

### ③ Merge Sort

- i)  $\Rightarrow$  Divide & Conquer Algorithm.
- ii)  $O(n \log n)$  time &  $O(n)$  Auxil space.
- iii) Stable Algorithm.
- iv) Best for Lists, but outperformed by quicksort in case of arrays.

#### eg: Merge 2 Sorted Arrays

$\Rightarrow$  Naive method  $\rightarrow$  create a new array, copy elem of both array in it & sort it.

$\Rightarrow$  2 pointer method.  $O(m+n)$  time.

$\Rightarrow$  Algo of 2 pointers

void 2pointMethod (int a[], int b[])

{ int i=0; j=0;

m  $\rightarrow$  length of a  
n  $\rightarrow$  length of b.

while (i < m && j < n)

{ if (a[i] <= b[j])

{ ~~cout~~ print(a[i]);  
i++;

}

else { print(b[j]);

j++;

}

}

while (i < m)

{ ~~cout~~ print(a[i]); i++;

}

while (j < n)

{ print(b[j]); j++;

}

## ✓ Merge fn of Merge Sort

from previous problem :- We can derive merge fn.

```
void merge (int a[], int low, int mid, int high)
```

```
{
    int n1 = mid - low + 1, n2 = high - mid;
    int left[n1], right[n2];
```

```
    for (int i = 0; i < n1; i++) { left[i] = a[low + i]; }
```

```
    for (int i = 0; i < n2; i++) { right[i] = a[mid + i + 1]; }
```

```
    int i = 0, j = 0, k = 0;
```

```
    while (i < n1 && j < n2)
```

```
    { if (left[i] <= right[j]) { a[k] = left[i];
        i++; k++; }
```

```
      else { a[k] = right[j];
        k++; j++; }
```

```
    }
```

```
    while (i < n1) { a[k] = left[i]; i++; k++; }
```

```
    while (j < n2) { a[k] = right[j]; j++; k++; }
```

```
}
```

Like our last  
ques but  
low, mid & high  
given  
low to mid  
→ arr 1  
i.e left  
mid to high  
→ arr 2  
i.e right



## vi) merge sorting algo

void MergeSort (int arr[], int l, int r)

{ if (r > l) → (check if atleast 2 elem)

int m = l + (r - l) / 2 → (same as  $\frac{r+l}{2}$ , but our formula avoids overflow.)

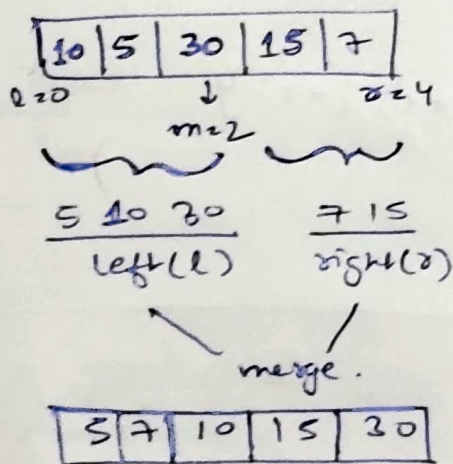
MergeSort (arr, l, m);

MergeSort (arr, m+1, r);

merge (arr, l, m, r); → (in next page).

}

}



Example(1):

$O(n)$  work at each level.

$O(\log n)$  level

→  $O(n \log n)$