# Week 11 – Big-Step Proofs



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Quiz

#### **Tuples**

(TU) 
$$\frac{e_1 \Rightarrow v_1 \quad \dots \quad e_k \Rightarrow v_k}{(e_1, \dots, e_k) \Rightarrow (v_1, \dots, v_k)}$$

Lists

(LI) 
$$\frac{e_1 \Rightarrow v_1 \qquad e_2 \Rightarrow v_2}{e_1 :: e_2 \Rightarrow v_1 :: v_2}$$

Global definitions

$$(\mathsf{GD}) \quad \frac{f = e \quad e \Rightarrow v}{f \Rightarrow v}$$

#### Local definitions

(LD) 
$$\frac{e_1 \Rightarrow v_1 \qquad e_0[v_1/x] \Rightarrow v_0}{\text{let } x = e_1 \text{ in } e_0 \Rightarrow v_0}$$

#### Function calls

(APP) 
$$\frac{e \Rightarrow \text{fun } x \rightarrow e_0 \quad e_1 \Rightarrow v_1 \quad e_0[v_1/x] \Rightarrow v_0}{e \ e_1 \ \Rightarrow \ v_0}$$

#### Pattern Matching

$$(PM) \quad \frac{e_0 \Rightarrow v' \equiv p_i[v_1/x_1, \dots, v_k/x_k] \qquad e_i[v_1/x_1, \dots, v_k/x_k] \Rightarrow v}{\operatorname{match} e_0 \operatorname{with} p_1 \rightarrow e_1 \mid \dots \mid p_m \rightarrow e_m \Rightarrow v}$$

#### **Built-in operators**

(OP) 
$$\frac{e_1 \Rightarrow v_1}{e_1 \operatorname{op} e_2 \Rightarrow v} \frac{e_2 \Rightarrow v_2}{e_1 \operatorname{op} e_2 \Rightarrow v}$$

Unary operators are treated analogously.

$$\lim_{R \to \infty} \frac{\partial}{\partial x} \left[ \frac{\partial}{\partial x} \left( \frac{\partial}{\partial x} \right) \left( \frac{\partial}{\partial x} \left( \frac{\partial}{\partial x} \right) \left( \frac{\partial}{\partial x} \right) \left( \frac{\partial}{\partial x} \left( \frac{\partial}{\partial x} \right) \right) \right] + \lim_{R \to \infty} \frac{\partial}{\partial x} \left( \frac{\partial}{\partial x} \left( \frac{\partial}{\partial x} \right) \left( \frac{\partial}{$$

$$\operatorname{LD} \frac{\operatorname{fun \ a \rightarrow \ [(a+1,a-1)] \Rightarrow fun \ a \rightarrow \ [(a+1,a-1)] \ pp}}{\operatorname{let \ f = \ fun \ a \rightarrow \ [(a+1,a-1)] \ in \ f \ 7 \Rightarrow [(8,6)]}}$$

$$\text{LD} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{APP'} \\ \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad 7 \Rightarrow 7 \quad \pi_0}{\text{(fun a } \rightarrow \text{[(a+1,a-1)])} \quad 7 \Rightarrow \text{[(8,6)]} } \\ \text{let f = fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{in f } 7 \Rightarrow \text{[(8,6)]}$$

$$\pi_0 = \text{LI} \frac{\text{OP} \frac{7 \Rightarrow 7 \quad 1 \Rightarrow 1 \quad 7 + 1 \Rightarrow 8}{7 + 1 \Rightarrow 8} \text{ OP} \frac{7 \Rightarrow 7 \quad 1 \Rightarrow 1 \quad 7 - 1 \Rightarrow 6}{7 - 1 \Rightarrow 6}}{(7 + 1, 7 - 1) \Rightarrow (8, 6) \quad [] \Rightarrow []}$$

```
LD = \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad APP, \\ \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad 7 \Rightarrow \text{7} \quad \pi_0}{\text{(fun a } \rightarrow \text{[(a+1,a-1)])} \quad 7 \Rightarrow \text{[(8,6)]}}
\text{let f = fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{in f } 7 \Rightarrow \text{[(8,6)]}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
APP = [3;6] = [3:6] match [3:6] with []=11 | X::XS=1X+9XS f [3:6] <math>\Rightarrow 9
```

$$\pi_f = \text{GD} \ \frac{\text{f = fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs}}{\text{f = fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs}} \frac{\text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs}}{\text{f = fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs}}$$

$$\pi_g = \text{GD} \ \frac{\text{g = fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs}}{\text{g = fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs}} \frac{\text{fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs}}{\text{g = fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs}}$$

