

Hw 1) Given 32 bit: 1010 1101 0001 0000 0000 0000 0000 0010  
 a) 2's comp signed int = -1391460350  
 b.) MIPS 

101011	01000	10000	0000 0000 0000 0010
op	rs	rt	imm

  

0x2b
8=\$t0
16=\$s0
2

sw \$s0, 2(\$t0)  
 # \$s0 → MEMORY[\$t0+2]

2. (4 pts) Determine the absolute value of a signed integer. Show the implementation of the following pseudo-instruction using three real instructions:  
 abs \$t1, \$t2

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addu $t1, $t2, $zero #interpretation of move pseudo instruction
bgez $t1, next       #if t1 >= 0 branch to next
subu $t1, $zero, $t1 #else if neg int take 0-(-int)= +int
next:...
```

3.)

```
#move $t1, $t2
addu $t1, $t2, $zero      #$t1 = $t2 + 0 therefore $t1 = $t2

#clear $t5
and $t5, $t5, $zero      #$t5 = $t5 AND 0 therefore $t5 = 0

#li $t5, imm32
lui $s1, (upper 16 bits)  #stores upper 16 bits
ori $s1, $s1, (lower 16 bits) #inserts the lower 16 bits

#addi $t5, $t3, imm32
lui $t5, (upper 16 bits)  #similar function as above
ori $t5, $t5, (lower 16 bits) #similar function as above
addi $t5, $t5, $t3        #$t5 = $t5 + $t3

#beq $t5, imm32, Label
lui $at, (upper 16 bits)  #stores upper 16 bits in $at
ori $at, $at, (lower 16 bits) #inserts lower 16 bits in $at
beq $t5, $at, Label       #if $t5 == $at branch to Label

#ble $t5, $t3, Label      #branch if $t5 <= $t3
slt $at, $t3, $t5         #if $t3 < $t5 $at=1
beq $at, $zero, Label     #if $at == 0 branch to Label
#if $t3 < $t5 then at=1 and $at != 0 ...no branch
#if $t5 <= $t3 then at=0 and $at == 0...branches

#bgt $t5, $t3, Label      #branch if $t5 > $t3
slt $at, $t3, $t5         #if $t3 < $t5 $at=1
bne $at, $zero, Label     #if $at != 0 branch to Label
#if $t3 < $t5 then $at=1 and $at != 0 ...branches

#bge $t5, $t3, Label      #branch if $t5 >= $t3
slt $at, $t5, $t3         #if $t5 < $t3 $at=1
beq $at, $zero, Label     #if $at == 0 branch to Label
#works similar to ble but changes places for rs and rt
```

4. (8 pts) Translate the following statements into MIPS assembly language. Assume that a, b, c, and d are allocated in \$s0, \$s1, \$s2, and \$s3. All values are signed 32-bit integers.

```
a) if ((a > b) || (b > c)) {d = 1;}
b) if ((a <= b) && (b > c)) {d = 1;}
```

```
#implementation for a: if (($s0 > $s1) || ($s1 > $s2)) {$s3 = 1;}
#pseudocode implementation
    bgt $s0, $s1, L1                #if yes jump to 'if' part (L1)
    ble $s1, $s2, next             #if no skip 'if' part
L1:  li $s3, 1                     #set $s3 to 1
next:...                           #branch here to skip 'if' inst
```

```
#implementation with no pseudocode
    slt $at, $s1, $s0      #if $s1 < $s0 $at=1
    bne $at, $zero, L1     #if $at != 0 branch to 'if' part (L1)
    slt $at, $s2, $s1      #if $s2 < $s1 $at=1
    beq, $at, $zero, next  #if $at == 0 skip 'if', branch to
next
L1:  ori $s3, $zero, 1     #OR inst used to set bits
next:...                  #branch here to skip 'if' inst
```

```
#implementation for b: if (($s0 <= $s1)&&($s1 > $s2)) {$s3=1;}
#pseudocode implementation
    bgt $s0, $s1, next          #if $s0>$s1 skip 'if'
    ble $s1, $s2, next          #if $s1<=$s2 skip 'if'
    li $s3, 1                   #set $s3 to 1
next:...                        #branch here to skip 'if' inst
```

```
#implementation with no pseudocode
    slt $at, $s1, $s0          #if $s1 < $s0 $at=1
    bne $at, $zero, next       #if $at != 0 skip 'if', branch to
next
    slt $at, $s2, $s1          #if $s2 < $s1 $at=1
    beq $at, $zero, next       #if $at == 0 skip 'if', branch to
next
L1: ori $s3, $zero, 1          #OR inst used to set bits
next:...                       #branch here to skip 'if' inst
```

5. (8 pts) Consider the following fragment of C code:

