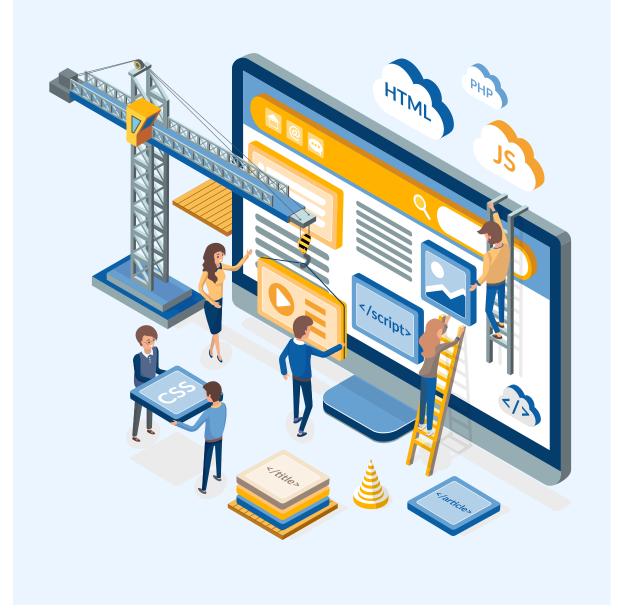


Task scheduling optimization problem

Multi-objective optimization





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Problem & Objective



Problem

- Task assignment problem in a computer
- Fog-cloud computing environment

Objective

 Assigning tasks to a machine located either in the fog or in the cloud while minimizing makespan and cost

2 Mathematical formulation



A Decision variable et constraints

Decision variables

• Xij, where i is the task number and j is the machine number (binary because if task i is assigned to machine j, 1; otherwise, 0).

Constraints

• Each task is associated with a single and unique VM. $\longrightarrow \Sigma_j X_{ij} = 1, \forall i$

• Each VM has a capacity constraint (Wi is the workload of task i and Cj is the processing capacity of VM j in $\Sigma_i W_i \cdot X_{ij} \leq C_j, \forall j \in \mathbb{N}$ 10s.)



Mathematical formulation



B Objective functions

 Minimize Makespan: minimize the total time taken to complete all tasks

 \longrightarrow Minimize $\Sigma_i \Sigma_j T_{ij} \cdot X_{ij}$

Minimize cost: minimize the total cost of VMs

$$\longrightarrow$$
 Minimize $\Sigma_i \Sigma_j (c_{ijp} + c_{ijm} + c_{ijb}) * X_{ij}$



Metaheuristic function: Bat Algorithm



Population Initialization



 Generating an initial population of bats, where each bats represents a potential solution according to its position

Parameters Initialization

 Defining parameters such as the number of bats, problem dimensions, and algorithmic parameters:

```
# Parameters

dim = num_tasks # Dimensionality of the tasks

epochs = 100 # Number of generations

pop_size = 100 # Population size

gam = 0.1 # Pulse rate increasing rate

qmin = 0 # Minimum frequency

qmax = 1 # Maximum frequency
```



Archive part

Metaheuristic function: Bat Algorithm



Search and Solution Update

```
solution = np.zeros((num_vms, num_tasks), dtype=int)
assigned_workloads = np.zeros(num_vms)

# Sort the tasks in ascending order based on the position of the bat
sorted_task_indices = np.argsort(position)
for task_index in sorted_task_indices:
    task_workload = workload[task_index]
    vm_candidates = np.where(np.array(capacities) >= task_workload)[0]
    for vm in vm_candidates:
        # Capacity constraint
        if assigned_workloads[vm] + task_workload <= capacities[vm]:
            solution[vm, task_index] = 1
            assigned_workloads[vm] += task_workload
            break # Assign task to the first available VM to make sure only one VM has this task
return solution</pre>
```

Function that returns a solution (2D matrix containing the X ij) respecting the constraints

def bat task assignment(position, workload, capacities):

