

# Category & Categorification

plan:

- Categories & Functors

- Adjoints

- Group actions on categories

- Categorifications

- Examples.

In Statistics / applied maths, quantities are concerned (i.e. find the probability or prove an equation holds)

In pure mathematics, structures are more often concerned, (i.e. finding two structures are related or the same).

And category theory is the formal way of carrying this out.

In short, pure mathematicians often works (at least) 1 categorical level higher than applied mathematicians.

What is a category:

Def<sup>n</sup>: A category  $\mathcal{C}$  consists of:

- objects:  $o \in \mathcal{C}$

- morphisms/arrows: Given  $C_1, C_2$  objects in  $\mathcal{C}$ ,  
 $f \in \text{Hom}_{\mathcal{C}}(C_1, C_2)$

such that morphisms can compose:  $f \in \text{Hom}(C_1, C_2)$ ,  $g \in \text{Hom}(C_2, C_3)$

$$\exists g \circ f \in \text{Hom}(C_1, C_3)$$

satisfies: associativity:  $(f_1 \circ f_2) \circ f_3 = f_1 \circ (f_2 \circ f_3)$

identity:  $\exists ! e_x \in \text{Hom}(x, x)$

s.t.  $e_x \circ f = f \quad \forall f \in \text{Hom}(y, x)$

$g \circ e_x = g \quad \forall g \in \text{Hom}(x, y)$

( $\equiv$  sign here means 'they are the same arrow')

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Example:

- $\text{Grp} = \text{category of groups}$ .

objects are groups (i.e. an object is a group)

morphisms are group morphisms

- $\text{Vect}_{\mathbb{K}}$

- $\text{Mod}_{\mathbb{R}}$

- $\text{Ring} = \text{category of rings}$ .

- $\text{Field} = \text{category of fields}$ .

(This category  $\text{Field}$  is a bit strange, as the only possible field maps are injections, i.e. field extensions)

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More on Categories:

- $\mathbb{K}$ -category: Hom-sets are  $\mathbb{K}$ -vector spaces.

Example:  $\text{Vect}_{\mathbb{K}}$

Non-Example:  $\text{Grp}$

• (pre-)Additive - category: Given  $A, B \in \mathcal{C}$ , possible to define  $A \oplus B$   
also  $f, g \in \text{Hom}(A, B)$ , possible to define  $f \circ g$

Example:  $\text{Mod}_R$

(Via uni property)

Non-Example:  $\text{Field}$ ,  $\text{Ring}$

• (pre-)Abelian - category: Additive, & given  $A \xrightarrow{f} B$ , possible to define ker  
& coker

Example:  $\text{Mod}_R$ .

Non-Example:  $\text{Ring}$ ,  $\text{Field}$  (has to send 1 to 1)

Example:  $\text{Shv}_X$

Non-Example:  $\text{VectBun}_X$  (finite rank)

$$0 \rightarrow \mathcal{O}(L) \xrightarrow{p} \mathcal{O} \rightarrow \text{sky scrapper} \rightarrow 0$$

(Side note on  $\text{VectBun}_X$ , by Serre-Swan 'theorem', Vector bundles correspond to f.g projective modules over  $\mathcal{O}(X)$ ,  $\text{ker}/\text{coker } (P_1 \rightarrow P_2)$  doesn't have to be proj again).

• Monoidal category: A category where 'tensor product' is possible,  
bi-additive  $\mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$

•  $D$ -Enriched Category: Each  $\text{Hom}(C_1, C_2)$  is an object of  $D$ .

Example:  $\mathbb{k}$ -Categories are exactly Vect $_{\mathbb{k}}$ -Enriched categories.

(pre-) Additive Categories are exactly Ab-Enriched categories

• 2-Category: ? (Enriched over 'cat')

i.e. morphisms are categories

## Functors:

Functors are maps between categories

$F: \mathcal{C} \rightarrow \mathcal{D}$  is an assignment

$F(x) \quad \forall x \in \mathcal{C} \text{ and } F(f) \quad \forall f \in \text{Hom}(x, y)$

st.  $F(e) = e$

$F(x \circ y) = F(x) \circ F(y)$

Example: Forgetful functor, Homology,  $\pi_1$  (not  $H^*$ )

- $\text{Cat} = \text{Category of all (small) categories}$   
morphisms are functors.

## Natural Transformations:

are maps between functors:  $F, G: \mathcal{C} \rightarrow \mathcal{D}$

$\eta: F \rightarrow G$  is an assignment

$\eta_x: F(x) \rightarrow G(x) \quad \text{st.}$

$$\begin{array}{ccc} F(x) & \xrightarrow{\eta_x} & G(x) \\ F(f) \downarrow & \square & \downarrow G(y) \\ F(y) & \xrightarrow{\eta_y} & G(y) \end{array}$$

(Think about a map between group representations)

- $\text{Func}(\mathcal{C}, \mathcal{D})$  is a category:

each object is a functor

morphisms are natural transformations,

Thus we see  $\text{Cat}$  is a 2-category: each hom set

is also a category. The morphisms of the hom sets are called 2-morphisms.

