Tutorial on Programming GPUs

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Abstract

The new NVIDIA Fermi GPU architecture has hardware for cache and fast native atomic operations. These features allow one to obtain decent performance with less programming effort than before. A short tutorial on programming GPUs will be presented, making use of simple examples. We will discuss the important concepts of data coalescence, warp divergence, and tiling. Examples will be presented in both Cuda Fortran and Cuda C. If time permits, we will illustrate how these ideas were applied to a Particle-in-Cell code.

Outline of Presentation

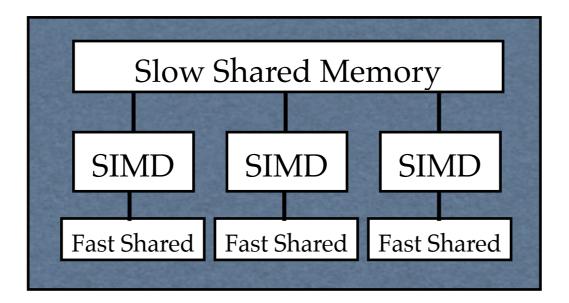
- Abstraction of future computer hardware
- Vector copy example (segmenting problem, data coalescence)
- Transpose example (tiling with shared memory)
- Sum reduction example (device functions)

All examples are in the file GPUTutorial.tar.gz README file describes contents There are 3 versions of each example in this file:

- Cuda C, with Fortran 2003 main program
- Cuda C, with C main program
- Cuda Fortran

Cuda Fortran will be used in this presentation

Simple abstraction of future hardware



- SIMD (vector) unit has multiple cores, each executing the same instruction
- Cores work in lockstep with fast shared memory and local synchronization
- Multiple SIMD units coupled via "slow" shared memory and global synchronization

Each compute node is a powerful computer by itself

• A supercomputer is a hierarchy of such powerful computers

This abstraction applies to both GPUs as well as upcoming multi-core Intel processors

• Programming languages generally differ, however



GPUs are graphical processing units which consist of:

- 12-30 SIMD multiprocessors, each with small (16-48KB), fast (4 clocks) shared memory
- Each multi-processor contains 8-32 processor cores
- Large (0.5-6.0 GB), slow (400-600 clocks) global shared memory, readable by all units
- No cache on some units
- Very fast (1 clock) hardware thread switching

GPU Technology has two special features:

- High bandwidth access to global memory (>100 GBytes/sec)
- Ability to handle thousands of threads simultaneously, greatly reducing memory "stalls"

NVIDIA M2090 has 512 cores!

Programming GPUs

Programming Massively Parallel Processors: A Hands-on Approach by David B. Kirk and Wen-mei W. Hwu [Morgan Kaufmann, 2010].

CUDA by Example: An Introduction to General-Purpose GPU Programming by Jason Sanders and Edward Kandrot [Addison-Wesley, 2011].

CUDA Application Design and Development by Rob Farber [Morgan Kaufmann, 2011].

CUDA Programming: A Developer's Guide to Parallel Computing with GPUs by Shane Cook [Morgan Kaufmann, 2012].

The CUDA Handbook: A Comprehensive Guide to GPU Programming by Nicholas Wilt [Addison-Wesley, 2013].

CUDA Fortran for Scientists and Engineers by Gregory Ruetsch and Massimiliano Fatica [Morgan Kaufmann, 2014].

http://developer.nvidia.com/cuda-downloads

Programming GPUs: Example 1

Let's consider a very simple example: copying one 1D array to another Each SIMD unit processes one block of the total problem

Original vector copy
Parallelized with OpenMP

```
subroutine copy0(a,b)
! simple 1d copy of length nx
! a = b
  implicit none
  real, dimension(:) :: a, b
  integer :: j, nx
  nx = min(size(a,1),size(b,1))

!$OMP PARALLEL DO PRIVATE(j)
  do j = 1, nx
      a(j) = b(j)
  enddo
!$OMP END PARALLEL DO
end subroutine
```

Stride 1 memory access important: Read memory locations in order Serial segmented 1d vector copy
Break up loops into blocks of size mx

```
subroutine copy1(a,b,mx)
! segmented 1d copy of length nx, with block size mx
! a = b
  implicit none
  integer :: mx
  real, dimension(:) :: a, b
  integer :: j, id, nx, nbx, joff
  nx = min(size(a,1), size(b,1))
  nbx = (nx - 1)/mx + 1
  do id = 1, nbx-1 ! outer loop over number of blocks
      joff = mx*id
     do j = 1, min(mx,nx-joff) ! loop over block
         a(j+joff) = b(j+joff)
      enddo
  enddo
  end subroutine
```

Programming GPUs: Copying one 1D array to another

Threads on each block (threadIdx) run on a single SIMD unit, execute the same instruction.

