Real-time Dynamic Function Call Detection

A heuristic-based approach

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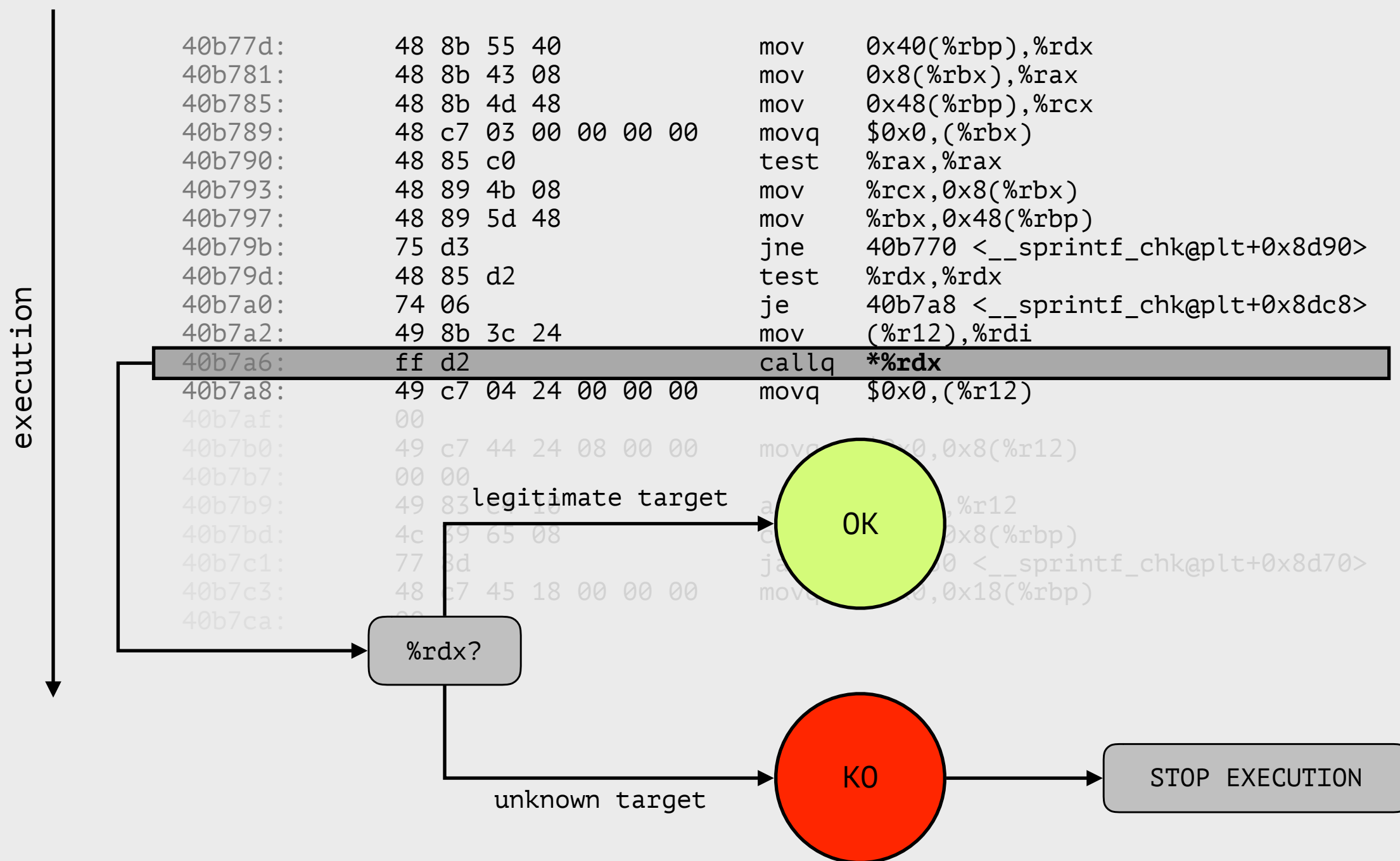Dennis Andriesse

VU Amsterdam

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UGA

INTRODUCTION

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let's try to keep everyone awake
for two minutes

