

## Balanced Parenthesis.

I/p: ([ )]

O/p: - Not valid

I/p: { } ( [ ( ) ] )

O/p: Valid.

Code:-

boolean matching(char a, char b)

```
{ return ((a == '(' && b == ')') ||  
          (a == '[' && b == ']') ||  
          (a == '{' && b == '}'));
```

}

boolean isBalanced(String str)

```
{ Deque<Character> s = new ArrayDeque<>();
```

```
for (int i = 0; i < str.length(); i++)
```

```
{ char x = str.charAt(i);
```

```
if (x == '(' || x == '[' || x == '{')
```

```
{ s.push(x); }
```

```
else { if (s.isEmpty() == false) { return false; }
```

```
else if (matching(s.peek(), x) == false) { return false; }
```

```
else { s.pop(); }
```

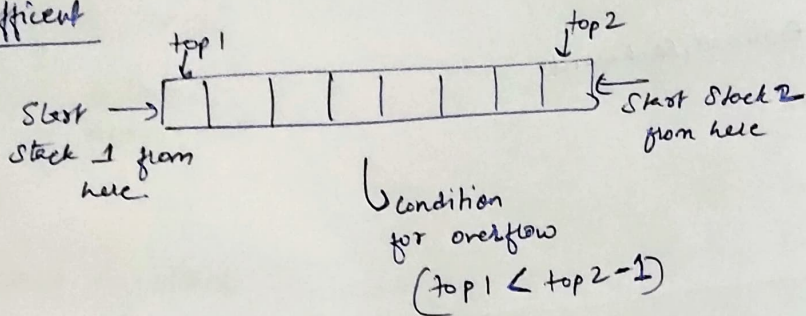
```
}  
return (s.isEmpty() == true);
```

→ if in ques this is valid too, then just do  $x++$  &  $x--$ ;  
at last  $x == 0$   
 $x++ \rightarrow (, [, {$   
 $x-- \rightarrow ), ], }$

2 or more Stacks in an array.

Naive → divide the array into 2 parts from middle.

Efficient



Stock Span Problem.

→ Span on a day → (including that day = 1) → (No. of days less <sup>with</sup> value than it on left side that are consecutive)

→ Eg:- [30, 20, 25, 20, 27, 29]  
          |  |  2  3  1  2

①

→ Naive Solution →  $O(n^2)$

→ Traversing left side of every array element.

② Code

```
void printSpan (int arr[], int n)
{
    for (int i = 0; i < n; i++)
    {
        int span = 1;
        for (int j = i - 1; j >= 0 && arr[j] <= arr[i]; j--)
            span++;
        print(span + " ");
    }
}
```



## ② Efficient solution

Void printSpan(int arr[], int n)

```
{ Stack S;
  S.push(0);
  print(1);
```

```
  for(int i=1; i<n; i++)
```

```
  { while (S.isEmpty() == false; && arr[S.top()] <= arr[i])
    { S.pop();
    }
    span = S.isEmpty() ? i+1 : i - S.top();
    print(span);
    S.push(i);
  }
}
```

removing  
all  
smaller  
items.

$$\text{Span of item} = \left( \text{Index of curr item} \right) - \left( \text{Index of prev greater elem.} \right)$$

Previous Greater Element. (Variation of Stock span).

I/p → [15, 10, 18, 12, 4, 6, 2, 8]

O/p → -1 15 -1 18 12 12 6 12

I/p → {8, 10, 12}

O/p → -1 -1 -1

P.7.0 →

Naive Sol  $\rightarrow$  Use two loops, traverse from right ~~for~~ each element.  
 $O(n^2)$

Efficient Sol  $\rightarrow$  printPrerGreater (int arr[], int n)

```
{ Stack <int> s;  
  s.push(arr[0]);
```

```
for (int i = 0; i < n; i++)
```

```
{ while (s.empty() == false && s.top() <= arr[i])  
  { s.pop();  
  }
```

```
int pg = (s.empty()) ? -1 : s.top();
```

```
cout << pg << " ";
```

```
s.push(arr[i]);
```

```
}  
}
```