Different blocks (blockIdx) run on different SIMD units

```
subroutine copy1(a,b,mx)
! a = b
  implicit none
  integer :: mx
  real, dimension(:) :: a, b
   integer :: j, id, nx, nbx, joff
  nx = min(size(a,1), size(b,1))
  nbx = (nx - 1)/mx + 1
  do id = 0, nbx-1
      joff = mx*id
      do j = 1, min(mx,nx-joff)
         a(j+joff) = b(j+joff)
      enddo
  enddo
  end subroutine
```

The loop parameters set by host calling function:

```
subroutine gpu copy1(a,b,mx)
! outer part of loop goes here
  implicit none
  integer :: mx, nx, nbx
  real, device, dimension(:) :: a, b
  type (dim3) :: dimBlock, dimGrid
  nx = min(size(a,1), size(b,1))
  nbx = (nx - 1)/mx + 1
  dimBlock = dim3(mx,1,1) ! size of block
  dimGrid = dim3(nbx, 1, 1) ! number of blocks
  call gcopy1<<<dimGrid,dimBlock>>>(a,b,nx)
  crc = cudaThreadSynchronize()
  end subroutine
The inner part of loop in GPU kernel function:
  attributes(global) subroutine gcopy1(a,b,nx)
! inner part of loops goes here
  implicit none
  integer, value :: nx
  real, dimension(nx) :: a, b
  integer :: j, js, id, mx
  mx = blockDim%x ! comes from dimBlock
  j = threadIdx%x ! comes from dimBlock
  id = blockIdx%x
                       ! comes from dimGrid
  js = j + mx*(id - 1)
  if (js \le nx) a(js) = b(js)
end subroutine
```

Programming GPUs: CUDA Fortran

In addition, you must initialize memory on the GPU

```
real, dimension(:), allocatable :: a, b
real, device, dimension(:), allocatable :: g_a, g_b
! allocate host data
  allocate(a(nx),b(nx),c(nx))
! allocate data on GPU, using Fortran90 array syntax
  allocate(g_a(nx),g_b(nx))
```

Copy from host to GPU

```
! Copy data to GPU, using Fortran90 array syntax g_b = b
```

Execute the subroutine:

```
! Execute on GPU: g_a = g_b call gpu_copy1(g_a,g_b,mx)
```

Copy from GPU back to host

```
! Copy data from GPU, using Fortran90 array syntax a = g_a
```

CUDA C is similar but more complex (no array syntax, separate memory spaces)

Programming GPUs: Copying one **2D array** to another

Serial segmented 2d vector copy

```
subroutine copy2(a,b,mx)
! segmented 2d copy of length nx, ny
! with block size mx
! a = b
  implicit none
  integer :: mx
  real, dimension(:,:) :: a, b
  integer :: j, k, nx, ny, js, id, nbx
  nx = min(size(a,1), size(b,1))
  ny = min(size(a,2), size(b,2))
  nbx = (nx - 1)/mx + 1
  do k = 1, ny
     do id = 0, nbx-1
         joff = mx*id
         do j = 1, min(mx,nx-joff)
            a(j+joff,k) = b(j+joff,k)
         enddo
      enddo
   enddo
  end subroutine
```

In Fortran, first index is adjacent in memory

Changes to host calling function: Each y value is a separate block, nbx*ny blocks

```
subroutine gpu_copy2a(a,b,mx)
! outer part of loop goes here
...
real, device, dimension(:,:) :: a, b
ny = min(size(a,2),size(b,2))
dimGrid = dim3(nbx,ny,1)

call gcopy2a<<<dimGrid,dimBlock>>>(a,b,nx,ny)
...
end subroutine
```

GPU kernel function:

```
attributes(global) subroutine gcopy2a(a,b,nx,ny)
implicit none
integer, value :: nx, ny
real, dimension(nx,ny) :: a, b
integer :: j, k, js, id, mx
mx = blockDim%x
j = threadIdx%x
id = blockIdx%x
k = blockIdx%y

js = j + mx*(id - 1)
if ((js <= nx).and.(k <= ny)) a(js,k) = b(js,k)
end subroutine</pre>
```