AN EXAMPLE TO ILLUSTRATE THE PROBLEM — CFI

```
40b77d:      48 8b 55 40               mov      0x40(%rbp),%rdx
40b781:      48 8b 43 08               mov      0x8(%rbx),%rax
40b785:      48 8b 4d 48               mov      0x48(%rbp),%rcx
40b789:      48 c7 03 00 00 00 00      movq     $0x0,(%rbx)
40b790:      48 85 c0                  test     %rax,%rax
40b793:      48 89 4b 08               mov      %rcx,0x8(%rbx)
40b797:      48 89 5d 48               mov      %rbx,0x48(%rbp)
40b79b:      75 d3                     jne      40b770 <__sprintf_chk@plt+0x8d90>
40b79d:      48 85 d2                  test     %rdx,%rdx
40b7a0:      74 06                     je       40b7a8 <__sprintf_chk@plt+0x8dc8>
40b7a2:      49 8b 3c 24               mov      (%r12),%rdi
                                       push     40b7a8
                                       jmp      *%rdx
40b7a6:      ff d2                     callq    *%rdx
40b7a8:      49 c7 04 24 00 00 00      movq     $0x0,(%r12)
40b7af:      00
40b7b0:      49 c7 44 24 08 00 00      movq     $0x0,0x8(%r12)
40b7b7:      00 00
40b7b9:      49 83 c4 10               add      $0x10,%r12
40b7bd:      4c 39 65 08               cmp      %r12,0x8(%rbp)
40b7c1:      77 8d                     ja       40b750 <__sprintf_chk@plt+0x8d70>
40b7c3:      48 c7 45 18 00 00 00      movq     $0x0,0x18(%rbp)
40b7ca:      00
```

execution

CFI **SHOULD** APPLY, BUT WILL IT?

THE PROBLEM

HOW TO INSTRUMENT **EVERY CALL**

DURING EXECUTION?

HOW TO DISTINGUISH **INTRA-PROCEDURAL JUMPS** FROM **JUMP-BASED CALLS**?

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THE REAL REAL PROBLEM

HOW TO DISTINGUISH **INTRA-PROCEDURAL**

JUMPS FROM **TAIL CALLS**?

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OUR PROPOSED SOLUTION: iCi

	-O0	-O1	-O2	-O3
	iCi	iCi	iCi	iCi
binutils	1,000	1,000	1,000	1,000
evince	0,991	0,985	0,985	0,986
coreutils	0,999	0,998	0,998	0,998
ffmpeg	0,997	0,997	0,997	0,997
SPEC2006 jcall worst	nc	nc	1,000	nc
SPEC2006 iCi worst	nc	nc	0,947	nc

