

图像灰度压缩

一幅4*4的图像, 灰度值序列如下. 请根据课堂上所讲代码, 写出构造解的S数组、l数组和b数组, 追踪解的S数组. 需要体现做题过程, 如: 每一轮i的循环写出内部循环的前两次和最后两次, 内部循环少于等于四次的需要全部的计算过程

灰度值 10 9 12 40 50 35 15 12 8 10 9 15 11 130 160 240

构造阶段的b数组

已知灰度值序列为 10 9 12 40 50 35 15 12 8 10 9 15 11 130 160 240 , 共 $n = 16$ 个元素。

根据 length 函数:

- $b[1] = \text{length}(10)$:
 - $i = 10 / 2 = 5$, $k = 2$;
 - $i = 5 / 2 = 2$, $k = 3$;
 - $i = 2 / 2 = 1$, $k = 4$;
 - $i = 1 / 2 = 0$, 返回 $k = 4$ 。
- $b[2] = \text{length}(9)$:
 - $i = 9 / 2 = 4$, $k = 2$;
 - $i = 4 / 2 = 2$, $k = 3$;
 - $i = 2 / 2 = 1$, $k = 4$;
 - $i = 1 / 2 = 0$, 返回 $k = 4$ 。
- $b[3] = \text{length}(12)$:
 - $i = 12 / 2 = 6$, $k = 2$;
 - $i = 6 / 2 = 3$, $k = 3$;
 - $i = 3 / 2 = 1$, $k = 4$;
 - $i = 1 / 2 = 0$, 返回 $k = 4$ 。
- $b[4] = \text{length}(40)$:
 - $i = 40 / 2 = 20$, $k = 2$;
 - $i = 20 / 2 = 10$, $k = 3$;
 - $i = 10 / 2 = 5$, $k = 4$;
 - $i = 5 / 2 = 2$, $k = 5$;
 - $i = 2 / 2 = 1$, $k = 6$;
 - $i = 1 / 2 = 0$, 返回 $k = 6$ 。
- $b[5] = \text{length}(50)$:
 - $i = 50 / 2 = 25$, $k = 2$;
 - $i = 25 / 2 = 12$, $k = 3$;
 - $i = 12 / 2 = 6$, $k = 4$;
 - $i = 6 / 2 = 3$, $k = 5$;

- $i = 3 / 2 = 1$, $k = 6$;
- $i = 1 / 2 = 0$, 返回 $k = 6$ 。

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- $b[16] = \text{length}(240)$:
 - $i = 240 / 2 = 120$, $k = 2$;
 - $i = 120 / 2 = 60$, $k = 3$;
 - $i = 60 / 2 = 30$, $k = 4$;
 - $i = 30 / 2 = 15$, $k = 5$;
 - $i = 15 / 2 = 7$, $k = 6$;
 - $i = 7 / 2 = 3$, $k = 7$;
 - $i = 3 / 2 = 1$, $k = 8$;
 - $i = 1 / 2 = 0$, 返回 $k = 8$ 。

以此类推, 计算出完整的 b 数组: $b = [0, 4, 4, 4, 6, 6, 6, 4, 4, 4, 4, 4, 4, 4, 8, 8, 8]$, 索引从 0 开始

构造阶段的 s 数组和 l 数组

最终结果:

$s[] = 15 \ 19 \ 23 \ 35 \ 41 \ 47 \ 53 \ 59 \ 65 \ 71 \ 77 \ 82 \ 86 \ 105 \ 113 \ 121$, 索引从 1 开始

$l[] : 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10 \ 11 \ 6 \ 7 \ 1 \ 2 \ 3$, 索引从 1 开始

步骤:

初始化: $L_{\max} = 256$, $header = 11$, $s[0] = 0$

- $i=1$ 时, $b[i]=4, b_{\max}=4, s[1]=s[0]+b_{\max}=4; l[1]=1; s[1] += header = 15$
 - $i=1 < j=2$ 不进入内层循环

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- $i=2$ 时, $b[i]=4, b_{\max}=4, s[2]=s[1]+b_{\max}=19; l[2]=1$
 - $j=2$
 - $b_{\max}=4 == b[i-j+1]=b[1]=4$ 不更新 b_{\max} ,
 - $s[2]=19 > s[0] + j * b_{\max}=0+2*4=8$,
 - 所以 $s[2]=s[0] + j * b_{\max}=2*4=8, l[2]=2$,
 - $s[2] += header = 19$