Programming GPUs: Copying one 2D array to another

Fewer blocks, more work per block

prior GPU kernel function:

Changes to host calling function: One block handles all x values

```
subroutine gpu_copy2b(a,b,mx)
! outer part of loop goes here
...
dimGrid = dim3(ny,1,1)

call gcopy2b<<<dimGrid,dimBlock>>>(a,b,nx,ny)
...
end subroutine
```

GPU kernel function:

end subroutine

Programming GPUs: Copying one 2D array to another

Doubly segmented serial 2d vector copy

```
subroutine copy3(a,b,mx,my)
! segmented 2d copy of length nx, ny
! with block size mx, my
! a = b
   implicit none
  integer :: mx, my
  real, dimension(:,:) :: a, b
   integer :: j, k, nx, ny, joff, koff
   integer :: idx, idy, nbx, nby
  nx = min(size(a,1), size(b,1))
  ny = min(size(a,2), size(b,2))
  nbx = (nx - 1)/mx + 1
  nby = (ny - 1)/my + 1
  do idy = 0, nby-1
  koff = my*idy
  do idx = 0, nbx-1
      joff = mx*idx
      do k = 1, min(my,ny-koff)
         do j = 1, min(mx,nx-joff)
            a(j+joff,k+koff) =
            b(j+joff,k+koff)
         enddo
      enddo
   enddo
  enddo
  end subroutine
```

Changes to host calling function: One block handles some x,y values

```
subroutine gpu_copy3(a,b,mx)
! outer part of loop goes here
...
real, device, dimension(:,:) :: a, b
dimBlock = dim3(mx,my,1)
dimGrid = dim3(nbx,nby,1)

call gcopy3<<<dimGrid,dimBlock>>>(a,b,nx,ny)
...
end subroutine
```

GPU kernel function:

```
attributes(global) subroutine gcopy3(a,b,nx,ny)
implicit none
integer, value :: nx, ny
real, dimension(nx,ny) :: a, b
integer :: j, k, js, ks, idx, idy, mx, my
mx = blockDim%x; my = blockDim%y
j = threadIdx%x; k = threadIdx%y
idx = blockIdx%x; idy = blockIdx%y

ks = k + my*(idy - 1)
js = j + mx*(idx - 1)
if ((js <= nx) .and. (ks <= ny)) then
    a(js,ks) = b(js,ks)
endif
end subroutine</pre>
```

Programming GPUs: Vector copy

Summary

- Processing must be segmented into independent blocks
- Inner loop runs on GPU, loop parameters are set on host
- Data coalescing is important (adjacent threads read adjacent memory locations)
- Each block should execute same instruction (avoid complex if statements)

```
Programming GPUs: CUDA Fortran
Sample output:
make cudaf
./fexample1
   0 :CUDA_DEVICE_NAME=Tesla M2090
  CUDA_MULTIPROCESSOR_COUNT= 16
  CUDA_GLOBAL_MEM_SIZE= 5636554752 ( 5.249451 GB)
  Capability= 20
using device j=0
Fortran empty kernel time= 2.2300000E-04
Fortran 1d copy time= 3.0000001E-06
GPU 1d copy time= 7.4000003E-05
1d copy maximum difference = 0.000000
Fortran 2d copy time= 3.6940000E-03
GPU 2d copy time= 4.3499999E-04
2d copy maximum difference = 0.000000
```

Homework: vary blocksize mx and data lengths nx, ny

Programming GPUs: Example 2
Transpose a 2D array:
First copy in to fast memory
Then copy out from fast memory
Copy to slow memory has stride 1

Original vector transpose

```
subroutine transpose0(a,b)
! simple 2d transpose of length nx,ny
! a = transpose(b)
  implicit none
  real, dimension(:,:) :: a, b
  integer :: j, k, nx, ny
  nx = min(size(a,2),size(b,1))
  ny = min(size(a,2),size(b,2))

do k = 1, ny
  do j = 1, nx
       a(k,j) = b(j,k)
  enddo
enddo
end subroutine
```