fscore

APPROACH

let's jump everywhere, but not in any ways

OVERVIEW

POSSIBLE APPROACHES

RESOLVE JUMPS

STATIC

FIND ENTRY POINTS

PROBLEM

POST-MORTEM

DYNAMIC

WITHOUT PRIOR KNOWLEDGE   iCi

REAL-TIME

WITH PRIOR KNOWLEDGE
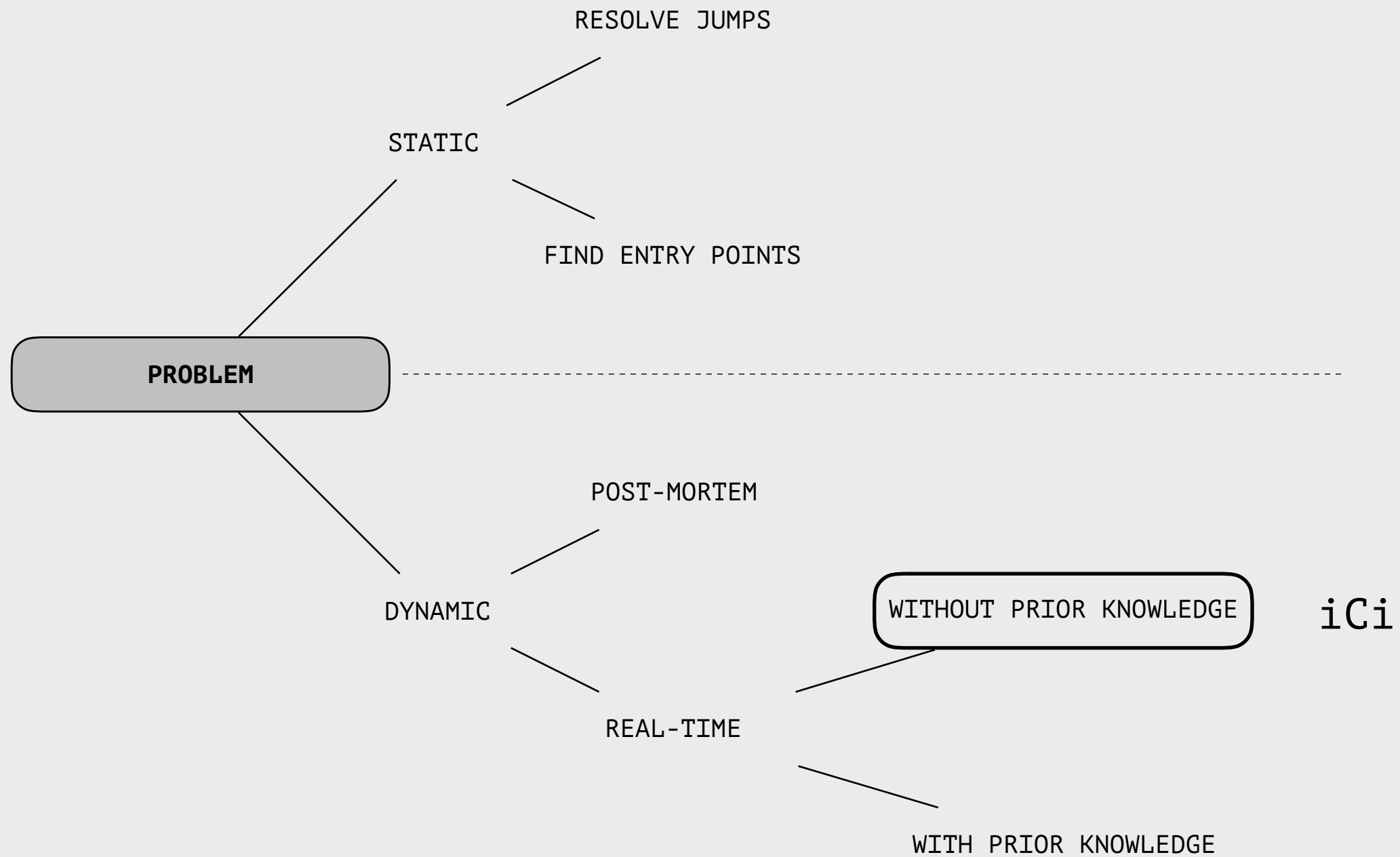
JMP SELECTION IN iCi
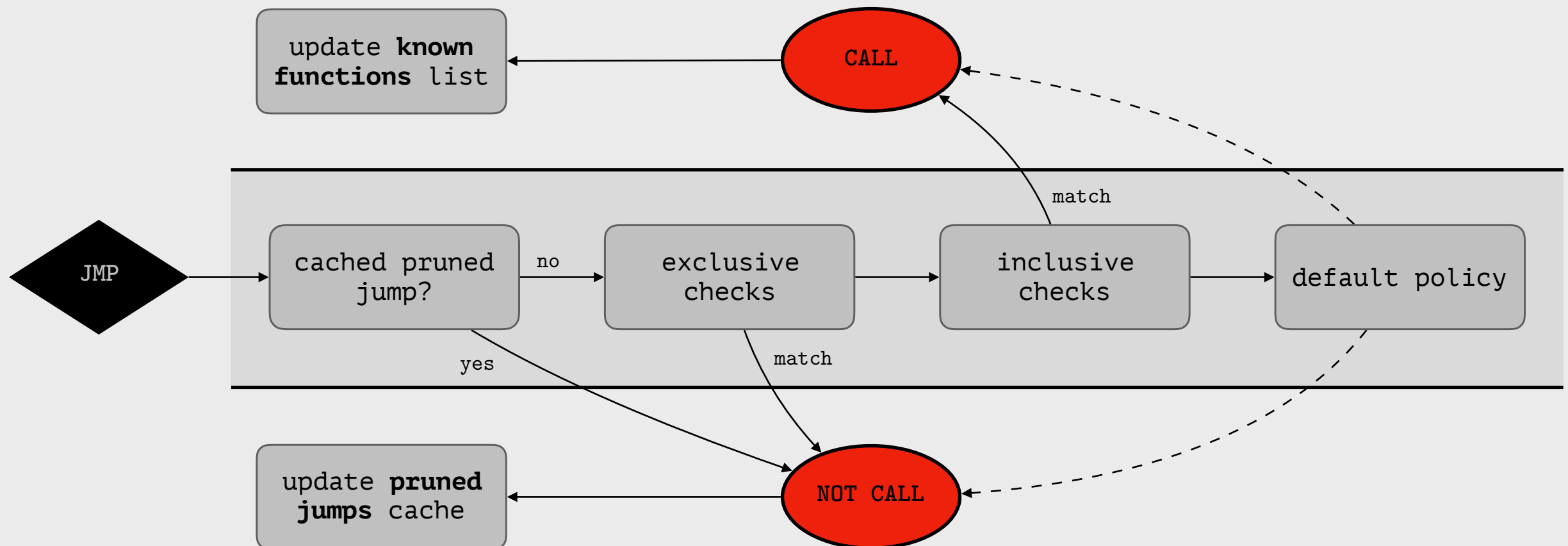
BACKGROUND INFORMATION

LIST OF ENTRY POINTS

Each time a function call is detected, its **target** is marked as an **entry point**.