Initialization

num_tasks = len(workload)
num vms = len(capacities)

```
for bat in population:
   # Get and evaluate a solution depending on position of the bat
   solution = bat task assignment(bat.pos, workload, capacity)
   bat fitness = [makespan(solution), total cost(solution)]
   bat.set fitness(bat fitness)
   bat.set sol(solution)
   if bat not in archive:
       is dominated = False
       for archived bat in archive:
           if bat fitness[0] < archived bat.fitness[0] and bat fitness[1] < archived bat.fitness[1]:
               # If the new bat dominates an archived bat, remove the archived bat
               archive.remove(archived bat)
           elif bat fitness[0] >= archived bat.fitness[0] and bat fitness[1] >= archived bat.fitness[1]:
               # If an archived bat dominates the new bat, it is dominated
               is dominated = True
       if not is dominated:
           # If the new bat is not dominated, add it to the archive
           archive.append(bat)
           # Selecting the last best bat as a reference
           best bat = bat
   else:
     # Checking the dominance in the archive
     for archived bat in archive:
       if bat.fitness[0]>=archived bat.fitness[0] and bat.fitness[1]>=archived bat.fitness[1] and bat.fitness!=archived bat.fitness:
         archive.remove(bat)
   makespan cost epoch.append(bat fitness)
makespan costs.append(makespan cost epoch)
```



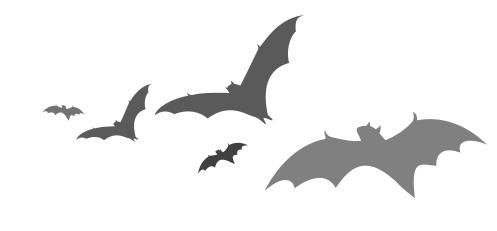
Metaheuristic function: Bat Algorithm

Search and Solution Update

- At each iteration, each bat conducts a local search around its current position
- The local search is guided by the quality of the best solution, and each bat updates its position accordingly
- We maintain an archive of the best solutions found, i.e. the Pareto optimal solutions

```
# Moving part
for bat in population:
 if bat not in archive:
   new pos = [0.00 for in range(dim)]
   new_vel = [0.00 for _ in range(dim)]
   # Frequence of the bat
   freq = random.uniform(qmin, qmax)
   # Random pulsation
   pulse chance = random.uniform(0, 1)
   for d in range(dim):
       # Updating velocity
       new_vel[d] = bat.vel[d] + (bat.pos[d] - best bat.pos[d]) * freq
       # If pulsation rate not loud enough
       if pulse_chance > bat.pulse_rate:
           # Stays close to the best bat
           new_pos[d] = best_bat.pos[d] + (random.uniform(-1, 1) * 0.1)
            # Moving toward the best bat and keeps exploring
           new_pos[d] = bat.pos[d] + new_vel[d]
   # Update
   bat.vel = new vel
   bat.pos = new pos
   bat.pulse rate = bat.max pulse rate * (1 - exp(-gam * e))
```

Metaheuristic function: Bat Algorithm



Termination Criterion:

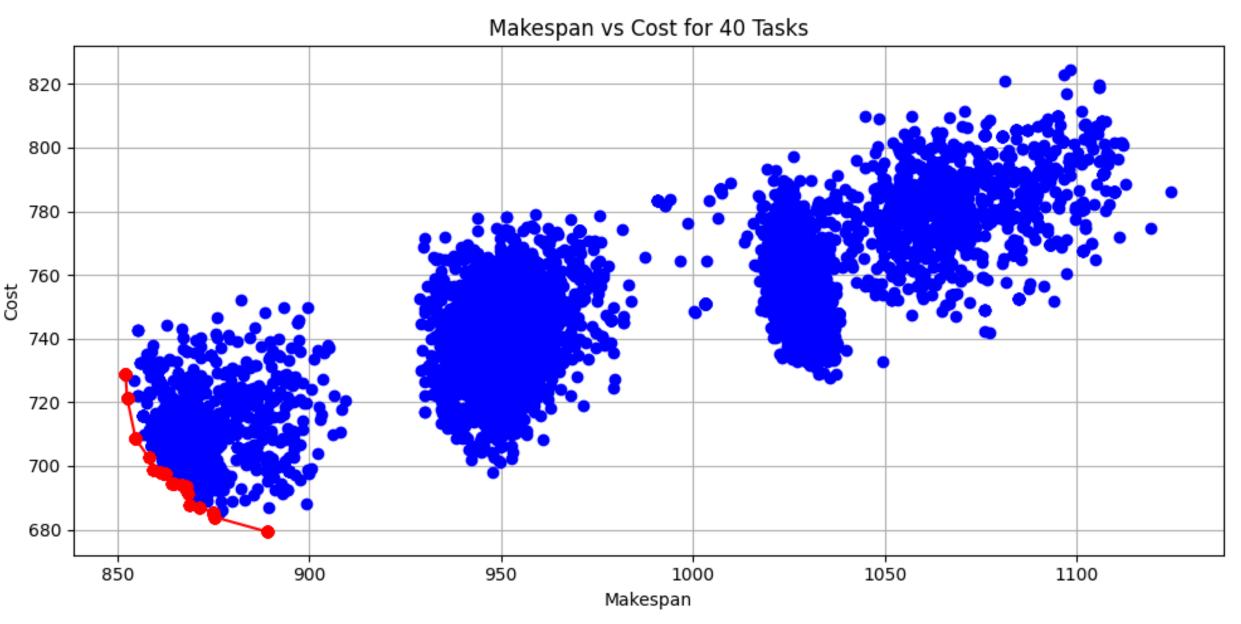
• The algorithm continues iterating until a predefined termination criterion is met, here it is the maximum number of iterations (epochs)



