

# Formal Languages and Compilers

21 October 2022

Using the JFLEX lexer generator and the CUP parser generator, realize a JAVA program capable of recognizing and executing the programming language described in the following.

## Input language

The input file is composed of two sections: *header* and *command* sections, separated by means of the sequence of characters “\$\$\$”. Comments are possible, and they are delimited by the starting sequence “{++” and by the ending sequence “++}”.

### Header section: lexicon

The *header* section can contain 3 types of tokens, each terminated with the character “;”:

- **<tk1>**: it begins with the character “?”, followed by a word composed of an even number (at least 6) of uppercase alphabetic letters. It is then followed by an octal number (possible digits are from 0 to 7) between –127 and 323, and optionally followed by 4 or more repetitions of the words “xx”, “yy”, or “zz” in any combination.
- **<tk2>**: it is composed by 2, 12, or 15 emails, where each email is a word composed of numbers, letters and characters “\_” and “.”, the character “@”, and a word composed of letters and numbers, a “.”, and the word “it”, “org”, or “com”. Each email is separated by the character “!” or “/”.
- **<tk3>**: it is the word “tk3”.

### Header section: grammar

In the *header section* the tokens **<tk1>** and **<tk3>** can appear in **any order and number (even 0 times)**, instead, **<tk2>** can appear only **0, 1, or 4 times**.

### Code section: grammar and semantic

The *command section* is composed of a list of **<commands>**. The list can be possibly **empty**, or with an **even** number of elements, **at least 4**. As a consequence, the list can be composed of 0, 4, 6, 8,... elements.

Two types of commands are possible:

- *Assignment*: it is a **<variable>** (same regular expression of C identifiers), followed by a “=”, and a **<bool\_expr>**. This command stores the result of the **<bool\_expr>** into an entry of a global symbol table with key **<variable>**. **This symbol table is the only global data structure allowed in all the examination, and it can be written only by means of an assignment command.** Each time an *assignment* command is executed, the command prints into the screen the **<variable>** name and the associated value.

- *CMP*: it has the following syntax:

**CMP <bool\_expr> <actions\_list>**

where **<bool\_expr>** represents the result of a boolean expression (i.e., a T (true) or a F (false) values). **<actions\_list>** is a list of at least one **<action>**, where an **<action>** is the word **WITH**, a **<bool\_expr.a>** (same regular expression of **<bool\_expr>**), the character “[”, a **<print>** instruction, and the character “]”. The **<print>** instruction is the word **print**, followed by a “(”, a *quoted*

*string*, a “)”, and a “;”. The `<print>` instruction is executed each time the result of `<bool_expr>` equals the result of `<bool_expr_a>`.

`<bool_expr>` can contain the following logical operators: AND, OR, NOT, and round brackets. Operands can be T (the true constant), F (the false constant), a `<variable>` (which represents the value stored in the symbol table by an *assignment* command), and the `fz_and()` function. The `fz_and()` function takes in input a list of `<bool_expr>` separated by a “,”, and returns the “logical and” of the results of the listed `<bool_expr>` (i.e., the `fz_and()` function returns T only if the results of all the listed `<bool_expr>` are T, otherwise it returns F).

## Goals

The translator must execute the language, and it must produce the output reported in the example. For any detail not specified in the text, follow the example.

## Example

### Input:

```
tk3 ;                                {++ tk3 ++}
name1.surname1@skenz.it/name2.surname2@abc.net; {++ tk2 ++}
?ABCFEF-36xyxyxyy ;                {++ tk1 ++}
tk3;                                {++ tk3 ++}
```

```
$$$ {++ division between header and command sections ++}
```

```
x1 = T;
x2 = NOT T AND NOT x1 ; {++ F AND F = F ++}
```

```
{++ fz_and(T, T, T, F) OR F = F OR F = F ++}
x3 = fz_and(T, T, fz_and(T, x1), F ) OR F;
```

```
CMP T AND F {++ T AND F = F ++}
WITH F OR F [ {++ executed ++}
    print("one");
]
WITH x1 [ {++ not executed ++}
    print("two");
]
WITH NOT x1 [ {++ executed ++}
    print("three");
]
```

### Output:

```
x1 T
x2 F
x3 F
"one"
"three"
```

**Weights:** Scanner 8/30; Grammar 9/30; Semantic 10/30