LIST OF RETURN LOCATIONS

Each time a function returns, the location of the last instruction is marked as a **return point**.

background information

AVAILABLE INFO FOR EACH JMP

information	accessible?
target of the JMP	ok (dynamic)
current %rip, %rsp	ok (dynamic)
current entry point	**~ok (depend on the known entry points)**
state of stack when last call occurred	**~ok (idem)**
return locations for the current function	**not at first call; then ~ok**

— precision over time

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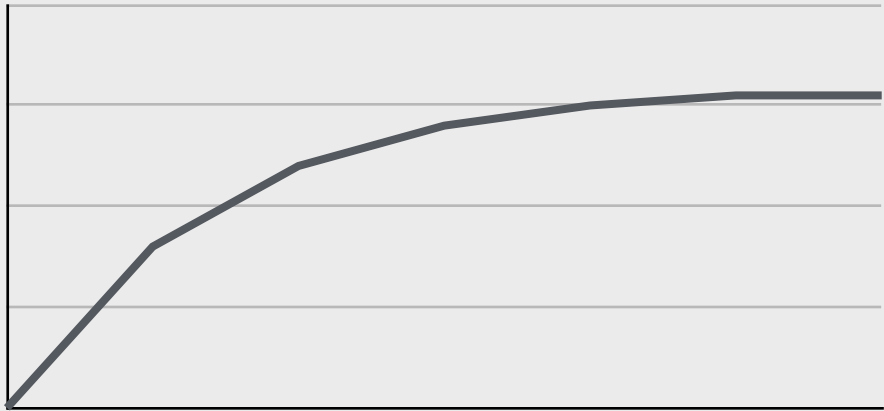
HEURISTICS

