The Art of Translation

- Person(p), which states that p is a person, and
- Loves(x, y), which states that x loves y,

write a sentence in first-order logic that means "everybody loves someone else."

```
\forall p. (Person(p) \rightarrow \exists q. (Person(q) \land p \neq q \land Loves(p, q))
```

How many of the following first-order logic statements are correct translations of "everyone loves someone else?"

```
∀p. (Person(p) →
∃q. (Person(q) ∧
Loves(p, q)
)
```

```
\forall p. (Person(p) \land \exists q. (Person(q) \land p \neq q \land Loves(p, q)
```

```
\forall p. (Person(p) \rightarrow \exists q. (Person(q) \land p \neq q \rightarrow Loves(p, q)
```

```
\exists p. (Person(p) \rightarrow \forall q. (Person(q) \land p \neq q \land Loves(p, q)
```

Answer at **PollEv.com/cs103** or text **CS103** to **22333** once to join, then **0**, **1**, **2**, **3**, or **4**.

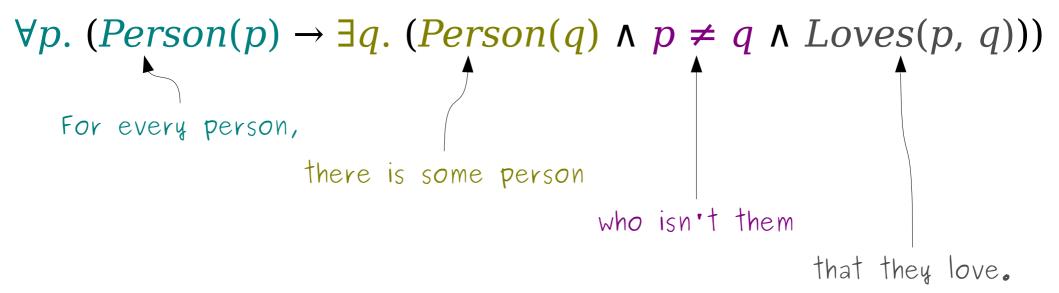
- Person(p), which states that p is a person, and
- Loves(x, y), which states that x loves y,

write a sentence in first-order logic that means "there is a person that everyone else loves."

```
\exists p. (Person(p) \land \forall q. (Person(q) \land p \neq q \rightarrow Loves(q, p))
```

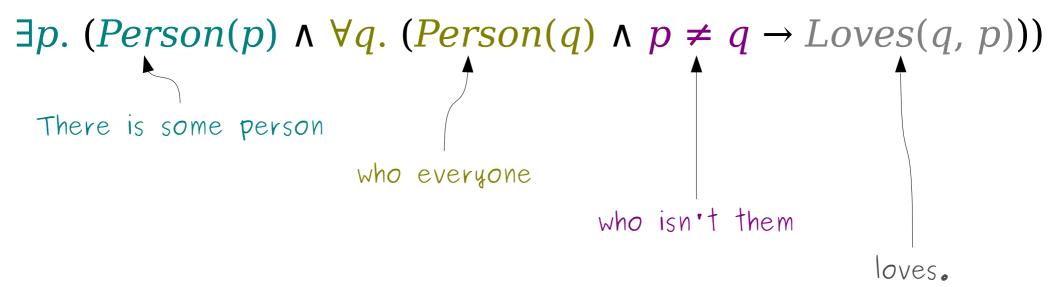
Combining Quantifiers

- Most interesting statements in first-order logic require a combination of quantifiers.
- Example: "Everyone loves someone else."