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- $i=3$ 时, $b[i]=4, b_{\max}=4, s[3]=s[2]+b_{\max}=19+4=23; l[3]=1$
 - $j=2$
 - $b_{\max}=4 == b[i-j+1]=b[2]=4$ 不更新 b_{\max} ,
 - $s[3]=23 == s[1] + j * b_{\max}=15+2*4=23$, 不更新 $s[3]$ 、 $l[3]$

- $j=3$
 - $bmax=4 == b[i-j+1]=b[1]=4$ 不更新 $bmax$,
 - $s[3]=23 > s[0] + j * bmax=0+2*4=8$,
 - 所以 $s[3]=s[0] + j * bmax=3*4=12$, $l[3]=3$
 - $s[3]+=header=12+11=23$
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- $i=4$ 时, $b[i]=6$, $bmax=6$, $s[4]=s[3]+bmax=23+6=29$; $l[4]=1$

- $j=2$
 - $bmax=6 > b[i-j+1]=b[3]=4$ 不更新 $bmax$,
 - $s[4]=29 < s[2] + j * bmax=19+2*6=41$, 不更新 $s[4]$ 、 $l[4]$
 - $j=3$
 - $bmax=6 > b[i-j+1]=b[2]=4$ 不更新 $bmax$,
 - $s[4]=29 > s[1] + j * bmax=15+3*6=33$,
 - 所以 $s[4]=s[1] + j * bmax=33$, $l[4]=3$
 - $j=4$
 - $bmax=6 > b[i-j+1]=b[1]=4$ 不更新 $bmax$,
 - $s[4]=33 > s[0] + j * bmax=0+2*4=8$,
 - 所以 $s[3]=s[0] + j * bmax=3*4=12$, $l[3]=3$
 - $s[3]+=header=12+11=23$
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- $i = 5$ 时, $b[5] = 6$, $bmax = 6$, $s[5] = s[4] + bmax = 35 + 6 = 41$, $l[5] = 1$

- $j=2$
 - $b[i - j + 1] = b[4] = 6 \rightarrow bmax = \max(6, 6) = 6$
 - $s[i - j] + j * bmax = s[3] + 2 * 6 = 23 + 12 = 35$
 - 由于 $35 < 41 \rightarrow$ 更新 $s[5] = 35$, $l[5] = 2$
- $j=3$
 - $b[i - j + 1] = b[3] = 4 \rightarrow bmax = \max(6, 4) = 6$
 - $s[i - j] + j * bmax = s[2] + 3 * 6 = 19 + 18 = 37$
 - 由于 $37 > 35 \rightarrow$ 不更新
- $j=4$
 - $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$
 - $s[i - j] + j * bmax = s[1] + 4 * 6 = 15 + 24 = 39$
 - 由于 $39 > 35 \rightarrow$ 不更新
- $j=5$
 - $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$
 - $s[i - j] + j * bmax = s[0] + 5 * 6 = 30$
 - 由于 $30 < 35 \rightarrow$ 更新 $s[5] = 30$, $l[5] = 5$

- $s[5] = 30 + 11 = 41$
 - $l[5] = 5$
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- $i = 6$ 时, $b[6] = 6$, $bmax = 6$, $s[6] = s[5] + bmax = 41 + 6 = 47$, $l[6] = 1$

◦ $j = 2$

- $b[i - j + 1] = b[5] = 6 \rightarrow bmax = \max(6, 6) = 6$
- $s[i - j] + j * bmax = s[4] + 2 * 6 = 35 + 12 = 47$
- 由于 $47 == 47 \rightarrow$ 不更新

◦ $j = 3$

- $b[i - j + 1] = b[4] = 6 \rightarrow bmax = \max(6, 6) = 6$
- $s[i - j] + j * bmax = s[3] + 3 * 6 = 23 + 18 = 41$
- 由于 $41 < 47 \rightarrow$ 更新 $s[6] = 41$, $l[6] = 3$

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◦ $j = 5$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$
- $s[i - j] + j * bmax = s[1] + 5 * 6 = 15 + 30 = 45$
- 由于 $45 > 41 \rightarrow$ 不更新

◦ $j = 6$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$
 - $s[i - j] + j * bmax = s[0] + 6 * 6 = 36$
 - 由于 $36 < 41 \rightarrow$ 更新 $s[6] = 36$, $l[6] = 6$
 - $s[6] = 36 + 11 = 47$
 - $l[6] = 6$
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- $i = 7$ 时, $b[7] = 4$, $bmax = 4$, $s[7] = s[6] + bmax = 47 + 4 = 51$, $l[7] = 1$