because... you know... provable approaches are overrated

exclusive checks

update **known functions** list

CALL

JMP

cached pruned jump? — no — exclusive checks — inclusive checks — default policy

yes

match

match

update **pruned jumps** cache

NOT CALL

exclusive checks
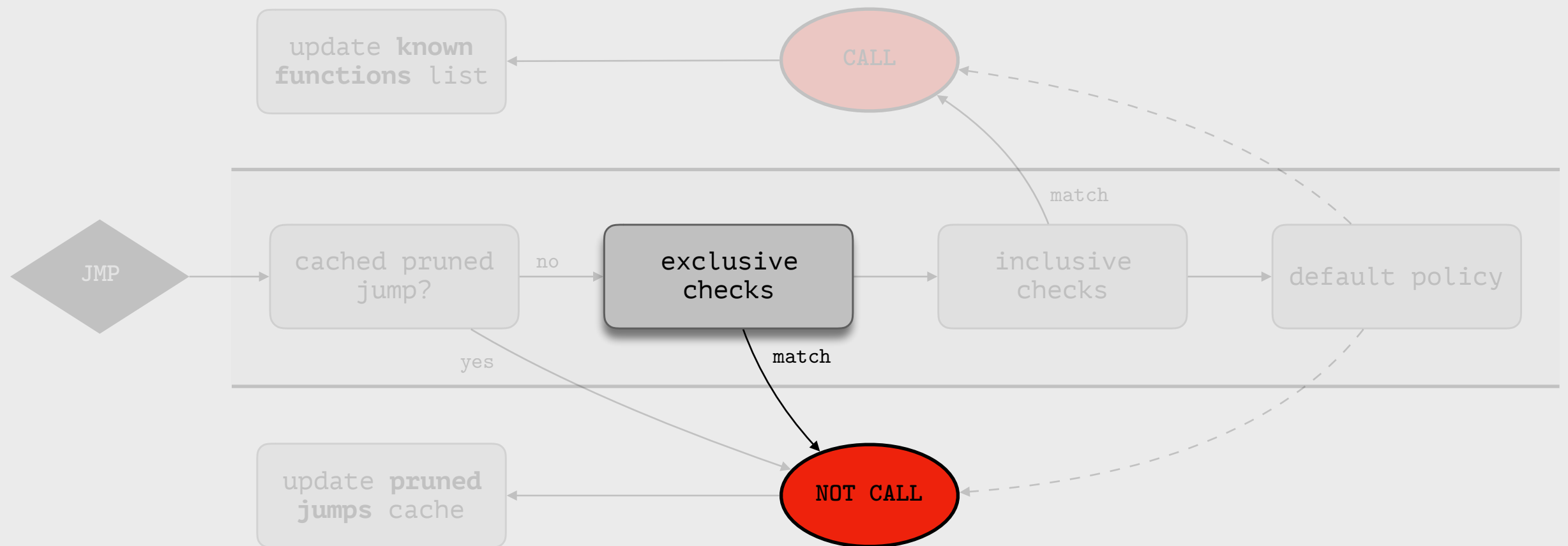
INTERNAL JUMP BEFORE

> A jump to a location that is between the current entry point and the current %rip **is not a call**.

```
…      :   …                         …
401e59:    c3                        retq
401e5a:    55                        push    %rbp          ← last known entry point
401e5b:    48 89 e5                  mov     %rsp,%rbp
401e5e:    48 83 ec 30               sub     $0x30,%rsp
401e62:    48 89 7d d8               mov     %rdi,-0x28(%rbp)
401e66:    48 8b 45 d8               mov     -0x28(%rbp),%rax
401e6a:    48 8b 80 b0 27 00 00      mov     0x27b0(%rax),%rax
…      :   …                         …       …
401e8a:    eb 2e                     jmp     401eba
401e8c:    48 8b 45 f8               mov     -0x8(%rbp),%rax
401e90:    8b 40 20                  mov     0x20(%rax),%eax
401e93:    85 c0                     test    %eax,%eax
401e95:    75 17                     jne     401eae
401e97:    48 8b 55 d8               mov     -0x28(%rbp),%rex
401e9b:    48 8d 4d e0               lea     -0x20(%rbp),%rcx
…      :   …                         …       …
401eb6:    48 89 45 f8               mov     %rax,-0x8(%rbp)
401eba:    48 83 7d f8 00            cmpq    $0x0,-0x8(%rbp)
401ebf:    75 cb                     jne     401e8c
401ec1:    90                        nop
401ec2:    c9                        leaveq
401ec3:    c3                        retq
401ec4:    55                        push    %rbp
…      :   …                         …       …
```
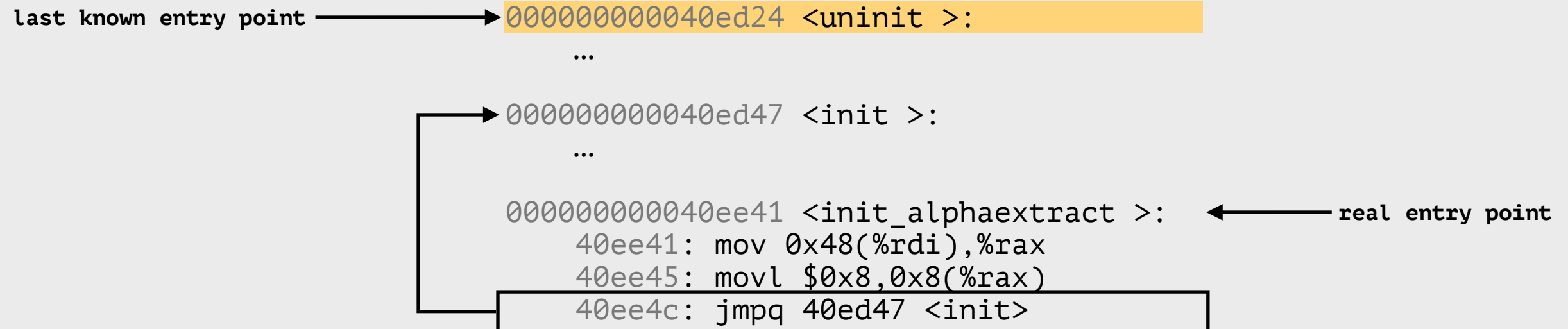
internal jump before

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ERROR EXAMPLE

last known entry point ⟶ `000000000040ed24 <uninit >:`
        …

`000000000040ed47 <init >:`
        …

`000000000040ee41 <init_alphaextract >:` ⟵ **real entry point**
        `40ee41: mov 0x48(%rdi),%rax`
        `40ee45: movl $0x8,0x8(%rax)`
        `40ee4c: jmpq 40ed47 <init>`