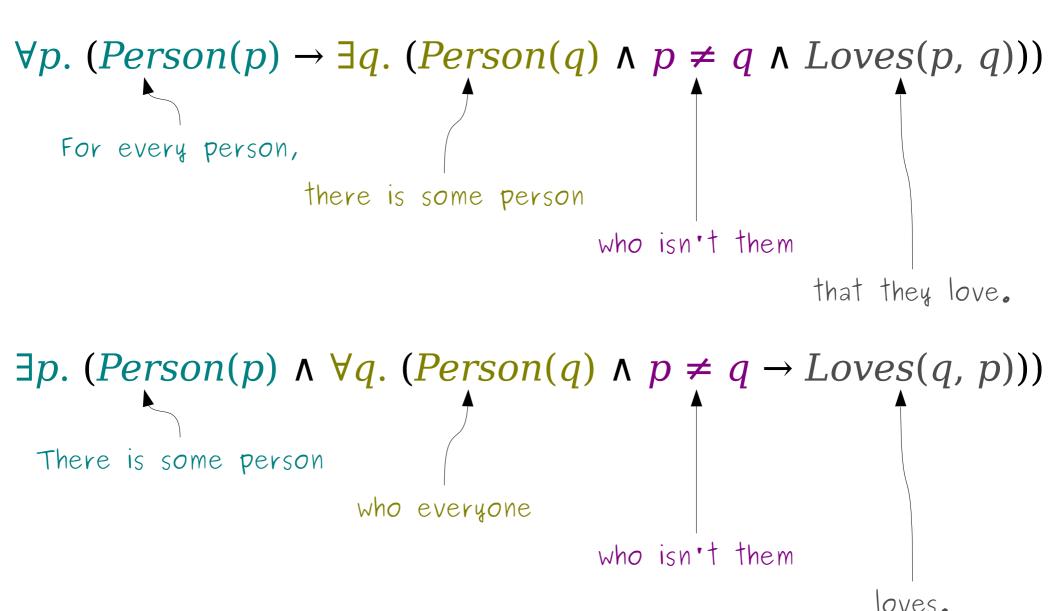


Combining Quantifiers

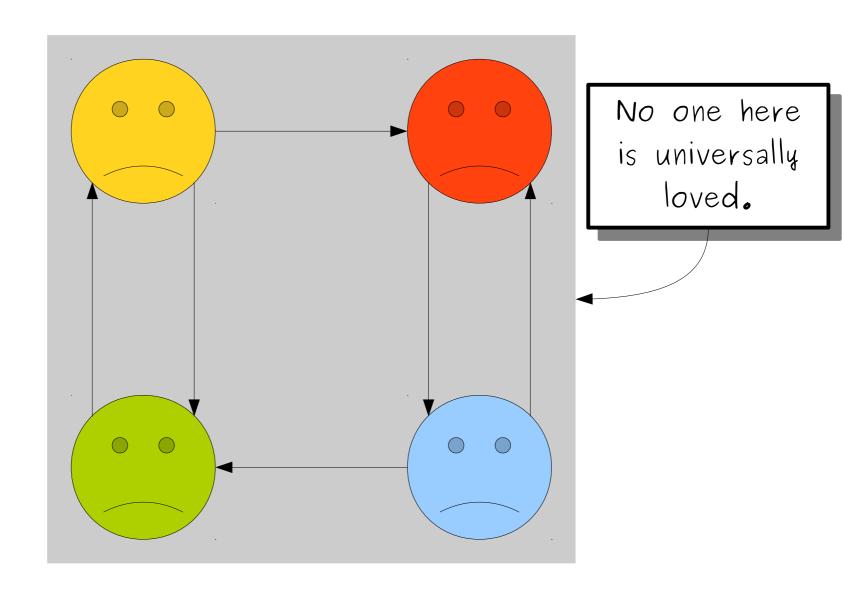
- Most interesting statements in first-order logic require a combination of quantifiers.
- Example: "There is someone everyone else loves."



