Week 11 – Big-Step Proofs



Tuples

$$(\mathsf{TU}) \quad \frac{e_1 \Rightarrow v_1 \quad \dots \quad e_k \Rightarrow v_k}{(e_1, \dots, e_k) \Rightarrow (v_1, \dots, v_k)}$$

Lists

(LI)
$$\frac{e_1 \Rightarrow v_1 \qquad e_2 \Rightarrow v_2}{e_1 :: e_2 \Rightarrow v_1 :: v_2}$$

Global definitions

$$(GD) \quad \frac{f = e \quad e \Rightarrow v}{f \Rightarrow v \quad \text{5.5}} = \frac{5+3}{5+3} = \frac{5+3}{5+3} = \frac{5}{5}$$

(e(x=543

Local definitions

(LD)
$$\frac{e_1 \Rightarrow v_1 \qquad e_0[v_1/x] \Rightarrow v_0}{\text{let } x = e_1 \text{ in } e_0 \Rightarrow v_0}$$

Function calls

(APP)
$$\frac{e \Rightarrow \text{fun } x \rightarrow e_0 \quad e_1 \Rightarrow v_1 \quad e_0[v_1/x] \Rightarrow v_0}{e \ e_1 \ \Rightarrow \ v_0}$$

Pattern Matching

$$(PM) \quad \frac{e_0 \Rightarrow v' \equiv p_i[v_1/x_1, \dots, v_k/x_k] \qquad e_i[v_1/x_1, \dots, v_k/x_k] \Rightarrow v}{\text{match } e_0 \text{ with } p_1 \rightarrow e_1 \mid \dots \mid p_m \rightarrow e_m \Rightarrow v}$$

Built-in operators

$$(\mathsf{OP}) \ \frac{e_1 \Rightarrow v_1 \qquad e_2 \Rightarrow v_2 \qquad v_1 \, \mathsf{op} \, v_2 \Rightarrow v}{e_1 \, \mathsf{op} \, e_2 \Rightarrow v}$$
 Unary operators are treated analogously.

lana>[(af1,a-1)] → lua>[(af1,a-1)]

```
\frac{(4 \ln a \rightarrow ((a + 1, a - 1)))}{\text{let f = fun a -> [(a+1, a-1)] in f 7 \Rightarrow [(8,6)]}
```

$$\widehat{\mathbb{I}_0} = \underbrace{\mathcal{E}(747, 7-1)3 \Rightarrow \mathcal{E}(8,6)}_{}$$

$$LD = \frac{\text{fun a -> [(a+1,a-1)]} \Rightarrow \text{fun a -> [(a+1,a-1)]}}{\text{fun a -> [(a+1,a-1)]}} + \frac{\text{fun a -> [(a+1,a-1)]}}{\text{(fun a -> [(a+1,a-1)])}} + \frac{\text{fun a -> [(a+1,a-1)]}}{\text{(fun a -> [(a+1,a-1)]}} + \frac{\text{fun a$$

 $LD = \frac{\text{fun a -> [(a+1,a-1)]} \Rightarrow \text{fun a -$

$$41 \frac{711 \Rightarrow 8 \quad (7-1) \Rightarrow (6)}{(7+1), 7-1) \Rightarrow (8),6}$$

$$\pi_0 = \text{LI} \frac{\text{OP} \frac{7 \Rightarrow 7 \quad 1 \Rightarrow 1 \quad 7 + 1 \Rightarrow 8}{7 + 1 \Rightarrow 8} \text{ OP} \frac{7 \Rightarrow 7 \quad 1 \Rightarrow 1 \quad 7 - 1 \Rightarrow 6}{7 - 1 \Rightarrow 6}}{(7 + 1, 7 - 1) \Rightarrow (8, 6) \quad [] \Rightarrow []}$$

$$[(7 + 1, 7 - 1)] \Rightarrow [(8, 6)]$$

$$\text{LD} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{APP'} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad 7 \Rightarrow 7 \quad \pi_0}{\text{(fun a } \rightarrow \text{[(a+1,a-1)])} \quad 7 \Rightarrow \text{[(8,6)]} } \\ \text{let f = fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{in f } 7 \Rightarrow \text{[(8,6)]}$$