```
for (i=0; i<=100; i=i+1) { a[i] = b[i] + c; }
```

Assume that *a* and *b* are arrays of words and the base address of *a* is in *\$a0* and the base address of *b* is in *\$a1*. Register *\$t0* is associated with variable *i* and register *\$s0* with *c*. Write the code in MIPS.

and \$t0, \$t0, \$zero	#clears \$t0 to use as iterator
addi \$t1, \$zero, 404	#sets \$t1 to 404(max iterator val)
loop: add \$t2, \$a1, \$t0	#\$t2 = B base addr + iterator
lw \$t2, 0(\$t2)	#\$t2 = B[i]
addu \$t2, \$t2, \$s0	#\$t2 = B[i] + C
add \$t3, \$a0, \$t0	#\$t3 = A base addr + iterator
sw \$t2, 0(\$t3)	#store sum of B[i] + C in A[i]
addi \$t0, \$t0, 4	#\$t0 = \$t0 + 4 (i++)
bne \$t0, \$t1, loop	#if iterator != 404 branch to loop

6. (8 pts) Add comments to the following MIPS code and describe in one sentence what it computes. Assume that *\$a0* is used for the input and initially contains *n*, a positive integer. Assume that *\$v0* is used for the output.

begin:	addi \$t0, \$zero, 0	#clears space for storage
	addi \$t1, \$zero, 1	#sets iterator to 1
loop:	slt \$t2, \$a0, \$t1	#is n < iterator?
	bne \$t2, \$zero, finish	#if n < i branch to finish
	add \$t0, \$t0, \$t1	#else store value of i in \$t0
	addi \$t1, \$t1, 2	#iterator = iterator + 2
	j loop	#repeat loop
finish:	add \$v0, \$t0, \$zero	#output = stored value of \$t0

This code starts with 1 and sums all the odd ints less than or equal to input *\$a0*.

7. (12 pts) The following code fragment processes an array and produces two important values in registers \$v0 and \$v1. Assume that the array consists of 5000 words indexed 0 through 4999, and its base address is stored in \$a0 and its size (5000) in \$a1. Describe what this code does. Specifically, what will be returned in \$v0 and \$v1?

		add \$a1, \$a1, \$a1	#a1 = 5,000 + 5,000 = 10,000
		add \$a1, \$a1, \$a1	#a1 = 10,000 + 10,000 = 20,000
		add \$v0, \$zero, \$zero	#clears \$v0 to use to sum values
		add \$t0, \$zero, \$zero	#clears \$t0 to use as iterator
<div style="display: inline-block; vertical-align: middle;"> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;"> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">1</div> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">outer:</div> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">{</div> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">inner:</div> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">{</div> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">skip:</div> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">}</div> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">}</div> <div style="display: inline-block; vertical-align: middle; margin-right: 10px;">}</div> </div> </div>		add \$t4, \$a0, \$t0	#t4 = A base address + iterator
		lw \$t4, 0(\$t4)	#t4 = A[i]
		add \$t5, \$zero, \$zero	#clears \$t5 to use as a counter
		add \$t1, \$zero, \$zero	#clears \$t1 to use as <u>iterator2</u>
		add \$t3, \$a0, \$t1	#t3 = A base address + <u>iterator2</u>
		lw \$t3, 0(\$t3)	#t3 = A[j]
		bne \$t3, \$t4, skip	#if A[j] != A[i] branch to skip
		addi \$t5, \$t5, 1	#else counter++
		addi \$t1, \$t1, 4	#point to next A[j]
		bne \$t1, \$a1, inner	#if <u>iterator2</u> != 20,000 loop to inner
		slt \$t2, \$t5, \$v0	#is counter < \$v0?
		bne \$t2, \$zero, next	#if \$v0 is greater jump to next
		add \$v0, \$t5, \$zero	#else \$v0 = counter value
		add \$v1, \$t4, \$zero	# \$v1 = A[i]
		addi \$t0, \$t0, 4	#point to next A[i]
	bne \$t0, \$a1, outer	#if iterator != 20,000 branch to outer	

- (1) Goes from A[0] to A[4999] using first iterator...A[i]
- (2) Goes from A[0] to A[4999] using a second iterator...A[j]
- (3) If elem in A[i] matches elem in A[j] increase counter, else j++ and check next element
- (4) If counter value for current check is greater than stored value: store the counter value and store the value for A[i], else i++ and check next element

\$v0 will return the highest counter value and \$v1 will return the index of the elem that generated the highest counter value. So the code is finding which element appears the most frequently and how many times it appears.