Evil Corp [UNSOLVED :-(]

Mathias Tausig (TwentyOneCool)

You were called by the Incident Response Team of Evil Corp, the urgently need your help. Somebody broke into the main server of the company, bricked the device and stole all the files! Nothing is left! This should have been impossible. The Hacker used some secret backdoor to bypass authentication. Without the knowledge of the secret backdoor other servers are at risk as well! The Incident Response Team has a full packet capture of the incident and performed an emergency cold boot attack on the server to retrieve the contents of the memory (its a really important server, Evil Corp is always ready for such kinds of incidents. However they were unable to retrieve much information from the RAM, what's left is only some parts of the "key_block" of the TLS server. Can you help Evil Corp to analyze the exploit the attacker used? (Flag is inside of the attackers' secret message).

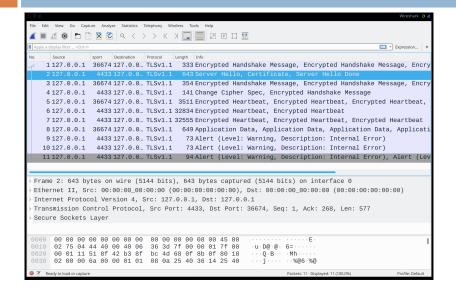
Task II

Keyblock

TT = Could not recover

File evilcorp.pcapng

Capture



Skills

Skills needed for this challenge:

- Basic cryptography knowledge
- Basic TLS knowledge
- RFC reading

Skills

Skills needed for this challenge:

- Basic cryptography knowledge
- Basic TLS knowledge
- RFC reading

RFC 4346¹

The Transport Layer Security (TLS) Protocol Version 1.1 April 2006

¹https://tools.ietf.org/html/rfc4346

What is a key_block?

6.3. Key Calculation

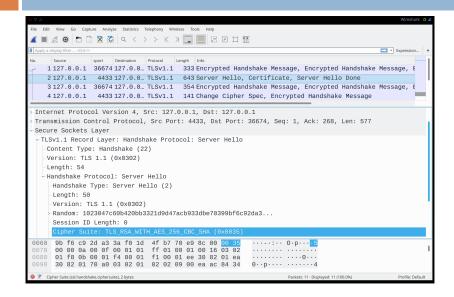
The Record Protocol requires an algorithm to generate keys, and MAC secrets from the security parameters provided by the handshake protocol.

The master secret is hashed into a sequence of secure bytes [...]

Then the **key_block** is partitioned as follows:

```
client_write_MAC_secret[SP.hash_size]
server_write_MAC_secret[SP.hash_size]
client_write_key[SP.key_material_length]
server_write_key[SP.key_material_length]
```

SecurityParameters



key_block I

Security Parameters

Cipher: AES-256 → 32 byte key

MAC: HMAC-SHA1 \rightarrow 20 byte hash size

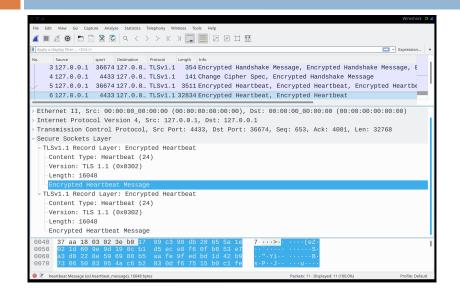
Available data

key_block II

Secrets

- client_write_MAC_secret: 6B 4F 93 6A TT TT TT TT TT TT 00 D9 F2 9B 4C B0 2D 88 36 CF
- server_write_MAC_secret: B0 CB F1 A6 7B 53 B2 00 B6 D9 DC EF 66 E6 2C 33 5D 89 6A 92
- server_write_key: 94 TT 0C EB 50 8D 81 C4 E4 40 B6 26 DF E3 40 9A 6C F3 95 84 E6 C5 86 40 49 FD 4E F2 A0 A3 01 06

Encrypted data



TLS records

```
struct {
    ContentType type;
    ProtocolVersion version;
    uint16 length;
    select (CipherSpec.cipher_type) {
        case stream: GenericStreamCipher;
        case block: GenericBlockCipher;
    } fragment;
} TLSCiphertext;
block-ciphered struct {
    opaque IV[CipherSpec.block_length];
    opaque content[TLSCompressed.length];
    opaque MAC[CipherSpec.hash_size];
    uint8 padding[GenericBlockCipher.padding_length];
    uint8 padding_length;
} GenericBlockCipher;
```

Brute force the server_write_key

Decrypt all messages sent by the server

- Brute force the server_write_key
 - Validate decryption
- Decrypt all messages sent by the server