In original, stride 1 memory access is obtained only when reading b array

```
subroutine transpose2(a,b,mx,my)
! segmented 2d transpose of length nx, ny
! with block size mx, my
! a = transpose(b)
  implicit none
  integer :: mx, my, j, k, nx, ny
  real, dimension(:,:) :: a, b
  integer :: idx, idy, joff, koff, nbx, nby
  real, dimension(mx+1,my) :: s
  nx = min(size(a,2), size(b,1))
  ny = min(size(a,1), size(b,2))
  nbx = (nx - 1)/mx + 1; nby = (ny - 1)/my + 1
  do idy = 0, nby-1
     koff = my*idy
     do idx = 0, nbx-1
         joff = mx*idx
         do k = 1, min(my,ny-koff)
            do j = 1, min(mx,nx-joff)
               s(j,k) = b(j+joff,k+koff)
            enddo
         enddo
         do j = 1, min(mx,nx-joff)
            do k = 1, min(my,ny-koff)
               a(k+koff,j+joff) = s(j,k)
            enddo
         enddo
     enddo
  enddo
  end subroutine
```

Programming GPUs: Transpose 2D array On GPU, mx = my

```
subroutine transpose2(a,b,mx,my)
! a = transpose(b)
  real, dimension(mx+1,my) :: s
  do idy = 0, nby-1
     koff = my*idy
      do idx = 0, nbx-1
         joff = mx*idx
         do k = 1, min(my,ny-koff)
            do j = 1, min(mx,nx-joff)
               s(j,k) =
               b(j+joff,k+koff)
            enddo
         enddo
         do j = 1, min(mx,nx-joff)
            do k = 1, min(my,ny-koff)
               a(k+koff,j+joff) =
               s(j,k)
            enddo
         enddo
      enddo
  enddo
  end subroutine
```

The inner part of loop in GPU kernel function:

```
attributes(global) subroutine
                      gtranspose2(a,b,nx,ny)
! inner part of loops go here
  implicit none
  integer, value :: nx, ny
  real, dimension(ny,nx) :: a
  real, dimension(nx,ny) :: b
  integer :: j, k, js, ks, jb, kb
  integer :: joff, koff, mx, mxv
  real, shared, dimension(*) :: s
  mx = blockDim%x; mxv = mx + 1
  j = threadIdx%x; k = threadIdx%y
  idx = blockIdx%x; idy = blockIdx%y
  koff = mx*(idy - 1)
  joff = mx*(idx - 1)
  ks = k + koff
  js = j + joff
  if ((js \le nx) .and. (ks \le ny)) then
      s(j+mxv*(k-1)) = b(js,ks)
  endif
  call syncthreads() ! synchronize threads
  is = k + ioff
  ks = j + koff
  if ((js \le nx) .and. (ks \le ny)) then
      a(ks,js) = s(k+mxv*(j-1))
  endif
  end subroutine
```

GPU kernel function:

```
attributes(global) subroutine
           gtranspose2(a,b,nx,ny)
real, shared, dimension(*) :: s
mx = blockDim%x; mxv = mx + 1
j = threadIdx%x; k = threadIdx%y
idx = blockIdx%x; idy = blockIdx%y
koff = mx*(idy - 1)
joff = mx*(idx - 1)
ks = k + koff
js = j + joff
if ((js \le nx).and.(ks \le ny)) then
   s(j+mxv*(k-1)) = b(js,ks)
endif
call syncthreads()
is = k + ioff
ks = j + koff
if ((js \le nx).and.(ks \le ny)) then
   a(ks,js) = s(k+mxv*(j-1))
endif
end subroutine
```

Host calling function similar to gpu_copy3

```
subroutine gpu_transpose2(a,b,mx)
...
real, device, dimension(:,:) :: a, b
...
dimBlock = dim3(mx,mx,1)
dimGrid = dim3(nbx,nby,1)
! calculate size of shared memory
ns = (mx + 1)*mx*sizeof(a(1,1))

call gtranspose2<<<dimGrid,dimBlock,ns>>>(a,b,nx,ny)
crc = cudaThreadSynchronize()
end subroutine
```

Programming GPUs: Transpose

Summary

- Stride 1 access is obtained to slow global memory
- Stride 1 access is not obtained to fast shared memory, but it is 100x faster
- Processing in pieces small enough to fit in fast memory (tiling) is important

Programming GPUs: CUDA Fortran

Sample output:

./fexample2

using device j= 0
Fortran empty kernel time= 2.2800000E-04
Fortran 2d transpose time= 4.2500001E-04
GPU 2d transpose time= 2.5099999E-04
2d transpose maximum difference = 0.000000

