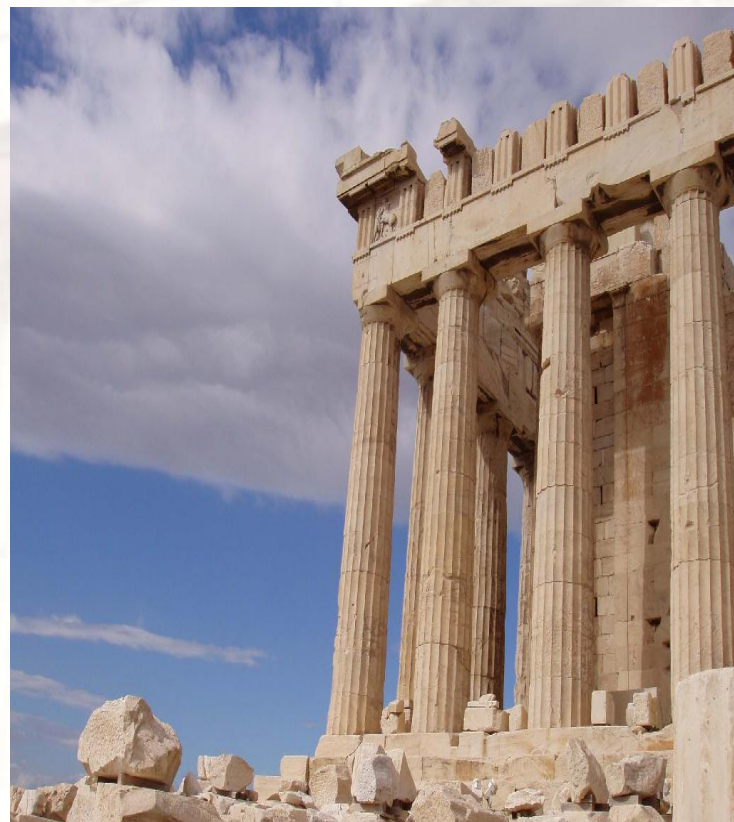


Ch3-- ILP & its exploration

Ch3-1

- Instruction Level Parallelism
- Basic Compiler Techniques for Exposing ILP
- Dynamic scheduling-- Scoreboard

3.1, 3.2, App C7





What is Instruction-Level Parallelism ?

❑ Instruction-level parallelism

➤ The potential overlap among instructions

❑ Basic Block ILP is quite small

- **Basic Block**: a straight-line code sequence with no branches in except to the entry and no branches out except at the exit
- average dynamic branch frequency 15% to 25%
=> 4 to 7 instructions execute between a pair of branches
- Plus instructions in BB likely to depend on each other



Recall from Pipelining Review

- ❑ When exploiting instruction-level parallelism, goal is to **maximize CPI**
- ❑ Pipeline $CPI = \text{Ideal pipeline CPI} + \text{Structural Stalls} + \text{Data Hazard Stalls} + \text{Control Stalls}$
 - Ideal pipeline CPI: measure of the maximum performance attainable by the implementation
 - Structural hazards: HW cannot support this combination of instructions
 - Data hazards: Instruction depends on result of prior instruction still in the pipeline
 - Control hazards: Caused by delay between the fetching of instructions and decisions about changes in control flow (branches and jumps)



How to exploit ILP?

□ there are two main approaches:

➤ Hardware-based dynamic approaches

- Used in server and desktop processors
- Not used as extensively in PMP processors

➤ Compiler-based static approaches

- Not as successful outside of scientific applications



Chapter 3

- ❑ ILP: Concepts and Challenge
- ❑ Basic compiler Techniques for exposing ILP
- ❑ Overcoming Data Hazards with Dynamic Scheduling
- ❑ Reducing Branch Costs with Dynamic Branch Prediction
- ❑ Hardware-based Speculation
- ❑ Exploiting ILP with Multiple Issue & Static Scheduling
- ❑ Exploiting ILP with Dynamic scheduling, Multiple issue, & speculation
- ❑ Advanced Techniques for instruction Delivery and speculation
- ❑ Multithreading: exploiting TLP improve uniprocessor throughput



Ideas to Reduce Stalls

	Technique	Reduces	Section
ChC	Forwarding and bypassing	Potential data hazard stalls	C.2,C.3
	Simple branch scheduling and prediction	Control hazard stalls	C.3
Ch3	Basic compiler pipeline scheduling	Data hazard stalls	C.2, 3.2
	Loop unrolling	Control hazard stalls	3.2
	Dynamic branch prediction	Control stalls	3.3
	Dynamic scheduling (scoreboard)	Data hazard stalls	C.7
	Dynamic scheduling (Tomasulo)	DH stalls from Anti and output dependences	3.4, 3.5
	Hardware-based speculation	Control stalls	3.6
	Issuing multi-instructions per cycle	Ideal CPI	3.7
ChH	Dynamic scheduling + multiple issue +Speculation	Data and control stalls	3.8
	Multi-threading	Data parallelism	3.11
	Compiler dependence analysis, software pipelining, trace schedule	Ideal CPI and data hazard stalls	H.2, H.3
	Hardware support for compiler speculation	Ideal CPI and data hazard stalls, branch hazard stalls	H.4, H.5



Instruction-Level Parallelism (ILP)

- ❑ To obtain substantial performance enhancements, we must exploit ILP across multiple basic blocks
- ❑ Simplest: loop-level parallelism to exploit parallelism among iterations of a loop
 - Vector & GPU is one way
 - If not vector, then either dynamic via branch prediction or static via loop unrolling by compiler