◦ $j = 2$

- $b[i - j + 1] = b[6] = 6 \rightarrow bmax = \max(4, 6) = 6$
- $s[i - j] + j * bmax = s[5] + 2 * 6 = 41 + 12 = 53$
- 由于 $53 > 51 \rightarrow$ 不更新

◦ $j = 3$

- $b[i - j + 1] = b[5] = 6 \rightarrow bmax = \max(4, 6) = 6$
- $s[i - j] + j * bmax = s[4] + 3 * 6 = 35 + 18 = 53$
- 由于 $53 > 51 \rightarrow$ 不更新

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◦ $j=6$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$
- $s[i - j] + j * bmax = s[1] + 6 * 6 = 15 + 36 = 51$
- 由于 $51 > 47 \rightarrow$ 不更新

◦ $j=7$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$
 - $s[i - j] + j * bmax = s[0] + 7 * 6 = 42$
 - 由于 $42 < 47 \rightarrow$ 更新 $s[7] = 42$, $l[7] = 7$
 - $s[7] = 42 + 11 = 53$
 - $l[7] = 7$
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- $i = 8$ 时, $b[8] = 4$, $bmax = 4$, $s[8] = s[7] + bmax = 53 + 4 = 57$, $l[8] = 1$

◦ $j=2$

- $b[i - j + 1] = b[7] = 4 \rightarrow bmax = \max(4, 4) = 4$
- $s[i - j] + j * bmax = s[6] + 2 * 4 = 47 + 8 = 55$
- 由于 $55 < 57 \rightarrow$ 更新 $s[8] = 55$, $l[8] = 2$

◦ $j=3$

- $b[i - j + 1] = b[6] = 6 \rightarrow bmax = \max(4, 6) = 6$
- $s[i - j] + j * bmax = s[5] + 3 * 6 = 41 + 18 = 59$
- 由于 $59 > 55 \rightarrow$ 不更新

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• $j=7$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$
- $s[i - j] + j * bmax = s[1] + 7 * 6 = 15 + 42 = 57$
- 由于 $57 > 53 \rightarrow$ 不更新

• $j=8$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$
 - $s[i - j] + j * bmax = s[0] + 8 * 6 = 48$
 - 由于 $48 < 53 \rightarrow$ 更新 $s[8] = 48$, $l[8] = 8$
 - $s[8] = 48 + 11 = 59$
 - $l[8] = 8$
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- $i = 9$ 时 $b[9] = 4$, $bmax = 4$, $s[9] = s[8] + bmax = 59 + 4 = 63$, $l[9] = 1$

◦ $j = 2$

- $b[i - j + 1] = b[8] = 4 \rightarrow bmax = \max(4, 4) = 4$
- $s[i - j] + j * bmax = s[7] + 2 * 4 = 53 + 8 = 61$
- 由于 $61 < 63 \rightarrow$ 更新 $s[9] = 61$, $l[9] = 2$

◦ $j = 3$

- $b[i - j + 1] = b[7] = 4 \rightarrow bmax = \max(4, 4) = 4$
- $s[i - j] + j * bmax = s[6] + 3 * 4 = 47 + 12 = 59$
- 由于 $59 < 61 \rightarrow$ 更新 $s[9] = 59$, $l[9] = 3$

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◦ $j = 8$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$
- $s[i - j] + j * bmax = s[1] + 8 * 6 = 15 + 48 = 63$
- 由于 $63 > 59 \rightarrow$ 不更新

◦ $j = 9$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$
- $s[i - j] + j * bmax = s[0] + 9 * 6 = 54$
- 由于 $54 < 59 \rightarrow$ 更新 $s[9] = 54$, $l[9] = 9$
- $s[9] = 54 + 11 = 65$
- $l[9] = 9$

- $i = 10$ 时, $b[10] = 4$, $bmax = 4$, $s[10] = s[9] + bmax = 65 + 4 = 69$, $l[10] = 1$

◦ $j = 2$

- $b[i - j + 1] = b[9] = 4 \rightarrow bmax = \max(4, 4) = 4$
- $s[i - j] + j * bmax = s[8] + 2 * 4 = 59 + 8 = 67$
- 由于 $67 < 69 \rightarrow$ 更新 $s[10] = 67$, $l[10] = 2$