exclusive checks
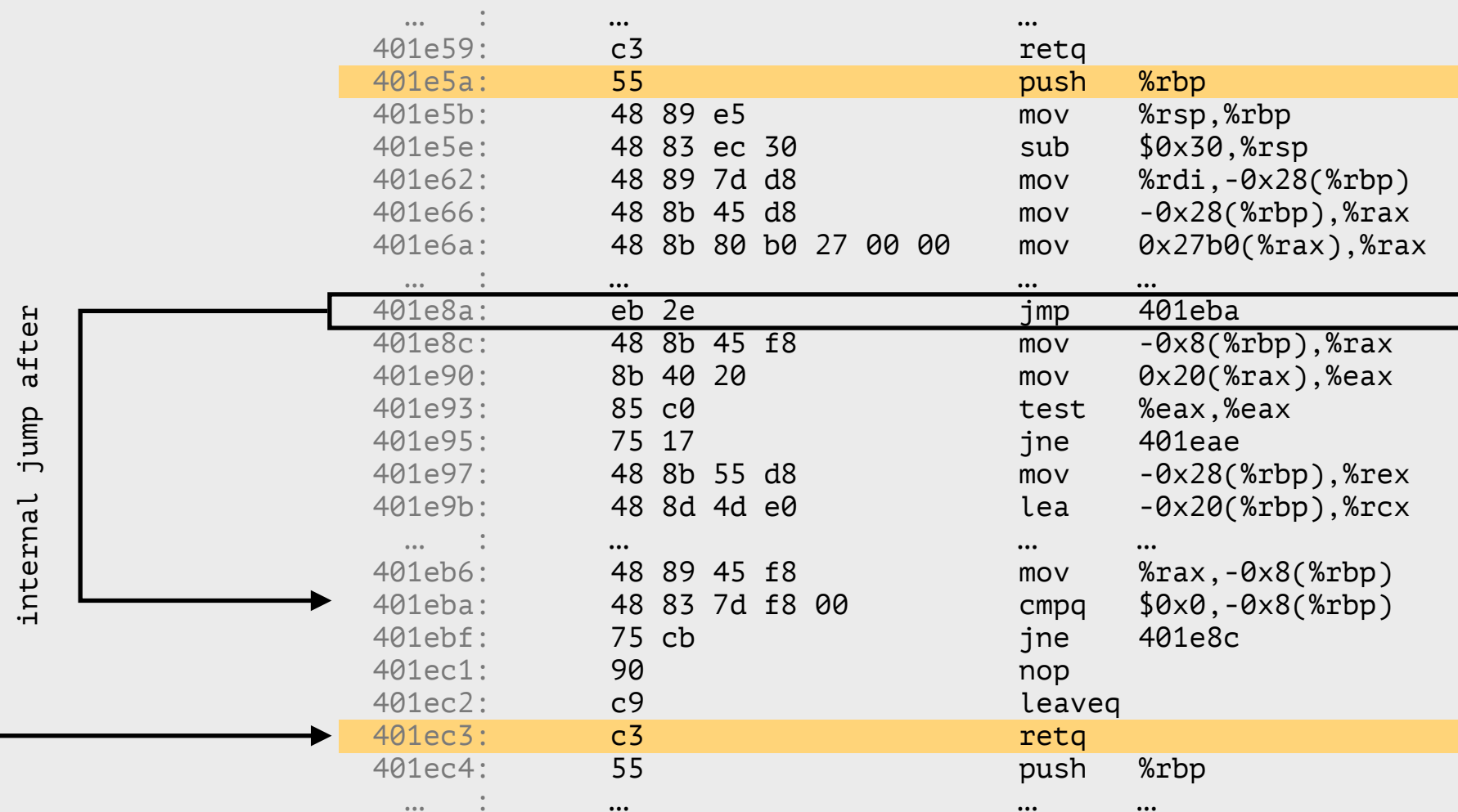
INTERNAL JUMP AFTER

A jump to a location that is between the current %rip and a known ret of the current function **is not a call**.

```
…    :      …                          …
401e59:      c3                        retq
401e5a:      55                        push    %rbp
401e5b:      48 89 e5                  mov     %rsp,%rbp
401e5e:      48 83 ec 30               sub     $0x30,%rsp
401e62:      48 89 7d d8               mov     %rdi,-0x28(%rbp)
401e66:      48 8b 45 d8               mov     -0x28(%rbp),%rax
401e6a:      48 8b 80 b0 27 00 00      mov     0x27b0(%rax),%rax
…    :      …                          …    …
401e8a:      eb 2e                     jmp     401eba
401e8c:      48 8b 45 f8               mov     -0x8(%rbp),%rax
401e90:      8b 40 20                  mov     0x20(%rax),%eax
401e93:      85 c0                     test    %eax,%eax
401e95:      75 17                     jne     401eae
401e97:      48 8b 55 d8               mov     -0x28(%rbp),%rex
401e9b:      48 8d 4d e0               lea     -0x20(%rbp),%rcx
…    :      …                          …    …
401eb6:      48 89 45 f8               mov     %rax,-0x8(%rbp)
401eba:      48 83 7d f8 00            cmpq    $0x0,-0x8(%rbp)
401ebf:      75 cb                     jne     401e8c
401ec1:      90                        nop
401ec2:      c9                        leaveq
401ec3:      c3                        retq
401ec4:      55                        push    %rbp
…    :      …                          …    …
```

internal jump after

known ret

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exclusive checks
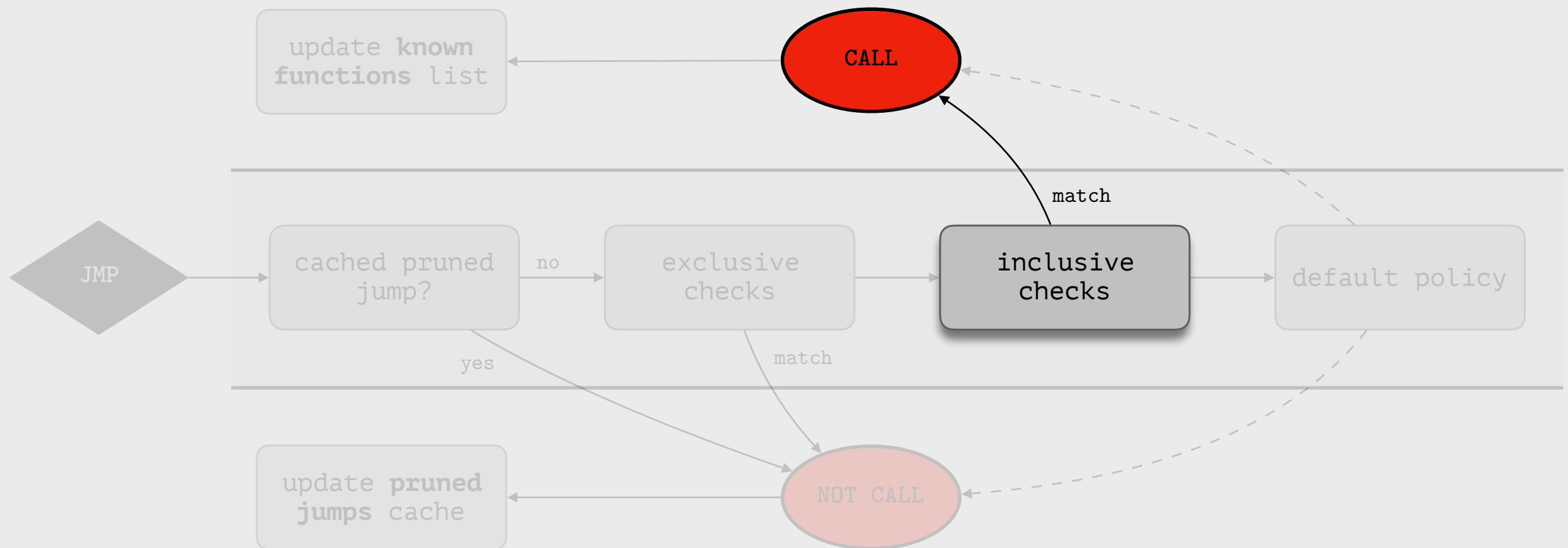
STACK INCONSISTENCY (IDEA)

