

CS2102: Database Systems

T01: Relational Model & Algebra

Question 1

Question 1

Recap

- Superkey

- Key

Relation

Question

Recap

Superkey

Definition

A **superkey** is a subset of attributes that uniquely identifies a tuple in a relation.

Terminology

- **Relation** A table *(as a set, the order of rows doesn't matter)*
- **Attribute** A column of the table
- **Tuple** A row on the table *(as a tuple, the order of values or columns matter)*

Question 1

Recap

- Superkey

- Key

Relation

Question

Recap

Superkey

Definition

A **superkey** is a subset of attributes that uniquely identifies a tuple in a relation.

Violations

Given a relation R , a set of attributes A is not a superkey if there are two tuples t_1 and t_2 such that

- $\forall a \in A : t_1.a = t_2.a$ *(i.e., all values in attributes in A are equal)*
- $\exists b \in \text{Attr}(R) - A : t_1.b \neq t_2.b$ *(i.e., at least one value in other attributes that is different)*

In this case, we have a *counter-example* where there is a value $t_1.b$ and $t_2.b$ that cannot be uniquely identified.

Question 1

Recap

- *Superkey*

- *Key*

Relation

Question

Recap

Key

Definition

A **key** is a **superkey** that is also minimal (i.e., no **proper subset** of the key is a superkey).

Important

| Minimal \neq Minimum

Terminology

- **Possible**
- **Definite**

There is no violation, at least in the current instance
No possible (*valid*) instance will violate this

Question 1

Recap
Relation
Question

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

How?

Any idea?

Question 1

Recap
Relation

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- Q1A
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- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
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1	1	2	0
0	0	1	2

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

How?

Steps

1. Pick a set of attributes
2. Try to find **violations!**
 - o If none found, then it is *possible* that the set is a superkey
 - o Otherwise, the set is *definitely* not a superkey

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Example

- Consider $\{B, C\}$

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

How?

Steps

1. Pick a set of attributes
2. Try to find **violations!**
 - If none found, then it is *possible* that the set is a superkey
 - Otherwise, the set is *definitely* not a superkey

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Example

- Consider $\{B, C\}$
- Look at row 2 & 3

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

How?

Steps

1. Pick a set of attributes
2. Try to find **violations!**
 - If none found, then it is *possible* that the set is a superkey
 - Otherwise, the set is *definitely* not a superkey

Question 1

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Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Example

- Consider $\{B, C\}$
- Look at row 2 & 3
- Different value for A!

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

How?

Steps

1. Pick a set of attributes
2. Try to find **violations!**
 - If none found, then it is *possible* that the set is a superkey
 - Otherwise, the set is *definitely* not a superkey

Question 1

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Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Example

- Consider $\{A, B, C\}$

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

How?

Steps

1. Pick a set of attributes
2. Try to find **violations!**
 - If none found, then it is *possible* that the set is a superkey
 - Otherwise, the set is *definitely* not a superkey

Question 1

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Relation

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- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Example

- Consider $\{A, B, C\}$
- All different values
 - No violation

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

How?

Steps

1. Pick a set of attributes
2. Try to find **violations!**
 - If none found, then it is *possible* that the set is a superkey
 - Otherwise, the set is *definitely* not a superkey

Question 1

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0	0	1	2

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

How?

Steps

1. Pick a set of attributes
2. Try to find **violations!**
 - o If none found, then it is *possible* that the set is a superkey
 - o Otherwise, the set is *definitely* not a superkey

Try it out (3 minutes)

Question 1

Recap
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- Q1B
- Q1C

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$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

Answer

- {A,C}
- {A,D}
- {A,B,C,D}
- {A,B,C}
- {A,B,D}
- {A,C,D}

Question 1

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$R(A,B,C,D)$

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A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

Answer

- {A,C}
- {A,D}
- {A,B,C,D}
- {A,B,C}
- {A,B,D}
- {A,C,D}

Question

Do you need check **ALL** possible subset of attributes? (*how many are there?*)

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

Answer

- $\{A,C\}$
- $\{A,D\}$
- $\{A,B,C,D\}$
- $\{A,B,C\}$
- $\{A,B,D\}$
- $\{A,C,D\}$

Improvement

- Consider $\{A,C\}$ (*it is a superkey*)

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1A

Question

Assuming that the given relation instance is *valid*, write down all the *possible* superkeys of R.

