

Constraint Satisfaction Problems

Artificial Intelligence

Constraint satisfaction problems (CSPs)

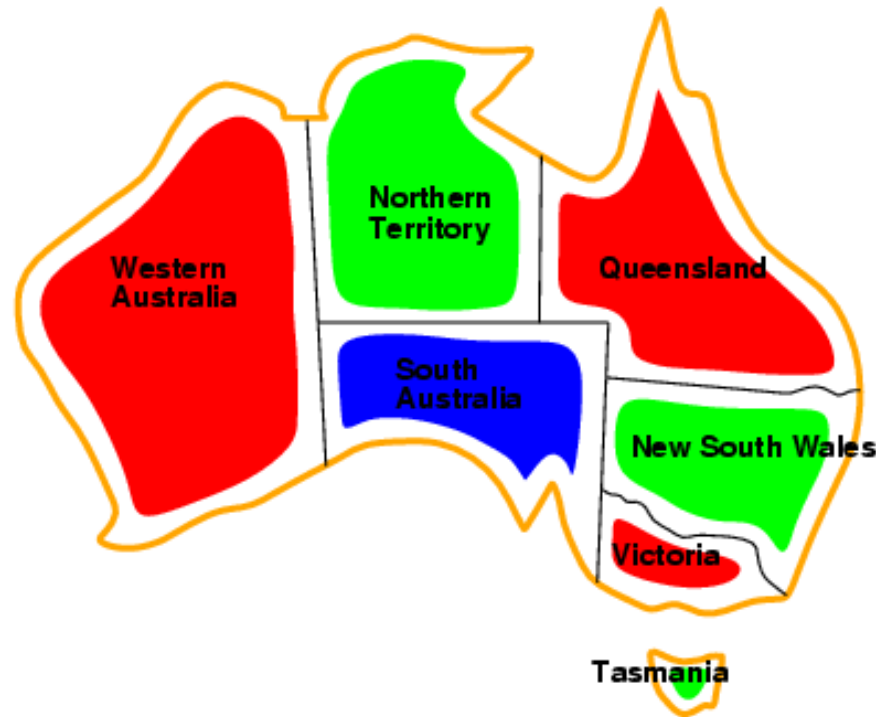
- Standard search problem: **state** is a "black box" – any data structure that supports successor function and goal test
CSP:
 - **state** is defined by **variables** X_i with **values** from **domain** D_i
 - **goal test** is a set of **constraints** specifying allowable combinations of values for subsets of variables
- Simple example of a **formal representation language**
- Allows useful **general-purpose** algorithms with more power than standard search algorithms

Example: Map-Coloring



- **Variables** WA, NT, Q, NSW, V, SA, T
- **Domains** $D_i = \{\text{red, green, blue}\}$
- **Constraints:** adjacent regions must have different colors
e.g., $WA \neq NT$, or (WA, NT) in $\{(\text{red, green}), (\text{red, blue}), (\text{green, red}), (\text{green, blue}), (\text{blue, red}), (\text{blue, green})\}$

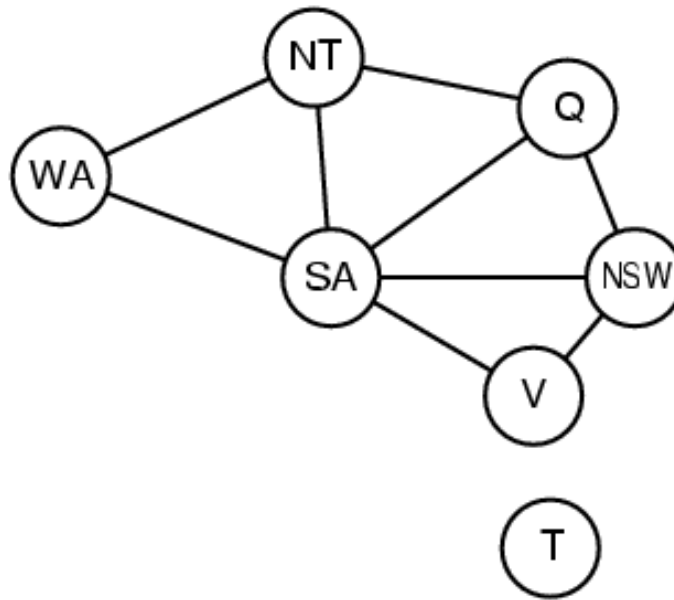
Example: Map-Coloring



- Solutions are complete and consistent assignments
- e.g., WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green

