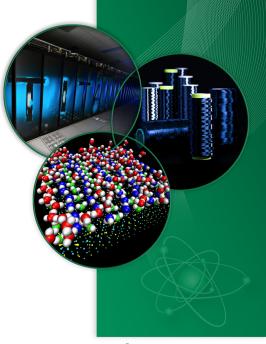
Parallel Algorithms for Monte Carlo Linear Solvers

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Motivation

- As we move towards exascale computing, the rate of errors is expected to increase dramatically
 - The probability that a compute node will fail during the course of a large scale calculation may be near 1
- Algorithms need to not only have increased concurrency/scalability but have the ability to recover from hardware faults
 - Lightweight machines
 - Heterogeneous machines
 - Both characterized by low power and high concurrency



Towards Exascale Concurrency and Resiliency

- Two basic strategies:
 - State with current "state of the art" methods and make incremental modifications to improve scalability and fault tolerance
 - Many efforts are heading in this direction, attempting to find additional concurrency to exploit
 - Start with methods having natural scalability and resiliency aspects and work at improving performance (e.g. Monte Carlo)
 - Soft failures a component of the tally variance
 - Hard failures potentially mitigated by replication
 - Concurrency enabled by several levels of parallelism



Outline

Monte Carlo Linear Solvers

- Domain Decomposition and Replication
- Scaling Studies
- Algorithm Variations
- Conclusions and Future Work



Monte Carlo Methods



Monte Carlo for Linear Systems

- ullet Suppose we want to solve $\mathbf{A}\mathbf{x} = \mathbf{b}$
- ullet If $ho({f I}-{f A})<1$, we can write the solution using the Neumann series

$$\mathbf{x} = \sum_{n=0}^{\infty} (\mathbf{I} - \mathbf{A})^n \mathbf{b} = \sum_{n=0}^{\infty} \mathbf{H}^n \mathbf{b}$$

where $\mathbf{H} \equiv (\mathbf{I} - \mathbf{A})$ is the Richardson iteration matrix

• Build the Neumann series stochastically

$$x_i = \sum_{k=0}^{\infty} \sum_{i_1}^{N} \sum_{i_2}^{N} \dots \sum_{i_k}^{N} h_{i,i_1} h_{i_1,i_2} \dots h_{i_{k-1},i_k} b_{i_k}$$

Define a sequence of state transitions

$$\nu = i \rightarrow i_1 \rightarrow \cdots \rightarrow i_{k-1} \rightarrow i_k$$



Forward Monte Carlo

- Typical choice (Monte Carlo Almost-Optimal):

$$\mathbf{P}_{ij} = rac{|\mathbf{H}_{ij}|}{\sum_{j=1}^{N} |\mathbf{H}_{ij}|}$$

- To compute solution component x_i :
 - Start a history in state *i* (with initial weight of 1)
 - ullet Transition to new state j based probabilities determined by ${f P}_i$
 - ullet Modify history weight based on corresponding entry in \mathbf{W}_{ij}
 - ullet Add contribution to ${f x}_i$ based on current history weight and value of ${f b}_j$
- \bullet A given random walk can only contribute to a single component of the solution vector with $\mathbf{x} \approx \mathbf{M_{MC}b}$



Sampling Example (Forward Monte Carlo)

Suppose

$$\mathbf{A} = \begin{bmatrix} 1.0 & -0.2 & -0.6 \\ -0.4 & 1.0 & -0.4 \\ -0.1 & -0.4 & 1.0 \end{bmatrix} \rightarrow \mathbf{H} \equiv (\mathbf{I} - \mathbf{A}) = \begin{bmatrix} 0.0 & 0.2 & 0.6 \\ 0.4 & 0.0 & 0.4 \\ 0.1 & 0.4 & 0.0 \end{bmatrix}$$

then

$$\mathbf{P} = \begin{bmatrix} 0.0 & 0.25 & 0.75 \\ 0.5 & 0.0 & 0.5 \\ 0.2 & 0.8 & 0.0 \end{bmatrix}, \quad \mathbf{W} = \begin{bmatrix} 0.0 & 0.8 & 0.8 \\ 0.8 & 0.0 & 0.8 \\ 0.5 & 0.5 & 0.0 \end{bmatrix}$$

• If a history is started in state 3, there is a 20% chance of it transitioning to state 1 and an 80% chance of moving to state 2



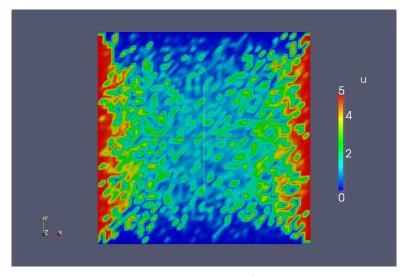


Figure : Forward solution. 2.5×10^3 total histories.



