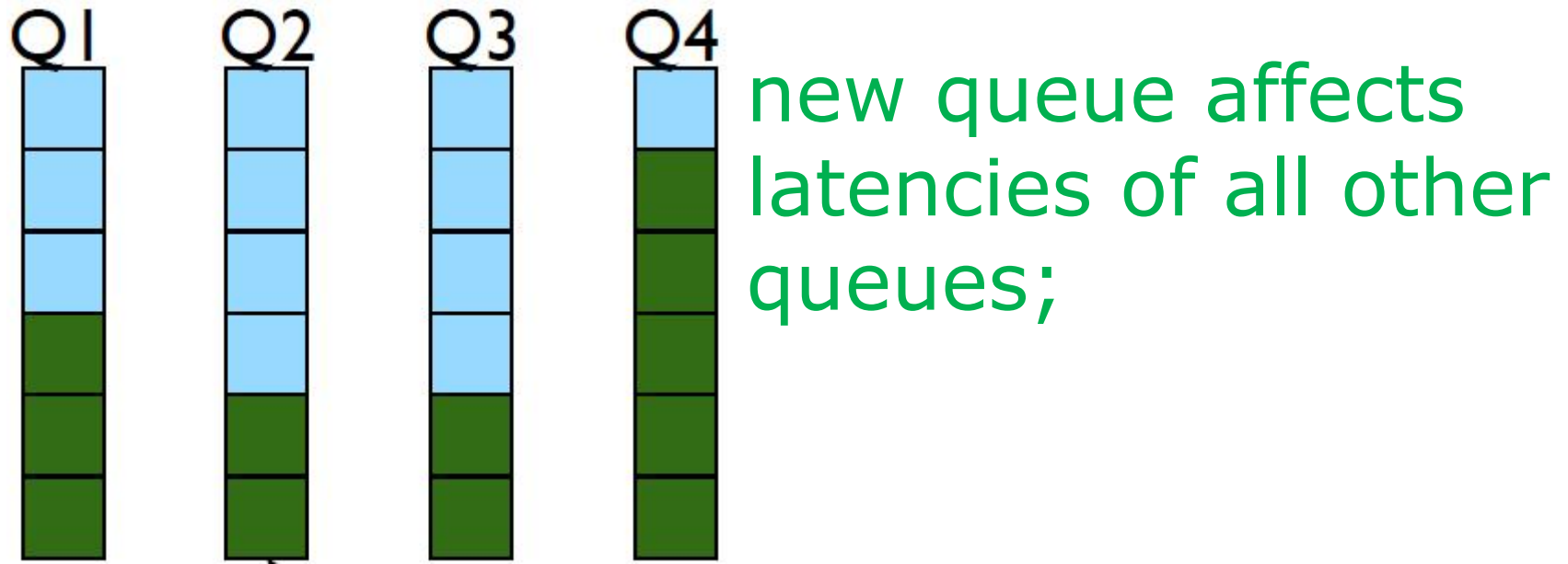
Low Latency Attack

- Tor router assigns each anonymous circuit its own queue
- Dequeue one packet from each queue in round-robin fashion

