

Project 1

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List a link to your github repository here!

PROBLEM 1

$$-\frac{d^2u}{dx^2} = f(x) \tag{1}$$

- source term: $f(x) = 100e^{-10x}$

- x range $x \in [0, 1]$

- boundary conditions: $u(0) = 0$ and $u(1) = 0$

$$u(x) = 1 - (1 - e^{-10})x - e^{-10x} \tag{2}$$

Checking analytically that an exact solution to Eq. 1 is given by Eq. 2.

$$\begin{aligned} \frac{du}{dx} &= 1 - e^{-10} + 10e^{-10x} \\ \frac{d^2u}{dx^2} &= -100e^{-10x} \\ -\frac{d^2u}{dx^2} &= 100e^{-10x} \\ -\frac{d^2u}{dx^2} &= f(x) \end{aligned}$$

PROBLEM 2

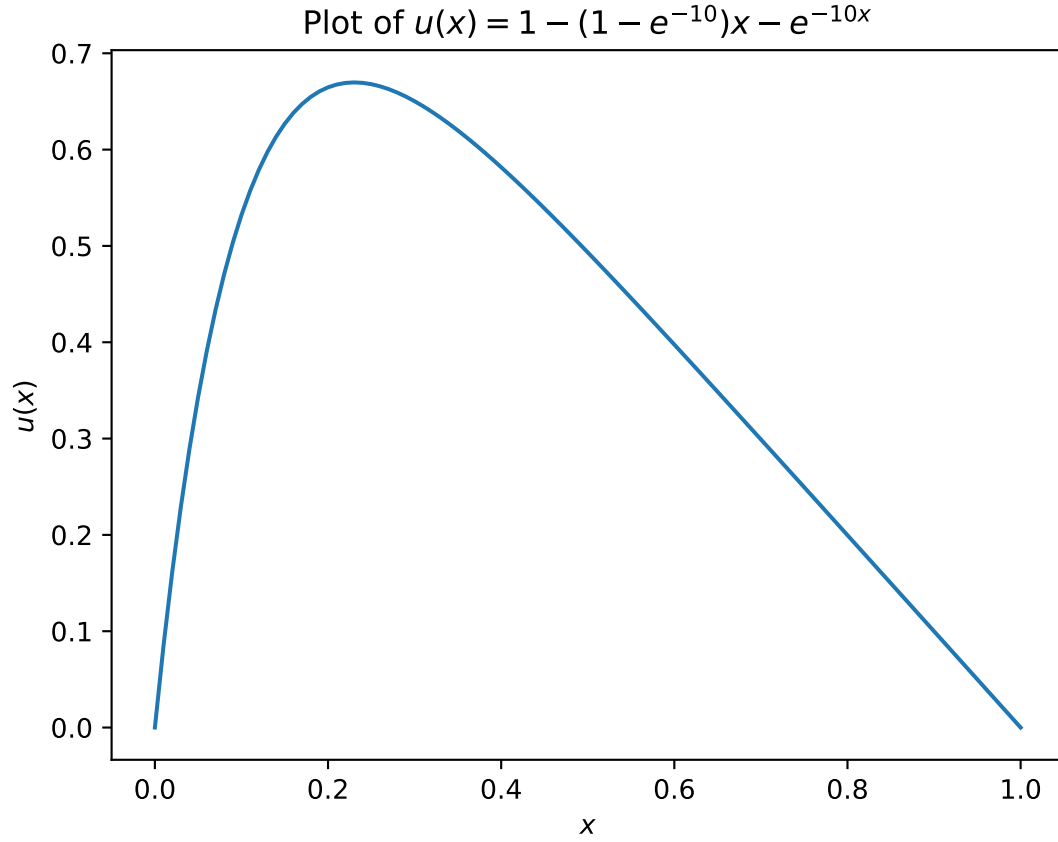


FIG. 1. Plot of $u(x)$.

PROBLEM 3

By using the Taylor approximation of the second derivative we can discretize the second derivative in the Poisson equation:

$$\frac{\partial^2 u}{\partial x^2} \approx \frac{u_{i+1} - 2u_i + u_{i-1}}{h^2} + O(h^2)$$

Here we have the stepsize h and the truncation error O . The one-dimensional Poisson equation can then be written for the approximated version of u as v like:

$$-\frac{v_{i+1} - 2v_i + v_{i-1}}{h^2} = f_i \quad (3)$$

PROBLEM 4

We can rewrite the discretized equation as a matrix equation for $n + 1$ number of points and $n - 1$ unknown points (v_0 and v_n are known) with the $n - 1 \times n - 1$ matrix \mathbf{A} . We rewrite the discretized Poisson function:

$$\begin{aligned} 2v_1 - v_2 &= f_1 h^2 \\ -v_1 + 2v_2 - v_3 &= f_2 h^2 \\ &\vdots \\ -v_{n-3} + 2v_{n-2} - v_{n-1} &= f_{n-2} h^2 \\ -v_{n-2} + 2v_{n-1} &= f_{n-1} h^2 \end{aligned}$$

This can be written in terms of the following matrix equation where we rewrite $f_i h^2$ for $i = 1, 2, \dots, n - 1$ as g_i

$$\begin{bmatrix} 2 & -1 & 0 & 0 & \dots & 0 \\ -1 & 2 & -1 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & -1 & 2 & -1 \\ 0 & 0 & \dots & 0 & -1 & 2 \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \\ \vdots \\ v_{n-2} \\ v_{n-1} \end{bmatrix} = \begin{bmatrix} g_1 \\ g_2 \\ \vdots \\ g_{n-2} \\ g_{n-1} \end{bmatrix}$$

PROBLEM 5

- a) Since the vector \vec{v}^* of length m represents a complete solution of the discretized Poisson equation it contains all values in \vec{v} in addition to the boundary conditions. The relation between n and m is therefore $m = n + 2$.
- b) Solving $\mathbf{A}\vec{v} = \vec{g}$ for \vec{v} gives us all but the first and last value in \vec{v}^* . So all but the boundary values.