Answer

- {A,C}
- {A,D}
- {A,B,C,D}
- {A,B,C}
- {A,B,D}
- {A,C,D}

Improvement

- Consider {A,C} (*it is a superkey*)
- All superset of {A,C} are also superkey! (*why?!*)

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1B

Question

Suppose that $\{A,C\}$ is definitely a superkey of R. Write down all the possible candidate keys of R.

How?

Do we need to recompute *everything*?

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1B

Question

Suppose that $\{A,C\}$ is definitely a superkey of R. Write down all the possible candidate keys of R.

Superkey

- $\{A,C\}$
- $\{A,D\}$
- $\{A,B,C,D\}$
- $\{A,B,C\}$
- $\{A,B,D\}$
- $\{A,C,D\}$

Steps

1. Highlight the sets that are *definitely* superkeys

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1B

Question

Suppose that $\{A,C\}$ is definitely a superkey of R. Write down all the possible candidate keys of R.

Superkey

- $\{A,C\}$
- $\{A,D\}$
- $\{A,B,C,D\}$
- $\{A,B,C\}$
- $\{A,B,D\}$
- $\{A,C,D\}$

Steps

1. Highlight the sets that are *definitely* superkeys
2. Remove all supersets of these sets as they are *definitely NOT* a key

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1B

Question

Suppose that $\{A,C\}$ is definitely a superkey of R. Write down all the possible candidate keys of R.

Superkey

- $\{A,C\}$
- $\{A,D\}$
- $\{A,B,C,D\}$
- $\{A,B,C\}$
- $\{A,B,D\}$
- $\{A,C,D\}$

Steps

1. Highlight the sets that are *definitely* superkeys
2. Remove all supersets of these sets as they are *definitely NOT* a key
3. The rest are the *possible* keys

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1B

Question

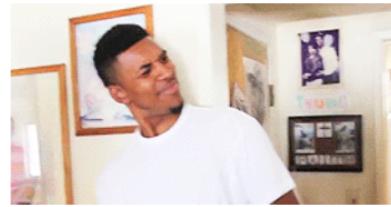
Suppose that $\{A,C\}$ is definitely a superkey of R. Write down all the possible candidate keys of R.

Superkey

- $\{A,C\}$
- $\{A,D\}$
- $\{A,B,C,D\}$
- $\{A,B,C\}$
- $\{A,B,D\}$
- $\{A,C,D\}$

Huh?

Why is $\{A,B,D\}$ still a possible key when we have $\{A,D\}$?



Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2
0	0	1	1

Q1B

Question

Suppose that $\{A,C\}$ is definitely a superkey of R. Write down all the possible candidate keys of R.

Superkey

- $\{A,C\}$
- $\{A,D\}$
- $\{A,B,C,D\}$
- $\{A,B,C\}$
- $\{A,B,D\}$
- $\{A,C,D\}$

Ans

The table can still be extended to make $\{A,D\}$ not a superkey.

Question 1

Recap
Relation

Question

- Q1A
- Q1B
- Q1C

Relation

$R(A,B,C,D)$

R

A	B	C	D
0	0	0	1
2	1	2	0
1	1	2	0
0	0	1	2

Q1C

Question

Which of these (*if any*) is likely to be the candidate key of R?

Discuss (3 minutes)

Question 2

Question 2

Recap
- Foreign Key
Relation
Question

Recap

Superkey

Definition

A **foreign key** is a subset of attributes of relation R_1 that refers to the primary key of relation R_2 .

Requirement: Each foreign key in R_1 must satisfy one of the following:

- Appear as **primary key** in R_2
- Be a NULL value (*or a tuple containing at least one NULL value*)

Terminology

- **Referencing**
- **Referenced**

R_1 (*check the value in this table*)

R_2 (*that the value appears in this table*)

Question 2

Recap
Relation
Question

Relation

$R(A,B)$

R

A	B
3	0
2	1
1	1
0	0

$S(W,X,Y,Z)$

S

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

A	B
3	0
2	1
1	1
0	0

$S(W,X,Y,Z)$

S

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R
(i.e., R.A).

How?

Any idea?

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B
3	0
2	1
1	1
0	0

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R
(i.e., R.A).

