# Week 11 – Big-Step Proofs

49369565



Tuples

(TU) 
$$\frac{e_1 \Rightarrow v_1 \quad \dots \quad e_k \Rightarrow v_k}{(e_1, \dots, e_k) \Rightarrow (v_1, \dots, v_k)}$$

Lists

$$(LI) \quad \frac{e_1 \Rightarrow v_1 \qquad e_2 \Rightarrow v_2}{e_1 :: e_2 \Rightarrow v_1 :: v_2}$$

Global definitions

$$(GD) \quad \frac{f = e \quad e \Rightarrow v}{f \Rightarrow v_{X=3+5}}$$

$$\downarrow \rightarrow \mathcal{E}$$

#### Local definitions

$$(\text{LD}) \quad \underbrace{\begin{array}{ccc} e_1 \Rightarrow v_1 & e_0[v_1/x] \Rightarrow v_0 \\ \text{let } x \neq e_1 \text{ in } e_0 & \Rightarrow v_0 \end{array}}_{\text{let } x \neq e_1 \text{ in } e_0 \Rightarrow v_0$$

#### **Function calls**

(APP) 
$$\frac{e \Rightarrow \text{fun } x \rightarrow e_0 \quad e_1 \Rightarrow v_1 \quad e_0[v_1/x] \Rightarrow v_0}{e \ e_1 \ \Rightarrow \ v_0}$$

$$(\mathsf{APP'}) \quad \frac{e_0 \Rightarrow \mathsf{fun} \; x_1 \ldots x_k \rightarrow e \quad e_1 \Rightarrow v_1 \ldots e_k \Rightarrow v_k \quad e[v_1/x_1, \ldots, v_k/x_k] \Rightarrow v}{e_0 \; e_1 \; \ldots \; e_k \; \Rightarrow \; v}$$

#### Pattern Matching

$$(PM) \quad \frac{e_0 \Rightarrow v' \equiv p_i[v_1/x_1, \dots, v_k/x_k] \qquad e_i[v_1/x_1, \dots, v_k/x_k] \Rightarrow v}{\text{match } e_0 \text{ with } p_1 \rightarrow e_1 \mid \dots \mid p_m \rightarrow e_m \Rightarrow v}$$

#### **Built-in operators**

 $fan a \ni [(a+1,a-1)]$   $= \int ka \cdot 2[(a+1,a-1)] \qquad (kn a \cdot 2[(a+1,a-1)]) + \int [(a+1,a-1)] \qquad (kn a \cdot 2[(a+1,a-1)]) \qquad (kn a \cdot 2[(a+1$ 

$$\pi_{o} = \frac{1}{[(741, 7-1)]} = [(8,6)]$$

```
 \text{LD} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \text{ APP'} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \text{ } 7 \Rightarrow 7 \text{ } \pi_0}{\text{(fun a } \rightarrow \text{[(a+1,a-1)])} \text{ } 7 \Rightarrow \text{[(8,6)]} } \\ \text{let f = fun a } \rightarrow \text{[(a+1,a-1)]} \text{ in f } 7 \Rightarrow \text{[(8,6)]}
```

$$\pi_0 = \text{LI} \frac{\text{OP} \frac{7 \Rightarrow 7 \quad 1 \Rightarrow 1 \quad 7 + 1 \Rightarrow 8}{7 + 1 \Rightarrow 8} \text{ OP} \frac{7 \Rightarrow 7 \quad 1 \Rightarrow 1 \quad 7 - 1 \Rightarrow 6}{7 - 1 \Rightarrow 6}}{(7 + 1, 7 - 1) \Rightarrow (8, 6) \quad [] \Rightarrow []}$$

$$[(7 + 1, 7 - 1)] \Rightarrow [(8, 6)]$$

```
 \text{LD} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{APP'} \\ \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad 7 \Rightarrow 7 \quad \pi_0}{\text{(fun a } \rightarrow \text{[(a+1,a-1)])} \quad 7 \Rightarrow \text{[(8,6)]} \\ \text{let f = fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{in f } 7 \Rightarrow \text{[(8,6)]}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l \rightarrow
  match l with [] -> 0 | x::xs -> x * f xs
```

