

## JUMP TO SOLUTION

Answer the following questions about the algorithm below, which attempts to find the median element of an odd-length array.

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Input: data: an array with an odd number of integers
Input: n: the length of data (odd)
Output: element m in data such that  $(n - 1)/2$  elements are smaller
        and  $(n - 1)/2$  elements are larger
1 Algorithm: BadMedian
2 med = data[1]
3 lo = hi = 0
4 for i = 1 to n - 1 do
5     if data[i + 1] < med then
6         if lo < i/2 then
7             lo = lo + 1
8         else
9             med = data[i + 1]
10            hi = hi + 1
11        end
12    else
13        if hi < i/2 then
14            hi = hi + 1
15        else
16            med = data[i + 1]
17            lo = lo + 1
18        end
19    end
20 end
21 return med
```

**Example.** If  $data = [2, 1, 4, 5, 3]$ ,  $med = 2$ ,  $lo = 0$ , and  $hi = 0$  before the first iteration. In the first iteration of the for loop,  $data[2] = 1 < 2$  and  $lo = 0 < 1/2$ , so  $lo$  becomes 1. In the second,  $data[3] = 4 > 2$  and  $hi = 0 < 2/2$ , so  $hi$  becomes 1. In the third,  $data[4] = 5 > 2$  and  $hi = 1 < 3/2$ , so  $hi = 2$ .

In the fourth and final iteration,  $data[5] = 3 > 2$  but  $hi = 2 \geq 4/2$ , so  $med$  becomes 3 and  $lo$  becomes 2. BadMedian returns 3, which is the median of *data*.

1. Prove that  $lo + hi = i$  after every iteration of the for loop.
2. Prove that BadMedian is incorrect.

# SOLUTION

## 1. Invariant

For each iteration two possible cases, each with two possible cases of their own:

Case 1:  $\text{data}[i+1]$  is less than the value of med

Case A: value of lo is less than  $i/2$  and lo increases by 1

Case B: otherwise med =  $\text{data}[i+1]$  and hi increases by 1

Case 2: otherwise

Case C: value of hi is less than  $i/2$  and hi increases by 1

Case D: otherwise med =  $\text{data}[i+1]$  and lo increases by 1

1st iteration: at start  $i = 1$  and lo and hi both = 0.  $\text{Data}[2]$  is either  $<$  med or not. If it is less than med lo will increase by 1 since lo's current value is 0 and 0 is  $< 1/2$ . If it is not less than med hi will increase by 1 since hi's current value is 0 and 0 is  $< 1/2$ . Therefore EITHER hi or lo will increase by 1 and as such  $\text{lo} + \text{hi} = i$  when  $i = 1$ .

$k+1$ st iteration: suppose that for  $\text{data}[k]$  where  $k=i$  that  $\text{lo} + \text{hi} = i$ . During iteration  $k+1$   $i$  is increased by 1. The value of  $\text{data}[i+1]$  will either be  $<$  med or not. If it is less than med the value of lo is either  $< i/2$  and lo increases by 1 otherwise hi increases by 1. If it is not less than med the value of hi is either less than  $i/2$  and hi increases by 1 otherwise lo increases by 1.

Therefore it must be true that each time  $i$  increases by 1 EITHER hi or lo will increase by 1 as well, hence the loop invariant holds that  $\text{lo} + \text{hi} = i$  for every iteration of the loop.

## 2. Counter Example

Consider data [0,0,1]

iteration 1:  $\text{data}[2] = 0$  -- case 2: 0 not less than med(0) -- case C:  $\text{hi}(0) < 1/2$   
hi increases by 1

iteration 2:  $\text{data}[3] = 1$  -- case 2: 1 not less than med(0) -- case D:  $\text{hi}(1) \text{ not } < 2/2$   
lo increases by 1 and med =  $\text{data}[i+1]$  making med = 1

Med return value is 1. Desired output should have  $(n-1)/2$  smaller elements and  $(n-1)/2$  larger elements, but output has 2 smaller elements and 0 larger elements.