Week 11 – Big-Step Proofs



Tuples

(TU)
$$\frac{e_1 \Rightarrow v_1 \quad \dots \quad e_k \Rightarrow v_k}{(e_1, \dots, e_k) \Rightarrow (v_1, \dots, v_k)}$$

Lists

(LI)
$$\frac{e_1 \Rightarrow v_1 \qquad e_2 \Rightarrow v_2}{e_1 :: e_2 \Rightarrow v_1 :: v_2}$$

Global definitions

$$(\mathsf{GD}) \quad \frac{f = e \quad e \Rightarrow v}{f \Rightarrow v}$$

Local definitions

(LD)
$$\frac{e_1 \Rightarrow v_1 \qquad e_0[v_1/x] \Rightarrow v_0}{\text{let } x = e_1 \text{ in } e_0 \Rightarrow v_0}$$

Function calls

(APP)
$$\frac{e \Rightarrow \text{fun } x \rightarrow e_0 \quad e_1 \Rightarrow v_1 \quad e_0[v_1/x] \Rightarrow v_0}{e \ e_1 \ \Rightarrow \ v_0}$$

Pattern Matching

$$(PM) \quad \frac{e_0 \Rightarrow v' \equiv p_i[v_1/x_1, \dots, v_k/x_k] \qquad e_i[v_1/x_1, \dots, v_k/x_k] \Rightarrow v}{\operatorname{match} e_0 \operatorname{with} p_1 \rightarrow e_1 \mid \dots \mid p_m \rightarrow e_m \Rightarrow v}$$

Built-in operators

(OP)
$$\frac{e_1 \Rightarrow v_1}{e_1 \operatorname{op} e_2 \Rightarrow v} \frac{e_2 \Rightarrow v_2}{e_1 \operatorname{op} e_2 \Rightarrow v}$$

Unary operators are treated analogously.

```
 \text{LD} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]}}{\text{let f = fun a } \rightarrow \text{[(a+1,a-1)]} \text{ in f } 7 \Rightarrow \text{[(8,6)]}}
```

```
\pi_0 = \begin{bmatrix} (7+1,7-1) \end{bmatrix} \Rightarrow \begin{bmatrix} (8,6) \end{bmatrix}
```

```
 \text{LD} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{APP'} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad 7 \Rightarrow 7 \quad \pi_0}{\text{(fun a } \rightarrow \text{[(a+1,a-1)])} \quad 7 \Rightarrow \text{[(8,6)]} } \\ \text{let f = fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{in f } 7 \Rightarrow \text{[(8,6)]}
```

$$\pi_0 = \operatorname{LI} \frac{\operatorname{OP} \frac{7 \Rightarrow 7 \ 1 \Rightarrow 1 \ 7 + 1 \Rightarrow 8}{7 + 1 \Rightarrow 8} \operatorname{OP} \frac{7 \Rightarrow 7 \ 1 \Rightarrow 1 \ 7 - 1 \Rightarrow 6}{7 - 1 \Rightarrow 6}}{(7 + 1, 7 - 1) \Rightarrow (8, 6) \ [] \Rightarrow []}$$

$$[(7 + 1, 7 - 1)] \Rightarrow [(8, 6)]$$

```
 \text{LD} \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{APP'} \\ \frac{\text{fun a } \rightarrow \text{[(a+1,a-1)]} \Rightarrow \text{fun a } \rightarrow \text{[(a+1,a-1)]} \quad 7 \Rightarrow 7 \quad \pi_0}{\text{(fun a } \rightarrow \text{[(a+1,a-1)])} \quad 7 \Rightarrow \text{[(8,6)]} \\ \text{let f = fun a } \rightarrow \text{[(a+1,a-1)]} \quad \text{in f } 7 \Rightarrow \text{[(8,6)]}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_f = \text{GD} \ \frac{\text{f} = \text{fun 1} \ -> \text{ match 1 with []} \ -> \ 1 \ | \ x::xs \ -> \ x+g \ xs \ \text{fun 1} \ -> \ \text{match 1 with []} \ -> \ 1 \ | \ x::xs \ -> \ x+g \ xs \ \Rightarrow \ \text{fun 1} \ -> \ \text{match 1 with []} \ -> \ 1 \ | \ x::xs \ -> \ x+g \ xs \ }
\pi_g = \text{GD} \ \frac{\text{g} = \text{fun 1} \ -> \ \text{match 1 with []} \ -> \ 0 \ | \ x::xs \ -> \ x*f \ xs \ \Rightarrow \ \text{fun 1} \ -> \ \text{match 1 with []} \ -> \ 0 \ | \ x::xs \ -> \ x*f \ xs \ \Rightarrow \ \text{fun 1} \ -> \ \text{match 1 with []} \ -> \ 0 \ | \ x::xs \ -> \ x*f \ xs \ }
```