Constraint graph

- **Binary CSP:** each constraint relates two variables
- **Constraint graph:** nodes are variables, arcs are constraints



Varieties of CSPs

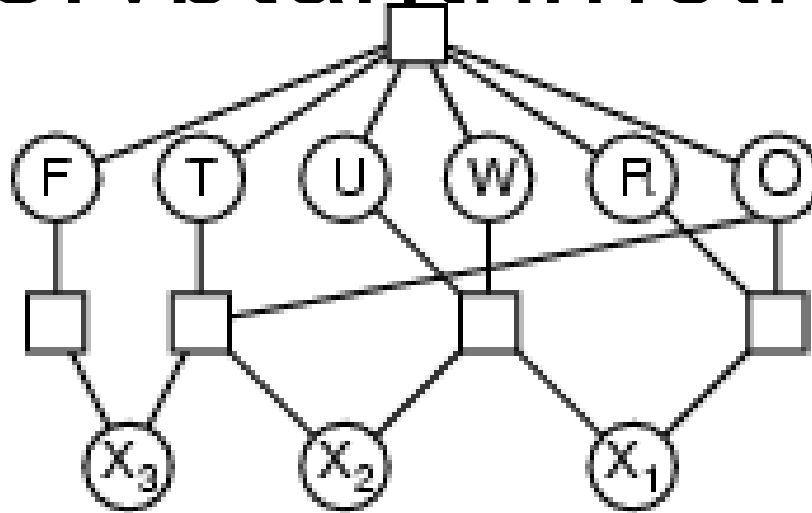
- Discrete variables
 - finite domains:
 - n variables, domain size $d \rightarrow O(d^n)$ complete assignments
 - e.g., Boolean CSPs, incl. Boolean satisfiability (NP-complete)
 - infinite domains:
 - integers, strings, etc.
 - e.g., job scheduling, variables are start/end days for each job
 - need a constraint language, e.g., $StartJob_1 + 5 \leq StartJob_3$
- Continuous variables
 - e.g., start/end times for Hubble Space Telescope observations
 - linear constraints solvable in polynomial time by LP

Varieties of constraints

- **Unary** constraints involve a single variable,
 - e.g., $SA \neq \text{green}$
- **Binary** constraints involve pairs of variables,
 - e.g., $SA \neq WA$
- **Higher-order** constraints involve 3 or more variables,
 - e.g., cryptarithmic column constraints

Example: Cryptarithmic

$$\begin{array}{r} \text{ T W O} \\ + \text{ T W O} \\ \hline \text{ F O U R} \end{array}$$



- **Variables:** $F T U W$
 $R O X_1 X_2 X_3$
- **Domains:** $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
- **Constraints:** $Alldiff(F, T, U, W, R, O)$

$$- O + O = R + 10 \cdot X_1$$

$$- X_1 + W + W = U + 10 \cdot X_2$$

$$- X_2 + T + T = O + 10 \cdot X_3$$

Real-world CSPs

- Assignment problems
 - e.g., who teaches what class
- Timetabling problems
 - e.g., which class is offered when and where?
- Transportation scheduling
- Factory scheduling
- Notice that many real-world problems involve real-valued variables

Backtracking search

- Variable assignments are **commutative**, i.e.,
[WA = red then NT = green] same as [NT = green then WA = red]
- => Only need to consider assignments to a single variable at each node
- Depth-first search for CSPs with single-variable assignments is called **backtracking** search
- Can solve n -queens for $n \approx 25$

