

## Homework 1.

- The file name of your homework (in PDF) should be in the format: “學號-ds-作業編號.pdf”. For example: **00757999-ds-hw1.pdf**
- Please submit your homework to Tronclass **before 23:59, October 19, 2025**

## Questions:

1. (10%) Order the following function by growth rate in increasing order:  
 $n^2$ ,  $n \log(\log n)$ ,  $n \log(n^2)$ ,  $n^n$

2. (10%) Ackerman's function  $A(m, n)$  is defined as:

$$A(m, n) = \begin{cases} n + 1, & \text{if } m = 0 \\ A(m - 1, 1), & \text{if } n = 0 \\ A(m - 1, A(m, n - 1)), & \text{otherwise} \end{cases}$$

What is the value of  $A(3, 4)$ ?

3. (12%) Show that the following statements are correct:
- (a)  $5n^2 - 6n = \Theta(n^2)$
  - (b)  $n! = O(n^n)$
4. (12%) Show that the following statements are incorrect:
- (a)  $10n^2 + 9 = O(n)$
  - (b)  $3^n = O(2^n)$
5. (16%) (from NTHU 106) Which ones of the following conjectures are correct?  
Please justify your answers:
- (a)  $3n^2 - 100n + 6 = O(n)$
  - (b)  $\log(n!) = \Theta(n \log n)$
  - (c)  $n! = O(n^n)$
  - (d)  $(n + a)^b = \Theta(n^b)$  where  $a$  and  $b$  are real constants
  - (e)  $(n + a)^b = \Omega(n^{b+1})$  where  $a$  and  $b$  are real constants

6. (20%) Fill the step count table (yellow regions) for the following program:

Statement	s/e	Frequency	Total Steps
<b>void</b> <i>mult</i> ( <b>int</b> <i>a</i> [][MAX_SIZE], <b>int</b> <i>b</i> [][MAX_SIZE], <b>int</b> <i>c</i> [][MAX_SIZE])	0	0	0
{	0	0	0
<b>int</b> <i>i, j, k</i> ;	0	0	0
<b>for</b> ( <i>i</i> = 0; <i>i</i> < MAX_SIZE; <i>i</i> ++)	1		
<b>for</b> ( <i>j</i> = 0; <i>j</i> < MAX_SIZE; <i>j</i> ++) {	1		
<i>c</i> [ <i>i</i> ][ <i>j</i> ] = 0;	1		
<b>for</b> ( <i>k</i> = 0; <i>k</i> < MAX_SIZE; <i>k</i> ++)	1		
<i>c</i> [ <i>i</i> ][ <i>j</i> ] += <i>a</i> [ <i>i</i> ][ <i>k</i> ] * <i>b</i> [ <i>k</i> ][ <i>j</i> ];	1		
}	0	0	0
}	0	0	0
Total			

7. (10%) Given an four-dimensional array A[100][200][300][400], if the address of A[3][20][32][45] is  $\alpha + X$ , what is the value of X ?

Note: (1)  $\alpha$  is the address of A[0][0][0][0]

(2) Only consider row major here.

8. (10%) (from NTCU 107) Consider the following function. Please derive the time complexity of function foo using recurrence relation. Please show your derivation step by step.

```

int i[n]; /* Assume that the values of all integers in i[n] have been initialized */
int foo(int a, int b, int c) {
    int d;
    if (a>b) return -1;
    d=(a+b)/2;
    if (c==i[d])
        return d;
    else if (c>i[d])
        return foo(d+1,b,c);
    else
        return foo(a, d-1,c);
}

```

**Answer:**

1.  $n^n$  = Biggest

$$\frac{n \log(n^2)}{n^2} = \frac{2n \log(n)}{n^2} \rightarrow 0 \quad \text{as } n \rightarrow \infty$$