[3;6] ] [3;6] match [3;6] muth [7=)7 | x:

f  $[3;6] \Rightarrow 9$ 

$$\pi_f = \operatorname{GD} \underbrace{ \begin{cases} f = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 1 \mid \text{x::xs} -> \text{x+g xs} \\ f \Rightarrow \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 1 \mid \text{x::xs} -> \text{x+g xs} \\ f \Rightarrow \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 1 \mid \text{x::xs} -> \text{x+g xs} \\ f \Rightarrow \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+g xs} \\ \end{cases}}$$

$$\pi_g = \operatorname{GD} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ g \Rightarrow \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{match } 1 \text{ with } [] -> 0 \mid \text{x::xs} -> \text{x+f xs} \\ \end{cases}} \underbrace{ \begin{cases} g = \text{fun } 1 -> \text{m$$

$$T_{e} = k_{n} \ell \operatorname{Juntch}(\operatorname{welh}()) \Rightarrow (\cdots)$$

$$T_{e} = f = T_{e} \qquad f \Rightarrow T_{e}$$

**Global Definitions** 

$$T_{2}$$
 = fun I -> match I with [] -> 1 | x::xs -> x + g xs

$$T_g$$
 = fun | -> match | with [] -> 0 | x::xs -> x \* f xs

$$\pi_{p} = G + \pi_{p} = \pi_{p} =$$

$$T_g = T_g \qquad T_g = T_g$$

$$g \Rightarrow T_g$$

**Global Definitions** 

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_{f} [3;6] \Rightarrow [3;6] PM \frac{[3;6] \ni [3;6]}{\text{match } [3;6] \text{ with } [] \rightarrow 1 \mid x::xs \rightarrow x+g xs \Rightarrow 9}
f [3;6] \Rightarrow 9
```

APP'

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
APP' \frac{\pi_{f} \text{ [3;6]} \Rightarrow \text{[3;6]} \text{ PM}}{\frac{3+g \text{ [6]} \Rightarrow 9}{\text{match [3;6] with [] -> 1 | x::xs -> x+g xs \infty 9}}{\text{f [3;6]} \Rightarrow 9}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

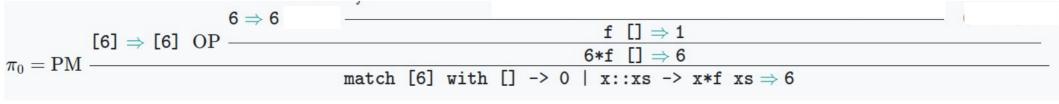
```
APP' \frac{\pi_{g} \ [6] \Rightarrow [6] \ \pi_{0}}{g \ [6] \Rightarrow 6}
\frac{3 \Rightarrow 3 \ \text{APP'}, \frac{\pi_{g} \ [6] \Rightarrow 6}{g \ [6] \Rightarrow 6}}{3+g \ [6] \Rightarrow 9}
\frac{3 \Rightarrow 3 \ \text{APP'}, \frac{\pi_{g} \ [6] \Rightarrow 6}{g \ [6] \Rightarrow 6}}{g \ [6] \Rightarrow 6}
\frac{3+g \ [6] \Rightarrow 9}{g \ [6] \Rightarrow 9}
\frac{3 \Rightarrow 3 \ \text{APP'}, \frac{\pi_{g} \ [6] \Rightarrow [6] \ \pi_{0}}{g \ [6] \Rightarrow 6}
\frac{3+g \ [6] \Rightarrow 9}{g \ [6] \Rightarrow 9}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_0 = \mathrm{PM} \ rac{6] \Rightarrow \texttt{[6]} \Rightarrow \texttt{[6]}}{\mathsf{match} \ \texttt{[6]} \ \mathsf{with} \ \texttt{[]} \Rightarrow \texttt{0} \ \texttt{x::xs} \Rightarrow \texttt{5}} \ \mathsf{x*f} \ \mathsf{xs} \Rightarrow \texttt{6}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```



```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_0 = \mathrm{PM} \ \frac{ \text{[6]} \Rightarrow \text{[6]} \ \mathrm{OP} }{ \begin{array}{c} 6 \Rightarrow 6 \ \mathrm{APP'} \\ \hline \\ \hline \\ \pi_0 = \mathrm{PM} \\ \hline \\ \hline \\ match \ [6] \ \text{with} \ [] \ -> \ 1 \ | \ x::xs \ -> \ x+g \ xs \Rightarrow 1 \\ \hline \\ \hline \\ 6*f \ [] \Rightarrow 6 \\ \hline \\ match \ [6] \ \text{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x*f \ xs \Rightarrow 6 \\ \hline \end{array}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_0 = \mathrm{PM} \ \frac{ \begin{array}{c} 6 \Rightarrow 6 \ \mathrm{APP} \end{array}, \frac{\pi_f \ \ \square \Rightarrow \square \ \ \mathrm{PM} \ \frac{\square \Rightarrow \square \ \ 1 \Rightarrow \square \ \ }{\mathrm{match} \ \ \square \ \ \mathrm{with} \ \ \square \Rightarrow \square \ \ }{\mathrm{match} \ \ \square \Rightarrow \square} \ \ \frac{1 \Rightarrow 1}{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{ \begin{array}{c} 6 \Rightarrow 6 \ \mathrm{APP} \end{array}, \frac{\pi_f \ \ \square \Rightarrow \square \ \ \mathrm{PM} \ \frac{\square \Rightarrow \square \ \ }{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{1 \Rightarrow \square \ \ }{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{1 \Rightarrow \square \ \ }{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{1 \Rightarrow \square \ \ }{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{1 \Rightarrow \square \ \ \square \ \ }{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \Rightarrow \square \ \ \square \ \ }{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \Rightarrow \square \ \ \square \ \square \ \square}{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \Rightarrow \square \ \ \square \ \square}{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \Rightarrow \square \ \ \square \ \square}{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square \ \square}{\mathrm{match} \ \ \square \Rightarrow \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square} \\ \pi_0 = \mathrm{PM} \ \frac{\square \ \ \square}{\mathrm{match} \ \ \square}
```

```
let rec f = fun l ->
                                                                                                                                                               T_{\parallel} = fun | -> match | with [] -> 1 | x::xs -> x + g xs
   match l with [] \rightarrow 1 \mid x::xs \rightarrow x + g xs
                                                                                                                                                              T_q = fun I -> match I with [] -> 0 | x::xs -> x * f xs
and q = fun l \rightarrow

\pi_{\varphi} = \frac{\oint = \Upsilon_{\ell}}{\oint \Rightarrow \Upsilon_{\ell}} \xrightarrow{\varphi} 

   match l with [] -> 0 | x::xs -> x * f xs
                                                                                                                                                         T_g = g = T_g \qquad T_g = T_g
g \Rightarrow T_g
                                        6\Rightarrow 6 \text{ APP'} \frac{\pi_f \text{ []} \Rightarrow \text{[] PM}}{\text{match [] with [] -> 1 | x::xs -> x+g xs \Rightarrow 1}} \qquad 6*1\Rightarrow 6
                                                                                                        f \mid \square \Rightarrow 1
            [6] ⇒ [6] OP —
```

 $6*f [] \Rightarrow 6$ 

match [6] with []  $\rightarrow$  0 | x::xs  $\rightarrow$  x\*f xs  $\Rightarrow$  6

$$\text{APP'} \frac{\pi_{g} \ [6] \Rightarrow [6] \ \pi_{0}}{\pi_{g} \ [6] \Rightarrow [6] \Rightarrow [6] \ \pi_{0}} \quad 3 + 6 \Rightarrow 9}{\text{match [3;6] with [] -> 1 | x::xs -> x+g xs \Rightarrow 9}}$$

$$\text{f [3;6] \ \Rightarrow 9}$$

 $\pi_0 = \mathrm{PM} -$ 

 $\pi_0 = \mathrm{APP'} \frac{\mathrm{fun} \ x \ -> \ x \ 3 \Rightarrow \mathrm{fun} \ x \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3 \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3}{\mathrm{(fun} \ y \ z \ -> \ z \ y) \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3}$ 

