

Binary Search Tree

* log operations when balanced.
 → If not balance all become $O(N)$

Why?

	Array (unsorted)	Array (sorted)	LL	BST	Hash Table
Search	$O(N)$	$O(\log N)$	$O(N)$	$O(\log N)$	$O(1)$
Insert	$O(1)$	$O(N)$	$O(1)$ & $(O(N) \text{ too sorted})$	$O(\log N)$	$O(1)$
Delete	$O(N)$	$O(N)$	$O(N)$	$O(\log N)$	$O(1)$
Find closest.	$O(N)$	$O(\log N)$	$O(N)$	$O(\log N)$	$O(N)$
Sorted traversal.	$O(N \log N)$	$O(N)$	$O(N \log N)$	$O(N)$	$O(N \log N)$

∴ If only Search Insert and delete then hashtable is best but additional one of the last two req, then bst can be used.

Intro

① → Tree + Binary Search. → BST.

→ It is implemented using linked list.

→ Implemented directly in Java → TreeSet & TreeMap.

Operations in BST

① Search operation in BST

not. \Rightarrow class considered \rightarrow class Node {
 int key;
 Node left, right;
 Node(int(x)) { key = x; }
}

Search operation. ($O(n)$).

\Rightarrow boolean Search (Node root, int x)

```
{  
    if (root == null)  
        return false;  
    else if (root.key == x)  
        return true;  
    else if (root.key < x)  
        return Search (root.right, x);  
    else return Search (root.left, x);  
}
```

(iterative).

```
while (root != null)  
{  
    if (root.key == x)  
        true.  
    else if (root.key > x)  
        root = root.left;  
    else ...  
}
```

② Insert operation in BST

\Rightarrow Very similar to search operation

\Rightarrow If key is there do nothing.

\Rightarrow If not present, then we end up with a leaf node.

Special case \rightarrow (root == null) \rightarrow root = x;

\hookrightarrow if (root == null) { return ^{new} Node(x); }

Node insert (Node root, int x)

```
{
    if (root == null)
        return new Node(x);

    if (root.key > x)
        root.left = insert (root.left, x);
    else if (root.key < x)
        root.right = insert (root.right, x);

    return root;
}
```

root, something
imp as it
connects root
& our new
node.

Insert

operation in BST (iterative).

→ Node insert (Node root, int x)

```
{
    Node temp = new Node(x);
    Node parent = null, curr = root;
```

```
while (curr != null)
```

```
{
    parent = curr;

    if (curr.key > x)
        curr = curr.left;
    else if (curr.key < x)
        curr = curr.right;
```

```
    else
        return root;
}
```

```
if (parent == null)
    return temp;
```

```
if (parent.key > x)
    parent.left = temp;
```

```
else { parent.right = temp; }
```

for
forming
connection

② Delete operation in BST

→ Search and delete.

→ if not present return false.
else delete and return true.

→ Cases → leaf node
internal node.
 → one child.
 → 2 children.

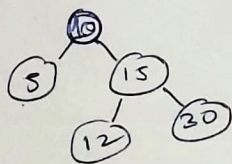
see code in Binary Tree section, same code.

③ Note:

Inorder traversal of BST is sorted.

③ Floor & Ceil Operations in BST.

→ I/P: $x = 14$
Root of tree.



O/P: 12 (closest lower value).

if not smaller then gives value return null.

→ We can easily see that smaller will lie in left subtree.

→ Node floor (root, x)

```
{ Node res = null;
while (root != null)
{ if (root.key == x)
return root;
else if (root.key > x)
root = root.left;
else
res = root;
root = root.right;
}
return res;
}
```

→ Similarly you can program for ceil too.

⑥ BST (Self Balancing)

* Idea for self balancing BST is to keep height in check so all operations stay as $O(\log n)$.

* If you know keys in adv, can simply sort and add values and it will be a balanced tree.

* After that we can do rotations to keep it balanced when nodes added. ($O(1)$ operation)

⑦ AVL & Red Black Trees

} All disturbed in Additional nodes.

Applications of BST.

- 1) Maintain sorted stream of data.
 - 2) Implementing doubly ended priority Queue.
 - 3) To solve problems like :-
 - a) Largest Smaller/Greater in a stream.
 - b) Floor/Ceiling in a stream.
-

TreeSet & TreeMap in Java

import java.util.*

public static void main (String[] args)

{ TreeSet<String> s = new TreeSet<String>();

s.add("gtg");

s.add("courses");

s.add("ide");

[also has lower,
higher, floor,
remove, ceiling].
size, isEmpty - O(1)]

System.out.println(s);

System.out.println(s.contains("ide"));

Iterator<String> i = s.iterator();

while (i.hasNext()) { System.out.println(i.next());

}

O/p: [courses gtg ide] will be in sorted order.

true

~~courses~~ courses

gtg

courses.

TreeMap

```
public static void main (String[] args)
```

```
{
```

```
    TreeMap<Integer, String> t = new TreeMap<Integer, String>();
```

```
    t.put(10, "geeks");
```

```
    t.put(15, "ide");
```

```
    t.put(5, "courses");
```

```
    System.out.println(t);
```

higher key
lower key
floor key
ceiling key

```
    for (Map.Entry<Integer, String> e : t.entrySet())
```

```
        System.out.println(e.getKey() + " " +  
                             e.getValue());
```

```
}
```

O/p:

{ 5 = courses, 10 = geeks, 15 = ide }

5 courses

10 geeks

15 ide

* Sorts acc to keys

Extra

To get a value acc to key

```
> t.get(key);
```

(t - higherEntry(10).getValue()); → gets closest higher value to key & print this value