Pontificia Universidad Católica del Perú - FCI

XieXieLucas Notebook - Froz/Phibrain/Ands

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1 Chinese Remainder Theorem

```
/**
 * Chinese remainder theorem.
 * Find z such that z % x[i] = a[i] for all i.
 * */
long long crt(vector<long long> &a, vector<long long> &x) {
 long long z = 0;
 long long n = 1;
 for (int i = 0; i < x.size(); ++i)
    n *= x[i];</pre>
```

```
for (int i = 0; i < a.size(); ++i) {
  long long tmp = (a[i] * (n / x[i])) % n;
  tmp = (tmp * mod_inv(n / x[i], x[i])) % n;
  z = (z + tmp) % n;
}
return (z + n) % n;</pre>
```

2 Cribas

3 Euler Totient

4 Inverso Modular

```
/** Inverso Modular **/
```

```
#define MAX 100
#define MOD 1000000009

long long inverso[MAX];

void inv(){
        inverso[1] = 1;
        REP(i,2,MAX) inverso[i] = ( (MOD-MOD/i) * inverso[MOD%i] ) % MOD;
}
```

5 Miller Rabin

```
//Miller-Rabin primality test
ll pow(ll a, ll b, ll c){
   11 \text{ ans} = 1;
   while(b){
       if(b\&1) ans = (1LL*ans*a)\%c;
       a = (1LL*a*a)%c;
       b >>=1;
   return ans;
bool miller(ll p, ll it = 10){
   if(p<2) return 0;</pre>
   if(p!=2 && (p&1) == 0) return 0;
   ll s=p-1;
   while((s&1) == 0) s>>=1;
   while(it--){
       ll a = rand()%(p-1)+1, temp = s;
       11 mod = pow(a,temp,p);
       while(temp!= p-1 && mod!=1 && mod!=p-1){
           mod = (1LL*mod*mod)%p;
           temp<<=1;
       }
       if(mod!=p-1 && (temp&1) == 0) return 0;
   return 1;
```

6 Number Theory

```
// Encuentra el menor positivo de la forma ax+ b , my+ n
// ( x,y enteros no necesariamente postivos)
// Si son positivos, hallar los coeficientes y sumar lo
// que falta para que de positivo
// d: mcd(a,b) (d > 0)
// x,y enteros tales que a*x + b*y = d
// las demas soluciones son ( x + (b/d)t , y - (a/d)t )
void gcdextend(ll a, ll b, ll &x, ll &y, ll &d){
   if(b == 0){
      if(a>0) x = 1, y = 0, d = a;
      else x = -1 , y = 0 , d = -a ;
      return ;
   gcdextend(b,a%b,x,y,d) ;
   11 x1 = y, y1 = x - (a/b)*y;
   x = x1, y = y1;
}
// menor positivo que es u modulo modPos
inline 11 ADDTOPOSITIVE(11 u, 11 modPOS){
   if(modPOS < 0) modPOS = -modPOS;</pre>
   if(u >= 0) return u%modPOS;
   u = -u:
   if(u%modPOS == 0) return 0;
   return modPOS*((u/modPOS)+1) - u ;
}
// Encuentra el menor positivo que es
// de la forma ax + b , my + n
// los demas son de la forma ans + ((a/d)*m)*t
// retorna -1 si no hay solucion ( mcd(a,m) no divide a n - b )
inline 11 FINDmenorLCS(11 a,11 b,11 m,11 n){
    //Cuidado con el caso a = 0 , m = 0 ,
          //porque es solo verificar b = n
    a = abs(a) ; m = abs(m) ;
    if( a == 0 ) {
       swap(a,m);
       swap(b,n);
    }
```

```
if(m == 0){
       if( (n-b) % a == 0) return n ;
       else return -1;
    ll x , y , d ;
     gcdextend(a,m,x,y,d);
    if((n-b)\%d != 0) return -1 ;
    ll temp = a*x*((n-b)/d) + b;
    temp = ADDTOPOSITIVE(temp,a*(m/d)) ;
    return temp ;
}
//Finds partition of n (number of ways to obtain n as a sum of positive
    numbers)
int partition(int n) {
 int[] dp = new int[n + 1];
 dp[0] = 1;
 for (int i = 1; i <= n; i++) {</pre>
   for (int j = 1, r = 1; i - (3 * j * j - j) / 2 >= 0; j++, r *= -1) {
     dp[i] += dp[i - (3 * j * j - j) / 2] * r;
     if (i - (3 * j * j + j) / 2 >= 0) {
       dp[i] += dp[i - (3 * j * j + j) / 2] * r;
   }
 }
 return dp[n];
```