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

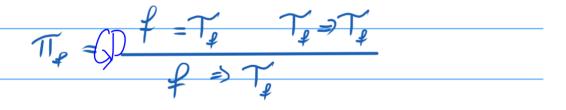
```
\pi_f = \operatorname{GD} \frac{\mathbf{f} = \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 1 \ | \ x::xs \ -> \ x+g \ xs \ \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 1 \ | \ x::xs \ -> \ x+g \ xs \ \to \ \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 1 \ | \ x::xs \ -> \ x+g \ xs \ \to \ \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+g \ xs \ \to \ \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \to \ \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \to \ \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{fun} \ 1 \ -> \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{match} \ 1 \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ xs \ \times \ \operatorname{with} \ [] \ -> \ 0 \ | \ x::xs \ -> \ x+f \ x+f
```

Te = fen (-) match (mill()-)? | y::XS-) X+ 9 X5

Global Definitions

$$T_{2}$$
 = fun I -> match I with [] -> 1 | x::xs -> x + g xs

$$T_g$$
 = fun | -> match | with [] -> 0 | x::xs -> x * f xs



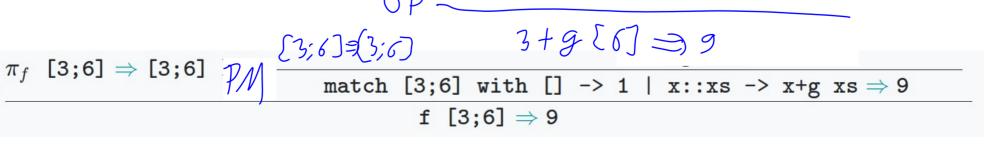
$$\frac{T_g}{g} = T_g \qquad T_g = T_g$$

$$\frac{g}{g} \Rightarrow T_g$$

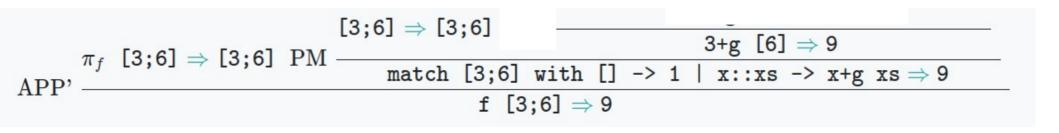
Global Definitions

APP'

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```



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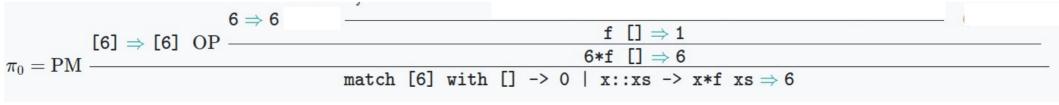
```
APP' \frac{\pi_{f} \ [3;6] \Rightarrow [3;6] \ PM}{\frac{[3;6] \Rightarrow [3;6] \ OP}{\frac{3 \Rightarrow 3 \ APP'}{\frac{\pi_{g} \ [6] \Rightarrow [6] \ \pi_{0}}{g \ [6] \Rightarrow 6}}{\frac{3+g \ [6] \Rightarrow 9}{g \ [6] \Rightarrow 9}}
f \ [3;6] \Rightarrow 9
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_0 = \mathrm{PM} \ rac{ 	ext{[6]} \Rightarrow 	ext{[6]} }{ 	ext{match [6] with [] -> 0 | x::xs -> x*f xs <math>\Rightarrow 6} }
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```



```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
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  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_0 = \mathrm{PM} \hspace{1cm} rac{6 \Rightarrow 6 \hspace{1cm} \mathrm{APP'}}{\mathrm{match} \hspace{1cm} [] \Rightarrow [] \hspace{1cm} rac{\mathrm{match} \hspace{1cm} [] \hspace{1cm} \hspace{1cm} \mathrm{with} \hspace{1cm} [] \hspace{1cm} \Rightarrow 1 \hspace{1cm} | \hspace{1cm} \mathrm{x::xs} \hspace{1cm} - \hspace{1cm} \rangle \hspace{1cm} \mathrm{x+g} \hspace{1cm} \mathrm{xs} \Rightarrow 1}{\mathrm{f} \hspace{1cm} [] \Rightarrow 1} \hspace{1cm} \pi_0 = \mathrm{PM} \hspace{1cm} \frac{6 \hspace{1cm} \hspace{1cm} \mathrm{f} \hspace{1cm} [] \Rightarrow 6}{\mathrm{match} \hspace{1cm} [6] \hspace{1cm} \mathrm{with} \hspace{1cm} [] \hspace{1cm} - \hspace{1cm} \rangle \hspace{1cm} 0 \hspace{1cm} | \hspace{1cm} \mathrm{x::xs} \hspace{1cm} - \hspace{1cm} \rangle \hspace{1cm} \mathrm{x+f} \hspace{1cm} \mathrm{xs} \Rightarrow 6}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
let rec f = fun l ->
                                                                                                                                                                      T_{\parallel} = fun | -> match | with [] -> 1 | x::xs -> x + g xs
   match l with [] \rightarrow 1 \mid x::xs \rightarrow x + g xs
                                                                                                                                                                     T_q = fun I -> match I with [] -> 0 | x::xs -> x * f xs
and q = fun l \rightarrow