 $\text{APP}, \frac{\pi_0}{--}$ 

 $\operatorname{APP}, \frac{\pi_0 \text{ fun w -> w+w} \Rightarrow \text{fun w -> w+w}}{\text{(fun w -> w+w) } 3 \Rightarrow 6}$   $\operatorname{(fun x -> x 3) (fun y z -> z y) (fun w -> w+w) \Rightarrow 6}$ 

 $APP' = \frac{\pi_0 \text{ fun w } -> \text{ w+w} \Rightarrow \text{fun w } -> \text{ w+w} \Rightarrow \text{fun$ 

 $APP' = \frac{\pi_0 \text{ fun w } -> \text{ w+w} \Rightarrow \text{fun w } -> \text{w+w} \Rightarrow \text{fun w } -> \text{ w+w} \Rightarrow \text{fun w } -> \text{ w+w} \Rightarrow \text{fun w } -> \text{ w+w} \Rightarrow \text{fun w$ 

$$\pi_0 = \text{APP'} \frac{\text{fun x -> x 3} \Rightarrow \text{fun x -> x 3 fun y z -> z y} \Rightarrow \text{fun z -> z 3} \Rightarrow \text{fun z -> z 3}}{(\text{fun y z -> z y})} \Rightarrow \text{fun z -> z 3}}$$

$$APP' = \frac{\pi_0 \text{ fun w -> w+w $\Rightarrow$ fun w -> w+w $\Rightarrow$ fun w -> w+w $\Rightarrow$ so } {(\text{fun w -> w+w}) \ 3 \Rightarrow 6} \frac{3 \Rightarrow 3 \ 3 \Rightarrow 3 \ 3 \Rightarrow 6}{3+3 \Rightarrow 6} \frac{(\text{fun w -> w+w}) \ 3 \Rightarrow 6}{(\text{fun w -> w+w}) \Rightarrow 6}$$

Prove that the function

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

terminates for all inputs  $a, b \geq 0$ .

let rec mul a b =
 match a with 0 -> 0 | \_ -> b + mul (a-1) b

Proof by Induction on n

Base: n = 0

To prove: if  $a \Rightarrow 0$  and  $b \Rightarrow m$ , then mul  $a \ b \Rightarrow 0 \cdot m$ 

let rec mul a b =
 match a with 0 -> 0 | \_ -> b + mul (a-1) b

Proof by Induction on n

Base: n = 0

To prove: if  $a\Rightarrow 0$  and  $b\Rightarrow m$ , then mul  $a\ b\Rightarrow 0\cdot m$ 

$$\text{APP'} \frac{\pi_{mul} \quad a \Rightarrow 0 \quad b \Rightarrow m \quad \text{PM} \; \frac{0 \Rightarrow \text{O} \quad \text{O} \Rightarrow \text{O}}{\text{match O with O -> O \mid \_ -> b + mul (O - 1)} \; b \Rightarrow 0}{\text{mul } a \; b \Rightarrow 0}$$

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if  $a \Rightarrow n$  and  $b \Rightarrow m$ , then  $\mathtt{mul}\ a\ b \Rightarrow n \cdot m$ 

Step:

To prove: if  $a\Rightarrow n+1$  and  $b\Rightarrow m$ , then  $\operatorname{\mathtt{mul}}\ a\ b\Rightarrow (n+1)\cdot m$ .

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if  $a \Rightarrow n$  and  $b \Rightarrow m$ , then  $\mathtt{mul}\ a\ b \Rightarrow n \cdot m$ 

Step:

To prove: if  $a\Rightarrow n+1$  and  $b\Rightarrow m$ , then  $\operatorname{\mathtt{mul}}\ a\ b\Rightarrow (n+1)\cdot m$ .