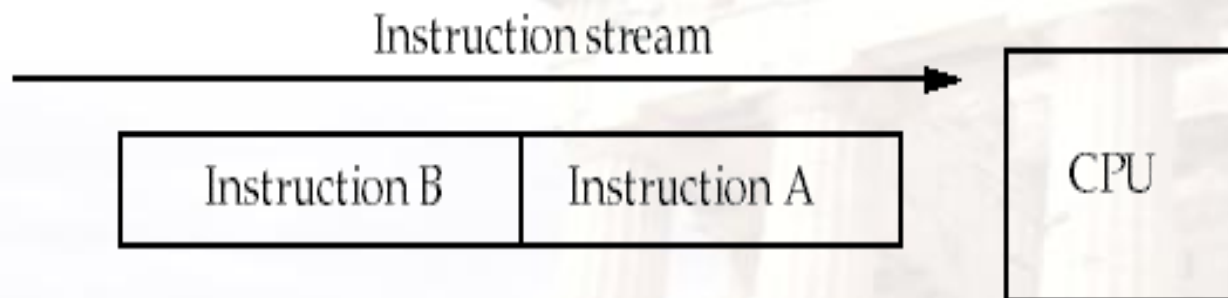
Data Dependence & hazard

- ❑ **Dependencies** are a **property of programs**, presence of dependence indicates potential for a hazard,
- ❑ Pipeline organization determines if dependence is detected and if it causes a stall, **actual hazard and length of any stall** is a **property of the pipeline**
- ❑ Data dependence conveys:
 - Possibility of a hazard (**register & memory location**)
 - Order in which results must be calculated
 - Upper bound on exploitable instruction level parallelism
- ❑ Dependencies that flow through memory locations are difficult to detect



Recall: Types of data hazards

- ❑ Consider two instructions, A and B. A occurs before B.



- ❑ RAW(Read after write) true dependence
 - Instruction A writes Rx, instruction B reads Rx
- ❑ WAW(Write after write) output dependence
 - Instruction A writes Rx, instruction B writes Rx
- ❑ WAR(Write after read) anti-dependence
 - Instruction A reads Rx, instruction B writes Rx
- ❑ Hazards are named according to the ordering that **MUST** be preserved by the pipeline



True Data Dependence and Hazards

□ True Data Dependence:

- Instr_J is **data dependent** on Instr_I
Instr_J tries to read operand before Instr_I writes it

```
    I: add r1, r2, r3  
    J: sub r4, r1, r3
```

- or Instr_J is data dependent on Instr_K which is dependent on Instr_I

□ Caused by a "**True Dependence**" (compiler term)

□ If **true** dependence caused a hazard in the pipeline, called a **Read After Write (RAW) hazard**



Name Dependence 1: Anti-dependence

- ❑ **Name dependence:** when 2 instructions use same register or memory location, called a **name**, but no flow of data between the instructions associated with that name;
- ❑ Instr_J writes operand before Instr_I reads it

```
    I: sub r4, r1, r3  
    J: add r1, r2, r3  
    K: mul r6, r1, r7
```

called an "**anti-dependence**" by compiler writers.
This results from reuse of the name "**r1**"

- ❑ If anti-dependence caused a hazard in the pipeline, called a **Write After Read (WAR) hazard**



Name Dependence 2: Output dependence

□ Instr_J writes operand before Instr_I writes it.

```
      I:  sub  r1, r4, r3  
    ↙   J:  add  r1, r2, r3  
      K:  mul  r6, r1, r7
```

□ Called an “output dependence” by compiler writers

This also results from the reuse of name “r1”

□ If anti-dependence caused a hazard in the pipeline, called a Write After Write (WAW) hazard



Name Dependence

❑ Two instructions **use the same name** but no flow of information

➤ Not a true data dependence, but is a problem when reordering instructions

➤ **Antidependence**: instruction j writes a register or memory location that instruction i reads

○ Initial ordering (i before j) must be preserved

• **ADD** x5, **x8**, x9
• **MUL** **x8**, x6, x10

➤ **Output dependence**: instruction i and instruction j write the same register or memory location

○ Ordering must be preserved

• **ADD** **x5**, x8, x10
• **MUL** **x5**, x6, x9



ILP and Data Hazards

- ❑ HW/SW must preserve **program order**:
order instructions would execute in if executed sequentially 1 at a time as determined by original source program
- ❑ HW/SW goal: exploit parallelism by preserving program order **only where it affects the outcome of the program**
- ❑ Instructions involved in a name dependence can execute simultaneously **if name used in instructions is changed** so instructions do not conflict
 - **Register renaming** resolves name dependence for registers
 - Either by **compiler** or by **HW**



Control Dependencies

- ❑ Every instruction is **control dependent** on some set of branches, and, in general, these control dependencies must be preserved to preserve program order

```
if p1 {  
    S1;  
};  
if p2 {  
    S2;  
}
```

- ❑ S1 is control dependent on p1, and S2 is control dependent on p2 but not on p1.



Control Dependence Ignored

- ❑ Control dependence need **not** be preserved
 - willing to execute instructions that should not have been executed, thereby violating the control dependences, **if** can do so without affecting correctness of the program
- ❑ Instead, 2 properties **critical to program correctness** are **exception behavior** and **data flow**



Examples

- Example 1:

- add x1,x2,x3
- **beq x4,x0,L**
- sub x1,x1,x6
- **L:** ...
- or x7,x1,x8

❑ or instruction dependent on add and sub

- Example 2:

- add x1,x2,x3
- beq x12,x0,skip
- **sub x4,x5,x6**
- add x5,x4,x9
- **skip:**
- or x7,x8,x9

❑ Assume x4 isn't used after skip

➤ Possible to move sub before the branch



Exception Behavior

❑ Preserving exception behavior

=> any changes in instruction execution order must not change how exceptions are raised in program

(=> no new exceptions)

➤ Example:

	DADDU	R2,R3,R4
	BEQZ	R2,L1
	LW	R1,0(R2)
L1:	

❑ Problem with moving LW before BEQZ?