Backtracking search

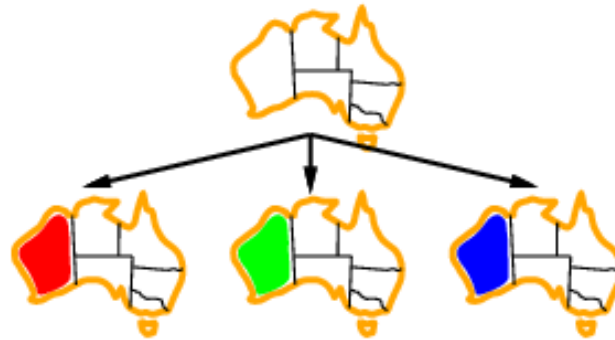
```
function BACKTRACKING-SEARCH(csp) returns a solution, or failure
  return RECURSIVE-BACKTRACKING({}, csp)

function RECURSIVE-BACKTRACKING(assignment, csp) returns a solution, or failure
  if assignment is complete then return assignment
  var ← SELECT-UNASSIGNED-VARIABLE(Variables[csp], assignment, csp)
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
    if value is consistent with assignment according to Constraints[csp] then
      add { var = value } to assignment
      result ← RECURSIVE-BACKTRACKING(assignment, csp)
      if result ≠ failure then return result
      remove { var = value } from assignment
  return failure
```