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if  $a \Rightarrow n$  and  $b \Rightarrow m$ , then  $\mathtt{mul}\ a\ b \Rightarrow n \cdot m$ 

Step:

To prove: if  $a \Rightarrow n+1$  and  $b \Rightarrow m$ , then  $\operatorname{mul}\ a\ b \Rightarrow (n+1) \cdot m$ .

mul 
$$((n+1) - 1)$$
  $m \Rightarrow n \cdot m$ 

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if  $a \Rightarrow n$  and  $b \Rightarrow m$ , then mul  $a \ b \Rightarrow n \cdot m$ 

Step:

To prove: if  $a\Rightarrow n+1$  and  $b\Rightarrow m$ , then  $\operatorname{\mathtt{mul}}\ a\ b\Rightarrow (n+1)\cdot m$ .

$$\operatorname{APP'} \frac{\pi_{mul} \quad a \Rightarrow n+1 \quad b \Rightarrow m \quad \operatorname{PM}}{m \Rightarrow m \quad \operatorname{by} I.H.} \frac{\operatorname{OP} \frac{n+1 \Rightarrow n+1 \quad 1 \Rightarrow 1 \quad (n + 1) - 1 \Rightarrow n}{(n + 1) - 1 \Rightarrow n} \quad m \Rightarrow m \\ \frac{m \Rightarrow m \quad \operatorname{by} I.H.}{mul \quad ((n+1) - 1) \quad m \Rightarrow n \cdot m} \quad m + (n \cdot m) \Rightarrow (n + 1) \cdot m}{m + mul \quad ((n+1) - 1) \quad m \Rightarrow (n+1) \cdot m}$$

$$\operatorname{Match} \quad n+1 \quad \text{with} \quad 0 \to 0 \quad | \quad -> m + mul \quad ((n+1) - 1) \quad m \Rightarrow (n+1) \cdot m}$$

$$\operatorname{Mul} \quad a \quad b \Rightarrow (n+1) \cdot m$$

Use big-step operational semantics to show that the function

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

terminates for all inputs and computes three times the sum of the input list's elements.

Proof by Induction on length of List

```
Base: I = []
```

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

APP

threesum  $[] \Rightarrow 0$ 

Proof by Induction on length of List

Base: I = []

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

```
\mathrm{APP} \ \frac{\pi_{ts} \ \texttt{[]} \Rightarrow \texttt{[]} \ \mathrm{PM} \ \frac{\texttt{[]} \Rightarrow \texttt{[]} \ \mathrm{O} \Rightarrow \texttt{0}}{\mathtt{match} \ \texttt{[]} \ \mathtt{with} \ \texttt{[]} -> \texttt{0} \ | \ \mathtt{x::xs} -> \ 3*\mathtt{x} \ + \ \mathtt{threesum} \ \mathtt{xs} \Rightarrow \texttt{0}}{\mathtt{threesum} \ \texttt{[]} \Rightarrow \texttt{0}}
```

Proof by Induction on length of List

Hypothisis:  $th_{SN} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$ 

Step: I = x :: xs

```
let rec threesum = fun l \rightarrow match \ l \ with [] \rightarrow 0 \ | \ x::xs \rightarrow 3*x + threesum xs
```

let rec threesum = fun l ->

Proof by Induction on length of List

Hypothisis:  $th_{501} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$ 

Step: I = x :: xs

match l with [] -> 0 | x::xs -> 3\*x + threesum xs

let rec threesum = fun l ->

match l with [] -> 0 | x::xs -> 3\*x + threesum xs

Proof by Induction on length of List

Hypothisis:  $th_{501} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$ 

Step: I = x :: xs

```
 \underbrace{x_{n+1} :: \mathtt{xs} \Rightarrow x_{n+1} :: \mathtt{xs} \Rightarrow x_{n+1}
```