**Global Definitions** 

$$\pi_f = \text{GD} \ \frac{\text{f} = \text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs \ \text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs \ \text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs}}{\text{f} \Rightarrow \text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs}}$$

$$\pi_g = \text{GD} \ \frac{\text{g} = \text{fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs \ \text{fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs}}{\text{g} \Rightarrow \text{fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs}}$$

$$T_{\mathcal{L}} = \text{fun } I \rightarrow \text{match } I \text{ with } [] \rightarrow 1 \mid x::xs \rightarrow x + g xs$$

$$T_g$$
 = fun I -> match I with [] -> 0 | x::xs -> x \* f xs

$$\Pi_{p} = \frac{f = T_{p}}{f} \xrightarrow{T_{p} \Rightarrow T_{p}}$$

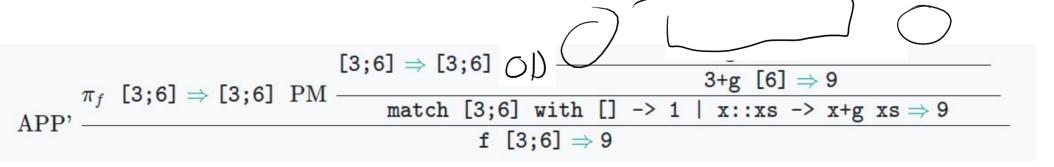
$$T_g = \frac{g = T_g}{g} \quad T_g \Rightarrow T_g$$

**Global Definitions** 

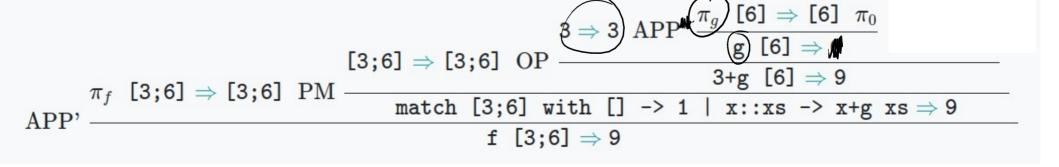
```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

$$\frac{\pi_{f} \ [3;6] \Rightarrow [3;6] \ )}{\text{APP'}} \frac{\left[3;6\right] \Rightarrow \left[3;6\right]}{\text{match} \ [3;6] \ \text{with} \ [] \ -> \ 1 \ | \ \text{x::xs} \ -> \ \text{x+g} \ \text{xs} \Rightarrow 9}{\text{f} \ [3;6] \Rightarrow 9}$$

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
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```



```
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```



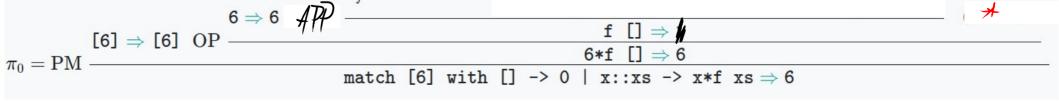
```
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  match l with [] -> 1 | x::xs -> x + g xs
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  match l with [] -> 0 | x::xs -> x * f xs
```

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```

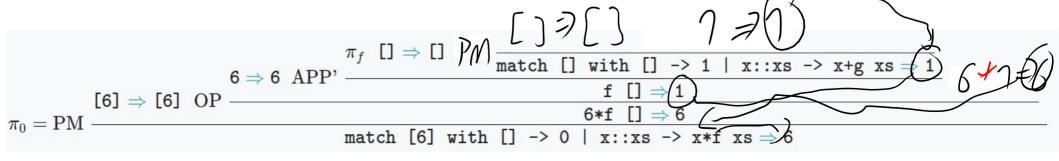


 $6*f [] \Rightarrow 6$  match [6] with [] -> 0 | x::xs -> x\*f xs \Rightarrow 6

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```



```
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```