Programming GPUs: Example 3

Sum reduction of 1D array

Original sum reduction Parallelized with OpenMP

```
subroutine sum0(a,sa)
! simple 1d sum reduction of length nx
! sa = sum(a)
  implicit none
  real :: sa
  real, dimension(:) :: a
  integer :: j

  sa = 0.0
!$OMP PARALLEL DO PRIVATE(j)
!$OMP& REDUCTION(+:sa)
  do j = 1, size(a,1)
      sa = sa + a(j)
  enddo
!$OMP END PARALLEL DO
  end subroutine
```

Stride 1 memory access important: Read memory locations in order Serial segmented 1d sum reduction Partial sums of blocks of size mx

```
subroutine sum1(a,sa,mx)
! 1d sum reductions, each of length mx
! sa = sum(a)
   implicit none
  integer :: mx
  real :: sa
  real, dimension(:) :: a
  integer :: j, js, jb, nx, nbx
  real :: t
  nx = size(a,1)
  nbx = (nx - 1)/mx + 1
  sa = 0.0
  do jb = 1, nbx
  t = 0.0
  do js = 1, min(mx,nx-mx*(jb-1))
      j = js + mx*(jb - 1)
     t = t + a(j)
  enddo
   sa = sa + t
  enddo
  end subroutine
```

Programming GPUs: Sum reduction of 1D array

On GPU, copy to shared memory one thread performs partial sum Not very efficient

```
subroutine sum1(a,sa,mx)
! sa = sum(a)
  real :: sa
  real, dimension(:) :: a
  sa = 0.0
  do jb = 1, nbx
  t = 0.0
  do js = 1, min(mx,nx-mx*(jb-1))
      j = js + mx*(jb - 1)
     t = t + a(j)
  enddo
  sa = sa + t
  enddo
  end subroutine
```

The inner part of loop in GPU kernel function:

```
attributes(global) subroutine gsum1(a,sa,nx)
! inner part of loop goes here
  implicit none
  integer, value :: nx
  real, dimension(:) :: a, sa
  integer :: j, js, jb, mx, joff, mxm
  real :: t
  real, shared, dimension(*) :: s
  mx = blockDim%x
  is = threadIdx%x
  jb = blockIdx%x
  joff = mx*(jb - 1)
  j = js + joff ! first copy to shared memory
  if (j \le nx) s(js) = a(j)
  call syncthreads()
  if (js==1) then ! one thread performs sum
     mxm = nx - joff
     if (mxm > mx) mxm = mx
     t = 0.0
     do j = 1, mxm
        t = t + s(j)
                         ! sum different blocks
     enddo
     t = atomicAdd(sa(1),t)
  endif
  end subroutine
```

Programming GPUs: Sum reduction of 1D array

GPU kernel function:

```
attributes(global) subroutine
           gsum1(a,sa,nx)
real, shared, dimension(*) :: s
mx = blockDim%x
is = threadIdx%x
jb = blockIdx%x
joff = mx*(jb - 1)
j = js + joff
if (j \le nx) s(js) = a(j)
call syncthreads()
if (js==1) then mxm = nx - joff
   if (mxm > mx) mxm = mx
   t = 0.0
   do j = 1, mxm
      t = t + s(j)
   enddo
   t = atomicAdd(sa(1),t)
endif
```

host calling function:

```
subroutine gpu_sum1(a,sa,mx)
implicit none
integer :: mx
real, device, dimension(:) :: a, sa
integer :: nx, nbx, ns
type (dim3) :: dimBlock, dimGrid
nx = size(a,1)
nbx = (nx - 1)/mx + 1
dimBlock = dim3(mx,1,1)
dimGrid = dim3(nbx,1,1)
sa(1) = 0.0
! calculate size of shared memory
ns = mx*sizeof(a(1))

call gsum1<<<dimGrid,dimBlock,ns>>>(a,sa,nx)
end subroutine
```

end subroutine

Parallel Sum Reduction:

0:	1	2	3	4	5	6	7	8
			—	V				<u> </u>
1:	6	8	10	12	5	6	7	8
2:	16	20	10	12	5	6	7	8
3:	36	20	10	12	5	6	7	8