Worst-Case

For each column, check **row-by-row**
(maybe skip any NULL value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B
3	0
2	1
1	1
0	0

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Example

- Consider S.W

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row** (maybe skip any NULL value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

A	B
3	0
2	1
1	1
0	0

$S(W,X,Y,Z)$

S

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Example

- Consider S.W
- Check **row 1**

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row** (maybe skip any **NULL** value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

A	B
3	0
2	1
1	1
0	0

$S(W,X,Y,Z)$

S

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Example

- Consider S.W
- Check **row 1 ; row 2**

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row** (maybe skip any **NULL** value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B	W	X	Y	Z
3	0	0	4	0	NULL
2	1	1	NULL	2	NULL
1	1	2	1	2	NULL
0	0	3	0	1	NULL

Example

- Consider S.W
- Check **row 1 ; row 2 ; row 3**

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row** (maybe skip any **NULL** value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B
3	0
2	1
1	1
0	0

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Example

- Consider S.W
- Check **row 1 ; row 2 ; row 3 ; row 4**

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row** (maybe skip any **NULL** value)

Question 2

Recap
Relation
Question
-Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B	W	X	Y	Z
3	0	0	4	0	NULL
2	1	1	NULL	2	NULL
1	1	2	1	2	NULL
0	0	3	0	1	NULL

Example

- Consider S.W
- Check **row 1 ; row 2 ; row 3 ; row 4**
- **Done** (*no violation, so it's possible*)

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R
(i.e., R.A).

Worst-Case

For each column, check **row-by-row**
(maybe skip any **NULL** value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B
3	0
2	1
1	1
0	0

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Example

- Consider S.X

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row** (maybe skip any NULL value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B
3	0
2	1
1	1
0	0

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Example

- Consider S.X
- Check **row 1**

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row**
(maybe skip any **NULL** value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B	W	X	Y	Z
3	0	0	4	0	NULL
2	1	1	NULL	2	NULL
1	1	2	1	2	NULL
0	0	3	0	1	NULL

Example

- Consider S.X
- Check **row 1**
- **Done** (*failed, so it's impossible*)

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row** (maybe skip any **NULL** value)

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

$S(W,X,Y,Z)$

S

A	B
3	0
2	1
1	1
0	0

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R (i.e., R.A).

Worst-Case

For each column, check **row-by-row** (maybe skip any **NULL** value)

Can you try the rest
(3 minutes)?

Question 2

Recap
Relation
Question
- Q2

Relation

$R(A,B)$

R

A	B
3	0
2	1
1	1
0	0

$S(W,X,Y,Z)$

S

W	X	Y	Z
0	4	0	NULL
1	NULL	2	NULL
2	1	2	NULL
3	0	1	NULL

Q2

Question

Write down all the possible foreign key in S that refer to attribute A in R
(i.e., R.A).

Answer

- $(S.W) \rightsquigarrow (R.A)$
- $(S.Y) \rightsquigarrow (R.A)$
- $(S.Z) \rightsquigarrow (R.A)$ (since all NULL)

Question 3

Question 3

Recap

- Unary

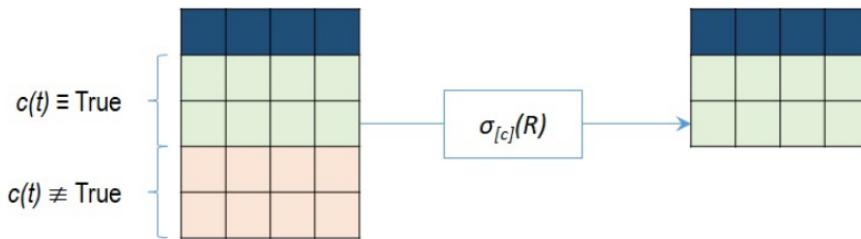
Schema

Question

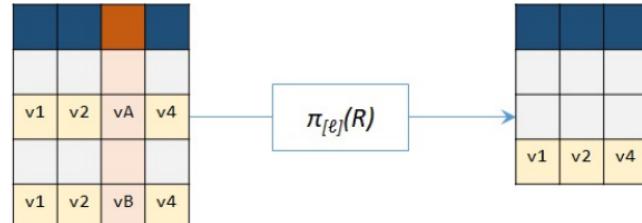
Recap

Unary

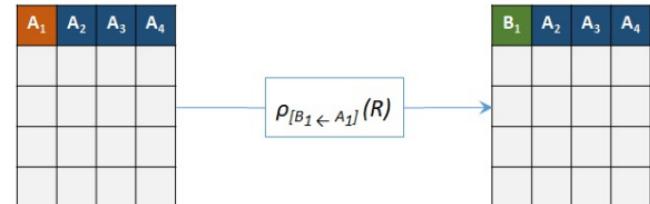
Selection



Projection



Renaming



Question 3

Recap
Schema
Question

Schema

$R(\underline{A}, C)$

R

A	C

$S(\underline{A}, D)$

S

A	D

$T(\underline{X}, Y)$

T

X	Y

Question 3

Recap
Schema

Question

- Q3(i)
- Q3(ii)
- Q3(iii)
- Q3(iv)

Question

Q3(i)

Question

Is $\pi_{[A]}(\sigma_{[A < 10]}(R)) \equiv \sigma_{[A < 10]}(\pi_{[A]}(R))$?