A short summary

□ ILP

- The potential overlap among instructions

□ Reduce stalls from

- Structural hazards
- Data hazards
- Control hazards

□ To keep the program correctness, we should

- Preserving Data flow
- Preserving exception behavior



Lecture for ILP: Software approaches

- ❑ Basic Compiler Technique for Exposing ILP
 - Loop unrolling
- ❑ Static Branch Prediction
- ❑ Static multiple Issue: VLIW
- ❑ Advanced Compiler Support for Exposing and Exploiting ILP
 - Software pipelining
 - Global Code scheduling
- ❑ Hardware Support for Exposing More Parallelism at compile time
 - Conditional or Predicated instructions
 - Compiler speculation with hardware support



FP Loop: Where are the Hazards?

```
Loop: LD    F0,0(R1) ;F0=vector element
      ADDD  F4,F0,F2 ;add scalar from F2
      SD    0(R1),F4 ;store result
      SUBI  R1,R1,8  ;decrement pointer 8B (DW)
      BNEZ  R1,Loop  ;branch R1!=zero
      NOP                    ;delayed branch slot
```

<i>Instruction producing result</i>	<i>Instruction using result</i>	<i>Execution in cycles</i>	<i>Latency in cycles</i>
FP ALU op	Another FP ALU op	4	3
FP ALU op	Store double	3	2
Load double	FP ALU op	1	1
Load double	Store double	1	0
Integer op	Integer op	1	0

Where are the stalls?



Specification for the latency

□ ALU F1, -, - : IF ID FD FD FD FD WB

□ ALU -, F1, -: IF ID s s s FD FD FD FD WB

□ ALU: IF ID FD FD FD FD WB

□ SW: IF ID s s EX DM

□ LW F1, - : IF ID EX DM WB

□ SW: F1, 8(R1): IF ID EX DM WB

MEM/WB.LDMR --→ DM input port

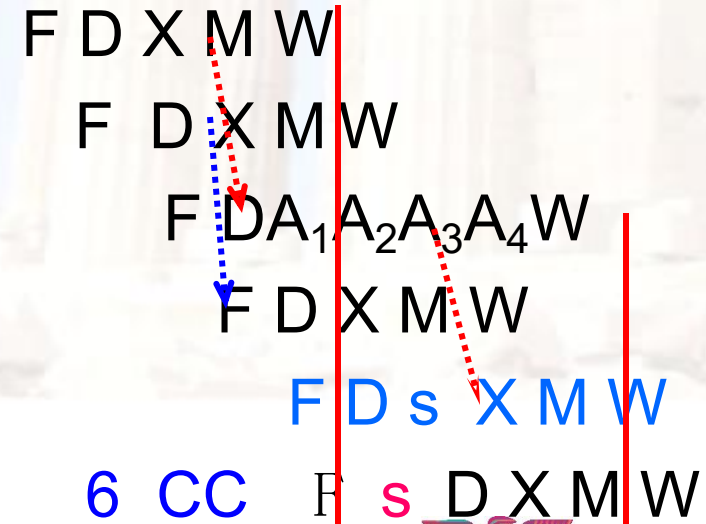


Reducing stalls from scheduling in BB and delayed branch

Loop: LD F0, 0(R1)
 ADDD F4, F0, F2
 SD 0(R1), F4
 SUBI R1, R1, #8
 BNEZ R1, Loop



Loop: LD F0, 0(R1)
 SUBI R1, R1, #8
 ADDD F4, F0, F2
 BNEZ R1, Loop
 SD +8(R1), F4





Unroll Loop Four Times (straightforward way)

```
1 Loop: LD      F0, 0(R1)
2      ADDD     F4, F0, F2
3      SD       0(R1), F4
4      LD       F6, -8(R1)
5      ADDD     F8, F6, F2
6      SD       -8(R1), F8
7      LD       F10, -16(R1)
8      ADDD     F12, F10, F2
9      SD       -16(R1), F12
10     LD       F14, -24(R1)
11     ADDD     F16, F14, F2
12     SUBI     R1, R1, #32
13     SD       +8(R1), F16
14     BNEZ     R1, LOOP
15     NOP
```

Annotations:

- 1 cycle stall (between lines 1 and 2)
- 2 cycles stall (between lines 2 and 3)
- drop SUBI & BNEZ (between lines 3 and 4)
- drop SUBI & BNEZ (between lines 6 and 7)
- 1 cycle stall (between lines 10 and 11)
- 1 cycle stall (waiting for F16) (between lines 11 and 12)
- alter to 4*8 (between lines 12 and 13)
- 1 cycle control stall (between lines 14 and 15)

Rewrite loop to
minimize stalls?

$14 + 3 \times (1+2) + 1 + 1 + 1 = 26$ clock cycles, or 6.5 per iteration
Assumes R1 is multiple of 4



Unrolled Loop That Minimizes Stalls

```
1 Loop: LD      F0, 0(R1)
2          LD      F6, -8(R1)
3          LD      F10, -16(R1)
4          LD      F14, -24(R1)
5          ADDD     F4, F0, F2
6          ADDD     F8, F6, F2
7          ADDD     F12, F10, F2
8          ADDD     F16, F14, F2
9          SD       0(R1), F4
10         SD       -8(R1), F8
11         SUBI     R1, R1, #32
12         SD       +16(R1), F12
13         BNEZ     R1, LOOP
14         SD       8(R1), F16 ; 8-32 = -24
```

□ What assumptions made when moved code?

- OK to move store past SUBI even though changes register
- OK to move loads before stores: get right data?
- When is it safe for compiler to do such changes?

14 clock cycles, or 3.5 per iteration



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Why Dynamic Scheduling ?

□ Example1 : Data Hazard

➤ DIVD F0,F2,F4
 ADDD F10,F0,F8
 SUBD F12,F8,F14

□ Example2: Structure Hazard

DIVD F2,F2,F4
ADDD F10,F0,F8 ; FP ADDer unpipelined
ADDD F12, F0,F4
MULD F16, F14, F4

- Problem: instruction (SUBD, MULD) stalled due to **irrelevant** forward instructions.