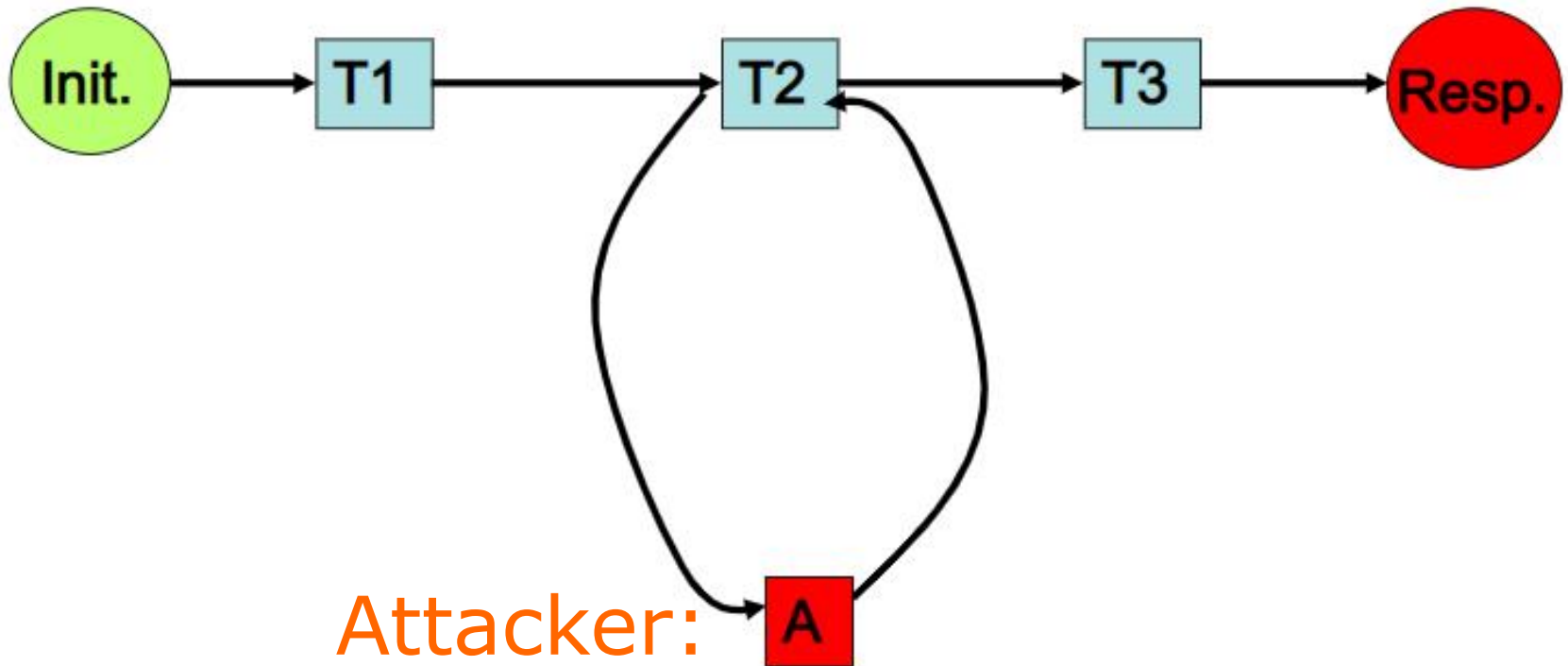


Low Latency Attack

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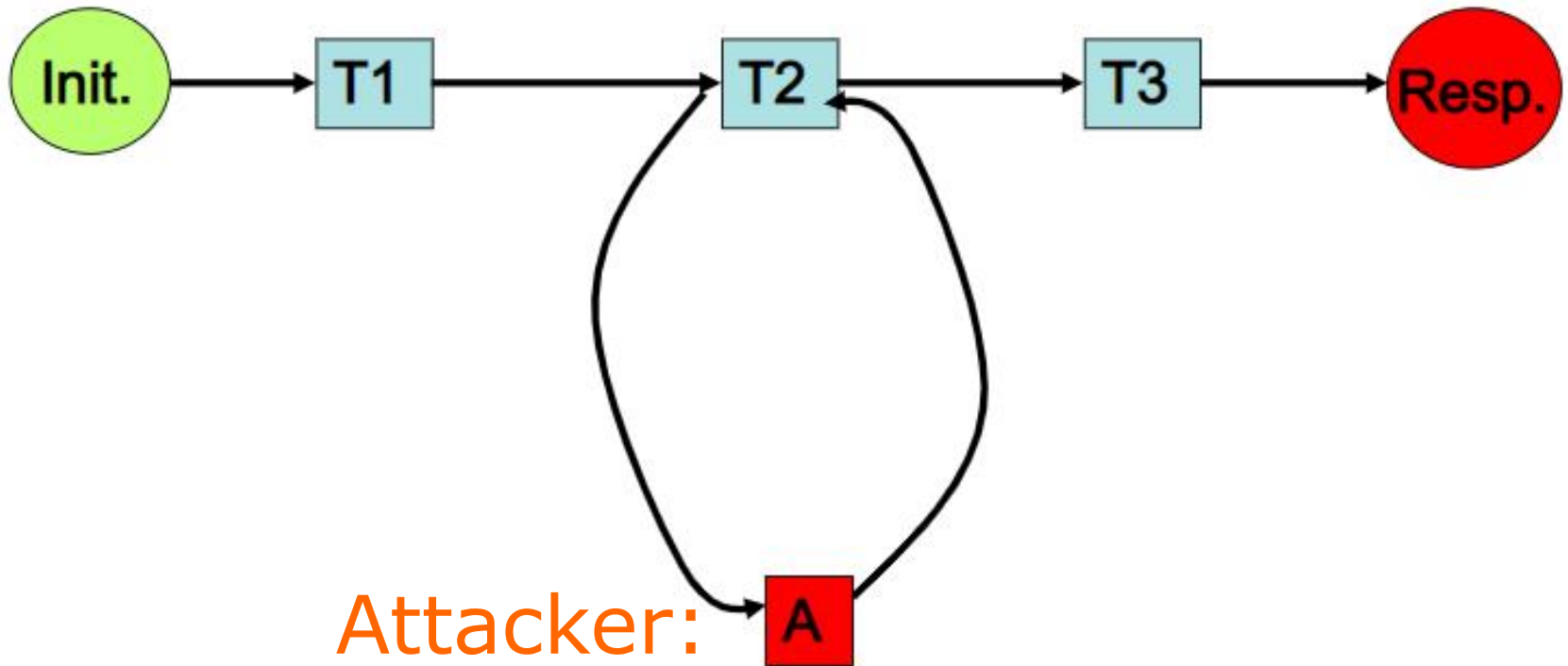


