

Noah Krill
Data Structures
Feb 4 2020

Algorithm 1 theoretical runtime is $O(N)$. This holds true with the empirical analysis because when the N was set to 1,000,000 the time was 0.001147 and when it was bumped up to 10,000,000 it took 0.014375. $1,000,000/0.001147 = 10,000,000/0.014375$. The first algorithm had one for loop to iterate through so its $O(N)$.

Algorithm 2 theoretical runtime is $O(N^2)$. This holds true with the empirical analysis because when the N was set to 10,000 the time was 0.145396 and when it was bumped up to 100,000 it took 14.5725. $10,000/0.145396 = 100,000/14.5725$. This algorithm had two sets of for loops causing this algorithm to be $O(N^2)$.

Algorithm 3 theoretical runtime is $O(N^3)$. This holds true with the empirical analysis because when the N was set to 200 the time was 0.014925 and when it was bumped up to 1,000 it took 2.65667. $200/0.014925 = 1,000/2.65667$. This algorithm had two sets of for loops causing this algorithm to be $O(N^3)$.

Algorithm 4 theoretical runtime is $O(N^2)$. Algorithm 2 theoretical runtime is $O(N^2)$. This holds true with the empirical analysis because when the N was set to 10,000 the time was 0.072802 and when it was bumped up to 100,000 it took 7.24494. $10,000/0.072802 = 100,000/7.24494$. This algorithm had two sets of for loops causing this algorithm to be $O(N^2)$.

Algorithm 5 theoretical runtime is $O(N^5)$. This holds true with the empirical analysis because when the N was set to 100 the time was 2.00145 and when it was bumped up to 200 it took 66.1488. $100/2.00145 = 200/66.1488$. This algorithm had three sets of for loops with a multiplication inside causing this algorithm to be $O(N^5)$.

Algorithm 6 theoretical runtime is $O(N^4)$. This holds true with the empirical analysis because when $N=1$ the time was 0.025785. When it was updated to 500 the time went up to 17.9082. $100/0.025785 = 500/17.9082$. This algorithm had three sets of for loops with a multiplication inside and an if statement causing this algorithm to be $O(N^4)$.

Exercise 2.6

A. on the n th day it would cost $((2^2)^{N-1})$

$$D = 2^{2^{N-1}}$$

$$\log D = 2^{N-1}$$

$$\log \log D + 1 = N$$

B. $O(\log \log N)$ days

2.11

A. linear = $.5 * 5 = 2.5$ ms

Noah Krill

Data Structures

Feb 4 2020

B. $O(N \log N) = 0.5 \times 5 \log_2 5 = 5.8 \text{ ms}$

C. quadratic $= 0.5 \times 5^2 = 12.5 \text{ ms}$

D. cubic $= 0.5 \times 5^3 = 62.5 \text{ ms}$