HW Schemes: Dynamic scheduling

- ❑ Key idea: **Allow instructions behind stall to proceed.**
Rearrange order of instructions to reduce stalls while maintaining data flow
- ❑ Enables **out-of-order execution**
and allows **out-of-order completion**
- ❑ Will distinguish when an instruction **begins execution** and when it **completes execution**; between 2 times, the instruction is **in execution**
- ❑ In a dynamically scheduled pipeline, all instructions pass through issue stage in order (**in-order issue**)



Adv. Of Dynamic Scheduling

- ❑ Handles cases when dependences unknown at compile time
 - (e.g., because they may involve a memory reference)
- ❑ It **simplifies** the compiler, Compiler doesn't need to have knowledge of microarchitecture
- ❑ Allows code that compiled for one pipeline to run efficiently on a different pipeline
- ❑ Hardware speculation, a technique with significant performance advantages, that builds on dynamic scheduling



Dynamic Scheduling Step 1

- ❑ Simple pipeline had 1 stage to check both structural and data hazards: Instruction Decode (ID), also called Instruction Issue
- ❑ Split the ID pipe stage of simple 5-stage pipeline into 2 stages:
- ❑ *Issue*—Decode instructions, check for structural hazards
- ❑ *Read operands*—Wait until no data hazards, then read operands



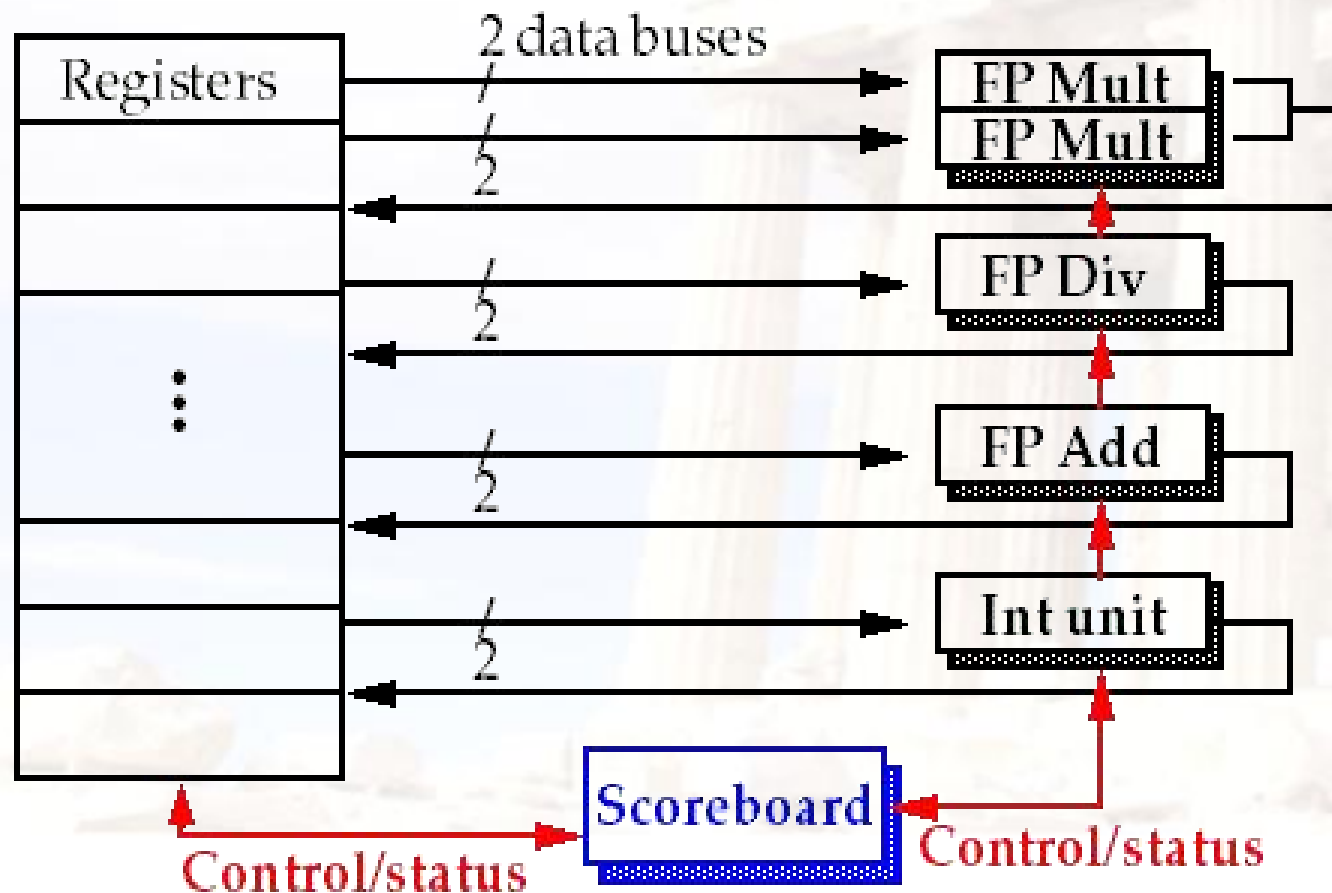
Dynamic Scheduling with a Scoreboard

□ Scoreboarding

- Named after CDC6600 scoreboard
- Allowing instructions to **execute out of order** when there are sufficient resources and no data dependences.
- In-order issue
- Out-of order completion
- Executing an instruction as early as possible