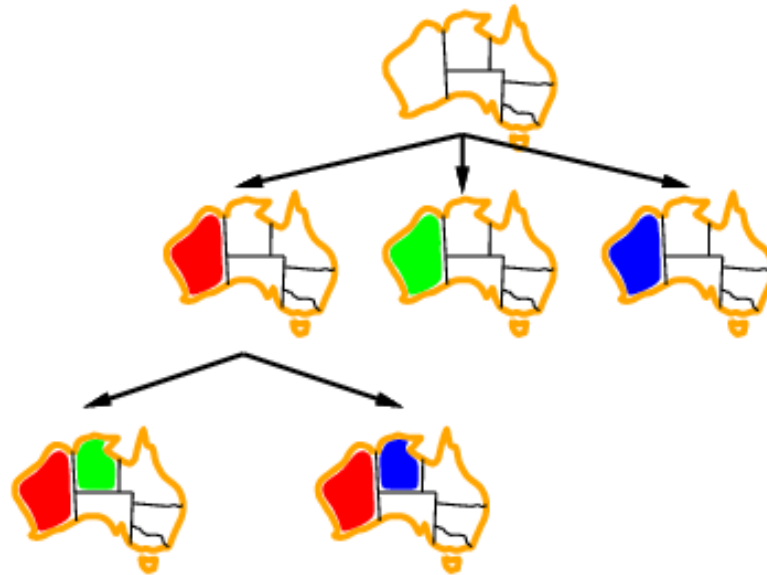
Backtracking example



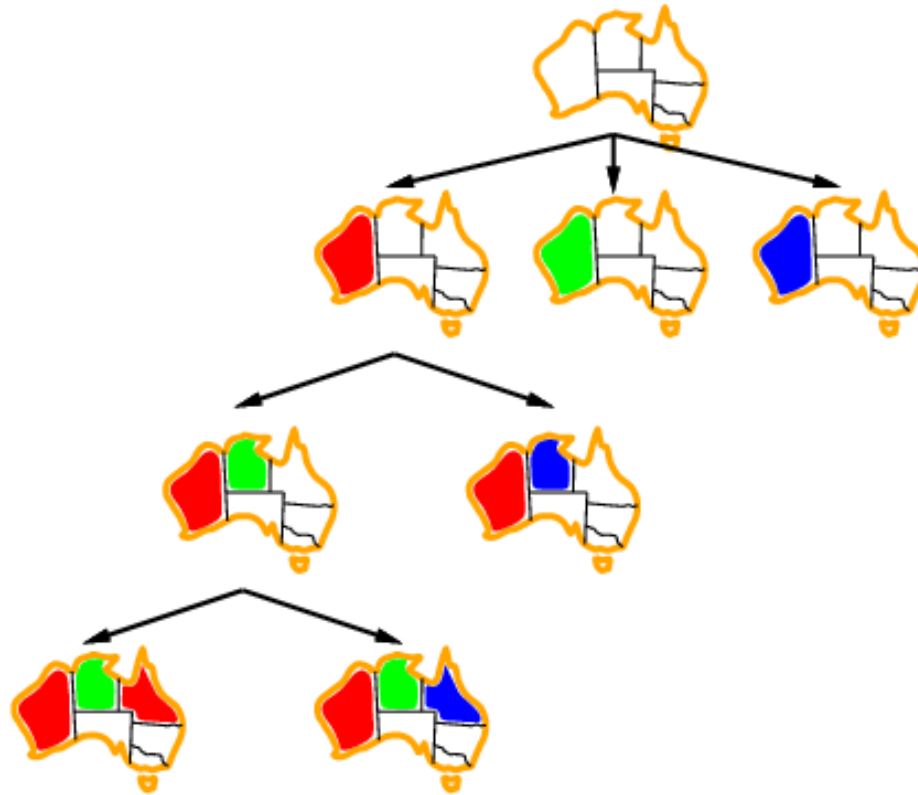
Backtracking example



Backtracking example



Backtracking example

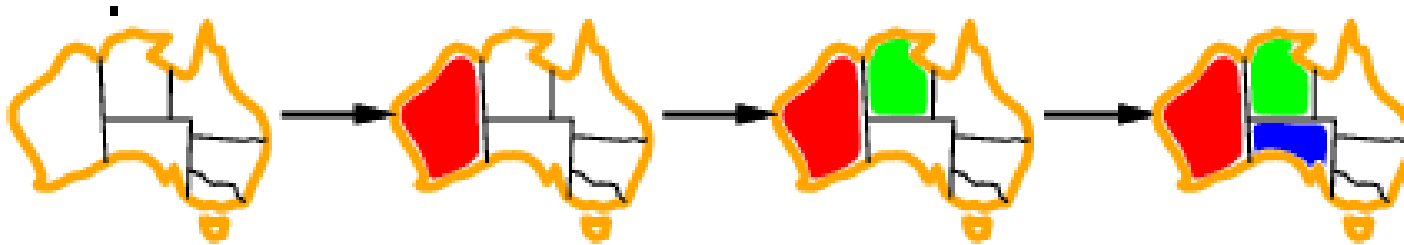


Improving backtracking efficiency

- **General-purpose** methods can give huge gains in speed:
 - Which variable should be assigned next?
 - In what order should its values be tried?
 - Can we detect inevitable failure early?

Most constrained variable

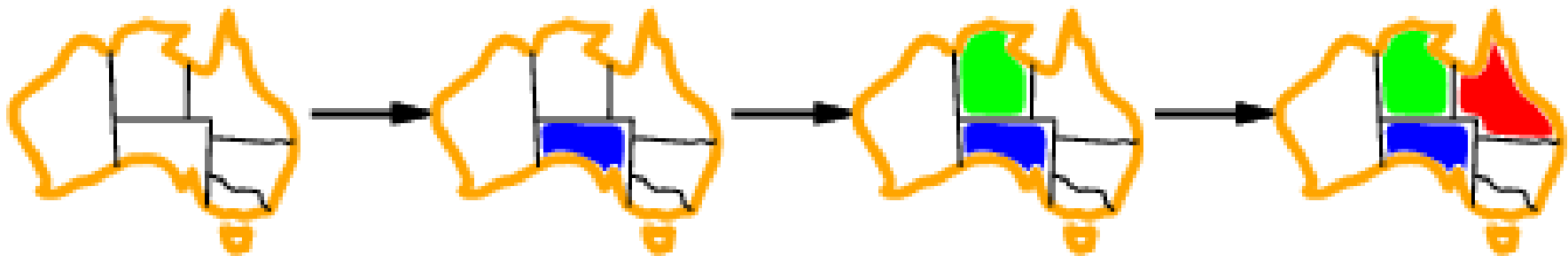
- Most constrained variable:
choose the variable with the fewest legal



