

## Pattern Matching

① i/p  $\rightarrow$  Geeks for Geeks , eKs  
Big string pattern (small string).

O/P  $\rightarrow$  2, 10

③ ABC ABCD  
ABD.

④ i/p  $\rightarrow$  A A A A A , AAA

O/P  $\rightarrow$  -1

O/P  $\rightarrow$  0, 1, 2

↳ it's like finding the pattern

in the given string.

# Overview of pattern Searching :

$m \rightarrow$  pattern length

$n \rightarrow$  string length.

in worst case  
hash fun. matches  
all times,  
it will perform  
poorer than  
naive

Naive

if all chars  
are distinct.  $\rightsquigarrow$

$$O((n-m+1) * m)$$

b. Take the pattern Compare every  
window of text one by one if matches  
print it.

$$O(n)$$

No preprocess

uses rolling hash to optimize  
the naive also

Rabin Karp

$$O((n-m+1) * m)$$

$\hookrightarrow$  it's like computing hash func  
then checking it

Kmp

$$O(n)$$

$\hookrightarrow$  preprocess pattern

Suffix tree

$$O(m)$$

Preprocess Text / String

, data structure

→ Naive.  
↳ we need to find / return all indexes where we find the pattern.

①  $\text{txt} = \text{ABABA}BCD \Rightarrow 0, 2$

$\text{ptrn} = \text{ABAB}$

②  $\text{AAAAA} \Rightarrow 0, 1, 2$   
 $\text{AAA}$

$\begin{array}{c} \text{A A A A A} \\ \hline \text{---} \\ \text{AAA} \end{array} \Rightarrow 5$   
 $\Rightarrow 3$

↳ need to check for inclusive.  
3 times  
ie.  $i \rightarrow [0, n-m]$

① check for next  $\frac{3}{m}$  chars

$\text{pattern}[j] = \text{string}[i+j]$

$\hookrightarrow j=0$

$\begin{array}{c} \text{A A A A A} \\ \hline \text{A A} \end{array}$

'if  $j = m$ :  
print(i)

$j=1$

$\begin{array}{c} \text{A A A A A} \\ \hline \text{A A A} \end{array}$

$\underline{\text{index}}$ .

$j=2$

$\begin{array}{c} \text{A A A A A} \\ \hline \text{A A A} \end{array}$

↳ if found 'j' as we found match.

for  $i \rightarrow 0 \rightarrow n-m+1$ :  
 $j=0$   
for  $j \rightarrow n \rightarrow 0 \rightarrow m$ :

$(n-m+1) * m$

$\lceil \frac{n}{m} \rceil = m$ :

# Improved Naive Algo For Distinct .

↳ in pattern .

0 1 2 3 4 5 6 7 8 9 10 11

ABC E ABE F ABCD

↑ ↑ ↑ ↑ ↑ ↑

ABc D

$n = 12$

$m = 4$

$i = 0$

while  $i \leq n - m + 1 :$

$i = 0 : i + = 3$

$i = 3 :$

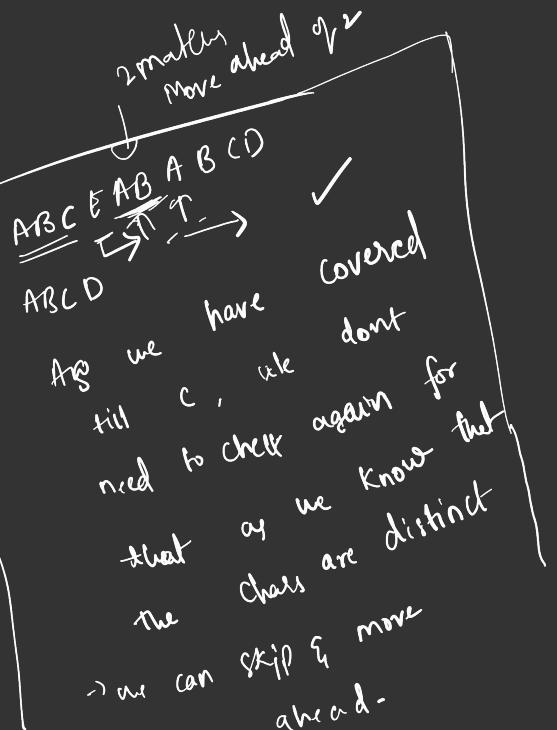
$i = 4 : i + = 2$

$i = 6 :$

$i = 7 :$

$i = 8 :$

Made with Goodnotes and vice versa



$J = 0$

while  $j \neq m :$

if  $\text{str}[i+j] \neq \text{pattern}[j] :$

break ;  
 $J++$

if  $j = m :$   
print (i)

↳ pattern found .

if  $J == 0 :$   
 $i + = 1$

else :

$i = i + j$  ↳ if we matched something

Rabin Karp Algo. :- rolling hash func.

- like naive algo, slides the pattern one by one in str / text
- Compare hash values of pattern & current text window. If hash values match  
then only compare individual chars.

a : 1

b : 2

c : 3

d : 4

e : 5

$P$  = Hash val of pattern

$t$  = Hash val of current window of text

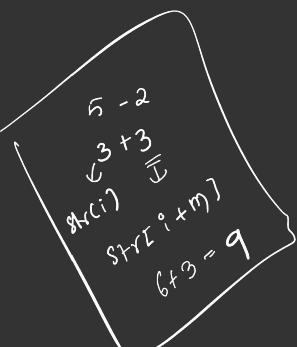
→ Simple Hash : sum of values

→ problem : Spurious hits.

god dog.  $\Rightarrow$  both give same value.

txt = 0 1 2 3 4 5 6 7 8 9  
0 1 2 3 4 5 6 7 8 9  
pat = abc

$$9 - 2 = 7$$



$P = 1+2+3 = 6$  = Computation of hash val of pattern

$i=0$	$t = 1+2+4 = 7$
$i=1$	$t = 2+4+1 = 7$
$i=2$	$t = 4+1+2 = 7$
$i=3$	$t = 1+2+3 = 6$ (match)
$i=4$	$t = 2+3+2 = 7$
$i=5$	$t = 3+2+1 = 6$ (spurious hit)
$i=6$	$t = 2+1+2 = 5$
$i=7$	$t = 1+2+3 = 6$ (match)

We can go with ascii vals also, we will get the spurious hits in that case also.