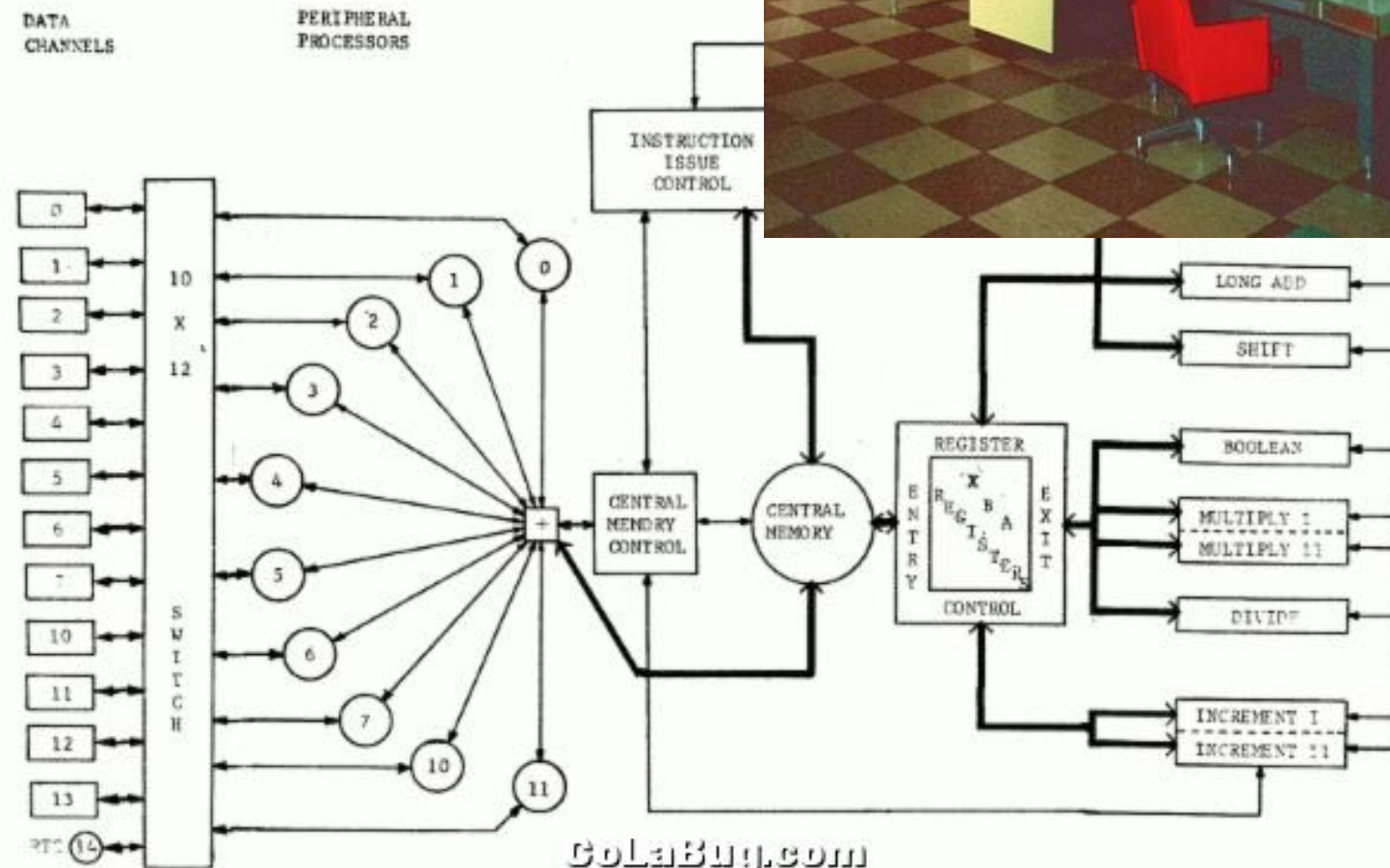
Basic structure of a pipelined processor with a scoreboard





