Recursion (Basic)

- 1) In colleng thelf directly on indirectly.
- 2) 2 Parts Base case.

 Recursive Call.
- 3) Every occurrênce problem can be solved with Plesation and vice versa.
- 4) Multiple Activation records are formed for securion due to for cells.

- Uses more space and memory.

- Also has greater time veg.

- Is cleen and simple.

pool occurre Search [int aver], book indalindrome (cher strl7, ints, inte) Port l, Port 8, if (s==e) of return true; 3 if (str[8] != str[e]) & return false; } 1 gloce) & return false; 3 g (arr[1] > 2x) & return toue; 3 ? (Sce) & setuon ? & Palindrome (str, s+1, e-1); glars [x] 2= x) { return true; } return true; setum recurrive Search (arr, l+1, r-1, 2); 7) Tail Recumen - when the recurrine all to at the end of the fin-3 carier I tompiles and doesn't use the current for anymore. to wreck -89: - Pactorial Pack (int N, int a) fact (PM N) & 9 (N1220) relation a; (if (N/220) seturn 1; return fact (N-1, N+a); setuon No feet (N-1); Non-Tail, Tail 1) Application

Divide and conquering (Deich Sort, Merge Sort)

Inheritaly - Tower of mener Recurring, Traversels OKSIBKS Johnpeus problem. Problems

Guess me output o void fun (int n)	
1 % (n = 20) retum;	
fun(n-1);	
cout ((n ccendl;	
fun (n-1);	
P	
int main(){	
fun(3);	
y setumo;	
de come de la come de	
3 fun(3) (
Copurps 1	
(s fun(s)) return. 1 2 1 3 1 3 1	
(s fun(o) (1) 3	
62 2	
12 6	
Always my to draw a tree 14.	
Heways my to draw a tree 22 to salve recursion problems.	
3	
7	
2 2	
Derint n to 1 winy recurrion.	
Point_n-to-1 ("int n) of ib (n: 0) of settern); Bell cutt.	
tout (n):	
$l_{n} + n_{-} + o_{-} = 1 (n - 1)$	
g Space → O(n) but we kill of	in the
	luce to

i.e use moderna? T can reduce to compiler 0(1)

04).

fint N (PAG n) 6 % (n 2 2 0) return 0; return n + FirstN(n-1); of find not fiboneci number (where n > 0) Fibo (Inf n) { (6 m, 5)

My

13 (n 2 e 1) return 13 - tral. if (a 2 2 n) return a; a = 16-10-1/a-2); in fib (int n) (13 (n 120) return 0; relum fib(n=1) + fib(n-2); Glalindrome ulready discussed.