Can we use the algebraic law?

Question 3

Recap
Schema

Question

- Q3(i)
- Q3(ii)
- Q3(iii)
- Q3(iv)

Question

Q3(i)

Question

Is $\pi_{[A]}(\sigma_{[A < 10]}(R)) \equiv \sigma_{[A < 10]}(\pi_{[A]}(R))$?

Can we use the algebraic law?

Reference

$$\pi_{[\ell]}(\sigma_{[\theta]}(R)) \not\equiv \sigma_{[\theta]}(\pi_{[\ell]}(R)) \quad (\text{unless } \theta \text{ uses only attributes in } \ell)$$

Question 3

Recap
Schema

Question

- Q3(i)
- Q3(ii)
- Q3(iii)
- Q3(iv)

Question

Q3(i)

Question

Is $\pi_{[A]}(\sigma_{[A < 10]}(R)) \equiv \sigma_{[A < 10]}(\pi_{[A]}(R))$?

Visual Proof

A	C
$\sigma_{[A < 10]}$	$\sigma_{[A < 10]}$
$\sigma_{[A \geq 10]}$	$\sigma_{[A \geq 10]}$

Question 3

Recap
Schema

Question

- Q3(i)
- Q3(ii)
- Q3(iii)
- Q3(iv)

Question

Q3(i)

Question

Is $\pi_{[A]}(\sigma_{[A < 10]}(R)) \equiv \sigma_{[A < 10]}(\pi_{[A]}(R))$?

Visual Proof

A	C
$\sigma_{[A < 10]}$	$\sigma_{[A < 10]}$
$\sigma_{[A \geq 10]}$	$\sigma_{[A \geq 10]}$

A	C
$\sigma_{[A < 10]}$	$\sigma_{[A < 10]}$

A
$\sigma_{[A < 10]}$
$\sigma_{[A \geq 10]}$

Question 3

Recap
Schema

Question

- Q3(i)
- Q3(ii)
- Q3(iii)
- Q3(iv)

Question

Q3(i)

Question

Is $\pi_{[A]}(\sigma_{[A < 10]}(R)) \equiv \sigma_{[A < 10]}(\pi_{[A]}(R))$?

Visual Proof

A	C
$\sigma_{[A < 10]}$	$\sigma_{[A < 10]}$
$\sigma_{[A \geq 10]}$	$\sigma_{[A \geq 10]}$

A	C
$\sigma_{[A < 10]}$	$\sigma_{[A < 10]}$

A
$\sigma_{[A < 10]}$
$\sigma_{[A \geq 10]}$

A
$\sigma_{[A < 10]}$

A
$\sigma_{[A < 10]}$

Question 3

Recap
Schema

Question

- Q3(i)
- **Q3(ii)**
- Q3(iii)
- Q3(iv)

Question

Q3(ii)

Question

Is $\pi_{[A]}(\sigma_{[C < 10]}(R)) \equiv \sigma_{[C < 10]}(\pi_{[A]}(R))$?

Can we use the algebraic law?

Question 3

Recap
Schema

Question

- Q3(i)
- **Q3(ii)**
- Q3(iii)
- Q3(iv)

Question

Q3(ii)

Question

Is $\pi_{[A]}(\sigma_{[C < 10]}(R)) \equiv \sigma_{[C < 10]}(\pi_{[A]}(R))$?

Can we use the algebraic law?

Reference

$$\pi_{[\ell]}(\sigma_{[\theta]}(R)) \not\equiv \sigma_{[\theta]}(\pi_{[\ell]}(R)) \quad (\text{unless } \theta \text{ uses only attributes in } \ell)$$

Question 3

Recap
Schema

Question

- Q3(i)
- **Q3(ii)**
- Q3(iii)
- Q3(iv)

Question

Q3(ii)

Question

Is $\pi_{[A]}(\sigma_{[C < 10]}(R)) \equiv \sigma_{[\text{C} < 10]}(\pi_{[\text{A}]}(R))$?

