

Advanced Topics in Machine Learning 2024

Amartya Sanyal

Home Assignment 5

Deadline: 23:59, Thursday, 24th October 2024

The assignments must be submitted individually – each student must write and submit a personal solution but we do not prevent you from discussing **high-level** ideas in small groups. If you use any LLM tool such as ChatGPT, please specify the **purpose** and **manner** in which you have utilized it.

We are interested in how you solved the problems, and not just in the final answers. Please explain your solutions and ideas as clearly as you can.

Late Penalty and multiple Submissions Late submissions will incur a penalty of 10% of the total marks for every hour of delay (rounded up) with a maximum allowed delay of 5 hours after which the submission server will close. If you submit multiple submissions, only the last submission will be considered relevant both for grading answers as well as late penalty.

Submission format: Please upload your answers in a single .pdf file. If you have created code please include at least the key parts in the PDF as well as a link to the full code (on Github, Colab, or similar).

We have studied the definition for Differential Privacy (DP)-Probably Approximately Correct (PAC) learning. We will now dive deep into it and design a generic DP-PAC learner. But first, let's define it again.

Definition 1. For any $\alpha, \beta, \epsilon > 0$, an algorithm $\mathcal{A}_{DP}^{\alpha, \beta, \epsilon} : (\mathcal{X} \times \mathcal{Y})^n \rightarrow \mathcal{H}^1$ is said to be an $(\alpha, \beta, \epsilon)$ -DP-PAC learner for hypothesis class \mathcal{H} with private sample complexity n , if for all $h^* \in \mathcal{H}$ and for all distributions \mathcal{D} over \mathcal{X} , the algorithm satisfies the following two properties.

- **Privacy** For any dataset S of size n sampled from \mathcal{D} and labelled by h^* , $\mathcal{A}_{DP}^{\alpha, \beta, \epsilon}$ is ϵ -DP with respect to S .
- **Utility** Define $\mathcal{R}(\hat{h}; \mathcal{D})$ to be the 0-1 error of \hat{h} on distribution \mathcal{D} labelled by h^* . Then,

$$\Pr_S \left[\mathbb{E}_{\hat{h} \sim \mathcal{A}_{DP}(S)} \left[\mathcal{R}(\hat{h}; \mathcal{D}) \right] \leq \alpha \right] \geq 1 - \beta$$

where the outer probability is over all datasets S of size n sampled from \mathcal{D} and labelled with h^* .

Exercise 1 (50 points). For any $d > 0$, let \mathcal{H}_d be any finite hypothesis class such that $|\mathcal{H}_d| \leq \exp(\text{poly}(d))$. For any such \mathcal{H}_d , design a generic $(\alpha, \beta, \epsilon)$ -DP-PAC learner with sample size polynomial in $d, \frac{1}{\alpha}, \frac{1}{\epsilon}, \log\left(\frac{1}{\beta}\right)$. Your solution does not need to be computationally efficient. You must write down the algorithm as well as the proof for its privacy and utility guarantee.

1. 10 points are reserved for a clear pseudocode of the algorithm.
2. 10 points are reserved for the proof of privacy.
3. 30 points are reserved for the proof utility. This proof contains two main steps which are described in the hint below and each carries equal weight of 15 points each.

¹For simplicity, we will shorthand $\mathcal{A}_{DP}^{\alpha, \beta, \epsilon}$ with \mathcal{A}_{DP} .

Hint: First step involves showing that no matter what \hat{h} is output, its error on the train set will be close to the error on the distribution. This argument involves a combination of Chernoff bound and union bound similar to what we did in class. Second step involves showing that we can privately find a \hat{h} with good training error. Each of these two steps carry 20 points each.

The above result dealt with the case of finite hypothesis class \mathcal{H} and any domain \mathcal{X} . Now we will look at **any** hypothesis class with finite VC dimension and finite input domain.

Exercise 2 (35 points). For any $d, p > 0$, let \mathcal{X}_p be a finite input domain such that $|\mathcal{X}_p| \leq \exp(p)$ and \mathcal{H}_d be any hypothesis class on $|\mathcal{X}_p|$ with VC dimension d . For any such $\mathcal{H}_d, \mathcal{X}_p$, design a generic $(\alpha, \beta, \epsilon)$ -DP-PAC learner with sample size polynomial in $d, p, \frac{1}{\alpha}, \frac{1}{\epsilon}, \log\left(\frac{1}{\beta}\right)$. Your solution does not need to be computationally efficient. You must write down the algorithm as well as the proof for its privacy and utility guarantee.

Hint: First step involves finding a finite set of hypothesis which is sufficient to output from. You will have to use the Sauer-Shelah lemma in this step. Second step is applying the previous theorem in Exercise 1.

Exercise 3 (15 points). Based on the above results and non-private learning theory, comment on the differences between private and non-private PAC learning from both statistical (amount of data required for small error) as well as computational (amount of time taken) perspective.

Do not write more than 100 words.

Good luck!

Amartya