- a.k.a. **minimum remaining values (MRV)**
heuristic

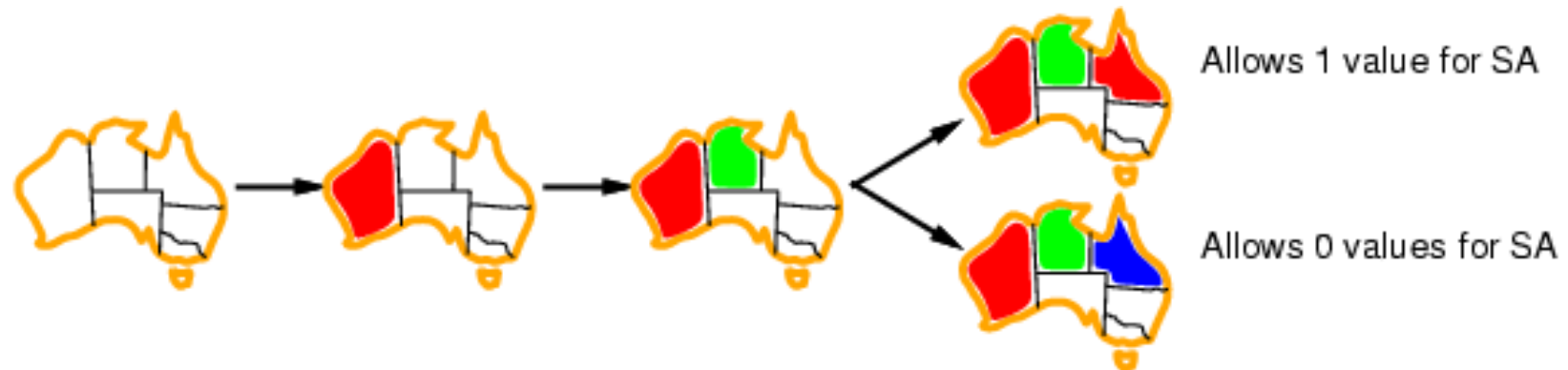
Most constraining variable

- A good idea is to use it as a tie-breaker among most constrained variables
- Most constraining variable:
 - choose the variable with the most constraints on remaining variables



Least constraining value

- Given a variable to assign, choose the least constraining value:
 - the one that rules out the fewest values in the remaining variables



- Combining these heuristics makes 1000 queens feasible

Forward checking

- Idea:
 - Keep track of remaining legal values for unassigned variables
 - Tern



/values



Forward checking

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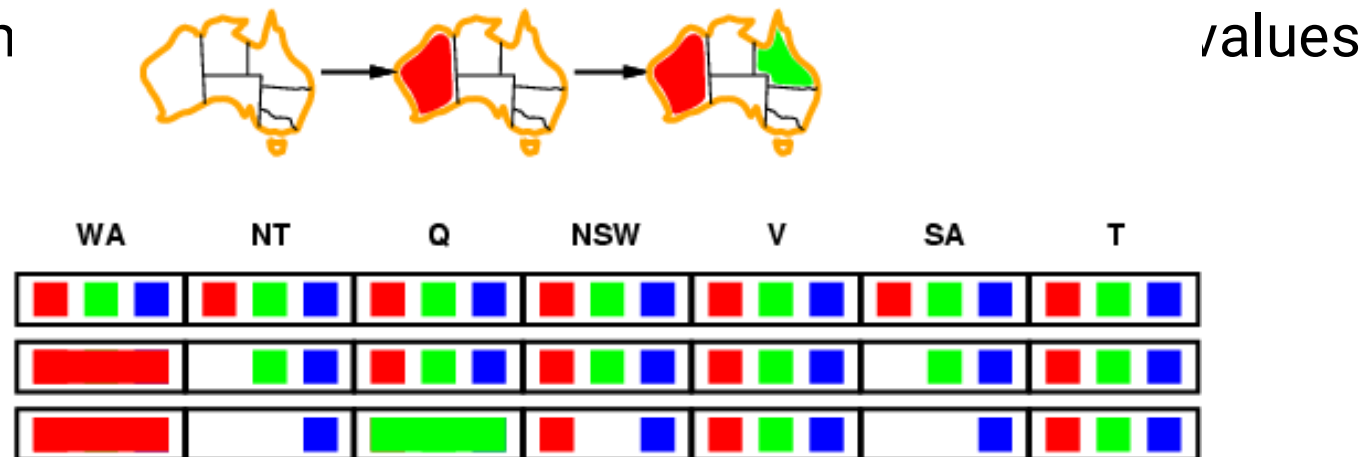