Low Latency Attack



Attacker: **A**
intends to infer Init's activity

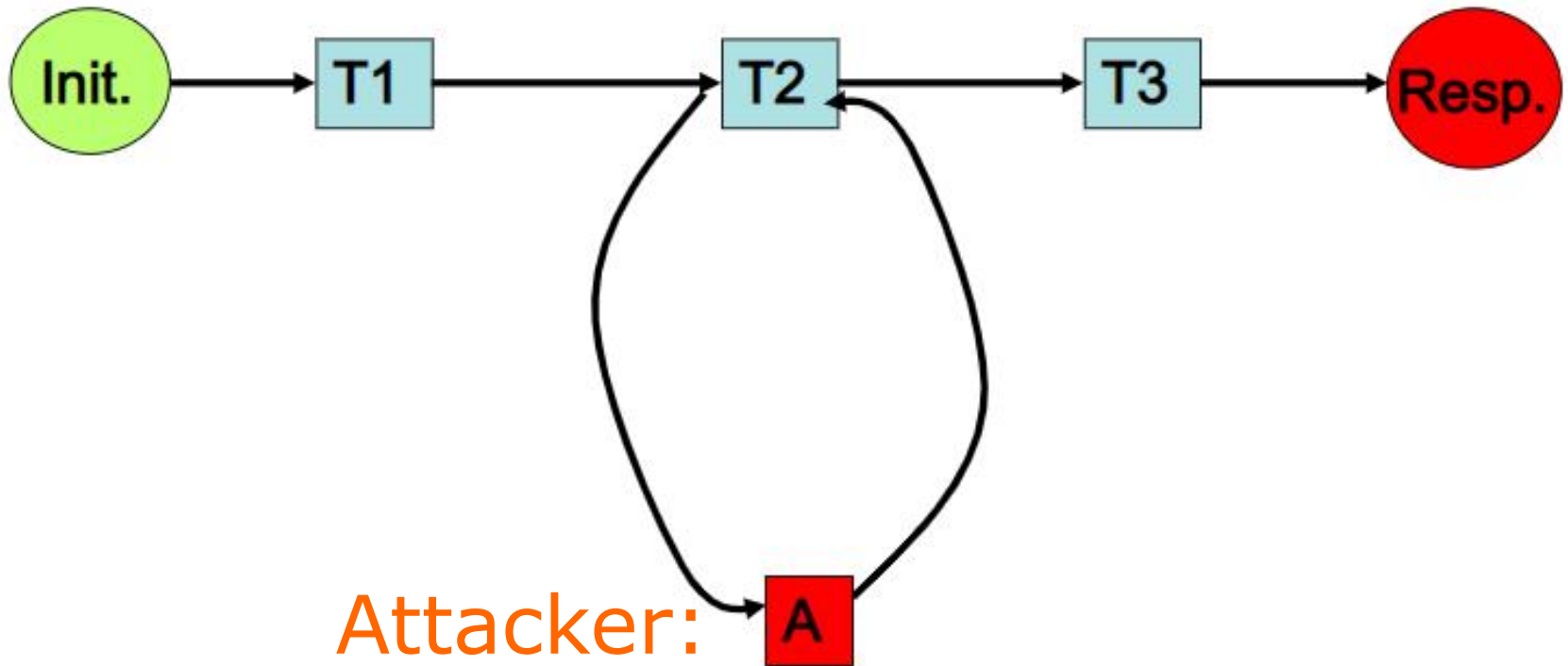
Low Latency Attack



Attacker:
intends to infer Init's activity

Assumption:
only Init and A occupy T2

Low Latency Attack



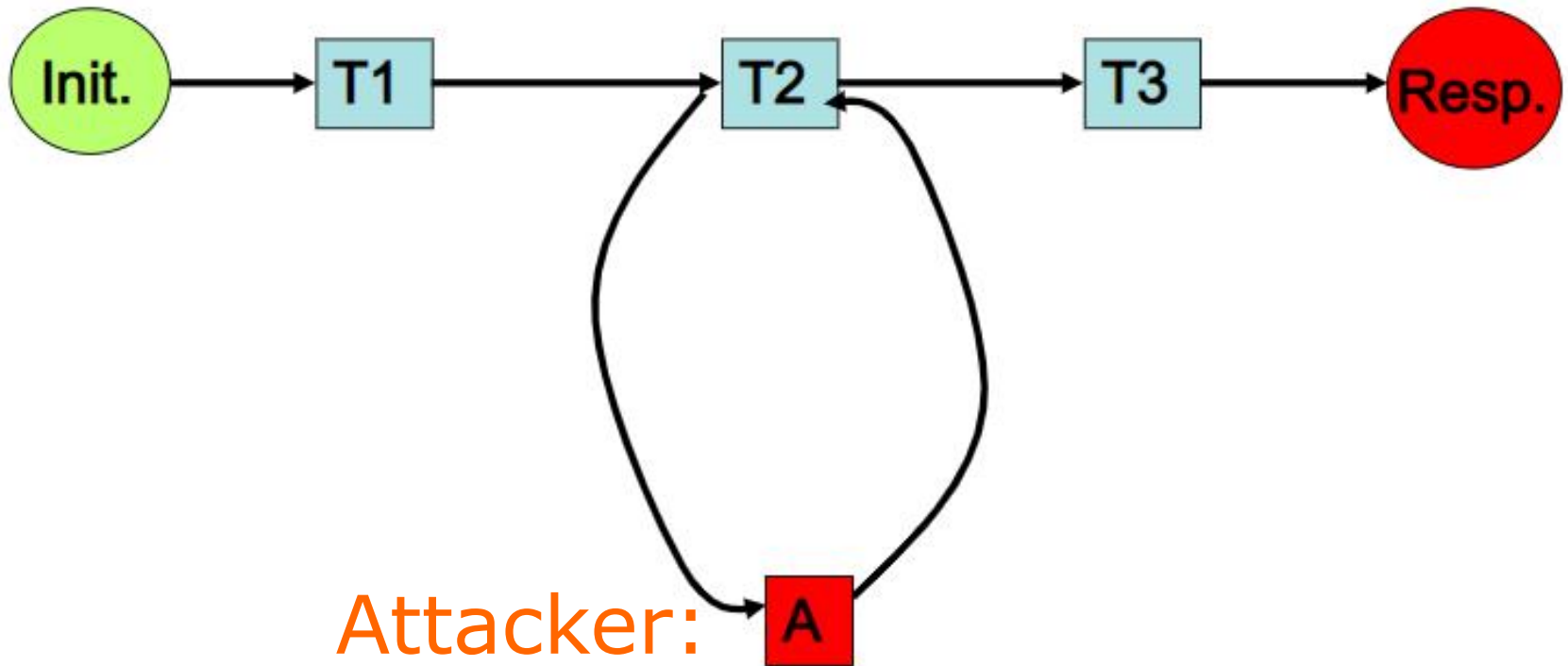
Attacker:

loop:

send packets to T2;
measure latency;

Low Latency Attack

larger latency indicate Init's traffic!



loop:

send packets to T2;
measure latency;

Cross Site Attack

Crawling:

- Deploy Tor routers
- Access darknet
- Crawl transaction information
- Extract Bitcoin accounts of interest

Correlation:

- Search the accounts on public websites



Readings

- Anonymous Communication
by Nick Mathewson
- Tor: The Second-Generation Onion Router (2012 DRAFT)
by Dingledine *et al.*

Disclaimer:

The content aims only for helping students understand principles of anonymous communication. It should not be used for abusive Internet activities.

Thank You