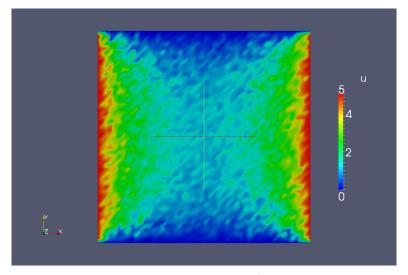


Figure : Forward solution. 2.5×10^4 total histories.



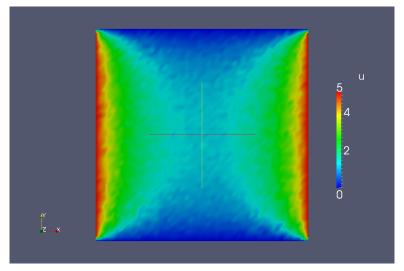


Figure : Forward solution. 2.5×10^5 total histories.



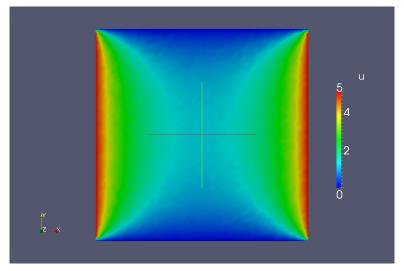


Figure : Forward solution. 2.5×10^6 total histories.



Domain Decomposition and Replication



Domain Decomposed Monte Carlo

- Each parallel process owns a piece of the domain (linear system)
- Random walks must be transported between adjacent domains through parallel communication
- Domain decomposition determined by the input system
- Load balancing not yet addressed

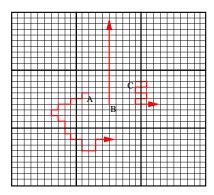
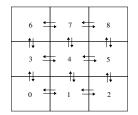
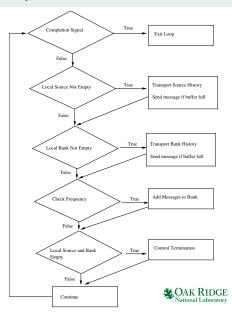


Figure: Domain decomposition example illustrating how domain-to-domain transport creates communication costs.

Asynchronous Monte Carlo Transport Kernel

- General extension of the Milagro algorithm (LANL)
- Asynchronous nearest neighbor communication of histories
- System-tunable communication parameters of buffer size and check frequency (performance impact)
- Need an asynchronous strategy for exiting the transport loop without a collective (running sum)





Exiting the Transport Loop without Collectives

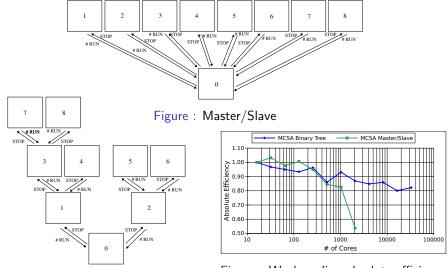


Figure: Binary Tree

Figure: Weak scaling absolute efficiency

Replication

Different batches of Monte Carlo samples can be combined in summation via superposition if they have different random number streams. For two different batches:

$$\mathbf{M_{MC}x} = \frac{1}{2}(\mathbf{M_1} + \mathbf{M_2})\mathbf{x}$$

Consider each of these batches independent *subsets* of a Monte Carlo operator where now the operator can be formed as a general additive decomposition of N_S subsets:

$$\mathbf{M_{MC}} = \frac{1}{N_S} \sum_{n=1}^{N_S} \mathbf{M_n}$$

We replicate the linear problem and form each subset on a different group of parallel processes. Applying the subsets to a vector requires an AllReduce to form the sum. Each subset is domain decomposed. *OAK RIDGE National Laboratory

Scaling Studies



Parallel Test - Simplified P_N (SP_N) Assembly Problem

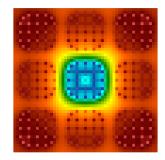


Figure : SP_N solution example

The (SP_N) equations are an approximation to the Boltzmann neutron transport equation used to simulate nuclear reactors

$$-\nabla \cdot \mathbb{D}_n \nabla \mathbb{U}_n + \sum_{i=1}^4 \mathbb{A}_{nm} \mathbb{U}_m = \frac{1}{k} \sum_{i=1}^4 \mathbb{F}_{nm} \mathbb{U}_m$$