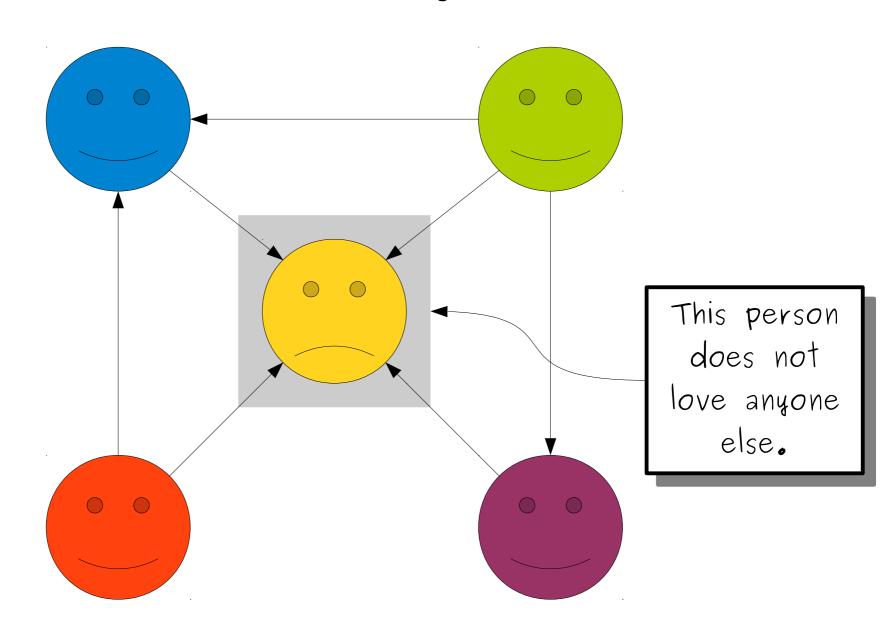
For Comparison



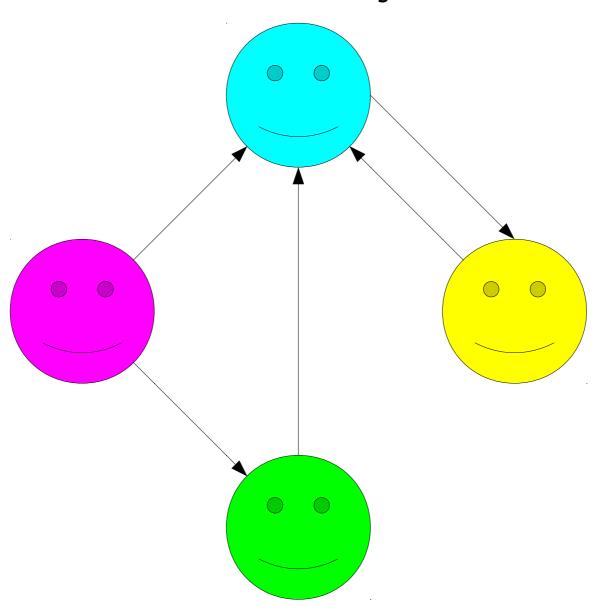
Everyone Loves Someone Else

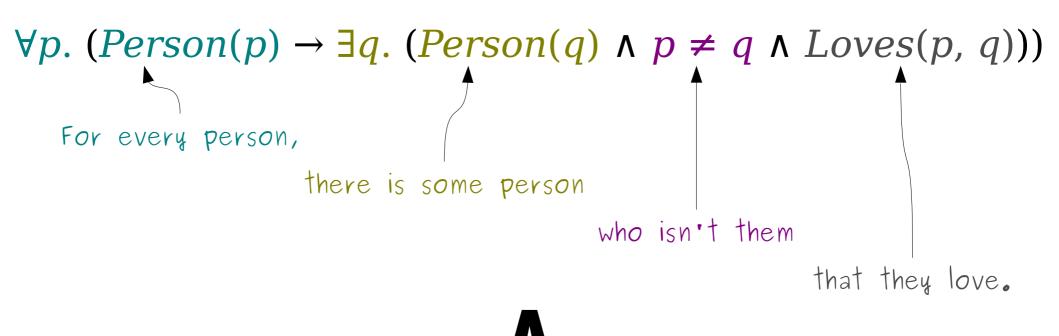


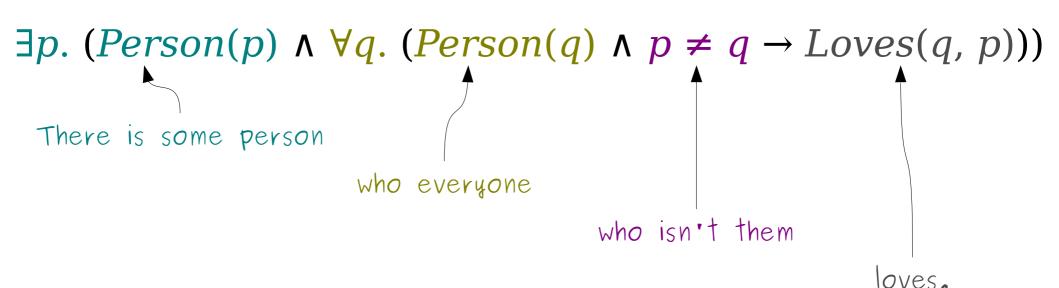
There is Someone Everyone Else Loves



Everyone Loves Someone Else *and*There is Someone Everyone Else Loves







Quantifier Ordering

The statement

$$\forall x. \exists y. P(x, y)$$

means "for any choice of x, there's some choice of y where P(x, y) is true."