Can we use the algebraic law?

Reference

$\pi_{[\ell]}(\sigma_{[\theta]}(R)) \not\equiv \sigma_{[\theta]}(\pi_{[\ell]}(R))$ (unless θ uses only attributes in ℓ)

Question 3

Recap
Schema

Question

- Q3(i)
- Q3(ii)
- **Q3(iii)**
- Q3(iv)

Question

Q3(iii)

Question

Is $\pi_{[D,Y]}(S \times T) \equiv \pi_{[D]}(S) \times \pi_{[Y]}(T)$?

Visual Proof

A	D	X	Y
Yellow	Yellow	Yellow	Yellow

Question 3

Recap
Schema

Question

- Q3(i)
- Q3(ii)
- **Q3(iii)**
- Q3(iv)

Question

Q3(iii)

Question

Is $\pi_{[D,Y]}(S \times T) \equiv \pi_{[D]}(S) \times \pi_{[Y]}(T) ?$

Visual Proof

A	D	X	Y
Yellow	Yellow	Grey	Grey
Grey	Grey	Yellow	Yellow

A	D	X	Y
Yellow	Yellow	Yellow	Yellow
Grey	Grey	Grey	Grey

D	Y
Yellow	Yellow
Grey	Grey

Question 3

Recap
Schema

Question

- Q3(i)
- Q3(ii)
- **Q3(iii)**
- Q3(iv)

Question

Q3(iii)

Question

Is $\pi_{[D,Y]}(S \times T) \equiv \pi_{[D]}(S) \times \pi_{[Y]}(T)$?

Visual Proof

D	Y
Yellow	Yellow
Grey	Grey
Yellow	Yellow
Grey	Grey

A	D
Yellow	Yellow
Grey	Grey
Yellow	Yellow

X	Y
Yellow	Yellow
Grey	Grey
Yellow	Yellow

D	Y
Yellow	Yellow
Grey	Grey
Yellow	Yellow
Grey	Grey

Question 3

Recap
Schema

Question

- Q3(iv)
- Q3(ii)
- Q3(iii)
- Q3(iv)

Question

Q3(iv)

Question

Is $\pi_{[D,Y]}(S \times T) \equiv \pi_{[D,Y]}(T \times S)$?

Try it out (3 minutes)

Question 4

Question 4

Question

- Q4(i)
- Q4(ii)
- Q4(iii)

Question

Q4(i)

Question

Find all pizzas that Moe likes but is not liked by Lisa.

Question 4

Question

- Q4(i)
- Q4(ii)
- Q4(iii)

Question

Q4(i)

Question

Find all pizzas that Moe likes but is not liked by Lisa.

Ask Yourself

- What are the required attributes?
- Which relation(s) hold the required attributes?
- What operation(s) do you need?
- What are the final attribute(s) needed?

Question 4

Question

- Q4(i)
- Q4(ii)
- Q4(iii)

Question

Q4(i)

Question

Find all pizzas that Moe likes but is not liked by Lisa.

Ask Yourself

- What are the required attributes?
cname and *pizza*
- Which relation(s) hold the required attributes?
Likes
- What operation(s) do you need?
Set difference
- What are the final attribute(s) needed?
pizza

Question 4

Question

- Q4(i)
- Q4(ii)
- Q4(iii)

Question

Q4(i)

Question

Find all pizzas that Moe likes but is not liked by Lisa.

Ask Yourself

- What are the required attributes?
cname and *pizza*
- Which relation(s) hold the required attributes?
Likes
- What operation(s) do you need?
Set difference
- What are the final attribute(s) needed?
pizza

Possible Answer

$$\pi_{[pizza]}(\sigma_{[cname='Moe']}(Likes)) - \pi_{[pizza]}(\sigma_{[cname='Lisa']}(Likes))$$

Question 4

Question

- Q4(i)
- Q4(ii)
- Q4(iii)

Question

Q4(i)

Question

Find all pizzas that Moe likes but is not liked by Lisa.

Alternative Step

Decomposition!

- Can you find all the pizzas that Moe likes?
- Can you find all the pizzas the Lisa likes?
- Can you subtract one from another?

Possible Answer

$$\pi_{[pizza]}(\sigma_{[cname='Moe']}(Likes)) - \pi_{[pizza]}(\sigma_{[cname='Lisa']}(Likes))$$

Question 4

Question

- Q4(i)
- **Q4(ii)**
- Q4(iii)

Question

Q4(ii)

Question

Find all customer-restaurant pairs (C, R) where C and R both located in the same area and C likes some pizza that is sold by R .