It is possible to define maps between natural transformations and 3-categories.

Another important example:

$\text{Mor}_{\text{(Morita)}}$ : objects are  $\text{Mod}_A$ ,  $A$  a ring

(1-)morphisms are given by bimodules  $A \otimes_B$

$$\text{Mod}_A \rightarrow \text{Mod}_B : - \otimes_A M_B$$

2-morphisms are isomorphisms (or homomorphisms) of bimodules

We will think about  $\text{Mor}$  a lot in this reading seminar (probably)

Adjoints:  $(F, G)$  between  $C$  &  $D$

$$C \begin{array}{c} \xrightarrow{F} \\ \xleftarrow{G} \end{array} D$$

are adjoints if  $\text{Hom}_D(FM, N) \cong \text{Hom}_C(M, GN)$

is a isomorphism functorial in both arguments.

Example: Tensor-Hom adjunction

(All adjunctions I can think of are more or less of this form)

$$\text{Hom}(Y \otimes X, Z) \cong \text{Hom}(Y, \text{Hom}(X, Z))$$

By produce is Ind-Res adjunction.

Relation to homological algebra:

right adjoints are always left exact

left adjoints are always right exact.

Therefore  $\otimes$  is right exact  
&  $\text{Hom}$  is left exact.

$F$  is left exact if  $F$  preserve limits

In particular preserve kernels. I.e.,

if  $0 \rightarrow M_1 \rightarrow M_2 \rightarrow M_3 \rightarrow 0$  is exact

then  $0 \rightarrow \tilde{F}(M_1) \rightarrow \tilde{F}(M_2) \rightarrow \tilde{F}(M_3) \rightarrow 0$  is exact

$G$  is right exact if  $G$

preserve colimits

Exact = left + right exact

Group actions on categories:

• Groupoid : a category s.t.  $\forall f: \text{Hom}(x, y), \exists! f^*: \text{Hom}(y, x)$

s.t.  $f \circ f^* = e_y$  (i.e. all morphisms are invertible)  
 $f^* \circ f = e_x$

Then  $\text{Hom}(x, x)$  is naturally a group

A groupoid is 'a group with multiple objects'.

Example of groupoid:  $\pi_1(X)$ .

objects : points

vers : paths / homotopy

$\pi_1(x, x) = \text{Hom}_{\pi_1(X)}(x, x)$ .

Rmk: One could instead form a 2-Cat where 2-morphisms are homotopies.

group actions on categories.

Naively:  $G \curvearrowright \mathcal{C}$  is

$F(g): \mathcal{C} \rightarrow \mathcal{C}$  a functor  $\forall g \in G$

$$\text{ie } F(gh) = F(g) \circ F(h).$$

But this is too strict: typically we won't get equality

General principle:  $= \rightsquigarrow \cong$

functors themselves live in a category, so we instead ask for:  $\eta_{g,h}: F(g)F(h) \cong F(gh)$

a natural isomorphism.

and  $F(g)F(h)F(k) \xrightarrow{\text{id}_{F(g)}\eta_{h,k}} F(g)F(hk)$

$$\downarrow \eta_{g,h} \text{id}_{F(k)} \qquad \qquad \qquad \downarrow \eta_{g,hk}$$

$$F(gh)F(k) \xrightarrow{\eta_{gh,k}} F(ghk) \quad \text{are equal.}$$

The naive action is the same as  $G \rightarrow \text{Aut}(\mathcal{C})$ , map of groups where  $\text{Aut}(\mathcal{C})$  is the group of automorphisms of  $\mathcal{C}$ .

The actual action is the same as  $G \rightarrow \text{Aut}(\mathcal{C})$ , a monoidal functor, where  $G$  is viewed as a monoidal category with objects elements of  $G$ ,  $\otimes = \text{group law}$ , arrows are identities,  $\text{Aut}(\mathcal{C})$  viewed as a monoidal cat with  $\otimes = \text{composition}$ , arrows are natural isomorphisms,

One can also view  $\mathcal{G}$  as a 2-Cat, with a signal object, arrows = group elements  
 arrow composition = group law, 2-morphisms = identity. View  $\text{Aut}(\mathcal{L})$  as a 2-Cat  
 as well, then the action is the same as a 2-functor from  $\mathcal{G}$  to  
 $\text{Aut}(\mathcal{L})$ .

The procedure of going from monoidal categories to 2-categories is called  
 degrouping. One can see 2-categories as 'monoidal categories with multiple objects'.

### Categorification:

Categorification is the process of replacing set-theoretic theorems  
 with category-theoretic analogues.

replaces	sets	with	categories
	functions		functors
	equations		natural isomorphisms

The opposite direction is called decategorification.

These are not precise procedures, and there can be many ways  
 of (de)categorifying.

Very often, categorification provides more structures and further  
 insights into the problem.

Inform example:  $\text{finVect}_K$  categorifies  $\mathbb{N}$ .