CDC6600 –First Supercomputer

top1 1964-1969





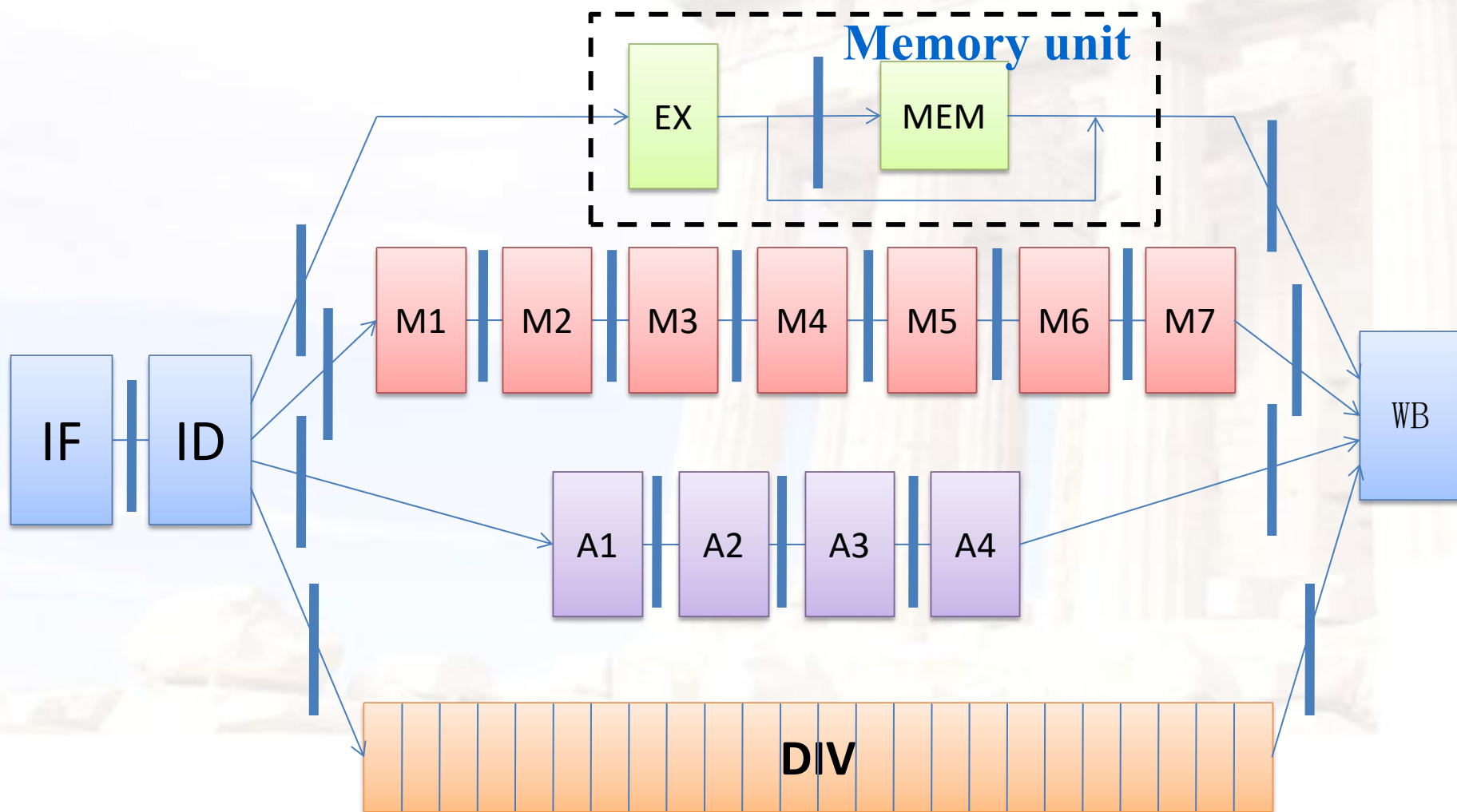
The pipeline stages with scoreboard

- ❑ The Five stages: IF, ID, EX, MEM, WB
 - IF: the same for all instructions
 - ID: split into two stages: issue and read operands
 - EX: no change
 - MEM: omitted for only concentrating on the FP operations
 - WB: no change
- ❑ So, the stages are: IF, IS, RO, EX, WB.

• Another way to look at missing MEM ?



Pipeline supports multiple outstanding FP operations





Scoreboard Pipeline stage description

❑ **Issue:** a instruction is issued when

- *The functional unit is available and*
- *No other active instruction has the same destination register.*
- Avoid **strutural** hazard and **WAW** hazard

❑ **Read Operands (RO)**

- *The read operation is delayed until both the operands are available.*
- *This means that no previously issued but ncompleted instruction has the operand as its destination.*
- This resolves **RAW** hazards dynamically

❑ **Execution (EX)**

- *Notify the scoreboard when completed so the functional unit can be reused.*

❑ **Write result (WB)**

- *The scoreboard checks for **WAR** hazards and stalls the completing instruction if necessary.*



The scoreboard algorithm

❑ Scoreboard-takes full responsibility for instruction issue and execution

- Create the dependence records
- Decide when to fetch the operand
- Decide when to enter execution
- Decide when the result can be written into the register file

❑ Three data structure

- Instruction status:
 - which of the four steps the instruction is in
- Functional unit status: $buzy, op, Fi, Fj, Fk, Qj, Qk, Rj, Rk$
- Register result status:
 - which functional unit will write that register



Example: Instruction status

LD F6, 34(R2)
LD F2, 45(R3)
MULTD F0, F2, F4
SUBD F8, F6, F2
DIVD F10, F0, F6
ADDD F6, F8, F2

Instruction	Instruction status			
	IS	RO	EX	WB
LD	✓	✓	✓	✓
LD	✓	✓	✓	
MULTD	✓			
SUBD	✓			
DIVD	✓			
ADDD				



Scoreboard Example

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2			
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Function Unit Status:

			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>			
Time	Name	Busy	Op	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
Clock cycle counter	Integer	No								
	Mult1	No								
	Mult2	No								
	Add	No								
	Divide	No								

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU									



Scoreboard Cycle 1

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2	✓		
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Function Unit Status:

			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>		
Time	Name	Busy	Op	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i> <i>Rk</i>
Clock cycle counter	Integer	Yes	load	F6	R2				Yes
	Mult1	No							
	Mult2	No							
	Add	No							
	Divide	No							

Register result status:

Clock
0

	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
FU	Int								





Scoreboard Cycle 2

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2	✓		
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Function Unit Status:

			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>		
Time	Name	Busy	Op	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>
Clock cycle counter	Integer	Yes	load	F6	R2				No
	Mult1	No							
	Mult2	No							
	Add	No							
	Divide	No							

Register result status:

Clock	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU Int								





Scoreboard Cycle 3

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2	<div>Calculate Address</div>		
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Function Unit Statous:

Clock cycle
counter

Time	Name	Busy	Op	des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>		
				<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
	Integer	Yes	load	F6	R2				No	
	Mult1	No								
	Mult2	No								
	Add	No								
	Divide	No								

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU	Int								



Scoreboard Cycle 4

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2	<div>✓</div>		
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Access Data Cache

Function Unit Status:

			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>		
Time	Name	Busy	Op	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>
Clock cycle counter	Integer	Yes	load	F6	R2				No
	Mult1	No							
	Mult2	No							
	Add	No							
	Divide	No							

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU	Int								





Scoreboard Cycle 5

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2			✓
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Function Unit Status:

			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>		
Time	Name	Busy	Op	<i>V_i</i>	<i>V_j</i>	<i>V_k</i>	<i>Q_j</i>	<i>Q_k</i>	<i>R_j</i> <i>R_k</i>
Clock cycle counter	Integer	No	load	F6	V _{R2}				No
	Mult1	No							
	Mult2	No							
	Add	No							
	Divide	No							

Register result status:

Clock	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU	M[R2+34]							



Scoreboard Cycle 6

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2	✓		
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Function Unit Status:

<i>Unit Statous:</i>			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>			
<i>Time</i>	<i>Name</i>	<i>Busy</i>	<i>Op</i>	<i>Vi</i>	<i>Vj</i>	<i>Vk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
e	Integer	Yes	load	F2	R3				Yes	
	Mult1	No								
	Mult2	No								
	Add	No								
	Divide	No								

