

Understanding Sorting and MergeSort

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What is Sorting?

Sorting is the process of arranging elements in a specific order—typically **ascending** or **descending**. It is a fundamental operation in computer science, used in searching, data analysis, and optimization problems.

Common Sorting Criteria:

- Numerical order (e.g., 1, 2, 3...)
- Lexicographical order (e.g., "apple", "banana", "cherry")
- Custom orderings (e.g., by date, priority, etc.)

Why Sorting Matters

Efficient sorting:

- Improves search performance (e.g., enables binary search)
- Helps in data normalization
- Is often a preprocessing step in complex algorithms (e.g., graph algorithms, database queries)

Types of Sorting Algorithms

Sorting algorithms are generally categorized by:

- Time complexity e.g., $O(n^2)$, $O(n \log n)$
- Space complexity
- Stability (whether equal elements retain their original order)
- In-place vs out-of-place (whether extra memory is used)

Examples:

Algorithm	Time Complexity	Stable	In-Place
Bubble Sort	$O(n^2)$	Yes	Yes
Insertion Sort	$O(n^2)$	Yes	Yes
MergeSort	$O(n \log n)$	Yes	No
QuickSort	$O(n \log n)$	No	Yes
HeapSort	$O(n \log n)$	No	Yes

MergeSort: A Divide and Conquer Algorithm

Concept

MergeSort is a classic example of the **divide and conquer** paradigm. It works by:

1. Dividing the array into two halves.
2. Recursively sorting each half.
3. Merging the two sorted halves into a single sorted array.

Key Idea

Instead of sorting the entire array at once, MergeSort breaks the problem into smaller subproblems, solves them independently, and then combines the results.

Step-by-Step Breakdown

Imagine sorting the array: [38, 27, 43, 3, 9, 82, 10]

1. **Divide:** Split into halves until each subarray has one element:
 - [38, 27, 43] and [3, 9, 82, 10]
 - Further split until: [38], [27], [43], [3], [9], [82], [10]
2. **Conquer:** Recursively sort each half (trivial for single-element arrays)
3. **Merge:** Combine sorted arrays:
 - Merge [38] and [27] → [27, 38]
 - Merge [27, 38] and [43] → [27, 38, 43]
 - Merge [3] and [9] → [3, 9], then merge with [82] → [3, 9, 82], and so on

4. Final merge gives the fully sorted array: [3, 9, 10, 27, 38, 43, 82]

Time and Space Complexity

- Time Complexity:
 - Best, Average, Worst: $O(n \log n)$
 - Because the array is split $\log n$ times and merging takes $O(n)$
- Space Complexity:
 - $O(n)$ due to the temporary arrays used during merging

Advantages of MergeSort

- Stable: Maintains the relative order of equal elements
- Predictable performance: Always $O(n \log n)$
- Good for linked lists and external sorting (e.g., sorting data on disk)

Disadvantages

- Not in-place: Requires additional memory
- Slower for small datasets compared to simpler algorithms like Insertion Sort

MergeSort Algorithm Pseudocode

Below are the pseudocodes for the Merge and MergeSort methods that make up the algorithm.

Merge Pseudocode

```
function Merge(left, right) is
    var result := empty vector
    while left is not empty and right is not empty do
    {
        if first(left) ≤ first(right) then
            add first(left) to result
            remove first from left
        else
            add first(right) to result
            remove first from right
    }

    // Either left or right may have elements left; consume them.
    // (Only one of the following loops will actually be entered.)
    while left is not empty do
    {
        add first(left) to result
        remove first from left
    }
    while right is not empty do
    {
        add first(right) to result
        remove first from right
    }
    return result
```

MergeSort Pseudocode

```
function MergeSort(vector m) is
    // Base case. A vector of zero or one elements is sorted, by definition.
    if length of m ≤ 1 then
        return m

    // Recursive case. First, divide the vector into equal-sized subvectors
    // consisting of the first half and second half of the vector.
    // This assumes vectors start at index 0.
    var left := empty vector
    var right := empty vector
    for i = 0 to length(m)/2 then
        add m[i] to left
    else
        add m[i] to right

    // Recursively sort both subvectors.
    left := MergeSort(left)
    right := MergeSort(right)

    // Then Merge the now-sorted subvectors.
    return Merge(left, right)
```

Quiz!

Here's a short quiz on the topic: [quiz](#)

Footer Separator

Markdown Viewer

How to view the markdown files in a browser...

- [Markdown Viewer](#)

Lecture Practices

Here are the lecture Practices...

- [Day 4](#)
- [Day 5](#)
- [Day 6](#)

Lecture Quizzes

Here are the lecture quizzes...

- [Day 4](#)
- [Day 5](#)
- [Day 6](#)

Weekly Topics

Here are the topics for the week...

- Recursion
- Pseudocode
- Sorting
- Searching
- Maps
- Time Complexity