$K^n$	$K^n \otimes K^m$	$K^n \otimes K^m$
$\uparrow$	$\uparrow$	$\uparrow$
$n$	$n+m$	$n \times m$

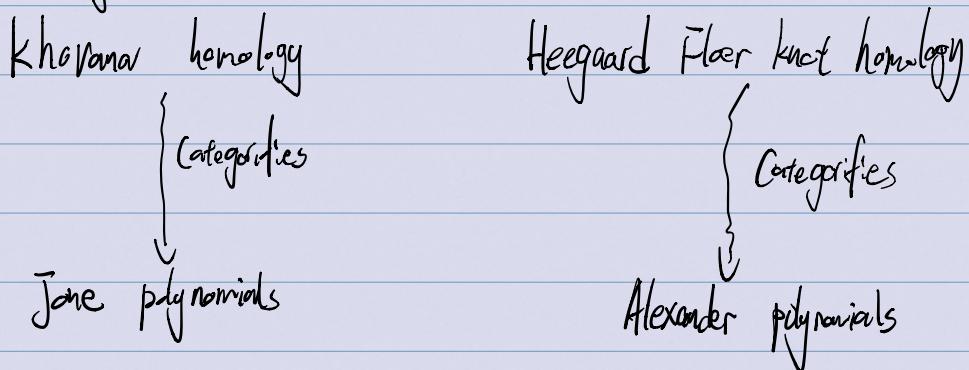
Monoid    Monoid

Decategorification is by taking dim/iso classes/Grothendieck groups.

$$\begin{aligned} \text{One sees that } n+m &= m+n \rightsquigarrow \mathbb{K}^n \oplus \mathbb{K}^m \cong \mathbb{K}^m \oplus \mathbb{K}^n \\ nm &= mn \rightsquigarrow \mathbb{K}^n \otimes \mathbb{K}^m \cong \mathbb{K}^m \otimes \mathbb{K}^n \end{aligned}$$

In a similar way, Graded vector<sub>K</sub> categorifies polynomials

Knot Theory:



A knot is  $i: S^1 \hookrightarrow \mathbb{R}^3 / S^3$

A link is a projection of image of  $i$  to  $\mathbb{R}^2$   
i.e.  $S^1 \hookrightarrow \mathbb{R}^3 \xrightarrow{?} \mathbb{R}^2$

A knot can give many link diagrams, they are related by Skein relations (Reidemeister moves)

Jone polynomials & Alexander polynomials are classical knot invariants (i.e. that  $k_1 \cong k_2 \Rightarrow P(k_1) = P(k_2)$ )

Khovanov homology of  $k$  is a graded vector space s.t. its Euler characteristic ( $\sum (-1)^i \dim V_i$ ) is Jone poly

Khovanov homology detects the unknot, it is not known if Jone poly does.

In Wei's talk, categorification of reps of  $Sl_2$  helped to construct a derived equivalence of:

$$D^b(T^*Gr(n, k)) \longrightarrow D^b(T^*Gr(n, n-k))$$

Given an (Artinian) abelian category  $\mathcal{C}$ , we can define its Grothendieck group  $k_0(\mathcal{C})$ :

It is an abelian group generated by  $[C]$ , where  $C \in ob(\mathcal{C})$ , with relations is  $0 \rightarrow C_1 \rightarrow C_2 \rightarrow C_3 \rightarrow 0$  is a SES, then  $[C_2] = [C_1] + [C_3]$

Example:  $\mathcal{C} = \text{Abelian groups}$ , then  $[\mathbb{Z}/p\mathbb{Z}] = 0 \in k_0(\mathcal{C})$   
as  $0 \rightarrow \mathbb{Z} \rightarrow \mathbb{Z} \rightarrow \mathbb{Z}/p\mathbb{Z} \rightarrow 0$

One can similarly define  $k_0(F)$ ,  $F$  a functor.

$k_0$  is a popular way of deategorification.

In Wei's talk,  $k_0(F)$  was used to construct the equivalence.

Very similarly, Chuang & Rouquier used a very similar categorification method to solve the abelian defect group conjecture for symmetric groups

A bit more on  $k_0$

As we said,

$$\begin{array}{c} A\text{-mod} \\ \left\{ \begin{array}{l} \text{Categories} \\ \vdots \\ k_0(A) \end{array} \right. \end{array}$$

However, sometimes we don't gain extra info.

For example, if  $A = \mathbb{C}G$ , then  $\mathbb{C}G\text{-mod} \cong \mathbb{C}H\text{-mod}$

iff  $k_0(\mathbb{C}G) \cong k_0(\mathbb{C}H)$  (This is a restatement of Masch  
Theorem of Character Table.)

But is it possible to recover  $G$  or  $H$  from  
 $\mathbb{C}G\text{-mod}/G\text{-rep}$  or  $k_0(\mathbb{C}G)$ .

The answer is Yes, and this is the context of  
Tannaka Duality (one form of it)

Let  $F: \mathbb{C}G\text{-mod} \rightarrow \text{Vect}$  be the forgetful functor,

then  $\text{Aut}(F) \cong G$

One can also view 'Riemann-Hilbert correspondence'  $\pi_1(X)\text{-rep} \leftrightarrow \text{Loc}(X)$   
as an instance of Tannaka duality.