Clock cycle
counter

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU		Int		M[R2+34]					





Scoreboard Cycle 7

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2			
LD	F2	45+	R3	✓		
MULTD	F0	F2	F4	✓		
SUBD	F8	F6	F2			
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Function Unit Status:

<i>Unit Statous:</i>			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>			
<i>Time</i>	<i>Name</i>	<i>Busy</i>	<i>Op</i>	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
	Integer	Yes	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4	Int		No	Yes
	Mult2	No								
	Add	No								
	Divide	No								

Register result status:

Clock

0

	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F31</i>
FU	Mult1	Int				M[R2+34]			



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Scoreboard Cycle 8

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2			
LD	F2	45+	R3		✓	Calculate Address
MULTD	F0	F2	F4	✓		
SUBD	F8	F6	F2	✓		
DIVD	F10	F0	F6			
ADDD	F6	F8	F2			

Function Unit Statous:

				des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>		
Time	Name	Busy	Op	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
Clock cycle counter	Integer	Yes	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4	Int		No	Yes
	Mult2	No								
	Add	Yes	Sub	F8	F6	F2		Int	Yes	No
	Divide	No								

Register result status:

Clock	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	<i>F30</i>
0	Mult1	Int.47	M[R2+34	Add				





Scoreboard Cycle 9

Instruction status:

Instruction	<i>j</i>	<i>k</i>	Issue	Read Oprand	Exec Comp	Write Result
LD	F6	34+	R2			
LD	F2	45+	R3		✓	
MULTD	F0	F2	F4	✓		
SUBD	F8	F6	F2	✓		
DIVD	F10	F0	F6	✓		
ADDD	F6	F8	F2			

Access Data Cache

Function Unit Statous:

Clock cycle
counter

Time	Name	Busy	Op	des <i>Fi</i>	<i>S1</i> <i>Fj</i>	<i>S2</i> <i>Fk</i>	<i>RS</i> <i>Qj</i>	<i>RS</i> <i>Qk</i>	<i>Rj</i>	<i>Rk</i>
	Integer	Yes	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4	Int		No	Yes
	Mult2	No								
	Add	Yes	Sub	F8	F6	F2		Int	Yes	No
	Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

Register result status:

Clock

0

	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
FU	Mult1	Int		M[R2+3]	Add	Div			





Example: Function unit status and Register status

Name	Functional unit status								
	Busy	Op	Fi	Fj	Fk	Qj	Qk	Rj	Rk
Integer	Yes	Load	F2	F3				No	
Mult1	Yes	Mult	F0	F2	F4	Integer		No	Yes
Mult2	No								
Add	Yes	Sub	F8	F6	F2		Integer	Yes	No
Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

	Register result status								
	F0	F2	F4	F6	F8	F10	F12	...	F30
FU	Mult1	Integer			Add	Divide		...	



Scoreboard Cycle 10

Instruction status:

Instruction	<i>j</i>	<i>k</i>	Issue	Read Oprand	Exec Comp	Write Result
LD	F6	34+	R2			
LD	F2	45+	R3			✓
MULTD	F0	F2	F4	✓		
SUBD	F8	F6	F2	✓		
DIVD	F10	F0	F6	✓		
ADDD	F6	F8	F2			

Function Unit Statous:

Clock cycle
counter

Time	Name	Busy	Op	des <i>Fi</i>	<i>S1</i> <i>Fj</i>	<i>S2</i> <i>Fk</i>	<i>RS</i> <i>Qj</i>	<i>RS</i> <i>Qk</i>	<i>Rj</i>	<i>Rk</i>
	Integer	No	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4	Int		Yes	Yes
	Mult2	No								
	Add	Yes	Sub	F8	F6	F2		Int	Yes	Yes
	Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

Register result status:

Clock

0

	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
FU	Mult1	1[R3+45]		M[R2+34]	Add	Div			





Scoreboard Cycle 11

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Op	Comp	Result
LD	F6	34+	R2			
LD	F2	45+	R3			
MULTD	F0	F2	F4	✓		
SUBD	F8	F6	F2	✓		
DIVD	F10	F0	F6	✓		
ADDD	F6	F8	F2			

Function Unit Statous:

Clock cycle
counter

<i>Unit Statous:</i>			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>			
<i>Time</i>	<i>Name</i>	<i>Busy</i>	<i>Op</i>	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
e	Integer	No	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4			No	No
	Mult2	No								
	Add	Yes	Sub	F8	F6	F2			No	No
	Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

Register result status:

Clock

0

	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
FU	Mult1	1[R3+45]		M[R2+34]	Add	Div			



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Scoreboard Cycle 12

Instruction status:

Instruction		<i>j</i>	<i>k</i>	Issue	Read Oprand	Exec Comp	Write Result
LD	F6	34+	R2				
LD	F2	45+	R3				
MULTD	F0	F2	F4			✓	
SUBD	F8	F6	F2			✓	
DIVD	F10	F0	F6	✓			
ADDD	F6	F8	F2				

Assume Mul takes 7 cycles

Assume Sub takes 3 cycles

Function Unit Statous:

Unit Statous:

			des	S1	S2	RS	RS			
Time	Name	Busy	Op	Fi	Fj	Fk	Qj	Qk	Rj	Rk
	Integer	No	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4			No	No
	Mult2	No								
	Add	Yes	Sub	F8	F6	F2			No	No
	Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

Register result status:

Clock

0

	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
FU	Mult1	4[R3+45]		M[R2+34]	Add	Div			



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Scoreboard Cycle 15

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Opnd	Comp	Result
LD	F6	34+	R2		✓	
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6	✓		
ADDD	F6	F8	F2			

Assume Mul takes 7 cycles

Assume Sub takes 3 cycles

Function Unit Statous:

Unit Statous:

			des	S1	S2	RS	RS			
Time	Name	Busy	Op	Fi	Fj	Fk	Qj	Qk	Rj	Rk
	Integer	No	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4			No	No
	Mult2	No								
	Add	No	Sub	F8	F6	F2			No	No
	Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU	Mult1	1[R3+45]		M[R2+34]	V-	Div			





