

MISCELLANEOUS SET 1

Permutation in String

Given two strings s1 and s2, write a function to return true if s2 contains the permutation of s1. In other words, one of the first string's permutations is the substring of the second string.

Example 1:

Input: s1 = "ab" s2 = "eidbaooo"

Output: True

Explanation: s2 contains one permutation of s1 ("ba").

```
public boolean checkInclusion(String s1, String s2) {
    int n = s1.length();
    int m = s2.length();
    if(n>m)
        return false;
    HashMap<Character, Integer> map = new HashMap<>();
    for(int i=0;i<n;i++){
        char c = s1.charAt(i);
        if(map.containsKey(c)){
            map.put(c, map.get(c)+1);
        }else
            map.put(c, 1);
    }
    int i=0;
    while(i<n){
        char c = s2.charAt(i++);
        if(map.containsKey(c)){
            map.put(c, map.get(c)-1);
        }
    }
    for(i=n;i<=m;i++){
        int flag = 0;
        for(Map.Entry<Character, Integer> obj:map.entrySet()){
            if(obj.getValue()!=0){
                flag = 1;
                break;
            }
        }
        if(flag==0)
            return true;
        else if(i==m)
            break;
        else{
            char c = s2.charAt(i-n);
```

```

        if(map.containsKey(c)){
            map.put(c, map.get(c)+1);
        }
        c = s2.charAt(i);
        if(map.containsKey(c)){
            map.put(c, map.get(c)-1);
        }
    }
}
return false;
}

```

Insertion Sort on linked list

```

public ListNode insertionSortList(ListNode head) {
    if(head!=null){
        ListNode curr = head.next,parent = head;
        while(curr!=null){
            ListNode temp = head,par = null;
            while(temp!=curr && temp.val<=curr.val){
                par = temp;
                temp = temp.next;
            }
            if(temp!=curr){
                parent.next = curr.next;
                if(par==null){
                    curr.next = head;
                    head = curr;
                }else{
                    curr.next = par.next;
                    par.next = curr;
                }
                curr = parent.next;
            }else{
                parent = curr;
                curr = curr.next;
            }
        }
    }
    return head;
}

```

Product of the Last K Numbers

Implement the class `ProductOfNumbers` that supports two methods:

1. `add(int num)`

- Adds the number `num` to the back of the current list of numbers.

2. `getProduct(int k)`

- Returns the product of the last `k` numbers in the current list.
- You can assume that always the current list has **at least** `k` numbers.

At any time, the product of any contiguous sequence of numbers will fit into a single 32-bit integer without overflowing.

Approach : Using Product Prefix array.

```
class ProductOfNumbers {
    List<Integer> list;
    public ProductOfNumbers() {
        list = new ArrayList<>();
        list.add(1);
    }

    public void add(int num) {
        if(num==0){
            list.clear();
            list.add(1);
        }else{
            list.add(list.get(list.size()-1)*num);
        }
    }

    public int getProduct(int k) {
        int pos = list.size() - k - 1;
        if(pos<0)
            return 0;
        return list.get(list.size()-1) / list.get(pos);
    }
}
```

Unique Binary Search Trees

Given n , how many structurally unique **BST's** (binary search trees) that store values $1 \dots n$?

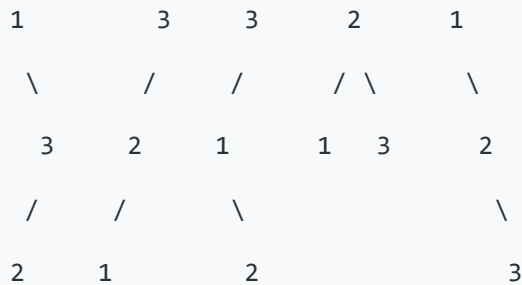
Example:

Input: 3

Output: 5

Explanation:

Given $n = 3$, there are a total of 5 unique BST's:



Approach : find catalan of n.

$C(0) = C(1) = 1;$

$C(2) = C(0) * C(1) + C(1) * C(0)$

$C(3) = C(0) * C(2) + C(1) * C(1) + C(2) * C(0)$

$C(4) = C(0) * C(3) + C(1) * C(2) + C(2) * C(1) + C(3) * C(0)$

.....