A jump occurring when %ebp ≠ %esp **is not a (tail) call**.

```
push    %rbp
mov     %rsp,%rbp
sub     $0x30,%rsp
push    …
pop     …
push    …
pop     …
push    …
jmp     XXX
```

```
push    %rbp
mov     %rsp,%rbp
sub     $0x30,%rsp
push    …
push    …
pop     …
pop     …
pop     …
jmp     XXX
```

4 push / 2 pop => NOT A CALL

3 push / 3 pop => MAYBE A CALL

inclusive checks

update **known functions** list

CALL

match

JMP

cached pruned jump?

no

exclusive checks

inclusive checks

default policy

yes

match

update **pruned jumps** cache

NOT CALL

KNOWN ENTRY POINT

A jump to a known entry point **is a call**.

EXTERNAL JUMP BEFORE

A jump to a location that is before the
current entry point **is a call**.

CROSS ENTRY POINT

A jump crossing a known entry point **is a call**.

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EXPERIMENTS

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introducing iCi

METHODOLOGY

ORACLE

INPUT: entry points of functions

ALGORITHM: for each non-sequential update of %rip, check its new value. If entry point => call.

REMARKS: existing corner cases, but best oracle we have

JCALL

ALGORITHM: CALL instructions + JMP to .plt => calls

REMARKS: low overhead, no false positive

JMP

ALGORITHM: CALL + JMP instructions => calls

REMARKS: high overhead, no false negative

BENCHMARK

BINARIES

- coreutils: 98 programs

- binutils: 13 programs

- evince: 1 program

- ffmpeg: 1 program

- SPEC CPU2006 (C/C++): 19 programs

OPTIMIZATION LEVEL

for *\{SPEC CPU2006}: -O0 to -O3 with gcc-6.0

COMPILATION

for coreutils: -O0 to -O3 with gcc-6.0 and clang-3.8

RESULTS

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F-SCORE

	-O0			-O1			-O2			-O3		
	jcall	**jmp**	**iCi**	**jcall**	**jmp**	**iCi**	**jcall**	**jmp**	**iCi**	**jcall**	**jmp**	**iCi**
binutils	1,000	0,223	1,000	1,000	0,222	1,000	0,999	0,222	1,000	1,000	0,221	1,000
evince	0,988	0,831	0,991	0,980	0,796	0,985	0,936	0,803	0,985	0,947	0,818	0,986
coreutils	0,998	0,518	0,999	0,998	0,352	0,998	0,946	0,405	0,998	0,997	0,351	0,998
ffmpeg	**0,948**	**0,227**	**0,997**	**0,919**	**0,183**	**0,997**	**0,874**	**0,182**	**0,997**	**0,884**	**0,181**	**0,997**
SPEC2006 jcall worst	nc	nc	nc	nc	nc	nc	0,848	0,432	1,000	nc	nc	nc
SPEC2006 iCi worst	nc	nc	nc	nc	nc	nc	0,906	0,260	0,947	nc	nc	nc