0 1 2 3 4 5 6 7

8 - 2 = 6

ABAB ABCD

AB

P = 3

t = 3

curr -  
ord(str(i))  
+  
next hash  $\Rightarrow$  ord(str(item))  
2

i = 0

i = 1

i = 2

i = 3

i = 4

i = 5

i = 6

t = 3  $\Rightarrow$  ✓

t = 3 - 1 + 1 = 3  $\Rightarrow$  ✗

t = 3 - 2 + 2 = 3  $\Rightarrow$  ✓

t = 3 - 2 + 2  $\Rightarrow$  ✗

3 - 2 + 2  $\Rightarrow$  ✓

3 - 1 + 3  $\Rightarrow$  5

5 - 2 + 4  $\Rightarrow$  7

$$t_{i+1} = t_i + \text{str}[i+m] - \text{str}[i].$$

Rolling Hash.

improved hash.

$$d = 5$$

we compute powers of  $d$  (weighted sum):

$\underline{\underline{=}} * \underline{\underline{=}}$

a : 1

b : 2

c : 3

d : 4

e : 5

$$h(abc) = 1 \times d^2 + 2 \times d + 3 \times d^0 = 25 + 10 + 3 = 38$$

$$h(dab) = 4 \times d^2 + 1 \times d + 2 \times d^0 = 107$$

Rolling hash func:

$$\frac{h(t_{i+1})}{t_{i+1}} = d(t_i - t \times t[i] \times d^{m-1}) + txt[i+m]$$

ex:

$$txt = \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ a & b & d & a & b & c & b & a & b & c \end{matrix}$$

$$\begin{aligned} pat &= abc \\ &\in 1 \times 5^2 + 2 \times 5 + 3 \\ &= 38 \end{aligned}$$

$$t_0 = 1 \times 5^2 + 2 \times 5 + 4 = 39$$

$$t_1 = 5 [ t_0 - 1 \times 5^2 ] + 1 = 71$$

$\hookrightarrow 5(39 - 25) + 1$

$$t_2 = 5(t_1 - 2 \times 5^2) + 2 = 107$$

$$t_{x1} = 132456$$

$$d = 10, m = 4$$

$$t_0 = 1324$$

$$t_1 = \underbrace{324 \times 10 + 5}_{3240 + 5} = 3245$$

$$= (1324 - 1 \times 10^3) \times 10 + 5$$

$$= (324 \times 10) + 5$$

$$= 3245$$

KMP  $\rightarrow$  Knuth - Morris - Pratt.

↳ Constructing LPS (longest proper prefix)

prefix      suffix      Array).

IP - ababc

0IP - LPS[ ] = [0, 0, 1, 2, 0]

• Proper prefixes of abcd

↳ "", a, ab, abc

• Suffixes of abcd

↳ "", d, cd, bcd, abcd

a

"

, a

" , b, ab

↓

0

ab

" , a

" , a, ab

↓

1

aba

" , a, ab

" , a, ba, aba

abab

" , a, ab, aba

" , b, ab, bab, abab

↓

longest one "1

2

ababc

" , a, ab, aba, abab

" , c, bc, abc, babc, ababc

↓

0

Common  
ones

are "

Made with Goodnotes  
length = 0

$\Rightarrow [0, 0, 1, 2, 0]$

'aaaq'

[ 0, 1, 2, 3 ]

'abcd'

[ 0, 0, 0, 0 ]

$\Rightarrow abac\ abad$

[ 0, 0, 1, 0, 1, 2, 3, 0 ]  
0 1 2 3 4

$a+b \rightarrow$ $\begin{matrix} 0 \\ 1 \end{matrix}$	$ab\bar{a}$	$abac$	$abaca$	$abacab$
$\begin{matrix} 0 \\ 1 \end{matrix}$	$\underline{a}, \underline{ab},$ $\underline{\bar{a}}, \underline{ba}$	$a, ab, aba$	$\underline{a}, \underline{ab}, \underline{aba}, \underline{abab}, \underline{bab}, \underline{babab}$	$\underline{a}, \underline{ab}, \underline{aba},$ $\underline{b}, \underline{ab}$
$\begin{matrix} 2 \\ 3 \end{matrix}$	$\underline{\bar{a}}$	$c, \underline{aca}, \underline{cab}, \underline{bac}$	$\underline{a}, \underline{ac}, \underline{aca}, \underline{acab}, \underline{aca}, \underline{cab}, \underline{bac}$	$\underline{\underline{a}}, \underline{\underline{ab}}, \underline{\underline{aba}},$ $\underline{\underline{b}}, \underline{\underline{ab}}$
$\begin{matrix} 4 \\ 5 \end{matrix}$	$\underline{\underline{\bar{a}}}$	$\underline{\underline{0}}$	$\underline{\underline{1}}$	$\underline{\underline{0}}$

$abaca$   
 $\underline{a}, \underline{ab}, \underline{aba}, \underline{abab}, \underline{bab}, \underline{babab}$

$abacab$   
 $\underline{a}, \underline{ab}, \underline{aba},$   
 $\underline{b}, \underline{ab}$

abb    abb