```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
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  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_0 = \mathrm{PM} \ \frac{ \begin{array}{c} 6 \Rightarrow 6 \ \mathrm{APP} \end{array}, \frac{\pi_f \ \ \square \Rightarrow \square \ \ \mathrm{PM} \ \frac{\square \Rightarrow \square \ \ 1 \Rightarrow \square}{\mathrm{match} \ \ \square \ \ \mathrm{with} \ \ \square \Rightarrow \square \ \ 1 \ \ \times : \times s \rightarrow x + g \ xs \Rightarrow 1} \\ \hline \pi_0 = \mathrm{PM} \ \frac{ \begin{array}{c} 6 \Rightarrow 6 \ \mathrm{APP} \end{array}, \frac{\pi_f \ \ \square \Rightarrow \square \ \ \mathrm{PM} \ \frac{\mathrm{match} \ \ \square \ \ \mathrm{with} \ \ \square \Rightarrow \square \ \ }{\mathrm{match} \ \ \square \Rightarrow 6} \\ \hline \pi_0 = \mathrm{PM} \ \frac{ \begin{array}{c} 6 \Rightarrow 6 \ \mathrm{APP} \end{array}, \frac{\pi_f \ \ \square \Rightarrow \square}{\mathrm{match} \ \ \square \Rightarrow \square} \ \ \frac{\mathrm{f} \ \ \square \Rightarrow \square}{\mathrm{match} \ \ \square \Rightarrow 6} \\ \hline \pi_0 = \mathrm{PM} \ \ \frac{\mathrm{f} \ \ \square \Rightarrow \square}{\mathrm{match} \ \ \square \Rightarrow 0 \ \ | \ x : : \times s \rightarrow x + g \ xs \Rightarrow 6} \\ \hline \end{array}} \ \ \frac{6 * 1 \Rightarrow 6}{\mathrm{f} \ \ \square \Rightarrow \square}
```

```
let rec f = fun l ->
match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
match l with [] -> 0 | x::xs -> x * f xs

\frac{f}{f} = fun | -> f(x) = fun | ->
```

 $f \mid \exists \Rightarrow 1$ 

 $6*f [] \Rightarrow 6$ 

match [6] with []  $\rightarrow$  0 | x::xs  $\rightarrow$  x\*f xs  $\Rightarrow$  6

$$APP' = \frac{\pi_{f} \ [3;6] \Rightarrow [3;6] \ PM}{\frac{[3;6] \Rightarrow [3;6] \ OP}{\frac{3 \Rightarrow 3 \ APP'}{\frac{\pi_{g} \ [6] \Rightarrow [6] \Rightarrow 9}{3+g \ [6] \Rightarrow 9}}}{\frac{3+g \ [6] \Rightarrow 9}{1 \ x::xs \rightarrow x+g \ xs \Rightarrow 9}}$$

[6] ⇒ [6] OP —

 $\pi_0 = \mathrm{PM}$  -

fun x -> x 3  $\Rightarrow$  fun x -> x 3 fun y z -> z y  $\Rightarrow$  fun y z -> z y (fun y z  $\rightarrow$  z y) 3  $\Rightarrow$  $\pi_0 = APP' -$ 

 $\pi_0 = \mathrm{APP'} \frac{\mathrm{fun} \ x \ -> \ x \ 3 \Rightarrow \mathrm{fun} \ x \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3 \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3}{\mathrm{(fun} \ y \ z \ -> \ z \ y) \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3}$ 

 $\text{APP}, \frac{\pi_0}{--}$ 

 $\operatorname{APP'} \frac{\pi_0 \text{ fun w -> w+w} \Rightarrow \text{fun w -> w+w}}{\text{(fun w -> w+w) } 3 \Rightarrow 6}$  (fun x -> x 3) (fun y z -> z y) (fun w -> w+w)  $\Rightarrow$  6

 $APP' \xrightarrow{\pi_0 \text{ fun w } -> \text{ w+w} \Rightarrow \text{ 3} \Rightarrow 3} \xrightarrow{3+3 \Rightarrow 6}$   $(\text{fun w } -> \text{ w+w}) \Rightarrow 6$ 

$$APP' = \frac{\pi_0 \text{ fun w -> w+w $\Rightarrow$ fun w -> w+w $\Rightarrow$ fun w -> w+w $\Rightarrow$ so } \frac{3 \Rightarrow 3 \ 3 \Rightarrow 3 \ 3 \Rightarrow 3 \ 3 \Rightarrow 6}{3+3 \Rightarrow 6} = \frac{(\text{fun w -> w+w}) \ 3 \Rightarrow 6}{(\text{fun x -> x 3}) \ (\text{fun y z -> z y}) \ (\text{fun w -> w+w}) \Rightarrow 6}$$

### Week 11 Tutorial 02 — Multiplication

Prove that the function

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

terminates for all inputs  $a, b \geq 0$ .

## Week 11 Tutorial 02 — Multiplication

let rec mul a b =  $match \ a \ with \ 0 \ -> \ 0 \ | \ \_ \ -> \ b \ + \ mul \ (a-1) \ b$ 

Proof by Induction on a

Base: a = 0

## Week 11 Tutorial 02 — Multiplication

let rec mul a b =
 match a with 0 -> 0 | \_ -> b + mul (a-1) b

Proof by Induction on a

Base: a = 0

$$ext{APP} = \frac{\pi_{mul} ext{ PM}}{ ext{match 0 with 0 -> 0 | _ -> b + mul (-1) b  $\Rightarrow$  0}}{ ext{mul 0 b  $\Rightarrow$  0}$$

## <u>Week 11 Tutorial 02 — Multiplication</u>

let rec mul a b =
 match a with 0 -> 0 | \_ -> b + mul (a-1) b

Proof by Induction on a

Hypothisis: mul a b = a\*b / q >Step: mul (a+1) b  $\Rightarrow$  (a+1)\*b

### <u>Week 11 Tutorial 02 — Multiplication</u>

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

Proof by Induction on a