/values

WA	NT	Q	NSW	V	SA	T
<div><div>red</div><div>green</div><div>blue</div></div>	<div><div>red</div><div>green</div><div>blue</div></div>	<div><div>red</div><div>green</div><div>blue</div></div>	<div><div>red</div><div>green</div><div>blue</div></div>	<div><div>red</div><div>green</div><div>blue</div></div>	<div><div>red</div><div>green</div><div>blue</div></div>	<div><div>red</div><div>green</div><div>blue</div></div>
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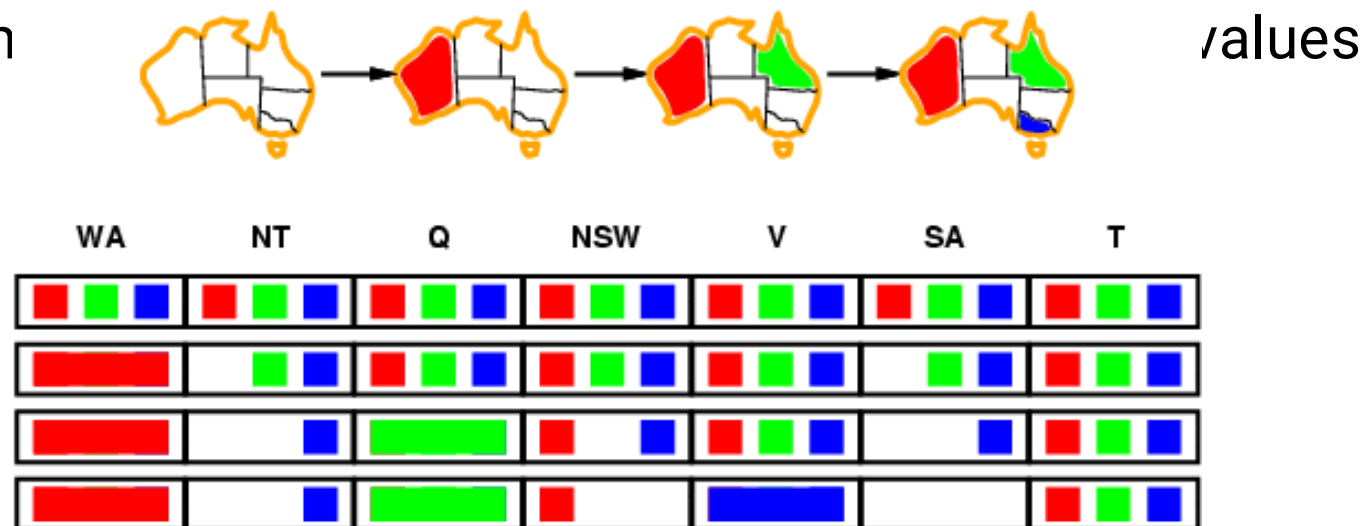
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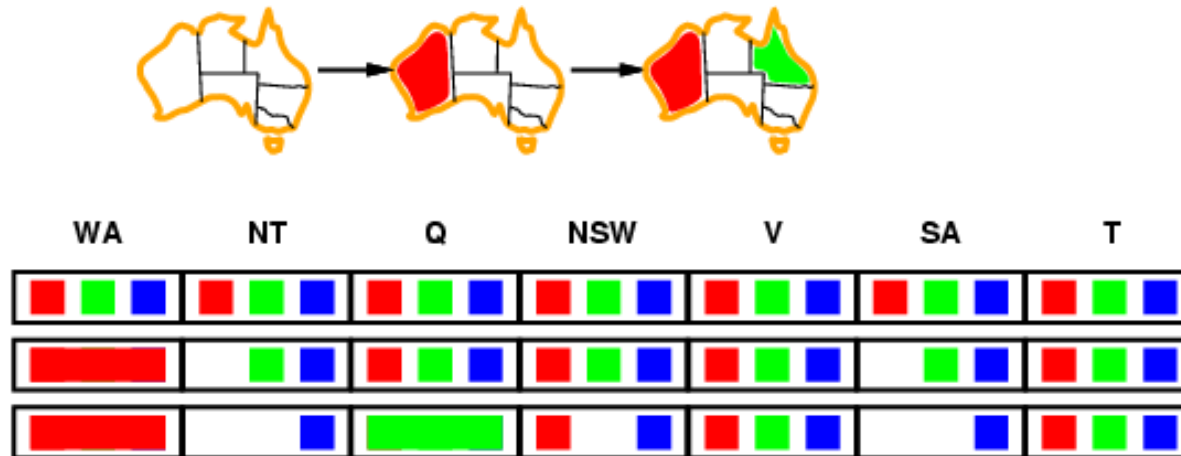
Forward checking