.....so on

```

public int numTrees(int n) {
    //calculate catalan(n)
    if(n==0)
        return 1;
    int dp[] = new int[n+1];
    dp[0] = dp[1] = 1;
    for(int i=2;i<=n;i++){
        dp[i] = 0;
        for(int j=0;j<i;j++){
            dp[i] += dp[j] * dp[i-j-1];
        }
    }
    return dp[n];
}

```

Delete Node in a BST

Given a root node reference of a BST and a key, delete the node with the given key in the BST. Return the root node reference (possibly updated) of the BST.

Basically, the deletion can be divided into two stages:

1. Search for a node to remove.
2. If the node is found, delete the node.

Note: Time complexity should be $O(\text{height of tree})$.

Example:

```
root = [5,3,6,2,4,null,7]
```

```
key = 3
```

```
    5
   / \
  3   6
 / \   \
2  4   7
```

```
private TreeNode smallCal(TreeNode root){
    if(root.left==null)
        return root;
    return smallCal(root.left);
}
public TreeNode deleteNode(TreeNode root, int key) {
    if(root==null)
        return null;
    TreeNode temp = root, par = null;
    while(temp!=null){
        if(temp.val==key){
            //leaf node
            if(temp.left==null && temp.right==null){
                if(par==null)
                    return null;
                else if(par.left == temp)
                    par.left = null;
                else
                    par.right = null;
                return root;
            }
            //two child node
            if(temp.left!=null && temp.right!=null){
                TreeNode small = smallCal(temp.right);
```

```

        int t = temp.val;
        temp.val = small.val;
        small.val = t;
        temp.right = deleteNode(temp.right, key);
        return root;
    }
    //one child node
    if(temp.left!=null){
        if(par==null)
            return temp.left;
        else if(par.left==temp)
            par.left = temp.left;
        else
            par.right = temp.left;
        return root;
    }
    if(temp.right!=null){
        if(par==null)
            return temp.right;
        else if(par.left==temp)
            par.left = temp.right;
        else
            par.right = temp.right;
        return root;
    }
} else {
    par = temp;
    if(key<temp.val)
        temp = temp.left;
    else
        temp = temp.right;
}
}
return root;
}

```

Insertion in AVL tree

We can maintain height of every node to make speed better to calculate height again.

```

class Solution {
    private int height(TreeNode root){
        if(root==null)

```

```

        return 0;
    return 1 + Math.max(height(root.left), height(root.right));
}
private TreeNode leftRotation(TreeNode x){
    TreeNode y = x.right;
    TreeNode T2 = y.left;

    y.left = x;
    x.right = T2;

    return y;
}
private TreeNode rightRotation(TreeNode x){
    TreeNode y = x.left;
    TreeNode T2 = y.right;

    y.right = x;
    x.left = T2;

    return y;
}
private TreeNode insert(TreeNode root, int key){
    if(root==null){
        return new TreeNode(key);
    }
    if(key<root.val)
        root.left = insert(root.left, key);
    else if(key>root.val)
        root.right = insert(root.right, key);
    else //duplicate node
        return null;

    //check for balance this node
    int LH = height(root.left);
    int rH = height(root.right);

    int balance = LH-rH;

    if(balance>1){ //left subtree height is more
        if(key<root.left.val){ //left - left case
            return rightRotation(root);
        }else{ //left - right case

```



```

        root.left = leftRotation(root.left);
        return rightRotation(root);
    }
} else if(balance<1){ //right subtree height is more
    if(key>root.right.val){ //right - right case
        return leftRotation(root);
    } else { //right - left case
        root.right = rightRotation(root.right);
        return leftRotation(root);
    }
}

return root;
}

public TreeNode sortedListToBST(ListNode head) {
    TreeNode root = null;
    while(head!=null){
        root = insert(root, head.val);
    }
    return root;
}
}

```

Add 1 to a number represented as linked list

A number **N** is represented in Linked List such that each digit corresponds to a node in linked list. You need to add 1 to it.