\pi_{\varphi} = \frac{\oint = \uparrow_{\ell} \qquad \uparrow_{\ell} \Rightarrow \uparrow_{\ell}}{\oint \Rightarrow \uparrow_{\ell}}

   match l with [] -> 0 | x::xs -> x * f xs
                                                                                                                                                                T_g = g = T_g \qquad T_g = T_g
g \Rightarrow T_g
                                          6\Rightarrow 6 \text{ APP'} \frac{\pi_f \text{ []} \Rightarrow \text{[] PM}}{\text{match [] with [] -> 1 | x::xs -> x+g xs \Rightarrow 1}} \qquad 6*1\Rightarrow 6
                                                                                                             f \mid \square \Rightarrow 1
             [6] ⇒ [6] OP —
```

 $6*f [] \Rightarrow 6$

match [6] with [] \rightarrow 0 | x::xs \rightarrow x*f xs \Rightarrow 6

$$\text{APP'} \frac{\pi_{g} \ [6] \Rightarrow [6] \ \pi_{0}}{\pi_{g} \ [6] \Rightarrow [6] \Rightarrow 6} \ 3+6 \Rightarrow 9}{\text{match [3;6] with [] -> 1 | x::xs -> x+g xs \Rightarrow 9}}$$

 $\pi_0 = \mathrm{PM} -$

 $\pi_0 = \mathrm{APP'} \frac{\mathrm{fun} \ x \ -> \ x \ 3 \Rightarrow \mathrm{fun} \ x \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3}{\mathrm{(fun} \ y \ z \ -> \ z \ y) \ 3 \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3}$

 $\text{APP}, \frac{\pi_0}{--}$

 $\operatorname{APP}, \frac{\pi_0 \text{ fun w -> w+w} \Rightarrow \text{fun w -> w+w}}{\text{ (fun w -> w+w) } 3 \Rightarrow 6}$ (fun x -> x 3) (fun y z -> z y) (fun w -> w+w) \Rightarrow 6

 $APP' = \frac{\pi_0 \text{ fun w } -> \text{ w+w} \Rightarrow \text{fun w } -> \text{ w+w} \Rightarrow \text{fun$

 $APP' \xrightarrow{\pi_0 \text{ fun } w \rightarrow w+w \Rightarrow \text{ fun } w \rightarrow w+w \Rightarrow$

$$\pi_0 = \text{APP'} \frac{\text{fun x -> x 3} \Rightarrow \text{fun x -> x 3 fun y z -> z y} \Rightarrow \text{fun z -> z 3} \Rightarrow \text{fun z -> z 3}}{(\text{fun y z -> z y})} \Rightarrow \text{fun z -> z 3}}$$

$$APP' = \frac{\pi_0 \text{ fun w -> w+w \Rightarrow 3 \Rightarrow 3 \Rightarrow 3 \Rightarrow 3 \Rightarrow 3 \Rightarrow 3 \Rightarrow 6}{(\text{fun w -> w+w}) 3 \Rightarrow 6}$$

Week 11 Tutorial 02 — Multiplication

Prove that the function

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

terminates for all inputs $a, b \geq 0$.

<u>Week 11 Tutorial 02 — Multiplication</u>

let rec mul a b =
 match a with 0 -> 0 | _ -> b + mul (a-1) b

Proof by Induction on n

Base: n = 0

To prove: if $a\Rightarrow 0$ and $b\Rightarrow m$, then $\mathtt{mul}\ a\ b\Rightarrow 0\cdot m$

Week 11 Tutorial 02 — Multiplication

let rec mul a b =
 match a with 0 -> 0 | _ -> b + mul (a-1) b

Proof by Induction on n

Base: n = 0

To prove: if $a\Rightarrow 0$ and $b\Rightarrow m$, then mul $a\ b\Rightarrow 0\cdot m$

$$\text{APP'} \frac{\pi_{mul} \quad a \Rightarrow 0 \quad b \Rightarrow m \quad \text{PM} \; \frac{0 \Rightarrow \text{O} \quad \text{O} \Rightarrow \text{O}}{\text{match O with O -> O \mid _ -> b + mul (O - 1)} \; b \Rightarrow 0}{\text{mul } a \; b \Rightarrow 0}$$

Week 11 Tutorial 02 — Multiplication

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if $a \Rightarrow n$ and $b \Rightarrow m$, then $\mathtt{mul}\ a\ b \Rightarrow n \cdot m$

Step:

<u>Week 11 Tutorial 02 — Multiplication</u>

