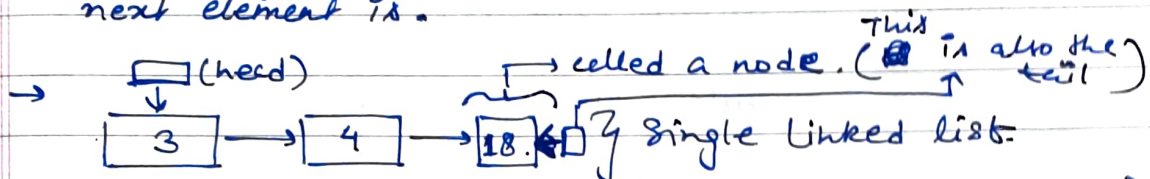


## Linked List

① Why? → In general, arrays are continuous (structures) data storing, & are fixed → Limitation.

② → Linked List is not continuous, but dynamic [not fixed]. That's why they also need addresses to point where next element is.



→ 4 doesn't know about 3, just about 18 (points to 18).

→ head has only address.

→ tail doesn't have address (points to null).

③ Syntax → 

```
class Node {  
    int val;  
    Node next;  
}
```

 } → All are private.  
→ data of LL Node.  
→ address of next LL node.

\* temp val is created everytime, i.e for anything to be done on linked list, temp val is created.

\* tail → generally found by an item having next = null. If not given, same used for finding size().

\* ④ In Java direct usage (using Collections).

→ `LinkedList<Integer> list = new LinkedList<>();`

has many functions →

a). `add(data)`;

e) `get`, `getFirst`, `getLast()`; → Index of also.

b). `add(index, data)`;

f) `clear()`; `clone()`;

c). `addFirst(data)`;

g) `peek()` → shows, doesn't remove

d). `addLast(data)`;

h) `poll()` → shows & removed.

i) `remove()`, `remove(object O)`, `remove(index)`, `removeFirstOccurrence()`;

j) `set(index, element)`

k) `size()`

l) `toArray()`,  
`toString()`;

⑤ Code & Internal working.

package com.Tyagi;

line ① public class LL {

line ② private class Node {

private int value;

private Node next;

public Node (int value) {

    this.value = value;  
}

public Node (int value, Node next) {

    this.value = value;

    this.next = next;

}

}

private Node head;

private Node tail;

private int size;

public LL() {

    this.size = 0;

}

} Constructor for LL.

---

package com.Tyagi

public class Main {

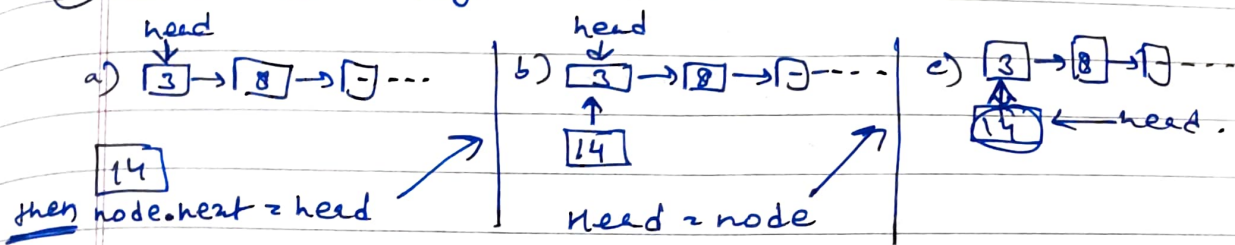
    public static void main (String[] args) {

        LL list = new LL();

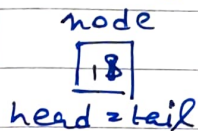
    }

}

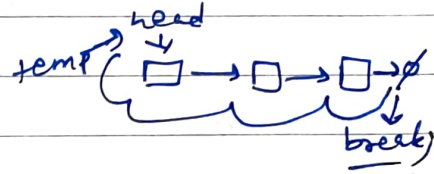
## 6) Insertion & Display.



## d) Inserting the first element.



## e) Display.



## Code [B/n Line ① & ②]

```
public void insertFirst(int val) {
```

```
    Node node = new Node(val); // adding new node.  
    node.next = head; // i.e null.
```

```
    head = node; // head equal to address of new node.
```

```
    if (tail == null) {
```

```
        tail = head;
```

```
    }
```

tells us if it is the first ele to be added.

```
    size += 1; —————> keeps track of size
```

```
}
```

```
public void display () {
```

```
    Node temp = head; // makes a new temp node.
```

```
    while (temp != null) {
```

```
        system.out.println(temp.value + " -> ");
```

```
        temp = temp.next;
```

// \n not needed.

```
    }
```

```
    System.out.print("END");
```

```
}
```

4) Insert at last

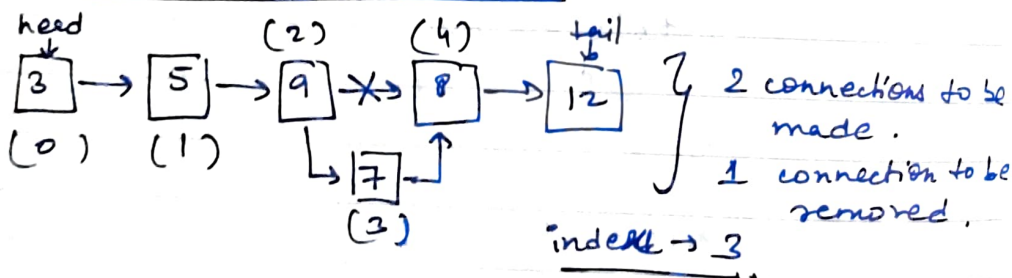
```
public void insertLast(int val) {
    if (tail == null) {
        insertFirst(val);
        return;
    }
```

```
    Node node = new Node(val);
    tail.next = node;
    tail = node;
    size++;
}
```

→ If tail not there then use fn in display() to go to last and insert the element. →  $O(n)$ .