Global Definitions

$$\pi_f = \text{GD} \ \frac{\text{f} = \text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs \ \text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs \ \text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs}}{\text{f} \Rightarrow \text{fun 1 -> match 1 with [] -> 1 \mid x::xs -> x+g \ xs}}$$

$$\pi_g = \text{GD} \ \frac{\text{g} = \text{fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs \ \text{fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs}}{\text{g} \Rightarrow \text{fun 1 -> match 1 with [] -> 0 \mid x::xs -> x*f \ xs}}$$

$$T_{\mathcal{L}} = \text{fun } I \rightarrow \text{match } I \text{ with } [] \rightarrow 1 \mid x::xs \rightarrow x + g xs$$

$$T_g$$
 = fun I -> match I with [] -> 0 | x::xs -> x * f xs

$$\pi_{\ell} = \frac{f = T_{\ell}}{f} \xrightarrow{T_{\ell} \Rightarrow T_{\ell}}$$

$$T_{g} = \frac{g = T_{g}}{g} = T_{g}$$

Global Definitions

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
	ext{APP'} = \frac{\pi_f \ [3;6] \Rightarrow [3;6]}{	ext{match } [3;6] \ 	ext{with } [] \ 	ext{->} 1 \ | \ 	ext{x::xs} \ 	ext{->} \ 	ext{x+g} \ 	ext{xs} \Rightarrow 9}{	ext{f } [3;6] \Rightarrow 9}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
APP' \frac{\pi_f \ [3;6] \Rightarrow [3;6] \ PM}{\frac{3+g \ [6] \Rightarrow 9}{match \ [3;6] \ with \ [] \ -> 1 \ | \ x::xs \ -> x+g \ xs \Rightarrow 9}{f \ [3;6] \Rightarrow 9}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

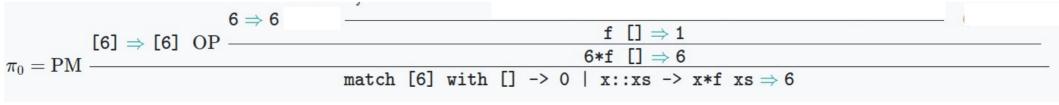
```
APP' \frac{\pi_{f} \ [3;6] \Rightarrow [3;6] \ PM}{\frac{[3;6] \Rightarrow [3;6] \ OP}{\frac{3 \Rightarrow 3 \ APP'}{\frac{\pi_{g} \ [6] \Rightarrow [6] \ \pi_{0}}{g \ [6] \Rightarrow 6}}{\frac{3+g \ [6] \Rightarrow 9}{g \ [6] \Rightarrow 9}}
f \ [3;6] \Rightarrow 9
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_0 = \mathrm{PM} \ rac{ 	ext{[6]} \Rightarrow 	ext{[6]} }{ 	ext{match [6] with [] -> 0 | x::xs -> x*f xs <math>\Rightarrow 6} }
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```