◦ $j = 3$

- $b[i - j + 1] = b[8] = 4 \rightarrow bmax = \max(4, 4) = 4$
- $s[i - j] + j * bmax = s[7] + 3 * 4 = 53 + 12 = 65$
- 由于 $65 < 67 \rightarrow$ 更新 $s[10] = 65$, $l[10] = 3$

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◦ $j = 9$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$

- $s[i - j] + j * bmax = s[1] + 9 * 6 = 15 + 54 = 69$

- 由于 $69 > 63 \rightarrow$ 不更新

- $j = 10$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$

- $s[i - j] + j * bmax = s[0] + 10 * 6 = 60$

- 由于 $60 < 63 \rightarrow$ 更新 $s[10] = 60$, $l[10] = 10$

- $s[10] = 60 + 11 = 71$

- $l[10] = 10$

- $i = 11$ 时, $b[11] = 4$, $bmax = 4$, $s[11] = s[10] + bmax = 71 + 4 = 75$, $l[11] = 1$

- $j = 2$

- $b[i - j + 1] = b[10] = 4 \rightarrow bmax = \max(4, 4) = 4$

- $s[i - j] + j * bmax = s[9] + 2 * 4 = 65 + 8 = 73$

- 由于 $73 < 75 \rightarrow$ 更新 $s[11] = 73$, $l[11] = 2$

- $j = 3$

- $b[i - j + 1] = b[9] = 4 \rightarrow bmax = \max(4, 4) = 4$

- $s[i - j] + j * bmax = s[8] + 3 * 4 = 59 + 12 = 71$

- 由于 $71 < 73 \rightarrow$ 更新 $s[11] = 71$, $l[11] = 3$

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- $j = 10$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$

- $s[i - j] + j * bmax = s[1] + 10 * 6 = 15 + 60 = 75$

- 由于 $75 > 67 \rightarrow$ 不更新

- $j = 11$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$

- $s[i - j] + j * bmax = s[0] + 11 * 6 = 66$

- 由于 $66 < 67 \rightarrow$ 更新 $s[11] = 66$, $l[11] = 11$

- $s[11] = 66 + 11 = 77$

- $l[11] = 11$

- $i = 12$ 时, $b[12] = 4$, $bmax = 4$, $s[12] = s[11] + bmax = 77 + 4 = 81$, $l[12] = 1$

- $j = 2$

- $b[i - j + 1] = b[11] = 4 \rightarrow bmax = \max(4, 4) = 4$

- $s[i - j] + j * bmax = s[10] + 2 * 4 = 71 + 8 = 79$

- 由于 $79 < 81 \rightarrow$ 更新 $s[12] = 79$, $l[12] = 2$

◦ $j=3$

- $b[i - j + 1] = b[10] = 4 \rightarrow bmax = \max(4, 4) = 4$

- $s[i - j] + j * bmax = s[9] + 3 * 4 = 65 + 12 = 77$

- 由于 $77 < 79 \rightarrow$ 更新 $s[12] = 77$, $l[12] = 3$

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◦ $j=11$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$

- $s[i - j] + j * bmax = s[1] + 11 * 6 = 15 + 66 = 81$

- 由于 $81 > 71 \rightarrow$ 不更新

◦ $j=12$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$

- $s[i - j] + j * bmax = s[0] + 12 * 6 = 72$

- 由于 $72 > 71 \rightarrow$ 不更新

- $s[12] = 71 + 11 = 82$

- $l[12] = 6$

- $i=13$ 时, $b[13] = 4$, $bmax = 4$, $s[13] = s[12] + bmax = 82 + 4 = 86$, $l[13] = 1$

◦ $j=2$

- $b[i - j + 1] = b[12] = 4 \rightarrow bmax = \max(4, 4) = 4$

- $s[i - j] + j * bmax = s[11] + 2 * 4 = 77 + 8 = 85$

- 由于 $85 < 86 \rightarrow$ 更新 $s[13] = 85$, $l[13] = 2$

◦ $j=3$

- $b[i - j + 1] = b[11] = 4 \rightarrow bmax = \max(4, 4) = 4$

- $s[i - j] + j * bmax = s[10] + 3 * 4 = 71 + 12 = 83$

- 由于 $83 < 85 \rightarrow$ 更新 $s[13] = 83$, $l[13] = 3$

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◦ $j=12$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(6, 4) = 6$