Example:

Input:

4
456
123
999
1879

Output:

457
124
1000
1880

```

private static int cal(Node head){
    if(head.next==null){

```

```

        if(head.data+1==10){
            head.data = 0;
            return 1;
        }else{
            head.data += 1;
            return 0;
        }
    }
    int val = cal(head.next);
    if(val==1){
        if(head.data+1==10){
            head.data = 0;
            return 1;
        }else{
            head.data += 1;
        }
    }
    return 0;
}

public static Node addOne(Node head) {
    if(head==null)
        return head;
    int val = cal(head);
    if(val==1){
        Node node = new Node(1);
        node.next = head;
        head = node;
    }
    return head;
}

```

M-Coloring Problem

Given an undirected graph and an integer **M**. The task is to determine if the graph can be colored with at most C colors such that no two adjacent vertices of the graph are colored with the same color. Here coloring of a graph means the assignment of colors to all vertices. Print 1 if it is possible to color vertices and 0 otherwise.

```

class solve {
    static boolean isSafe(List<Integer>[] G, int color[], int s, int c){
        List<Integer> l = G[s];
        for(int i=0;i<l.size();i++){
            int a = l.get(i);

```

```

        if(color[a]==c)
            return false;
        }
        return true;
    }
    // initial value s = 0 and color[i] = 0
    public static boolean graphColoring(List<Integer>[] G, int[] color, int s, int C) {
        if(s==G.length){
            return true;
        }
        for(int i=1;i<=C;i++){
            if(isSafe(G, color, s, i)){
                color[s] = i;
                if(graphColoring(G, color, s+1, C))
                    return true;
                color[s] = 0; //Backtrack
            }
        }
        return false;
    }
}

```

Prerequisite Tasks

There are a total of N tasks, labeled from 0 to N-1. Some tasks may have prerequisites, for example to do task 0 you have to first complete task 1, which is expressed as a pair: [0, 1]

Given the total number of tasks and a list of prerequisite pairs, find if it is possible to finish all tasks.

Note: If there is any **cycle** exist in graph, then no topological sort can be find i.e. impossible to complete all tasks. If order of tasks has to print, then print **topological sort** of graph.

```

class Solution {
    boolean isCycle(int s, ArrayList<ArrayList<Integer>> graph, boolean vis[],
boolean temp[]){
        if(temp[s])
            return true;
        if(vis[s])
            return false;
        vis[s] = true;
        temp[s] = true;
        ArrayList<Integer> l = graph.get(s);
        for(int i=0;i<l.size();i++){
            if(isCycle(l.get(i), graph, vis, temp))

```

```

        return true;
    }
    temp[s] = false;
    return false;
}
public boolean canFinish(int n, int[][] pre){
    ArrayList<ArrayList<Integer>> graph = new ArrayList<>();
    for(int i=0;i<n;i++){
        graph.add(new ArrayList<>());
    }
    int m = pre.length;
    for(int i=0;i<m;i++){
        graph.get(pre[i][1]).add(pre[i][0]);
    }
    boolean vis[] = new boolean[n];
    boolean temp[] = new boolean[n];
    for(int i=0;i<n;i++){
        if(!vis[i]){
            if(isCycle(i, graph, vis, temp))
                return false;
        }
    }
    return true;
}
}

```

Word Search

Given a 2D board of letters and a word. Check if the word exists in the board. The word can be constructed from letters of adjacent cells only. ie - horizontal or vertical neighbors. The same letter cell can not be used more than once.

```

class Solution {
    boolean dfs(int x,int y,int n,int m,char board[][], boolean vis[][], String word, int l){
        if(l==word.length()){
            return true;
        }
        vis[x][y] = true;
        int row[] = {-1,1,0,0};
        int col[] = {0,0,-1,1};
        for(int k=0;k<4;k++){
            int i = x + row[k];

```

```

        int j = y + col[k];
        if(i>=0 && j>=0 && i<n && j<m && !vis[i][j] &&
word.charAt(l)==board[i][j]){
            if(dfs(i,j,n,m,board,vis,word,l+1))
                return true;
        }
    }
    vis[x][y] = false; //Backtrack
    return false;
}
public int wordSearch(char[][] board, String word) {
    int n = board.length;
    int m = board[0].length;
    boolean vis[][] = new boolean[n][m];
    for(int i=0;i<n;i++){
        for(int j=0;j<m;j++){
            if(word.charAt(0)==board[i][j]){
                if(dfs(i,j,n,m,board,vis,word,1))
                    return 1;
            }
        }
    }
    return 0;
}
}

```