7 Pollard Rho

```
#define MAXL (50000>>5)+1
#define GET(x) (mark[x>>5]>>(x&31)&1)
#define SET(x) (mark[x>>5] |= 1<<(x&31))

int mark[MAXL];
int P[50000], Pt = 0;
void sieve() {
    register int i, j, k;
    SET(1);
    int n = 46340;
    for (i = 2; i <= n; i++) {</pre>
```

```
if (!GET(i)) {
           for (k = n/i, j = i*k; k >= i; k--, j -= i)
               SET(j);
           P[Pt++] = i;
       }
   }
}
11 mul(unsigned 11 a, unsigned 11 b, unsigned 11 mod) {
    11 \text{ ret} = 0:
    for (a %= mod, b %= mod; b != 0; b >>= 1, a <<= 1, a = a >= mod ? a -
       if (b&1) {
           ret += a;
           if (ret >= mod) ret -= mod;
       }
   }
    return ret;
}
void exgcd(ll x, ll y, ll &g, ll &a, ll &b) {
    if (y == 0)
       g = x, a = 1, b = 0;
    else
       exgcd(y, x\%y, g, b, a), b = (x/y) * a;
}
ll inverse(ll x, ll p) {
   ll g, b, r;
    exgcd(x, p, g, r, b);
   if (g < 0) r = -r;
   return (r\%p + p)\%p;
}
ll mpow(ll x, ll y, ll mod) { // mod < 2^32
   ll ret = 1:
    while (y) {
       if (v&1)
           ret = (ret * x) \% mod;
       y >>= 1, x = (x * x) \mod;
   }
    return ret % mod;
}
11 mpow2(11 x, 11 y, 11 mod) {
    ll ret = 1;
    while (y) {
```

```
if (y&1)
           ret = mul(ret, x, mod);
       v >>= 1. x = mul(x. x. mod):
   return ret % mod;
}
int isPrime(ll p) { // implements by miller-babin
   if (p < 2 || !(p&1))
                             return 0;
   if (p == 2)
                                     return 1;
   11 q = p-1, a, t;
   int k = 0, b = 0;
   while (!(q\&1)) q >>= 1, k++;
   for (int it = 0; it < 2; it++) {</pre>
       a = rand()\%(p-4) + 2;
       t = mpow2(a, q, p);
       b = (t == 1) \mid \mid (t == p-1);
       for (int i = 1; i < k && !b; i++) {</pre>
           t = mul(t, t, p);
           if (t == p-1)
              b = 1;
       }
       if (b == 0)
           return 0;
   return 1;
}
ll pollard_rho(ll n, ll c) {
   11 x = 2, y = 2, i = 1, k = 2, d;
   while (true) {
       x = (mul(x, x, n) + c);
       if (x >= n) x -= n;
       d = \_gcd(x - y, n);
       if (d > 1) return d;
       if (++i == k) v = x, k <<= 1;
   }
   return n;
void factorize(int n, vector<ll> &f) {
   for (int i = 0; i < Pt && P[i]*P[i] <= n; i++) {</pre>
       if (n%P[i] == 0) {
               while (n\%P[i] == 0)
                      f.push_back(P[i]), n /= P[i];
```

```
}
   if (n != 1) f.push back(n):
void llfactorize(ll n, vector<ll> &f) {
   if (n == 1)
       return :
   if (n < 1e+9) {
       factorize(n, f);
       return :
   }
   if (isPrime(n)) {
       f.push_back(n);
       return ;
   }
   11 d = n:
   for (int i = 2; d == n; i++)
       d = pollard_rho(n, i);
   llfactorize(d, f);
   llfactorize(n/d, f);
}
vector<ll> f;
map<ll, int> r;
int main() {
   sieve();
   11 n;
   scanf("%11d", &n);
   llfactorize(n, f):
   for (auto &x : f) r[x]++;
   ll last:
   for (auto it = r.begin(); it != r.end(); it++) {
       if (it != r.begin()) printf(" ");
       last = it -> first;
       printf("%lld", last);
       if (it->second > 1){
           for( int i = 0 ; i < it->second - 1 ; ++i ) printf(" %lld",last
               );
       }
   return 0;
```