- $s[i - j] + j * bmax = s[1] + 12 * 6 = 15 + 72 = 87$

- 由于 $87 > 75 \rightarrow$ 不更新

◦ $j=13$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(6, 4) = 6$

- $s[i - j] + j * bmax = s[0] + 13 * 6 = 78$
- 由于 $78 > 75 \rightarrow$ 不更新
- $s[13] = 75 + 11 = 86$
- $l[13] = 7$

- $i=14$ 时, $b[14] = 8$, $bmax = 8$, $s[14] = s[13] + bmax = 86 + 8 = 94$, $l[14] = 1$

◦ $j=2$

- $b[i - j + 1] = b[13] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[12] + 2 * 8 = 82 + 16 = 98$
- 由于 $98 > 94 \rightarrow$ 不更新

◦ $j=3$

- $b[i - j + 1] = b[12] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[11] + 3 * 8 = 77 + 24 = 101$
- 由于 $101 > 94 \rightarrow$ 不更新

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◦ $j=13$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[1] + 13 * 8 = 15 + 104 = 119$
- 由于 $119 > 94 \rightarrow$ 不更新

◦ $j=14$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[0] + 14 * 8 = 112$
- 由于 $112 > 94 \rightarrow$ 不更新
- $s[14] = 94 + 11 = 105$
- $l[14] = 1$

- $i = 15$ 时, $b[15] = 8$, $bmax = 8$, $s[15] = s[14] + bmax = 105 + 8 = 113$, $l[15] = 1$

◦ $j=2$

- $b[i - j + 1] = b[14] = 8 \rightarrow bmax = \max(8, 8) = 8$
- $s[i - j] + j * bmax = s[13] + 2 * 8 = 86 + 16 = 102$
- 由于 $102 < 113 \rightarrow$ 更新 $s[15] = 102$, $l[15] = 2$

◦ $j=3$

- $b[i - j + 1] = b[13] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[12] + 3 * 8 = 82 + 24 = 106$

- 由于 $106 > 102 \rightarrow$ 不更新

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◦ $j = 14$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[1] + 14 * 8 = 15 + 112 = 127$
- 由于 $127 > 102 \rightarrow$ 不更新

◦ $j = 15$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[0] + 15 * 8 = 120$
- 由于 $120 > 102 \rightarrow$ 不更新
- $s[15] = 102 + 11 = 113$
- $l[15] = 2$

- $i=16$ 时, $b[16] = 8, bmax = 8, s[16] = s[15] + bmax = 113 + 8 = 121, l[16] = 1$

◦ $j = 2$

- $b[i - j + 1] = b[15] = 8 \rightarrow bmax = \max(8, 8) = 8$
- $s[i - j] + j * bmax = s[14] + 2 * 8 = 105 + 16 = 121$
- 由于 $121 == 121 \rightarrow$ 不更新

◦ $j = 3$

- $b[i - j + 1] = b[14] = 8 \rightarrow bmax = \max(8, 8) = 8$
- $s[i - j] + j * bmax = s[13] + 3 * 8 = 86 + 24 = 110$
- 由于 $110 < 121 \rightarrow$ 更新 $s[16] = 110, l[16] = 3$

◦ $j = 15$

- $b[i - j + 1] = b[2] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[1] + 15 * 8 = 15 + 120 = 135$
- 由于 $135 > 110 \rightarrow$ 不更新

◦ $j = 16$

- $b[i - j + 1] = b[1] = 4 \rightarrow bmax = \max(8, 4) = 8$
- $s[i - j] + j * bmax = s[0] + 16 * 8 = 128$
- 由于 $128 > 110 \rightarrow$ 不更新
- $s[16] = 110 + 11 = 121$
- $l[16] = 3$

追踪解的s数组

已知结果

- l 数组: $[0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 6, 7, 1, 2, 3]$

计算过程

从 `n = 16` 开始调用 `Traceback` 函数：

1. **初始调用：** `Traceback(16, i, s, 1)`
 - 递归调用： `Traceback(16 - 1[16], i, s, 1) → Traceback(13, i, s, 1)`
 - 记录： `s[i++] = 13`
2. **递归调用：** `Traceback(13, i, s, 1)`
 - 递归调用： `Traceback(13 - 1[13], i, s, 1) → Traceback(6, i, s, 1)`
 - 记录： `s[i++] = 6`
3. **递归调用：** `Traceback(6, i, s, 1)`
 - 递归调用： `Traceback(6 - 1[6], i, s, 1) → Traceback(0, i, s, 1)`
 - 记录： `s[i++] = 0`
4. **终止条件：** `n == 0`，返回