- Brute force the server_write_key
 - Validate decryption
 - Validate padding
 - Verify MAC
- Decrypt all messages sent by the server

- Brute force the server_write_key
 - Validate decryption
 - Validate padding
 - Verify MAC
- Decrypt all messages sent by the server

6.2.3. Record Payload Protection

The MAC is generated as:

```
HMAC_hash(MAC_write_secret, seq_num +
TLSCompressed.type + TLSCompressed.version +
TLSCompressed.length + TLSCompressed.fragment));
```

- Brute force the server_write_key
 - Validate decryption
 - Validate padding
 - Verify MAC
- Decrypt all messages sent by the server

6.2.3. Record Payload Protection

```
The MAC is generated as:
```

```
HMAC_hash(MAC_write_secret, seq_num +
TLSCompressed.type + TLSCompressed.version +
TLSCompressed.length + TLSCompressed.fragment));
```

Problem: seq_num is implicit, not sent over the wire.

- Brute force the server_write_key
 - Validate decryption
 - Validate padding
 - Verify MAC
- Decrypt all messages sent by the server

Pitfall

TLS padding \neq PKCS#7 padding

Pitfall

TLS padding \neq PKCS#7 padding

PKCS#7 padding

Oc ff 3d 5b \rightarrow 0c ff 3d 5b 04 04 04 04

Pitfall

TLS padding \neq PKCS#7 padding

PKCS#7 padding

Oc ff 3d 5b \rightarrow 0c ff 3d 5b 04 04 04 04

TLS padding

Oc ff 3d 5b \rightarrow 0c ff 3d 5b 03 03 03

```
def remove_tls_padding(data):
    """Remove TLS padding which contains n bytes of
    value n-1 (n >= 1)"""
    padding_length = data[-1]
    padding = data[-(padding_length+1):]
    for padding_byte in padding:
        if(padding_byte != padding_length):
            raise ValueError("Bad padding")
    unpadded = data[:-(padding_length+1)]
    return unpadded
```

```
def decrypt_block(secret, block):
    try:
        iv = block[:16]
        ciphertext = block[16:]
        cipher = AES.new(secret, AES.MODE_CBC, iv)
        decryption = cipher.decrypt(ciphertext)
        plaintext = remove_tls_padding(decryption)
        # Last 20 bytes of the plaintext are the MAC
        plaintext = plaintext[:-20]
        # No compression is used
        return plaintext
    except ValueError:
        return None
```

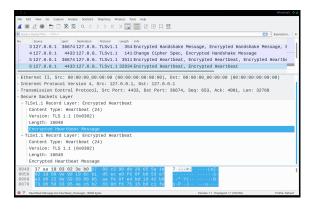
```
decrypted = {}
for key_byte in range(256):
    key_cand = key[:byte_to_bruteforce] +
        bytes([key_byte])+key[byte_to_bruteforce+1 :]
    block = application_data[0]
    plaintext = decrypt_block(key_cand, block)
    if(plaintext):
        decrypted[key_candidate] = []
        decrypted[key_candidate].append(plaintext)
        # Valid decryption for first block
        # -> check if we can decrypt the next blocks
        for i in range(1, len(application_data)):
            block = application_data[i]
            plaintext = decrypt_block(key_cand,block)
            if(plaintext):
                decrypted[key_cand].append(plaintext)
```

```
# The correct key should decrypt all blocks
valid_keys = list(filter(
    lambda k:len(decrypted[k]) == len(application_data)
    decrypted.keys()))
if(len(valid_keys) != 1):
    print('ERROR: Key not found.')
    exit(1)
key = valid_keys[0]
print("Decrpytion worked for key '{0}':".format(key))
for i, plaintext in enumerate(decrypted[key]):
    print('## Record {0} plaintext ##\n{1}')
        .format(i, plaintext))
```