Ask Yourself

- What are the required attributes?
- Which relation(s) hold the required attributes?
- What operation(s) do you need?
- What are the final attribute(s) needed?

Question 4

Question

- Q4(i)
- **Q4(ii)**
- Q4(iii)

Question

Q4(ii)

Question

Find all customer-restaurant pairs (C, R) where C and R both located in the same area and C likes some pizza that is sold by R .

Ask Yourself

- What are the required attributes?
cname, rname, area and pizza
- Which relation(s) hold the required attributes?
Customers, Restaurants, Sells (is it enough?)
- What operation(s) do you need?
at least equi-join (can we do better?)
- What are the final attribute(s) needed?
cname and rname

Question 4

Question

- Q4(i)
- **Q4(ii)**
- Q4(iii)

Question

Q4(ii)

Question

Find all customer-restaurant pairs (C, R) where C and R both located in the same area and C likes some pizza that is sold by R .

Ask Yourself

- What are the required attributes?
cname, rname, area and pizza
- Which relation(s) hold the required attributes?
Customers, Restaurants, Sells, and Likes
- What operation(s) do you need?
natural join!
- What are the final attribute(s) needed?
cname and rname

Question 4

Question

- Q4(i)
- **Q4(ii)**
- Q4(iii)

Question

Q4(ii)

Question

Find all customer-restaurant pairs (C, R) where C and R both located in the same area and C likes some pizza that is sold by R .

Ask Yourself

- What are the required attributes?
cname, rname, area and pizza
- Which relation(s) hold the required attributes?
Customers, Restaurants, Sells, and Likes
- What operation(s) do you need?
natural join!
- What are the final attribute(s) needed?
cname and rname

Possible Answer

```
 $\pi_{[cname, rname]}($   
    Customers  $\bowtie$   
    Restaurants  $\bowtie$   
    Sells  $\bowtie$   
    Likes  
)
```

Can you find other answers?

Question 4

Question

- Q4(i)
- Q4(ii)
- **Q4(iii)**

Question

Q4(iii)

Question

Suppose the relation *Likes* contains all information about all customers. In other words, if the pair $(cname, pizza)$ is not in the relation *Likes*, it means that the customer *cname* dislikes the pizza *pizza*. Write a relational algebra expression to find for all customers, the pizza that they dislike. The result should be of the form $(cname, pizza)$.

Decomposition

Let's apply decomposition. What simpler operations can you do?

Question 4

Question

- Q4(i)
- Q4(ii)
- **Q4(iii)**

Question

Q4(iii)

Question

Suppose the relation *Likes* contains all information about all customers. In other words, if the pair $(cname, pizza)$ is not in the relation *Likes*, it means that the customer *cname* dislikes the pizza *pizza*. Write a relational algebra expression to find for all customers, the pizza that they dislike. The result should be of the form $(cname, pizza)$.

Decomposition

Let's apply decomposition. What simpler operations can you do?

- Can you find **ALL** possible *cname* and *pizza* pair?

Question 4

Question

- Q4(i)
- Q4(ii)
- **Q4(iii)**

Question

Q4(iii)

Question

Suppose the relation *Likes* contains all information about all customers. In other words, if the pair $(cname, pizza)$ is not in the relation *Likes*, it means that the customer *cname* dislikes the pizza *pizza*. Write a relational algebra expression to find for all customers, the pizza that they dislike. The result should be of the form $(cname, pizza)$.

Decomposition

Let's apply decomposition. What simpler operations can you do?

- Can you find **ALL** possible *cname* and *pizza* pair?
- Then we simply subtract the pizza that the customer likes.

Question 4

Question

- Q4(i)
- Q4(ii)
- **Q4(iii)**

Question

Q4(iii)

Question

Suppose the relation *Likes* contains all information about all customers. In other words, if the pair $(cname, pizza)$ is not in the relation *Likes*, it means that the customer *cname* dislikes the pizza *pizza*. Write a relational algebra expression to find for all customers, the pizza that they dislike. The result should be of the form $(cname, pizza)$.

Decomposition

Let's apply decomposition. What simpler operations can you do?

- Can you find **ALL** possible *cname* and *pizza* pair?
- Then we simply subtract the pizza that the customer likes.

Can you write this
solution down?

```
postgres=# exit
```

```
Press any key to continue . . .
```