→ If tail there →  $O(1)$

9) Insert at a particular Index



Code

```
public void insert(int val, int index) {
    if (index == 0) {
        insertFirst(val);
        return;
    }
    if (index == size) {
        insertLast(val);
        return;
    }
```



```

Node temp = head;
for (int i=1; i<index; i++) {
    temp = temp.next; // [9] will be temp finally.
}

Node node = new Node(val, temp.next);
temp.next = node;

size++;
}

```

★★ Remember [9] is temp at that stage, so new node created will point to temp.next i.e [7].

∴ then temp.next = node ([9] → [7])

### ⑦ Deletion (First, Last & index).

```

public int deleteFirst() {
    int val = head.value; // imagine as (node.next).value
    head = head.next;
    if (head == null) {
        tail = null;
    }
    size--;
    return val; // returning deleted element.
}

```

```

public int deleteLast() {

```

↳ thinking here is we iterate to n-2 and then make its address point to null.

Last element will go into garbage value.

Getting value at the index given [prereq for delete]  
last & for any value also.

```
→ public Node get(int index){  
    Node node = head;  
    for (int i=0; i<index; i++){  
        node = node.next;  
    }  
    return node;  
}
```

```
public int deleteLast(){  
    if (size <= 1){  
        deleteFirst();  
    }
```

```
    Node secondLast = get(size-2);
```

```
    int val = tail.value;
```

```
    tail = secondLast;
```

```
    tail.next = null;
```

```
    return null;
```

```
}
```

if tail not there  
use a temp  
node.

(optional step)  
storing value to be  
deleted to return

Deleting any element

// 2 temp values also be used

```
public int delete(int index){
```

```
    if (index == 0) { deleteFirst(); }
```

```
    if (index == size-1) { deleteLast(); }
```

```
    Node prev = get(index-1);
```

```
    int val = prev.next.value;
```

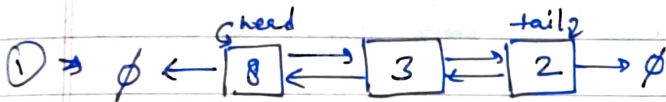
```
    prev.next = prev.next.next;
```

```
    return val;
```

### Bonus fn.

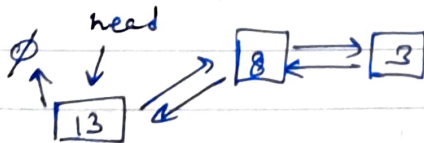
```
public Node find (int value) {  
    Node node = head  
    while (node != null)   
        can also be used.  
        (* write node instead of i here.)  
    for (int i = 0; i < size; i++) {  
        if (node != null) { if (node.value == value) {  
            return node;  
        }  
        node = node.next;  
    }  
    return null; // node not found.  
}
```

### Doubly LL & Circular LL



Syntax  $\rightarrow$  class Node {  
 int val;  
 Node next;  
 Node prev;  $\rightarrow$  Extra added.  
}

### ② $\Rightarrow$ Insertion



from ① to ②

Create new node.

node.next = head

node.prev = null.

check for null pointer exception  
[ head.prev = node  
head = node

$\rightarrow$  if (head != null) {  
 head.prev = node;  
}

- ③ Insert at last of  
Display remains the same. Deletion remains same.
- ④ To reverse linked list, start from tail & node.prev, use a temp variable. till temp  $\neq \emptyset$ .

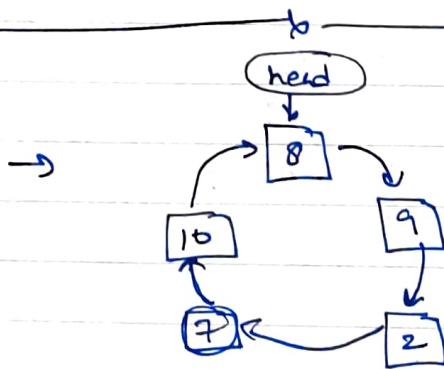
### ⑤ Insertion in b/n

- Same as SLL if index given.
- If Question is Insert node after a given value.  
i.e. Insert after 5.

Sol. create temp reach at 5.  
Create the new node.

→ node.next = temp.next;  
temp.next = node;    node  
~~temp~~.next.prev = node;  
etc.

### Circular LL



### Only difference:

- things are not null, until list is empty.

tail.next = node  
node.next = head  
tail = node.

} insertion

if (head ==  $\emptyset$ ) {  
head = node;  
tail = node;

}

do while loop for display → do {

System.out.println(node.val);

node = node.next;

} while (node != head);