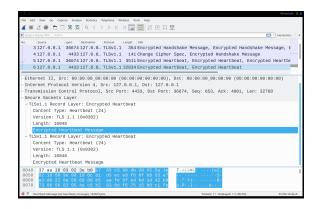
What's in it

 $[...]00 \times 00x / 00x /$ \x00\x00----BEGIN RSA PRIVATE KEY----\nMIIB3gIBAA J1AK2H8Iak4azSVdHXcySgXqfSUPKF86beNbnwfF0IOt1RZmd0J bgz\nUyglXntWL5RNVcVv8IT0MW/cnj9bAJ/v1lAVpcoijJTj/T XGq6g+p0IIAKNFSKo2\npdQOPHSWxlvbyGTo8WECAwEAAQJkJj9 5P2QmLb5qlgbj5SXH1zufBeWKb7Q4qVQd\nRTAkMVXYuWK7UZ9W a9nYulyjvg9RoWOO+SaDNqhiTWKosQ+ZrvG3A1TDMcVZSkPx\nb XCuhhRpp4j0T9levQi0s8tR1YuFzVFi8QIzANNLrgK2Y0JiDlyu 78t/eVbBey4m\nuh2xaxvEd8xGX4bIBlTuWlKIqwPNxE8fygmv4 uHFAjMAOj7Uk1ThY+UCYdeCm4/P\neVqkPYu7jNTHG2TGr/B6hs txyFpXBlq6MJQ/qPdRXLkLFuOCMwCf/OLCTQPpBiQn\ny5HoPRp MNW4m0M4F46vdN5MaCoMUU+pvbpbXfYI3/BrTapeZZCNfnQIzAJ 7 XzW9K\nj8cTPIuDcS/qpQvAiZneOmKaV5vAtcQzYb75cgu3BUz NuyH8v2P/Br+RJmm5AjMA\njp9N+xdEm4dW51lyUp6boVU6fxZi $mfYRfYANU2bVFmbsSAU9jzjWbOBuXexKKcX7 \nXGo = \n----EN$ D RSA PRIVATE KEY----\n\x00\x00\x00\x00\x00\x00 0\x00\x00[...]

Why is it there?



Why is it there?

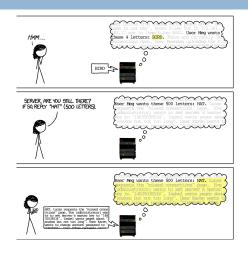




Heartbleed I

HOW THE HEARTBLEED BUG WORKS: SERVER, ARE YOU STILL THERE? IF SO, REPLY "POTATO" (6 LETTERS). Jser Meg wants these 6 letters: POTATO. da wants pages about "irl games". Unlock ecure records with master key 5130985733 Jser Meg wants these 6 letters: POTATO. SERVER, ARE YOU STILL THERE? IF SO, REPLY "BIRD" (4 LETTERS). 7 are in /tmp/files-3843. User Meg wants these 4 letters: BIRD. There are currently

Heartbleed II



Source: https://xkcd.com/1354/

Heartbleed III

Facts

- Heartbeat is a TLS extension²
- Meant to allow the usage of keep-alive functionality without performing a renegotiation
- Buffer overread vulnerablility in the openssl implementation
- Published: April 2014
- Fixed: openssl 1.0.1g, 1.0.2-beta2
- Led to numerous openssl forks/reimplementations (libressl, BoringSSL, ...)
- CVE-2014-0160

²RFC 6520: Transport Layer Security (TLS) and Datagram Transport Layer Security (DTLS) Heartbeat Extension (2012)

And now???



 $Source: CC-BY-SA, Takkk, \\ https://de.wikipedia.org/wiki/Datei:Basalt_house-wall_-_Matraszentistvan, _Hungary.jpg$

Ideas I

Key Exchange

- RSA key exchange was used
- Private key is key for certificate used in the handshake
- This should allow us to decrypt whole conversation
- Unfortunately, we need the ClientKeyExchange message from the TLS handshake, and it's not in the capture

Ideas II

Missing keys

Searching for the remaining unknown bytes of the *client_write_key* in the decrypted data was not successful.

Secure Renegotiation

- A second TLS extension is used in the handshake: Secure Renegotiation (RFC 5746: Transport Layer Security (TLS) Renegotiation Indication Extension, 2010)
- That's why the handshake messages sent by the client (containing the ClientKeyExchange) in the first frame are encrypted
- No decrypted handshake messages found
- Maybe the previous session keys are in the decrypted data?

Ideas IV

Finished

- We can decrypt the Finished message sent by the server at the end of the handshake
- It contains only the verify_data (RFC, 7.4.9, Finished)
- verify_data = PRF(master_secret, finished_label,
 MD5(handshake_messages) +
 SHA-1(handshake_messages)) [0..11];

Ideas V

Other data

- One more block with readable content found in the plaintext
 - (Part of) A certificate?
 - Looking for a DER encoded sequence (0x8081, 0x8082) reveals another 800 bit public key

Problems & Fixes

- \blacksquare Not upgrading openssl (and the rest of the system?) for 5 years \to upgrade system
- \blacksquare Use of deprecated RSA key exchange in TLS handshake \to use ECDHE
- (Use of deprecated TLS version \rightarrow use TLS v1.3)
- \blacksquare (Use of deprecated cipher suite \to use modern cipher suite with AEAD)

The End!