Scoreboard Cycle 16

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2	✓	✓	
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6	✓	✓	
ADDD	F6	F8	F2	✓		

Assume Mul takes 7 cycles

Function Unit Statous:

<i>Unit Statous:</i>			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>			
<i>Time</i>	<i>Name</i>	<i>Busy</i>	<i>Op</i>	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
	Integer	No	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4			No	No
	Mult2	No								
	Add	Yes	ADD	F6	F8	F2			Yes	Yes
	Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU	Mult1	1[R3+45]		Add	V-	Div			





Scoreboard Cycle 17

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2	✓	✓	
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6	✓	✓	
ADDD	F6	F8	F2			

Assume Mul takes 7 cycles

Function Unit Statous:

<i>Unit Statous:</i>			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>			
<i>Time</i>	<i>Name</i>	<i>Busy</i>	<i>Op</i>	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
	Integer	No	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4			No	No
	Mult2	No								
	Add	Yes	ADD	F6	F8	F2			No	No
	Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU	Mult1	1[R3+45]		Add	V-	Div			





Scoreboard Cycle 18

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2		✓	
LD	F2	45+	R3			
MULTD	F0	F2	F4			
SUBD	F8	F6	F2			
DIVD	F10	F0	F6	✓	✓	
ADDD	F6	F8	F2			

Last cycle of Mul

Function Unit Statous:

<i>Unit Statous:</i>			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>			
<i>Time</i>	<i>Name</i>	<i>Busy</i>	<i>Op</i>	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
	Integer	No	load	F2	R3				No	
	Mult1	Yes	Mul	F0	F2	F4			No	No
	Mult2	No								
	Add	Yes	ADD	F6	F8	F2			No	No
	Divide	Yes	Div	F10	F0	F6	Mult1		No	Yes

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU	Mult1	1[R3+45]		Add	V-	Div			





Scoreboard Cycle 19

Instruction status:

				Read	Exec	Write
Instruction	<i>j</i>	<i>k</i>	Issue	Oprand	Comp	Result
LD	F6	34+	R2			
LD	F2	45+	R3			
MULTD	F0	F2	F4			✓
SUBD	F8	F6	F2			
DIVD	F10	F0	F6	✓		
ADDD	F6	F8	F2		✓	

Function Unit Statous:

<i>Unit Statous:</i>			des	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>			
<i>Time</i>	<i>Name</i>	<i>Busy</i>	<i>Op</i>	<i>Fi</i>	<i>Fj</i>	<i>Fk</i>	<i>Qj</i>	<i>Qk</i>	<i>Rj</i>	<i>Rk</i>
	Integer	No	load	F2	R3				No	
	Mult1	No	Mul	F0	F2	F4			No	No
	Mult2	No								
	Add	Yes	ADD	F6	F8	F2			No	No
	Divide	Yes	Div	F10	F0	F6	Mult1		Yes	Yes

Register result status:

Clock		<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
0	FU	V*	1[R3+45]		Add	V-	Div			





Examples: Dynamic Scheduling

□ Example 1: RAW

fdiv.d **f0**,f2,f4

fadd.d f10,**f0**,f8

fsub.d f12,f8,f14

- fsub.d is not dependent, issue before fadd.d

□ Example 2:

fdiv.d f0,f2,f4

fmul.d f6,**f0**,f8

fadd.d **f0**,f10,f14

- fadd.d is not dependent, but the antidependence makes it impossible to issue earlier without register renaming

□ Example 1: structural hazard

fdiv.d f2,f2,f4

Fadd.d f10,f0, f8 ;unpipelined
;Adder

Fadd.d f12, f0,f4

fMul.d f16, f14, f4

- fMul.d is not dependent, issue before fadd.d

□ Example 2:

fdiv.d f2,f2,f4

fdiv.d f6,**f0**,f8

fadd.d **f0**,f10,f14

- fadd.d is not dependent, but the antidependence makes it impossible to issue earlier without register renaming



Limitations of Scoreboard-1

□ ILP

- If we can't find independent instructions to execute, scoreboard (or any dynamic scheduling scheme for that matter) helps very little.

□ Size of the "issued" queue

- This determines how far ahead the CPU can look for instructions to execute in parallel.
- It's called the **window**.
- For now, we assume that a **window** can **not** span a branch.
- In other words, the window includes instructions only within basic blocks.



Limitations of Scoreboard-2

- ❑ Number, types, and speed of the functional units
 - This determines how often a structural hazard results in stall.
- ❑ The presence of anti-dependences and output dependences
 - **WAR** and **WAW** hazards limit the scoreboard more than **RAW** hazards, lead to **WAR** and **WAW** stalls.
 - **RAW** hazards are problems for any technique.
 - But **WAR** and **WAW** hazards can be solved in ways other than scoreboards.



Register Renaming

----solve WAW、WAR

□ Example

```
fdiv.d f0,f2,f4  
fadd.d f10,f0,f8  
fsd f10,0(x1)  
fsub.d f12,f10,f14  
fmul.d f14, f2, f16 ;f14 WAR
```

Diagram showing data flow: Red arrows from **f0** to **f10** and from **f10** to **fsd**. A blue arrow from **f14** to **fmul.d**.

- Though multiplier can finish execution far earlier before sub, but it can write result into **f14** due to WAR dependence on sub.

□ After Reg Rename

```
fdiv.d f0,f2,f4  
fmul.d S,f0,f8  
fsd S,0(x1)  
fsub.d T,S,f14  
fmul.d U, f2, f16
```

Diagram showing data flow: Red arrows from **f0** to **S** and from **S** to **fsd**. A red arrow from **S** to **fsub.d**. A green arrow from **U** to **fmul.d**.

- Now only RAW hazards remain, which can be strictly ordered. WAR disappear.