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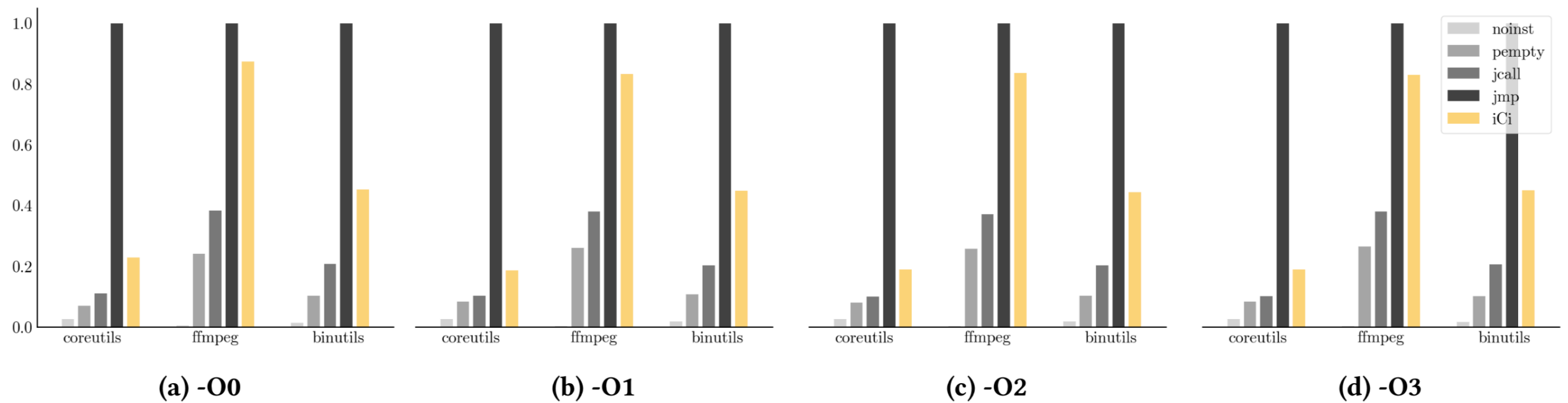
OVERHEAD

Figure 2: Overhead of each analysis in comparison with `jmp` for `-O2`

COMPILATION

	-O0			-O1			-O2			-O3		
	jcall	jmp	iCi	jcall	jmp	iCi	jcall	jmp	iCi	jcall	jmp	iCi
gcc-6.3	0,998	0,518	0,999	0,998	0,352	0,998	0,946	0,405	0,998	0,997	0,351	0,998
clang-3.8	0,997	0,360	0,998	0,997	0,359	0,998	0,997	0,359	0,998	0,997	0,359	0,998

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FINAL WORDS

time to wake everyone up

PROBLEM NOT ADDRESSED BEFORE (TO OUR KNOWLEDGE)

– **TOOLS**: No high-level API to instrument every call
(PIN, DynamoRIO, etc.)

– **RESEARCH**: A lot of papers about tail calls… but none
about detecting them dynamically

RESULTS INTERPRETATION

iCi is cool! It achieves great f-scores on C/C++

programs, and works with optimization levels and
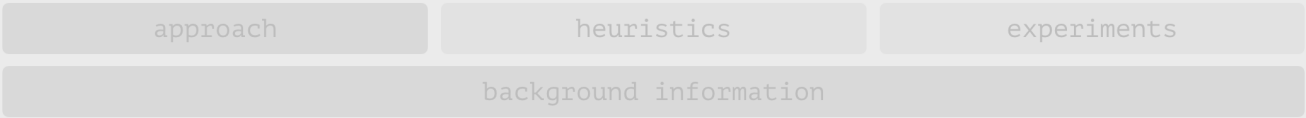
different compilers.

BUT...

RESULTS INTERPRETATION

Is it really worth it?

A simpler implementation (jcall) performs almost as

good, with a lower overhead.

YES if it is critical to catch all calls

NO otherwise

OPENING

background information

AVAILABLE INFO FOR EACH JMP

information	accessible?
target of the JMP	ok (dynamic)
current %rip, %rsp	ok (dynamic)
current entry point	~ok (depend on the known entry points)
state of stack when last call occurred	~ok (idem)
return locations for the current function	not at first call; then ~ok

— precision over time
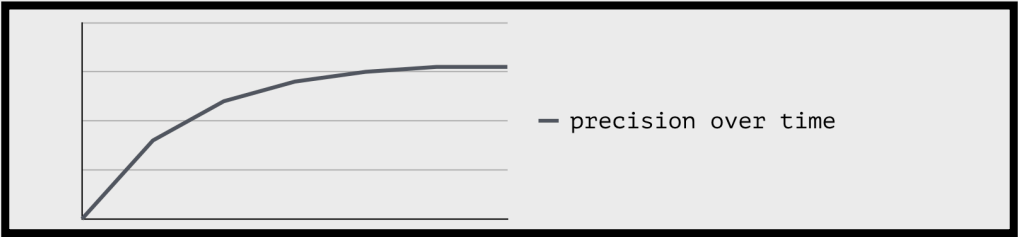
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+ overhead decreases over execution

OPENING

BETTER OVER TIME: keep learnt info from an execution to

another > should reduce overhead significantly

MIXED ANALYSIS: first static analysis (e.g. Nucleus)

https://github.com/Frky/iCi

IMPLEMENTATION	BENCHMARK (EXCEPT SPEC)

EVERYTHING NEEDED TO REPLAY