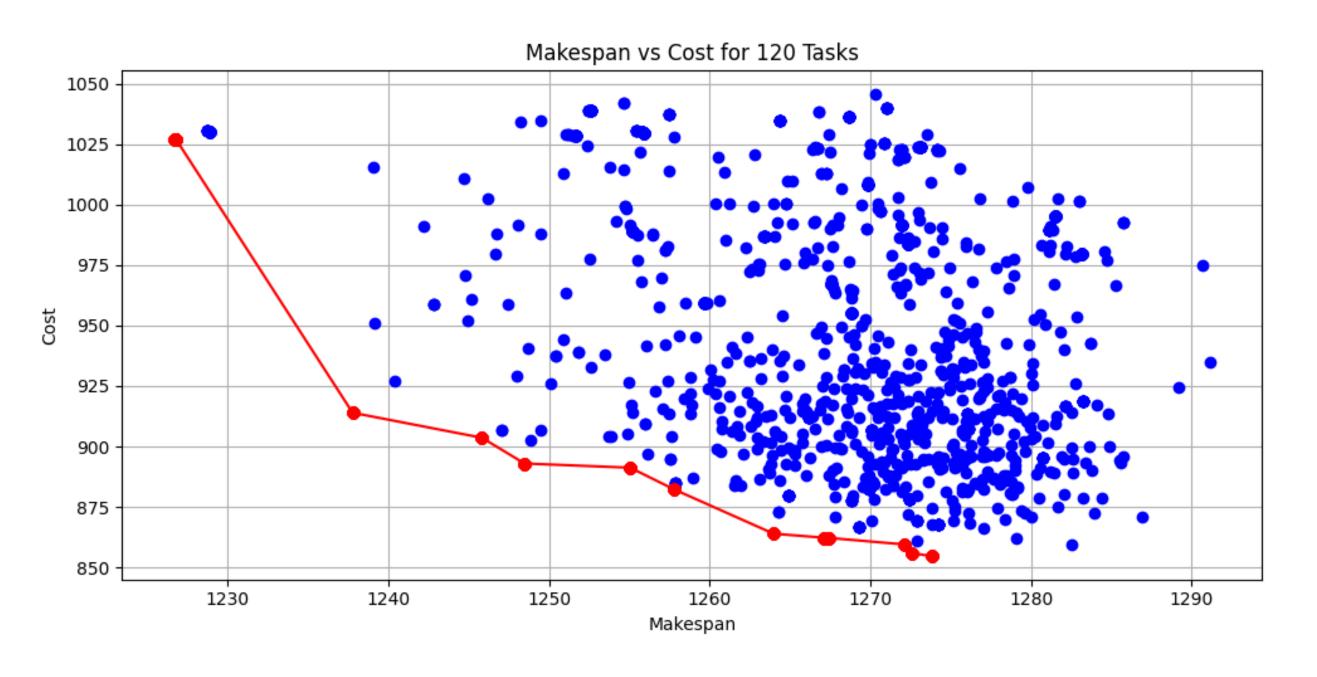


- O(enm): e the number of epochs, n the number of tasks and m the number of bats.
- e and m fixed leads to an execution time linear in n.
- O(ln(n)) possible?
- 30 seconds: solve a task scheduling problem with 40 tasks in 100 epochs with a population of 100 bats.











- Comparison with the solution of the group of Justine and David
- Task scheduling problem with a genetic algorithm but with different constraints
- Memory capacity constraint rather than a CPU capacity and a different calculation for the makespan
- Not completely comparable but same execution time
- Both performant and solutions graphically sufficient

5 Conclusion



- A good adaptation of the BAT algorithm to make it multi-objective in our problem
- Convincing results