PROBLEM 6

- a) We can use the Thomas algorithm to row reduce the matrix \mathbf{A} to give us a solution of the equation $\mathbf{A}\vec{v} = \vec{g}$. This is done in two steps. We first define the three diagonals as three vectors. The sub diagonal \vec{a} the diagonal \vec{b} and the superdiagonal \vec{c} where the i element in these vectors corresponds to the i row of the matrix \mathbf{A} . We define n unknowns and the matrix \mathbf{A} as $n \times n$

- i) The first step is forwards substitution. We define a number w for each step and overwrite the i element of both b and g starting at index 2. This means overwrite the values in the original vectors b and g , but at the same time we don't have to define new vectors to store new values. For large n this will reduce both the computation time and memory usage for the data machine

for $i = 2, \dots, n$:

$$\begin{aligned} w &= \frac{a_i}{b_{i-1}} \\ b_i &= b_i - w c_{i-1} \\ g_i &= g_i - w g_{i-1} \end{aligned}$$

- ii) The second and last step is back substitution where we find an expression for \vec{v} . We start at our last element and work our way backwards:

$$\begin{aligned} v_n &= \frac{g_n}{b_n} \\ v_i &= \frac{g_i - c_i v_{i+1}}{b_i} \quad \text{for } i = n - 1, \dots, 1 \end{aligned}$$

- b) We find the number of FLOPs for this algorithm by counting the number of floating point operations the computer has to do. For the first step we have 3 FLOPs (1 subtraction 1 division and one multiplication) for defining \tilde{b}_i and \tilde{g}_i each. Since we loop this operation $n - 1$ times we end up with a total number of $6(n - 1)$ FLOPs.

For back substitution we have 1 FLOP calculating v_n , and three FLOPs calculating v_i $n - 1$ times. This gives us a total of $3(n - 1) + 1$ FLOPs

The total FLOPs of the general algorithm is $9n - 8$

PROBLEM 7

PROBLEM 8

PROBLEM 9

For the special case we dont need to do new computations for every i element of the vectors \vec{a} , \vec{b} and \vec{c} .

- i) The first step is forwards substitution. We have to define the first index and then loop over the rest. We define $w = b - a\tilde{c}$ to simplify the expressions:

$$\begin{aligned}\tilde{c}_1 &= \frac{c_1}{b_1} \\ \tilde{c}_i &= \frac{c_i}{w} \quad \text{for } i = 2, \dots, n - 1\end{aligned}$$

$$\begin{aligned}\tilde{g}_1 &= \frac{g_1}{b_1} \\ \tilde{g}_i &= \frac{g_i - a_i \tilde{g}_{i-1}}{w} \quad \text{for } i = 2, \dots, n\end{aligned}$$

- ii) Second and last step is back substitution where we find an expression for \vec{v} . We start at our last element and work our way backwards:

$$\begin{aligned}v_n &= \tilde{g}_n \\ v_i &= \tilde{g}_i - \tilde{c}_i v_{i+1} \quad \text{for } i = n - 1, \dots, 1\end{aligned}$$

PROBLEM 10

We write equations using the LaTeX `equation` (or `align`) environments. Here is an equation with numbering

$$\mathbf{F} = \frac{d\mathbf{p}}{dt}, \tag{4}$$

and here is one without numbering:

$$\oint_C \mathbf{F} \cdot d\mathbf{r} = 0.$$

Sometimes it is useful to refer back to a previous equation, like we're demonstrating here for equation 4.

We can include figures using the `figure` environment. Whenever we include a figure or table, we *must* make sure to actually refer to it in the main text, e.g. something like this: "In figure 2 we show ...". Also, note the LaTeX code

FIG. 2. Write a descriptive caption here that explains the content of the figure. Note the font size for the axis labels and ticks — the size should approximately match the document font size.

we used to get correct quotation marks in the previous sentence. (Simply using the " key on your keyboard will give

the wrong result.) Figures should preferably be vector graphics (e.g. a `.pdf` file) rather than raster graphics (e.g. a `.png` file).

By the way, don't worry too much about where LaTeX decides to place your figures and tables — LaTeX knows more than we do about proper document layout. As long as you label all your figures and tables and refer to them in the text, it's all good. Of course, in some cases it can be worth trying to force a specific placement, to avoid the figure/table appearing many pages away from the main text discussing it, but this isn't something you should spend time on until the very end of the writing process.

Next up is a table, created using the `table` and `tabular` environments. We refer to it by table [I](#).

| Number of points | Output |
|------------------|--------|
| 10 | 0.3086 |
| 100 | 0.2550 |

TABLE I. Write a descriptive caption here, explaining the content of your table.

Finally, we can list algorithms by using the `algorithm` environment, as demonstrated here for algorithm [1](#).

| | |
|---|--------------------|
| Algorithm 1 Some algorithm | |
| Some maths, e.g $f(x) = x^2$. | ▷ Here's a comment |
| for $i = 0, 1, \dots, n - 1$ do | |
| Do something here | |
| while Some condition do | |
| Do something more here | |
| Maybe even some more math here, e.g $\int_0^1 f(x)dx$ | |