Scaling problem -1×1 to 17×17 array of fuel assemblies with 289 pins each resolved by a 2×2 spatial mesh and 200 axial zones

- 7 energy groups, 1 angular moment, 1.6M to 273.5M degrees of freedom
- 64 to 10,800 computational cores via domain decomposition
- We are usually interested in solving generalized eigenvalue problem - we use the operator from these problems to test the kernel scaling

Monte Carlo Communication Parameters

		Message Check Frequency				
		128	256	512	1024	
	256	1.054	1.061	1.076	1.076	
	512	1.103	1.146	1.211	1.270	
	1024	1.062	1.088	1.133	1.176	
Mesage Buffer Size	2048	1.030	1.042	1.072	1.107	
	4096	1.010	1.012	1.025	1.050	
	8192	1.001	1.000	1.008	1.018	
	16384	1.017	1.003	1.010	1.009	

- 64 cores, 1.6M DOFs, history length of 15, 3 histories per DOF
- Buffer size dominates the trend
- 27% decrease in runtime observed for bad parameter choices
- Worth the time to do this parameter study when running on new hardware

Monte Carlo Strong Scaling



Monte Carlo Weak Scaling



Monte Carlo Replication Scaling



Algorithm Variations



Monte Carlo Synthetic Acceleration

- Devised by Evans and Mosher in the 2000's as an acceleration scheme for radiation diffusion problems (LANL)
- Can be abstracted as a general linear solver with Monte Carlo as a preconditioner
- Combine with Richardson iteration as a "smoother" in between Monte Carlo steps:

$$\begin{aligned} \mathbf{r}^k &= \mathbf{b} - \mathbf{A}\mathbf{x}^k \\ \mathbf{x}^{k+1/2} &= \mathbf{x}^k + \mathbf{r}^k \\ \mathbf{r}^{k+1/2} &= \mathbf{b} - \mathbf{A}\mathbf{x}^{k+1/2} \\ \mathbf{x}^{k+1} &= \mathbf{x}^{k+1/2} + \mathbf{M}_{\mathbf{MC}}\mathbf{r}^{k+1/2} \end{aligned}$$



Matrix-Free Algorithm

- ullet At each application of $\mathbf{M}_{\mathbf{MC}}$, execute the Monte Carlo process
- Must perform Monte Carlo every time you want to apply with a better approximation requiring more time and more operations
- Variations in random number streams are amortized over iterations



Stochastic Approximate Inverse Algorithm

- ullet Construct ${f M_{MC}}$ as a sparse matrix by executing the Monte Carlo process once and tallying the row entries
- Use this operator as a stochastic approximation to the inverse
- A better approximation to the inverse requires more setup time and more storage
- Will investigate a drop tolerance strategy

L	N_S	NNZ	NNZ Ratio	MC Time (s)	Setup Time (s)	MCSA Iters
3	2	484,714	0.41	0.104	0.671	255
3	3	622,123	0.52	0.145	0.705	255
5	2	783,153	0.66	0.158	0.737	185
5	3	1,032,573	0.87	0.237	0.831	171
5	4	1,241,442	1.05	0.302	0.906	171
10	3	1,969,540	1.66	0.433	1.061	95
10	4	2,416,572	2.04	0.570	1.214	95
15	3	2,867,005	2.42	0.645	1.317	132
15	4	3,544,181	2.99	0.833	1.534	67
15	5	4,157,269	3.50	1.029	1.765	66

Table : MCSA Performance. A had 115,600 rows and 1,186,464 non-zero entries.

Unpreconditioned Algorithm Comparison



Conclusions and Future Work



Conclusions

- Monte Carlo methods offer great potential for both resilient and highly parallel solvers
- A fully asynchronous algorithm provides a scheme without collectives
- Good scaling demonstrated so far on reasonably load balanced problems
- Replication a potential resiliency strategy with some overhead
- Matrix-free and stochastic approximate inverse algorithms are complementary - trade operations for storage



Conclusions

 Vectorization an area of active research with a focus on heterogenous architectures (Titan and Summit)

• Multiple threading models are being explored (Kokkos, HPX, etc.)

• Extending methods to broader problem areas is significant algorithmic challenge and an attractive area for continued research

Performance modeling and resiliency simulations this FY

