

Week 9 Problem Set

Running Time of Programs

[Show with no answers] [Show with all answers]

1. (Asymptotic running times)

Suppose you have the choice between three algorithms:

- Algorithm A solves your problem by dividing it into five subproblems of half the size, recursively solving each subproblem, and then combining the solutions in linear time.
- Algorithm B solves problems of size n by recursively solving two subproblems of size $n - 1$ and then combining the solutions in constant time.
- Algorithm C solves problems of size n by dividing them into nine subproblems of size $\frac{n}{3}$, recursively solving each subproblem, and then combining the solutions in $\mathcal{O}(n^2)$ time.

Estimate the running times of each of these algorithms. Which one would you choose?

[show answer]

2. (Recurrences for algorithm analysis)

Recall the recurrence for Mergesort: $T(1) = 0$; $T(n) = 2T(\frac{n}{2}) + (n - 1)$ for $n > 1$.

Prove by induction that $T(n) = n \cdot (\log_2 n - 1) + 1$ for all $n = 2^k$ (with $k \geq 1$).

[show answer]

3. (Recursive algorithms)

- Analyse the complexity of the following recursive algorithm to test whether a number x occurs in an *unordered* list $L = [x_1, x_2, \dots, x_n]$ of size n . Take the cost to be the number of list element comparison operations.

Search($x, L = [x_1, x_2, \dots, x_n]$):

if $x_1 = x$ **then return yes**

else if $n > 1$ **then return** **Search**($x, [x_2, \dots, x_n]$)

else return no

- Analyse the complexity of the following recursive algorithm to test whether a number x occurs in an *ordered* list $L = [x_1, x_2, \dots, x_n]$ of size n . Take the cost to be the number of list element comparison operations.

BinarySearch($x, L = [x_1, x_2, \dots, x_n]$):

if $n = 0$ **then return no**

else if $x_{\lceil \frac{n}{2} \rceil} > x$ **then return** **BinarySearch**($x, [x_1, \dots, x_{\lceil \frac{n}{2} \rceil - 1}]$)

else if $x_{\lceil \frac{n}{2} \rceil} < x$ **then return** **BinarySearch**($x, [x_{\lceil \frac{n}{2} \rceil + 1}, \dots, x_n]$)

else return yes

[show answer]

4. Challenge Exercise

Without using the Master Theorem, give tight big-Oh upper bounds for the divide-and-conquer recurrence $T(1) = 1$; $T(n) = T(\frac{n}{2}) + g(n)$, for $n > 1$, where

a. $g(n) = 1$

b. $g(n) = 2n$

c. $g(n) = n^2$

[\[show answer\]](#)