8 Simplex Method

```
// Two-phase simplex algorithm for solving linear programs:
      maximize c^T x
11
      subject to Ax <= b
                 x >= 0
// INPUT: A -- an m x n matrix
        b -- an m-dimensional vector
         c -- an n-dimensional vector
        x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
         above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
 int m, n;
 VI B, N;
 VVD D;
       LPSolver(const VVD &A, const VD &b, const VD &c) :
              m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VD(n+2))  {
              for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
                     D[i][i] = A[i][i];
              for (int i = 0; i < m; i++)</pre>
                     B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i];
              for (int j = 0; j < n; j++)
                     N[i] = i; D[m][i] = -c[i];
              N[n] = -1; D[m+1][n] = 1;
       }
       void Pivot(int r, int s) {
              for (int i = 0; i < m+2; i++) if (i != r)
                     for (int j = 0; j < n+2; j++) if (j != s)
                             D[i][j] = D[r][j] * D[i][s] / D[r][s];
              for (int j = 0; j < n+2; j++) if (j != s)
                     D[r][i] /= D[r][s];
```

```
for (int i = 0; i < m+2; i++) if (i != r)</pre>
              D[i][s] /= -D[r][s];
       D[r][s] = 1.0 / D[r][s];
       swap(B[r], N[s]);
}
bool Simplex(int phase) {
       int x = phase == 1 ? m+1 : m;
       while (true) {
              int s = -1;
               for (int i = 0: i <= n: i++) {
                      if (phase == 2 && N[j] == -1) continue;
                      if (s == -1 || D[x][j] < D[x][s] ||
                               D[x][j] == D[x][s] && N[j] < N[s]) s
                                    = j;
               }
              if (D[x][s] >= -EPS) return true;
               int r = -1;
               for (int i = 0; i < m; i++) {</pre>
                      if (D[i][s] <= 0) continue;</pre>
                      if (r == -1 ||
                               D[i][n+1] / D[i][s] < D[r][n+1] / D[r</pre>
                                   ][s] ||
                       D[i][n+1] / D[i][s] == D[r][n+1] / D[r][s
                            7 &&
                               B[i] < B[r]) r = i;
}
               if (r == -1) return false;
               Pivot(r, s);
       }
DOUBLE Solve(VD &x) {
       int r = 0:
       for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r =
       if (D[r][n+1] <= -EPS) {</pre>
               Pivot(r, n);
               if (!Simplex(1) || D[m+1][n+1] < -EPS)</pre>
                      return -numeric_limits<DOUBLE>::infinity();
               for (int i = 0; i < m; i++) if (B[i] == -1) {
                      int s = -1:
                      for (int j = 0; j <= n; j++)</pre>
                              if (s == -1 || D[i][j] < D[i][s] ||</pre>
```

```
D[i][j] == D[i][s] && N[j] < N
                                                   [s]) s = i;
                              Pivot(i, s);
                      }
               }
               if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
               x = VD(n);
               for (int i = 0; i < m; i++) if (B[i] < n)
                      x[B[i]] = D[i][n+1]:
               return D[m][n+1];
       }
};
int main() {
  const int m = 4:
  const int n = 3;
 DOUBLE A[m][n] = {
   \{ 6, -1, 0 \},
   \{-1, -5, 0\},\
   { 1, 5, 1 },
   \{-1, -5, -1\}
  DOUBLE _b[m] = \{ 10, -4, 5, -5 \};
 DOUBLE _c[n] = \{ 1, -1, 0 \};
  VVD A(m):
  VD b(_b, _b + m);
  VD c(_c, _c + n);
  for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);</pre>
  LPSolver solver(A, b, c);
  VD x;
  DOUBLE value = solver.Solve(x);
  cerr << "VALUE: "<< value << endl;</pre>
  cerr << "SOLUTION:";</pre>
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
  cerr << endl:
 return 0;
```

9 Teorema de Lucas