Pontificia Universidad Católica del Perú - FCI

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1 Bipartite Matching

2 Dinic Flow

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
//
// Running time:
// O(|V|^2 |E|)
//
// INPUT:
// - graph, constructed using AddEdge()
// - source and sink
//
// OUTPUT:
// - maximum flow value
// - To obtain actual flow values, look at edges with capacity > 0
// (zero capacity edges are residual edges).
```

```
typedef long long LL;
struct Edge {
 int u, v;
 LL cap, flow;
 Edge() {}
 Edge(int u, int v, LL cap): u(u), v(v), cap(cap), flow(0) {}
struct Dinic {
 int N:
 vector<Edge> E;
 vector<vector<int>> g;
 vector<int> d, pt;
 Dinic(int N): N(N), E(O), g(N), d(N), pt(N) {}
 void AddEdge(int u, int v, LL cap) {
   if (u != v) {
     E.emplace_back(Edge(u, v, cap));
     g[u].emplace_back(E.size() - 1);
     E.emplace_back(Edge(v, u, 0));
     g[v].emplace_back(E.size() - 1);
 }
 bool BFS(int S, int T) {
   queue<int> q({S});
   fill(d.begin(), d.end(), N + 1);
   d[S] = 0;
   while(!q.empty()) {
     int u = q.front(); q.pop();
     if (u == T) break;
     for (int k: g[u]) {
       Edge &e = E[k];
      if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
        d[e.v] = d[e.u] + 1;
        q.emplace(e.v);
      }
     }
   return d[T] != N + 1;
 LL DFS(int u, int T, LL flow = -1) {
```

```
if (u == T || flow == 0) return flow;
   for (int &i = pt[u]; i < g[u].size(); ++i) {</pre>
     Edge &e = E[g[u][i]];
     Edge &oe = E[g[u][i]^1];
     if (d[e.v] == d[e.u] + 1) {
       LL amt = e.cap - e.flow;
       if (flow != -1 && amt > flow) amt = flow;
       if (LL pushed = DFS(e.v, T, amt)) {
         e.flow += pushed;
         oe.flow -= pushed;
         return pushed;
     }
   return 0;
 LL MaxFlow(int S, int T) {
   LL total = 0;
   while (BFS(S, T)) {
     fill(pt.begin(), pt.end(), 0);
     while (LL flow = DFS(S, T))
       total += flow;
   }
   return total;
 }
};
```

3 Edmonds Blossom

```
// Maximum general matching (not necessarily bipartite)
// Make sure to set N in main()
// Claimed O(N^4) running time

int N; // the number of vertices in the graph
typedef vector<int> vi;
typedef vector< vector<int> > vvi;
vi match;
vi vis;

void couple(int n, int m) { match[n]=m; match[m]=n; }
```

```
// True if augmenting path or a blossom (if blossom is non-empty).
// the dfs returns true from the moment the stem of the flower is
// reached and thus the base of the blossom is an unmatched node.
// blossom should be empty when dfs is called and
// contains the nodes of the blossom when a blossom is found.
bool dfs(int n, vvi &conn, vi &blossom) {
 vis[n]=0;
 REP(i, 0, N) if(conn[n][i]) {
   if(vis[i]==-1) {
     vis[i]=1;
     if(match[i] == -1 || dfs(match[i], conn, blossom)) {
                      couple(n,i);
                     return true;
              }
   }
   if(vis[i]==0 || SZ(blossom)) { // found flower
     blossom.pb(i); blossom.pb(n);
     if(n==blossom[0]) { match[n]=-1; return true; }
     return false:
   }
 }
 return false;
}
// search for an augmenting path.
// if a blossom is found build a new graph (newconn) where the
// (free) blossom is shrunken to a single node and recurse.
// if a augmenting path is found it has already been augmented
// except if the augmented path ended on the shrunken blossom.
// in this case the matching should be updated along the
// appropriate direction of the blossom.
bool augment(vvi &conn) {
       REP(m, 0, N) if(match[m] == -1) {
              vi blossom:
              vis=vi(N,-1);
              if(!dfs(m, conn, blossom)) continue;
              if(SZ(blossom)==0) return true; // augmenting path found
// blossom is found so build shrunken graph
              int base=blossom[0], S=SZ(blossom);
              vvi newconn=conn;
              REP(i, 1, S-1) REP(j, 0, N)
                     newconn[base][j]=newconn[j][base]|=conn[blossom[i]][
                          i];
              REP(i, 1, S-1) REP(j, 0, N)
```

```
newconn[blossom[i]][j]=newconn[j][blossom[i]]=0;
             newconn[base][base]=0; // is now the new graph
             if(!augment(newconn)) return false;
             int n=match[base];
// if n!=-1 the augmenting path ended on this blossom
   if(n!=-1) REP(i, 0, S) if(conn[blossom[i]][n]) {
     couple(blossom[i], n);
     if(i&1) for(int j=i+1; j<S; j+=2)</pre>
                    couple(blossom[j],blossom[j+1]);
     else for(int j=0; j<i; j+=2)</pre>
                    couple(blossom[j],blossom[j+1]);
     break;
   return true;
 return false;
// conn is the NxN adjacency matrix
// returns the number of edges in a max matching.
int edmonds(vvi &conn) {
 int res=0;
 match=vi(N.-1):
 while(augment(conn)) res++;
 return res:
}
set<pair<int,int> > used;
int main(){
 int n:
 cin >> n;
 N = n:
 vvi conn;
 vi tmp;
 tmp.assign(n,0);
 REP(i, 0, n) conn.push_back(tmp);
 int u, v;
 while(cin >> u >> v){
   u--; v--;
   if(u > v) swap(u,v);
   if(used.count(make_pair(u,v))) continue;
   used.insert(make_pair(u,v));
   conn[u][v] = conn[v][u] = 1;
```