```
Hypothisis:
mul a b = a*b
Step:
```

$$APP = \frac{\text{APP } \frac{\text{by I.H.}}{\text{mul } (a+1-1) \ b \Rightarrow a*b} \ b \ + \ (a*b) \Rightarrow (a+1)*b}{\text{b + mul } (a+1-1) \ b \Rightarrow (a+1)*b}$$

$$APP = \frac{\pi_{mul} \ PM}{\text{match } a+1 \ \text{with } 0 \ -> \ 0 \ | \ \_ \ -> \ b \ + \ \text{mul } (a+1-1) \ b \Rightarrow (a+1)*b}{\text{mul } (a+1) \ b \Rightarrow (a+1)*b}$$

Use big-step operational semantics to show that the function

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

terminates for all inputs and computes three times the sum of the input list's elements.

Proof by Induction on length of List

```
Base: I = []
```

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

APP

threesum  $[] \Rightarrow 0$ 

Proof by Induction on length of List

Base: I = []

```
let rec threesum = fun l -> match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

Proof by Induction on length of List

Hypothisis:  $th_{SN} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$ 

Step: I = x :: xs

```
let rec threesum = fun l \rightarrow match \ l \ with [] \rightarrow 0 \ | \ x::xs \rightarrow 3*x + threesum xs
```

### <u>Week 11 Tutorial 03 — Threesum</u>

let rec threesum = fun l ->

match l with [] -> 0 | x::xs -> 3\*x + threesum xs

Proof by Induction on length of List

Hypothisis:  $th_{501} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$ 

Step: I = x :: xs

```
x_{n+1}::xs \Rightarrow x_{n+1}::xs
```

### <u>Week 11 Tutorial 03 — Threesum</u>

let rec threesum = fun l ->

match l with [] -> 0 | x::xs -> 3\*x + threesum xs

Proof by Induction on length of List

Hypothisis:  $th_{501} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$ 

Step: I = x :: xs

```
 \text{APP} \frac{ \text{OP} \frac{3 \Rightarrow 3 \ x_{n+1} \Rightarrow x_{n+1} \ 3 * x_{n+1} \Rightarrow 3x_{n+1}}{3 * x_{n+1} \Rightarrow 3x_{n+1}} \ \text{APP} \frac{\text{by I.H.}}{\text{threesum } xs \Rightarrow 3\sum_{i=1}^{n} x_i} \ 3x_{n+1} + 3\sum_{i=1}^{n} x_i \Rightarrow 3\sum_{i=1}^{n+1} x_i}{3 * x_{n+1} \Rightarrow 3x_{n+1}} \\ \text{APP} \frac{x_{ts} \ x_{n+1} :: xs \Rightarrow x_{n+1} :: xs \Rightarrow x_{n+1} :: xs \Rightarrow x_{n+1} :: xs \text{ with } [] \rightarrow 0 \ | \ x :: xs \rightarrow 3 * x + \text{ threesum } xs \Rightarrow 3\sum_{i=1}^{n+1} x_i}{\text{threesum } (x_{n+1} :: xs) \Rightarrow 3\sum_{i=1}^{n+1} x_i}
```