```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
\pi_0 = \mathrm{PM} \hspace{1cm} rac{6 \Rightarrow 6 \hspace{1cm} \mathrm{APP'}}{\mathrm{match} \hspace{1cm} [] \Rightarrow [] \hspace{1cm} rac{\mathrm{match} \hspace{1cm} [] \hspace{1cm} \hspace{1cm} \mathrm{with} \hspace{1cm} [] \hspace{1cm} 	o \hspace{1cm} 1 \hspace{1cm} \hspace{1cm} \hspace{1cm} \hspace{1cm} \mathrm{x::xs} \hspace{1cm} \hspace{1cm}
```

```
let rec f = fun l ->
  match l with [] -> 1 | x::xs -> x + g xs
and g = fun l ->
  match l with [] -> 0 | x::xs -> x * f xs
```

```
let rec f = fun l ->
                                                                                                                                                                      T_{\parallel} = fun | -> match | with [] -> 1 | x::xs -> x + g xs
   match l with [] \rightarrow 1 \mid x::xs \rightarrow x + g xs
                                                                                                                                                                     T_q = fun I -> match I with [] -> 0 | x::xs -> x * f xs
and q = fun l \rightarrow

\pi_{\varphi} = \frac{\oint = \uparrow_{\ell} \qquad \uparrow_{\ell} \Rightarrow \uparrow_{\ell}}{\oint \Rightarrow \uparrow_{\ell}}

   match l with [] -> 0 | x::xs -> x * f xs
                                                                                                                                                                T_g = g = T_g \qquad T_g = T_g
g \Rightarrow T_g
                                          6\Rightarrow 6 \text{ APP'} \frac{\pi_f \text{ []} \Rightarrow \text{[] PM}}{\text{match [] with [] -> 1 | x::xs -> x+g xs \Rightarrow 1}} \qquad 6*1\Rightarrow 6
                                                                                                             f \mid \square \Rightarrow 1
             [6] ⇒ [6] OP —
```

 $6*f [] \Rightarrow 6$

match [6] with [] \rightarrow 0 | x::xs \rightarrow x*f xs \Rightarrow 6

$$\text{APP'} \frac{\pi_{g} \ [6] \Rightarrow [6] \ \pi_{0}}{\pi_{g} \ [6] \Rightarrow [6] \Rightarrow 6} \ 3+6 \Rightarrow 9}{\text{match [3;6] with [] -> 1 | x::xs -> x+g xs \Rightarrow 9}}$$

 $\pi_0 = \mathrm{PM} -$

fun x -> x 3 \Rightarrow fun x -> x 3 fun y z -> z y \Rightarrow fun y z -> z y (fun y z \rightarrow z y) 3 \Rightarrow $\pi_0 = APP' -$

 $\pi_0 = \mathrm{APP'} \frac{\mathrm{fun} \ x \ -> \ x \ 3 \Rightarrow \mathrm{fun} \ x \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ y \ z \ -> \ z \ y \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3}{\mathrm{(fun} \ y \ z \ -> \ z \ y) \ 3 \Rightarrow \mathrm{fun} \ z \ -> \ z \ 3}$

 $\text{APP}, \frac{\pi_0}{--}$

 $\operatorname{APP}, \frac{\pi_0 \text{ fun w -> w+w} \Rightarrow \text{fun w -> w+w}}{\text{(fun w -> w+w) } 3 \Rightarrow 6}$ $\operatorname{(fun x -> x 3) (fun y z -> z y) (fun w -> w+w) \Rightarrow 6}$

 $APP' \xrightarrow{\pi_0 \text{ fun } w \rightarrow w+w \Rightarrow 3 \Rightarrow 3} \xrightarrow{3+3 \Rightarrow 6}$ $(\text{fun } w \rightarrow w+w) \Rightarrow 3 \Rightarrow 6$ $(\text{fun } w \rightarrow w+w) \Rightarrow 6$

 $APP' \xrightarrow{\pi_0 \text{ fun } w \rightarrow w+w \Rightarrow \text{ fun } w \rightarrow w+w \Rightarrow$

$$\pi_0 = \text{APP'} \frac{\text{fun x -> x 3} \Rightarrow \text{fun x -> x 3 fun y z -> z y} \Rightarrow \text{fun z -> z 3} \Rightarrow \text{fun z -> z 3}}{(\text{fun y z -> z y})} \Rightarrow \text{fun z -> z 3}}$$

$$APP' = \frac{\pi_0 \text{ fun w -> w+w \Rightarrow fun w -> w+w \Rightarrow fun w -> w+w \Rightarrow so } \frac{3 \Rightarrow 3 \ 3 \Rightarrow 3 \ 3 \Rightarrow 3 \ 3 \Rightarrow 6}{3+3 \Rightarrow 6} = \frac{(\text{fun w -> w+w}) \ 3 \Rightarrow 6}{(\text{fun x -> x 3}) \ (\text{fun y z -> z y}) \ (\text{fun w -> w+w}) \Rightarrow 6}$$

Prove that the function

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

terminates for all inputs $a, b \geq 0$.

let rec mul a b =
 match a with 0 -> 0 | _ -> b + mul (a-1) b

Proof by Induction on n

Base: n = 0

To prove: if $a \Rightarrow 0$ and $b \Rightarrow m$, then mul $a \ b \Rightarrow 0 \cdot m$

let rec mul a b =
 match a with 0 -> 0 | _ -> b + mul (a-1) b

Proof by Induction on n

Base: n = 0

To prove: if $a\Rightarrow 0$ and $b\Rightarrow m$, then mul $a\ b\Rightarrow 0\cdot m$

$$\text{APP'} \frac{\pi_{mul} \quad a \Rightarrow 0 \quad b \Rightarrow m \quad \text{PM} \; \frac{0 \Rightarrow \texttt{0} \quad \texttt{0} \Rightarrow \texttt{0}}{\mathsf{match} \; 0 \; \mathsf{with} \; \texttt{0} \; \texttt{>} \; \texttt{0} \; | \; _ \; \texttt{>} \; b \; + \; \mathsf{mul} \; \; (0 \; \texttt{-} \; \texttt{1}) \; \; b \Rightarrow \texttt{0}}{\mathsf{mul} \; \; a \; b \Rightarrow \texttt{0}}$$

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if $a \Rightarrow n$ and $b \Rightarrow m$, then $\mathtt{mul}\ a\ b \Rightarrow n \cdot m$

Step:

```
let rec mul a b =
  match a with 0 -> 0 | _ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if $a \Rightarrow n$ and $b \Rightarrow m$, then mul $a \ b \Rightarrow n \cdot m$

Step:

$$\text{APP'} \frac{\pi_{mul} \quad a \Rightarrow n+1 \quad b \Rightarrow m \quad \text{PM}}{m+1 \Rightarrow n+1} \quad \text{OP} \frac{m \Rightarrow m}{m + \text{mul } ((n+1) - 1) \quad m \Rightarrow (n+1) \cdot m}{m + \text{mul } ((n+1) - 1) \quad m \Rightarrow (n+1) \cdot m} \\ \text{match } n+1 \text{ with } 0 \rightarrow 0 \mid_{-} \rightarrow m + \text{mul } ((n+1) - 1) \quad m \Rightarrow (n+1) \cdot m}$$

```
let rec mul a b = match a with 0 -> 0 \mid \_ -> b + mul (a-1) b
```

Proof by Induction on n

Hypothisis:

if $a \Rightarrow n$ and $b \Rightarrow m$, then $\mathtt{mul}\ a\ b \Rightarrow n \cdot m$

Step:

mul
$$((n+1) - 1)$$
 $m \Rightarrow n \cdot m$

let rec mul a b =
 match a with 0 -> 0 | _ -> b + mul (a-1) b

Proof by Induction on n

Hypothisis:

if $a \Rightarrow n$ and $b \Rightarrow m$, then mul $a \ b \Rightarrow n \cdot m$

Step:

$$\operatorname{APP'} \frac{\pi_{mul} \quad a \Rightarrow n+1 \quad b \Rightarrow m \quad \operatorname{PM}}{m \Rightarrow m \quad \operatorname{by} I.H.} \frac{\operatorname{OP} \frac{n+1 \Rightarrow n+1 \quad 1 \Rightarrow 1 \quad (n+1) - 1 \Rightarrow n}{(n+1) - 1 \Rightarrow n} \quad m \Rightarrow m}{mul \quad ((n+1) - 1) \quad m \Rightarrow n \cdot m} \quad m + (n \cdot m) \Rightarrow (n+1) \cdot m}{m + mul \quad ((n+1) - 1) \quad m \Rightarrow (n+1) \cdot m}$$

$$\operatorname{Match} \quad n+1 \quad \text{with} \quad 0 \to 0 \quad | \quad -> m + mul \quad ((n+1) - 1) \quad m \Rightarrow (n+1) \cdot m}{mul \quad a \quad b \Rightarrow (n+1) \cdot m}$$

Use big-step operational semantics to show that the function

```
let rec threesum = fun l ->
  match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

terminates for all inputs and computes three times the sum of the input list's elements.

Proof by Induction on length of List

Base: I = []

```
let rec threesum = fun l -> match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

APF

threesum $[] \Rightarrow 0$

Proof by Induction on length of List

```
Base: I = []
```

```
let rec threesum = fun l -> match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

Proof by Induction on length of List

Hypothisis: $th_{501} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$

Step: I = x :: xs

```
let rec threesum = fun l \rightarrow match \ l \ with [] \rightarrow 0 \ | \ x::xs \rightarrow 3*x + threesum xs
```

Proof by Induction on length of List

Hypothisis: $th_{501} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$

Step: I = x :: xs

```
let rec threesum = fun l ->
match l with [] -> 0 | x::xs -> 3*x + threesum xs
```

```
\pi_{ts} \ x_{n+1} :: \mathtt{xs} \Rightarrow x_
```

let rec threesum = fun l ->

match l with [] -> 0 | x::xs -> 3*x + threesum xs

Proof by Induction on length of List

Hypothisis: $th_{501} \downarrow \Rightarrow \sum_{i=1}^{n} x_i$

Step: I = x :: xs

```
 \text{APP} \frac{ \text{OP} \frac{3 \Rightarrow 3 \ x_{n+1} \Rightarrow x_{n+1} \ 3 * x_{n+1} \Rightarrow 3x_{n+1}}{3 * x_{n+1} \Rightarrow 3x_{n+1}} \ \text{APP} \frac{ \text{by I.H.}}{\text{threesum } xs \Rightarrow 3\sum_{i=1}^{n} x_i} \ 3x_{n+1} + 3\sum_{i=1}^{n} x_i \Rightarrow 3\sum_{i=1}^{n+1} x_i }{3 * x_{n+1} :: xs \ \text{PM}} 
 \frac{x_{n+1} :: xs \Rightarrow x_{n+1} :: xs \Rightarrow x_{n+1} :: xs \ \text{OP} \frac{3 \Rightarrow 3 \ x_{n+1} \Rightarrow 3x_{n+1}}{3 * x_{n+1} \Rightarrow 3x_{n+1}} \ \text{APP} \frac{\text{by I.H.}}{\text{threesum } xs \Rightarrow 3\sum_{i=1}^{n} x_i} \ 3x_{n+1} + 3\sum_{i=1}^{n} x_i \Rightarrow 3\sum_{i=1}^{n+1} x_i }{3 * x_{n+1} \Rightarrow 3x_{n+1} \Rightarrow 3x_{n+1} \ \text{Threesum } xs \Rightarrow 3\sum_{i=1}^{n+1} x_i } 
 \text{match } x_{n+1} :: xs \ \text{with } [] \rightarrow 0 \ | \ x :: xs \rightarrow 3 * x \ + \text{threesum } xs \Rightarrow 3\sum_{i=1}^{n+1} x_i } 
 \text{threesum } (x_{n+1} :: xs) \Rightarrow 3\sum_{i=1}^{n+1} x_i
```