```
int res = edmonds(conn);

cout<<res*2<<endl;

REP(i, 0, n) {
    if(match[i] > i){
        cout<<ii+1<<" "<<match[i] + 1<<endl;
    }
}

return 0;
}
</pre>
```

4 Min Cost Max Flow

```
const int MAXN = 5010;
const 11 INF = 1e15;
struct edge { int dest;ll origcap, cap; ll cost; int rev; };
struct MinCostMaxFlow {
   vector<edge> adj[MAXN];
   11 dis[MAXN], cost;
   int source, target, iter;
   11 cap;
   edge* pre[MAXN];
   int queued[MAXN];
   MinCostMaxFlow (){}
   void AddEdge(int from, int to, 11 cap, 11 cost) {
       adj[from].push_back(edge {to, cap, cap, cost, (int)adj[to].size()})
       adj[to].push_back(edge {from,0, 0, -cost, (int)adj[from].size()
            - 1});
   }
   bool spfa() {
       REP(i,0,MAXN) queued[i] = 0;
       fill(dis, dis + MAXN, INF);
       queue<int> q;
       pre[source] = pre[target] = 0;
       dis[source] = 0;
       q.emplace(source);
       queued[source] = 1;
```

```
while (!q.empty()) {
           int x = q.front();
           11 d = dis[x];
           q.pop();
           queued[x] = 0;
           for (auto& e : adj[x]) {
              int y = e.dest;
              11 w = d + e.cost;
              if (e.cap < 1 || dis[y] <= w) continue;</pre>
              dis[v] = w:
              pre[y] = &e;
              if(!queued[y]){
                      q.push(y);
                      queued[y] = 1;
              }
           }
       }
       edge* e = pre[target];
       if (!e) return 0;
       while (e) {
           edge& rev = adj[e->dest][e->rev];
           e \rightarrow cap -= cap;
           rev.cap += cap;
           cost += cap * e->cost;
           e = pre[rev.dest];
       }
       return 1;
   pair<11,11> GetMaxFlow(int S, int T) {
       cap = 1, source = S, target = T, cost = 0;
       while(spfa()) {}
       11 totflow = 0;
       for(auto e: adj[source]){
           totflow += (e.origcap - e.cap);
       return make_pair(totflow, cost);
};
```

5 Push Relabel Max Flow

```
// Fast O(|V|^3) flow, works for n ~ 5000 with no problem
// Actual flow values in edges with cap > 0 (0 cap = residual)
typedef long long LL;
struct Edge {
 int from, to, cap, flow, index;
 Edge(int from, int to, int cap, int flow, int index) :
   from(from), to(to), cap(cap), flow(flow), index(index) {}
};
struct PushRelabel {
 int N:
 vector<vector<Edge> > G;
 vector<LL> excess:
 vector<int> dist, active, count;
 queue<int> Q;
 PushRelabel(int N) :
       N(N), G(N), excess(N), dist(N), active(N), count(2*N) {}
  void AddEdge(int from, int to, int cap) {
   G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
   if (from == to) G[from].back().index++;
   G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
 }
       void Enqueue(int v) {
              if (!active[v] && excess[v] > 0) {
                      active[v] = true; Q.push(v);
              }
       }
  void Push(Edge &e) {
   int amt = int(min(excess[e.from], LL(e.cap - e.flow)));
   if (dist[e.from] <= dist[e.to] || amt == 0) return;</pre>
   e.flow += amt:
   G[e.to][e.index].flow -= amt:
   excess[e.to] += amt;
   excess[e.from] -= amt;
   Enqueue(e.to);
 }
```

```
void Gap(int k) {
 for (int v = 0: v < N: v++) {
   if (dist[v] < k) continue;</pre>
   count[dist[v]]--;
   dist[v] = max(dist[v], N+1);
   count[dist[v]]++;
   Enqueue(v);
 }
}
void Relabel(int v) {
 count[dist[v]]--;
 dist[v] = 2*N:
  for (int i = 0; i < G[v].size(); i++)</pre>
   if (G[v][i].cap - G[v][i].flow > 0)
     dist[v] = min(dist[v], dist[G[v][i].to] + 1);
 count[dist[v]]++:
  Enqueue(v);
     void Discharge(int v) {
     for (int i = 0; excess[v] > 0 && i < G[v].size(); i++)</pre>
             Push(G[v][i]):
     if (excess[v] > 0) {
                if (count[dist[v]] == 1) Gap(dist[v]);
                else Relabel(v):
     }
LL GetMaxFlow(int s, int t) {
  count[0] = N-1:
  count[N] = 1;
  dist[s] = N:
  active[s] = active[t] = true;
  for (int i = 0; i < G[s].size(); i++) {</pre>
   excess[s] += G[s][i].cap;
   Push(G[s][i]);
  while (!Q.empty()) {
   int v = Q.front();
   ()qoq. []
   active[v] = false;
   Discharge(v);
```

```
}
LL totflow = 0;
for (int i = 0; i < G[s].size(); i++) totflow += G[s][i].flow;
return totflow;
}
};</pre>
```