$$n \log(n^2) < n^2 < n^n$$

$$\log n < n^2$$

$$n \log(\log n) < n \log(n^2) < n^2 < n^n$$

---

2.  $A(0, n) = n + 1$

$$A(1, n) = A(0, A(1, n-1)) = A(1, n-1) + 1$$

$$A(1, 0) = A(0, 1) = 2$$

$$A(1, 1) = A(0, A(1, 0)) = 2 + 1 = 3$$

$$A(1, 2) = A(0, A(1, 1)) = 3 + 1 = 4$$

$$A(1, n) = n + 2$$

$$A(2, n) = A(1, A(2, n-1)) = A(2, n-1) + 2$$

$$A(2, 0) = A(1, 1) = 3$$

$$A(2, 1) = A(1, A(2, 0)) = 5$$

$$A(2, 2) = A(1, A(2, 1)) = 7$$

$$A(2, n) = 2n + 3$$

$$A(3, n)$$

$$A(3, 0) = A(2, 1) = 2 \cdot 1 + 3 = 5 \quad \xrightarrow{+8} = 8 - 3$$

$$A(3, 1) = A(2, A(3, 0)) = 2(5) + 3 = 13 \quad \xrightarrow{+16} = 16 - 3$$

$$A(3, 2) = A(2, A(3, 1)) = 2(13) + 3 = 29 = 32 - 3$$

$$A(3, n) = 2^{3+n} - 3$$

$$A(3, 4) = 2^7 - 3 = 128 - 3 = 125$$

3. a.  $5n^2 - 6n = \Theta(n^2)$

$$c_1 g(n) \leq f(n) \leq c_2 g(n)$$

$$(c_1=4 \quad c_2=5) \checkmark \quad n_0=6$$

$$4n^2 \leq 5n^2 - 6n \leq 5n^2 \quad \text{for all } n \geq 6$$

Correct

b.  $n! = O(n^n)$

$$n! \leq C n^n \quad C=1$$

$$n! \leq 1 \cdot n^n \quad \text{for all } n \geq 1$$

Correct

4. a.  $10n^2 + 9 \leq Cn$

$$10n + \frac{9}{n} \leq C$$

$$\text{as } n \rightarrow \infty$$

$$\infty \leq C \quad \text{is } \underline{\text{wrong}}$$

b.  $3^n \leq C \cdot 2^n$

$$\frac{3^n}{2^n} \leq C$$

$$\left(\frac{3}{2}\right)^n \leq C$$

$$\text{as } n \rightarrow \infty$$

$$\infty \leq C \quad \text{is } \underline{\text{wrong}}$$

5. a. False  
 $\exists n - 100 + \frac{6}{n} \leq C$

a)  $n \rightarrow \infty$

$\infty \leq C$  is wrong

b. True

$C_1 \cdot n \cdot \log n \leq \log(n!)$   $\leq C_2 \cdot n \cdot \log n$

Upper bound  $C_2 = 1$

$\log(n!) = \log(1 \cdot 2 \cdots n)$

$= \log(1) + \log(2) + \cdots + \log(n)$

$= \sum_{k=1}^n \log(k) \leq n \log(n)$

lower bound  $C_1 = \frac{1}{2}$

$\log(n!) \geq \sum_{k=n/2}^n \log(k)$

$\sum_{k=n/2}^n \log(k) \geq n/2 \cdot \log(n/2)$

$\log(n!) \geq n/2 \cdot \log(n/2)$

$= \frac{n}{2} (\log(n) - \log(2))$

drop the constants

$\log(n!) \geq n \log(n)$

Upper and lower bound

Proved

5. c. True

$n! \leq C \cdot n^n$

$C > 1$

$n! \leq n^n$  for all  $n \geq 1$

D. True

$C_1 n^b \leq (n+a)^b \leq C_2 n^b$

Substitute

$C_1 \leq \left(\frac{n+a}{n}\right)^b \leq C_2$

$C_1 \leq \left(1 + \frac{a}{n}\right)^b \leq C_2$  exists for all  $n \geq 1$

as  $n \rightarrow \infty$

$C_1 \leq 1 \leq C_2$

5. e. False

$(n+a)^b \geq C \cdot n^{b+1}$

$\left(\frac{n+a}{n}\right)^b \geq C \cdot n$

$\left(1 + \frac{a}{n}\right)^b \geq C \cdot n$

as  $n \rightarrow \infty$

$1 \geq \infty$  is wrong

6. (20%) Fill the step count table (yellow regions) for the following program: `let n = MAX_SIZE`

Statement	s/e	Frequency	Total Steps
<b>void mult(int a[][MAX_SIZE], int b[][MAX_SIZE], int c[][MAX_SIZE])</b>	0	0	0
<b>{</b>	0	0	0
<b>int i, j, k;</b>	0	0	0
<b>for (i = 0; i &lt; MAX_SIZE; i++)</b>	1	$n+1$	$n+1$
<b>for (j = 0; j &lt; MAX_SIZE; j++) {</b>	1	$n^2+n$	$n^2+n$
<b>c [ i ][ j ] = 0;</b>	1	$n^2$	$n^2$
<b>for (k = 0; k &lt; MAX_SIZE; k++)</b>	1	$n^3+n^2$	$n^3+n^2$
<b>c [ i ][ j ] += a [ i ][ k ] * b [ k ][ j ];</b>	1	$n^3$	$n^3$
<b>}</b>	0	0	0
<b>}</b>	0	0	0
Total	$O(n^3)$		$2n^3+3n^2+2n+1$

7.  $\alpha + x$

$$\begin{aligned}
 x &= 3 \times 200 \times 300 \times 400 + 20 \times 300 \times 400 + 32 \times 400 + 45 \\
 &= \underline{74,412,845}
 \end{aligned}$$

$$8. \quad T(a, b, c) = \begin{cases} -1 & \text{if } a > b \\ d & \text{if } c = i[(a+b)/2] \\ T((a+b)/2 + 1, b, c) & \text{if } c > i[(a+b)/2] \\ T(a, (a+b)/2 - 1, c) & \text{if } c < i[(a+b)/2] \end{cases}$$

The function does constant-time work like comparison and assignments, then the function calls itself recursively on half of the array so each recursions halves the size.

$$T(n) = T(n/2) + C$$

$$T(n/2) = T(n/4) + 2C$$

$$T(n/4) = T(n/8) + 3C$$

$$\text{So, } T(n) = T(n/2^k) + kC$$

The recursion stops when base case is reached which is when  $T(n/2^k)$  takes  $T(1)$ .

$$\text{So, } n/2^k = 1$$

$$n = 2^k$$

$$k = \log_2 n$$

So,

$$T(n) = T(1) + C \log n$$

$$\underline{\underline{T(n) = O(\log n)}}$$