```
let rec mul a b = match a with 0 -> 0 \mid \_ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if $a \Rightarrow n$ and $b \Rightarrow m$, then **mul** $a \ b \Rightarrow n \cdot m$

Step:

Week 11 Tutorial 02 — Multiplication

```
let rec mul a b = match a with 0 -> 0 \mid \_ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if $a \Rightarrow n$ and $b \Rightarrow m$, then $\mathtt{mul}\ a\ b \Rightarrow n \cdot m$

Step:

mul
$$((n+1) - 1)$$
 $m \Rightarrow n \cdot m$

<u>Week 11 Tutorial 02 — Multiplication</u>

let rec mul a b =
 match a with 0 -> 0 | _ -> b + mul (a-1) b

Proof by Induction on n

Hypothisis:

if $a \Rightarrow n$ and $b \Rightarrow m$, then **mul** $a \ b \Rightarrow n \cdot m$

Step:

$$\operatorname{APP'} \frac{\pi_{mul} \quad a \Rightarrow n+1 \quad b \Rightarrow m \quad \operatorname{PM}}{m \Rightarrow m \quad \operatorname{by} I.H.} \frac{\operatorname{OP} \frac{n+1 \Rightarrow n+1 \quad 1 \Rightarrow 1 \quad (n+1) - 1 \Rightarrow n}{(n+1) - 1 \Rightarrow n} \quad m \Rightarrow m}{m \quad ((n+1) - 1) \quad m \Rightarrow n \cdot m} \quad m + (n \cdot m) \Rightarrow (n+1) \cdot m}{m \quad \operatorname{match} \quad n+1 \quad \operatorname{with} \quad 0 \to 0 \quad | \quad -> m + \operatorname{mul} \quad ((n+1) - 1) \quad m \Rightarrow (n+1) \cdot m}$$

$$\operatorname{mul} \quad a \quad b \Rightarrow (n+1) \cdot m$$

Use big-step operational semantics to show that the function

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

terminates for all inputs and computes three times the sum of the input list's elements.

Proof by Induction on length of List

```
Base: I = []
```

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

APP

threesum $[] \Rightarrow 0$

Proof by Induction on length of List

Base: I = []

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

```
\mathrm{APP} \ \frac{\pi_{ts} \ \texttt{[]} \Rightarrow \texttt{[]} \ \mathrm{PM} \ \frac{\texttt{[]} \Rightarrow \texttt{[]} \ \mathrm{O} \Rightarrow \texttt{0}}{\mathtt{match} \ \texttt{[]} \ \mathtt{with} \ \texttt{[]} -> \texttt{0} \ | \ \mathtt{x::xs} -> \ \mathtt{3*x} \ + \ \mathtt{threesum} \ \mathtt{xs} \Rightarrow \texttt{0}}{\mathtt{threesum} \ \texttt{[]} \Rightarrow \texttt{0}}
```

Proof by Induction on length of List

Hypothisis: $th_{SN} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$

Step: I = x :: xs

```
let rec threesum = fun l \rightarrow match \ l \ with [] \rightarrow 0 \ | \ x::xs \rightarrow 3*x + threesum xs
```

Proof by Induction on length of List

Hypothisis: $th_{SN} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$

Step: I = x :: xs

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

```
\pi_{ts} \ x_{n+1} \colon \colon \mathsf{xs} \Rightarrow x_{n+1} \colon \mathsf{xs} \to \mathsf{xs} \Rightarrow x_{n+1} \colon \mathsf{xs} \to \mathsf
```

let rec threesum = fun l ->

match l with [] -> 0 | x::xs -> 3*x + threesum xs

Proof by Induction on length of List

Hypothisis: $th_{501} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$

Step: I = x :: xs

```
 \underbrace{x_{n+1} :: \mathtt{xs} \Rightarrow x_{n+1} \Rightarrow 3x_{n+1} \\ \underbrace{x_{n+1} :: \mathtt{xs} \Rightarrow x_{n+1} :: \mathtt{xs} \Rightarrow x_{n
```