• The choice of *y* can be different every time and can depend on *x*.

Quantifier Ordering

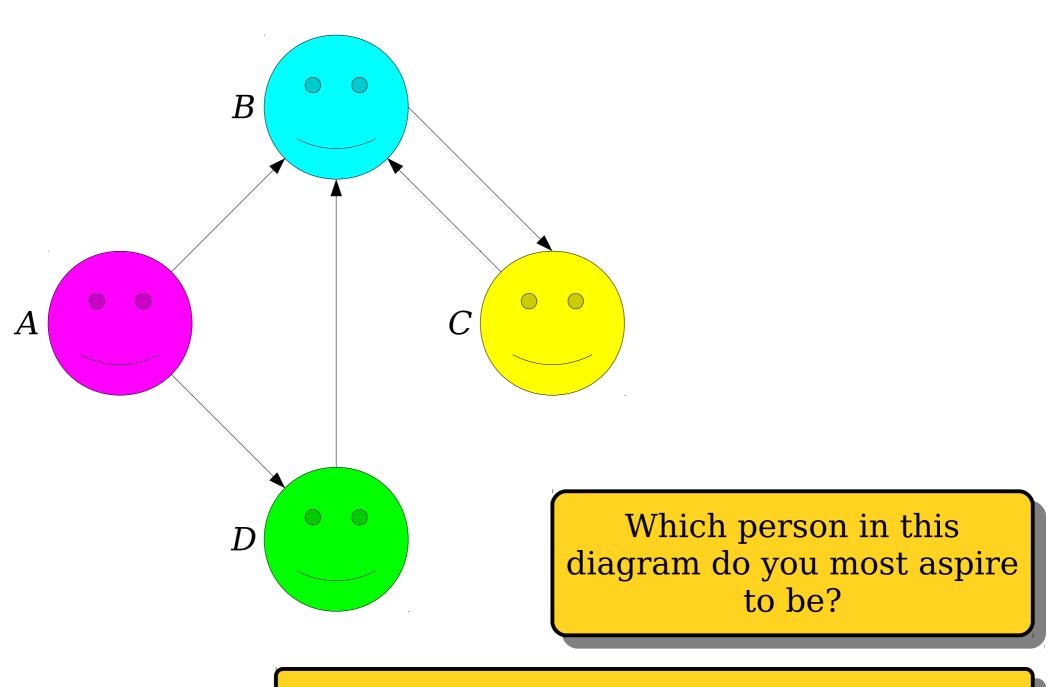
The statement

$$\exists x. \ \forall y. \ P(x, y)$$

means "there is some x where for any choice of y, we get that P(x, y) is true."

• Since the inner part has to work for any choice of *y*, this places a lot of constraints on what *x* can be.

Order matters when mixing existential and universal quantifiers!



Answer at **PollEv.com/cs103** or text **CS103** to **22333** once to join, then **A**, **B**, **C**, or **D**.

Set Translations

- Set(S), which states that S is a set, and
- $-x \in y$, which states that x is an element of y,

write a sentence in first-order logic that means "the empty set exists."

First-order logic doesn't have set operators or symbols "built in." If we only have the predicates given above, how might we describe this?

 $\exists S. (Set(S) \land \neg \exists x. x \in S)$

 $\exists S. (Set(S) \land \forall x. x \notin S)$

Both of these translations are correct.

Just like in propositional logic, there are many different equivalent ways of expressing the same statement in first—order logic.

- Set(S), which states that S is a set, and
- $-x \in y$, which states that x is an element of y,

write a sentence in first-order logic that means "two sets are equal if and only if they contain the same elements."

```
\forall S. (Set(S) \rightarrow \forall T. (Set(T) \rightarrow (S = T \leftrightarrow \forall x. (x \in S \leftrightarrow x \in T)))
```

```
\forall S. (Set(S) \rightarrow \forall T. (Set(T) \rightarrow (S = T \leftrightarrow \forall x. (x \in S \leftrightarrow x \in T)))
```

You sometimes see the universal quantifier pair with the ↔ connective. This is especially common when talking about sets because two sets are equal when they have precisely the same elements.