

# SEARCHING.

o) We already know linear search.

## ① Binary Search.

a)  $\Rightarrow$  I/p: arr[] = [10, 20, 30, 40, 50]

x = 20

O/p: 1.

b)  $\Rightarrow$  Normally we can use linear search, but it takes  $O(n)$  time in worst case (element not there in array).

## c) Corner Cases in Binary Search.

$\Rightarrow$  I/p: arr[] = {10, 15}

I/p: arr[] = {10, 10}

O/p: -1 x = 20

O/p: 0 (or) 1

d) Idea of binary search is to use the fact that array is sorted. and cut the array by half every time

Code:

```
int BSearch (int arr[], int n, int x)
```

```
{ int low = 0, high = n-1;
```

```
  while (low <= high)
```

```
  { int mid = (low + high) / 2;
```

```
    if (arr[mid] == x)
```

```
      { return mid; }
```

```
    else if (arr[mid] > x)
```

```
      { low high = mid - 1; }
```

```
    else { low = mid + 1; }
```

```
  }
```

```
  return -1;
```

Note: 2 elem

$(0 + 1/2) = 0$

is midpoint.

## ② Recursive Code for Binary Search.

① →

```
int bSearch (int arr[], int low, int high, int x)
{
    if (low > high) return -1;

    int mid = (low + high) / 2;

    if (arr[mid] == x) { return mid; }
    else if (arr[mid] > x)
    { return bSearch(arr, low, mid-1, x); }
    else
    { return bSearch(arr, mid+1, high, x); }
}
```

## ③ Analysis of Binary Search

→ In general iterative and recursive req.  $O(\log N)$  time.  
recursive also req.  $\rightarrow O(\log N)$  extra space.

→ till 16 ele  $\rightarrow$  4 iterations.  $\rightarrow$  loop is divided by  
constant amount  
i.e 2  
17 to 32 ele  $\rightarrow$  5 iterations.  
33 to 64 ele  $\rightarrow$  6 iterations.