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Constraint propagation

- Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:

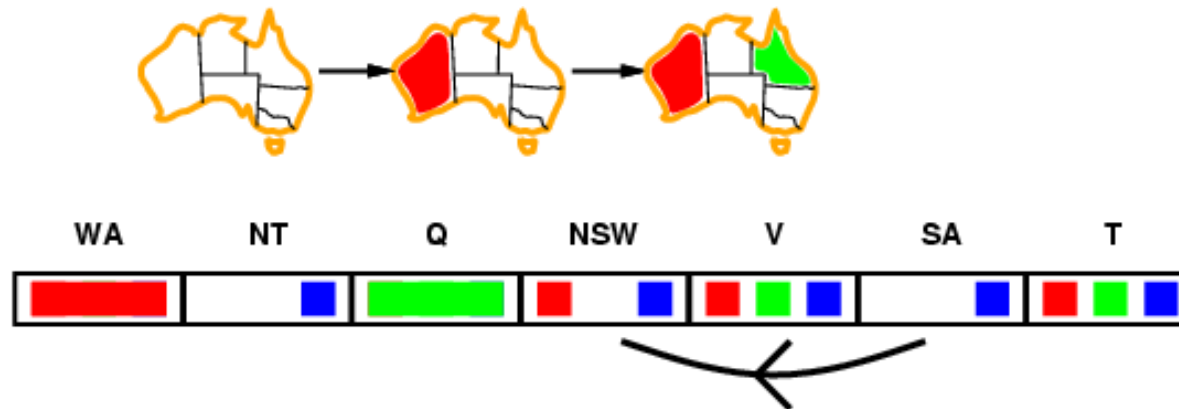


- NT and SA cannot both be blue!
- Constraint propagation algorithms repeatedly enforce constraints locally...

Arc consistency

- Simplest form of propagation makes each arc **consistent**
- $X \rightarrow Y$ is consistent iff