Programming GPUs: Sum reduction of 1D array

Implementation of local parallel sum reduction device functions are inlined, sdata is in shared memory

```
attributes(device) subroutine lsum2(sdata,n)
! finds local sum of n data items shared by threads
! using binary tree method. input is modified
  implicit none
  real, dimension(*) :: sdata
  integer, value :: n
  integer :: 1, k
  real :: s
  1 = threadIdx%x
  k = blockDim%x/2
  s = 0.0
  if (1 \le n) s = sdata(1)
  do while (k > 0)
      if (1 \le k) then
         if ((1+k) \le n) then
            s = s + sdata(1+k)
            sdata(1) = s
         endif
      endif
     call syncthreads()
     k = k/2
  enddo
  end subroutine
```

Programming GPUs: Sum reduction of 1D array

Prior GPU kernel function:

```
attributes(global) subroutine
           qsum1(a,sa,nx)
real, shared, dimension(*) :: s
mx = blockDim%x
js = threadIdx%x
jb = blockIdx%x
joff = mx*(jb - 1)
j = js + joff
if (j \le nx) s(js) = a(j)
call syncthreads()
if (js==1) then
   mxm = nx - joff
   if (mxm > mx) mxm = mx
   t = 0.0
   do j = 1, mxm
      t = t + s(j)
   enddo
   t = atomicAdd(sa(1),t)
endif
end subroutine
```

Multiple parallel sum reductions: write out partial sum of each block

GPU kernel function:

```
attributes(global) subroutine gsum2(a,d,nx)
implicit none
integer, value :: nx
real, dimension(:) :: a, d
integer :: j, js, jb, mx, joff, mxm
real, shared, dimension(*) :: s
mx = blockDim%x
is = threadIdx%x
jb = blockIdx%x
joff = mx*(jb - 1)
j = js + joff
if (j \le nx) s(js) = a(j)
call syncthreads()
mxm = nx - joff
if (mxm > mx) mxm = mx
call lsum2(s,mxm)
if (js==1) d(jb) = s(1)
end subroutine
```

Programming GPUs: Sum reduction of 1D array

Multiple parallel sum reductions: write out partial sum of each block

GPU kernel function:

```
attributes(global) subroutine
           qsum2(a,d,nx)
implicit none
integer, value :: nx
real, dimension(:) :: a, d
integer :: j, js, jb, mx, joff, mxm
real, shared, dimension(*) :: s
mx = blockDim%x
js = threadIdx%x
jb = blockIdx%x
joff = mx*(jb - 1)
j = js + joff
if (j \le nx) s(js) = a(j)
call syncthreads()
mxm = nx - joff
if (mxm > mx) mxm = mx
call lsum2(s,mxm)
if (js==1) d(jb) = s(1)
```

host calling function:

```
subroutine gpu sum2(a,d,mx)
  implicit none
  integer :: mx
  real, device, dimension(:) :: a, d
  integer :: nx, nbx, ns
  type (dim3) :: dimBlock, dimGrid
  nx = size(a,1)
  nbx = (nx - 1)/mx + 1
  dimBlock = dim3(mx,1,1)
  dimGrid = dim3(nbx,1,1)
! calculate size of shared memory
  ns = mx*sizeof(a(1))
  call gsum2<<<dimGrid,dimBlock,ns>>>(a,d,nx)
  crc = cudaThreadSynchronize()
  end subroutine
```

end subroutine

Programming GPUs: Sum reduction of 1D array

Many times only local sums are required If global sum is needed, one can combine these two procedures

```
Perform parallel sums for each block:

call gpu_sum2(g_a,g_d,mx)

Then add the partial sums with serial algorithm:

call gpu_sum1(g_d,g_s,mx)
```

If array is long, one can iterate gpu_sum2 several times. This is done by the host procedure gpu_sum3: call gpu_sum3(g_a,g_d,g_s,mx)

Sample output:

./fexample3

Fortran empty kernel time= 5.1999999E-05 Fortran 1d sum time= 4.0000000E-06 GPU 1d sum time= 2.5300001E-04 1d sum maximum difference = 0.000000 s,t = 4501500. 4501500.

Programming GPUs: Conclusions

- Vector algorithms are relatively easy
- Processing has to be done in small blocks
- Tiling algorithms are well known from cache based machines
- Sum reductions are harder
- Irregular problems such as reordering data can be very hard
- Very useful to develop a serial version before implementing a Cuda version
- Libraries such as BLAS, FFTs, are available to avoid reinventing the wheel
- New languages such as OpenACC (similar to OpenMP) are evolving