追踪解的 `s` 数组

最终 `s` 数组存储的是每个分段的起始位置（按调用顺序逆序记录）：

```
s = [0, 6, 13]
```

结果解释

- `s[0] = 0`：第一个分段从像素 1 开始（0 是前一个位置）
- `s[1] = 6`：第二个分段从像素 7 开始（6 是前一个位置）
- `s[2] = 13`：第三个分段从像素 14 开始（13 是前一个位置）

验证分段方案

- **分段 1：** 像素 1-6（长度 6，`bmax = 6`）
- **分段 2：** 像素 7-13（长度 7，`bmax = 4`）
- **分段 3：** 像素 14-16（长度 3，`bmax = 8`）

Code

```
#include <iostream>
#include <vector>
using namespace std;

// 计算整数的二进制位数
int length(int i)
{
    int k = 1;
    i /= 2;
    while (i > 0)
    {
        k++;
        i /= 2;
    }
    return k;
}
```

```

// 图像压缩函数
void Compress(int n, int p[], int s[], int l[], int b[])
{
    int Lmax = 256, header = 11;
    s[0] = 0;
    for (int i = 1; i <= n; i++)
    {
        b[i] = length(p[i]);
        int bmax = b[i];
        s[i] = s[i - 1] + bmax;
        l[i] = 1;
        for (int j = 2; j <= i && j <= Lmax; j++)
        {
            if (bmax < b[i - j + 1])
                bmax = b[i - j + 1];
            if (s[i] > s[i - j] + j * bmax)
            {
                s[i] = s[i - j] + j * bmax;
                l[i] = j;
            }
        }
        s[i] += header;
    }
}

// 回溯函数，确定最优分段
void Traceback(int n, int &i, int s[], int l[])
{
    if (n == 0)
        return;
    Traceback(n - l[n], i, s, l);
    s[i++] = n - l[n];
}

// 输出压缩结果
void Output(int s[], int l[], int b[], int n)
{
    cout << "图像压缩后的最小空间为: " << s[n] << endl;
    int m = 0;
    Traceback(n, m, s, l);
    s[m] = n;
    cout << "将原灰度序列分成" << m << "段序列段" << endl;
    for (int j = 1; j <= m; j++)
    {
        l[j] = l[s[j]];
        b[j] = b[s[j]];
    }
    for (int j = 1; j <= m; j++)
    {
        cout << "段" << j << ": 长度=" << l[j] << ", 存储位数=" << b[j] << endl;
    }
}

int main()
{
    // 示例灰度值序列 (4x4 图像)

```

```

    int p[] = {0, 10, 9, 12, 40, 50, 35, 15, 12, 8, 10, 9, 15, 11, 130, 160,
240};
    int n = 16; // 图像像素数

    // 初始化数组
    int *s = new int[n + 1];
    int *l = new int[n + 1];
    int *b = new int[n + 1];

    // 执行压缩
    Compress(n, p, s, l, b);
    // 输出b,s,l数组
    cout << "b数组:" << endl;
    for (int i = 1; i <= n; i++)
    {
        cout << b[i] << " ";
    }
    cout << endl;

    cout << "s数组:" << endl;
    for (int i = 1; i <= n; i++)
    {
        cout << s[i] << " ";
    }
    cout << endl;
    cout << "l数组:" << endl;
    for (int i = 1; i <= n; i++)
    {
        cout << l[i] << " ";
    }
    cout << endl;

    // 输出结果
    Output(s, l, b, n);

    // 释放内存
    delete[] s;
    delete[] l;
    delete[] b;

    return 0;
}

```

Out Put

```
b数组：
4 4 4 6 6 6 4 4 4 4 4 4 4 8 8 8
s数组：
15 19 23 35 41 47 53 59 65 71 77 82 86 105 113 121
l数组：
1 2 3 4 5 6 7 8 9 10 11 6 7 1 2 3
图像压缩后的最小空间为： 121
将原灰度序列分成3段序列段
段1： 长度=6， 存储位数=6
段2： 长度=7， 存储位数=4
段3： 长度=3， 存储位数=8
```

最终分段

- 段 1： 像素 1-6 (长度 6, b=6)
- 段 2： 像素 7-13 (长度 7, b=4)
- 段 3： 像素 14-16 (长度 3, b=8)

总位数： 每段头信息11位

$$(6 \times 6 + 7 \times 4 + 3 \times 8) + 3 \times 11 = 121 \text{ 位}$$