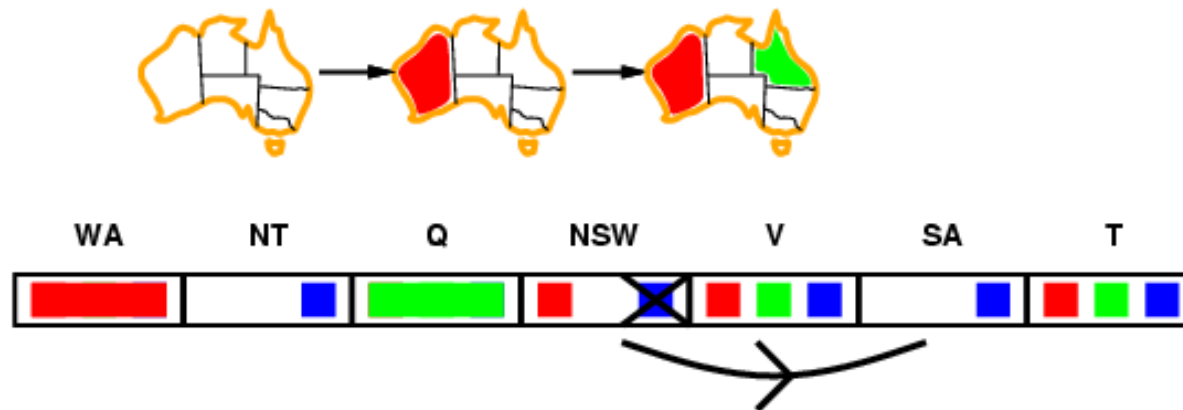
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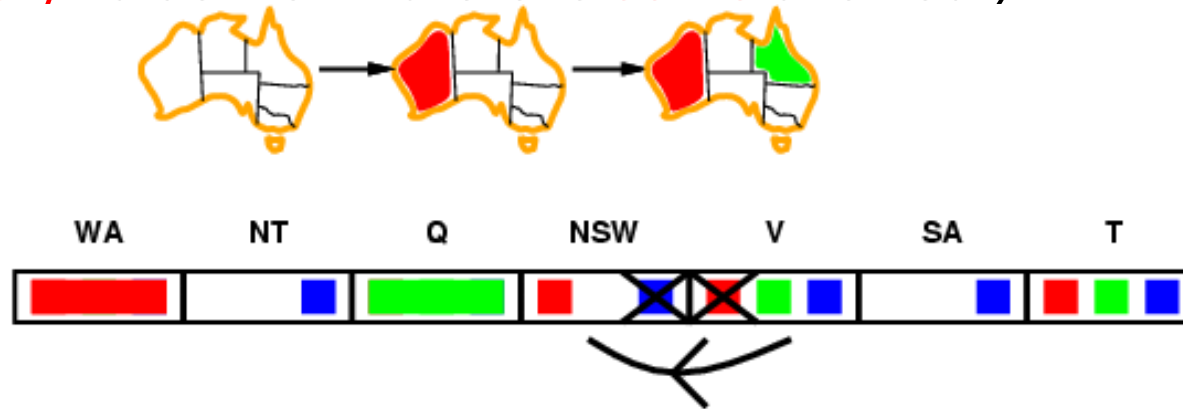
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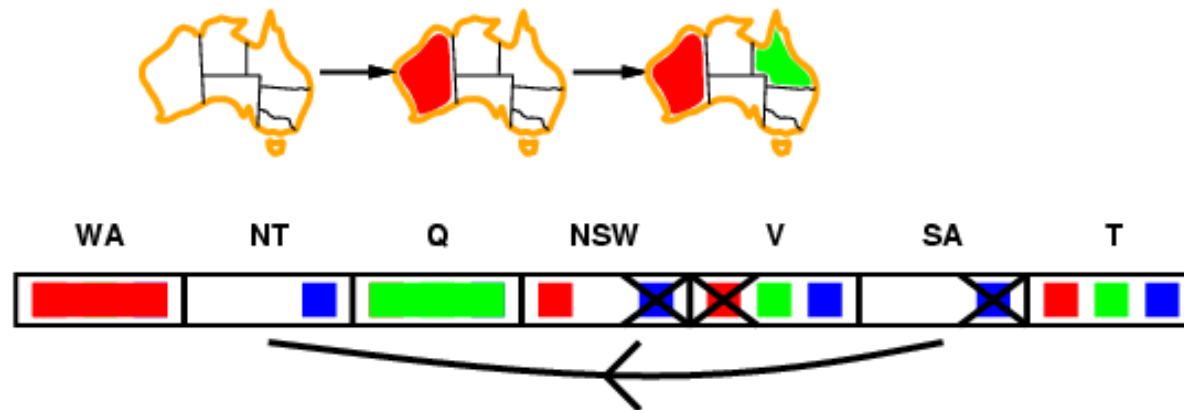


- If X loses a value, neighbors of X need to be rechecked

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- If X loses a value, neighbors of X need to be rechecked
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment