

Relational Schema Design

Elmasri and Navathe 9.1 – 9.2

Outline

- ER to Relational Mapping
 - Steps 1 - 7

Step 1 – Regular Entity Types

- For each regular entity type E
 - create relation R that includes all simple attributes of E and simple component attributes of composite attributes of E
 - choose one of the key attributes of E as the primary key

Step 2 – Weak Entity Types

- For each weak entity type W with owner entity type E
 - create relation R including all simple attributes and simple component attributes of W
 - include primary key of E 's corresponding relation as a foreign key in W
 - primary key of R is combination of foreign key from E and W 's partial key

Step 3 –Binary 1:1 Relationship Types

- For each binary 1:1 relationship type R
 - identify relations corresponding to participating entities, say T and S
 - if there is total participation, make that S
 - include primary key of T as a foreign key in S
 - include all simple attributes and simple components of composites of R as attributes of S

Step 4

Regular 1:N Binary Relationship Types

- For each regular 1:N binary relationship type R
 - identify relations, say S and T, corresponding to participating entity types
 - if S is at “N side” of relationship type then include primary key of T in S
 - include any simple attributes or simple components of R as attributes of S

Step 5 – Binary M:N Relationship Types

- For each binary N:M relationship type K
 - create new relation R to represent K
 - if S and T are the relations corresponding to the participating entity types then include their primary keys as foreign keys in R
 - combination of foreign keys will be primary key of R
 - include any simple attributes or simple components of K as attributes of R

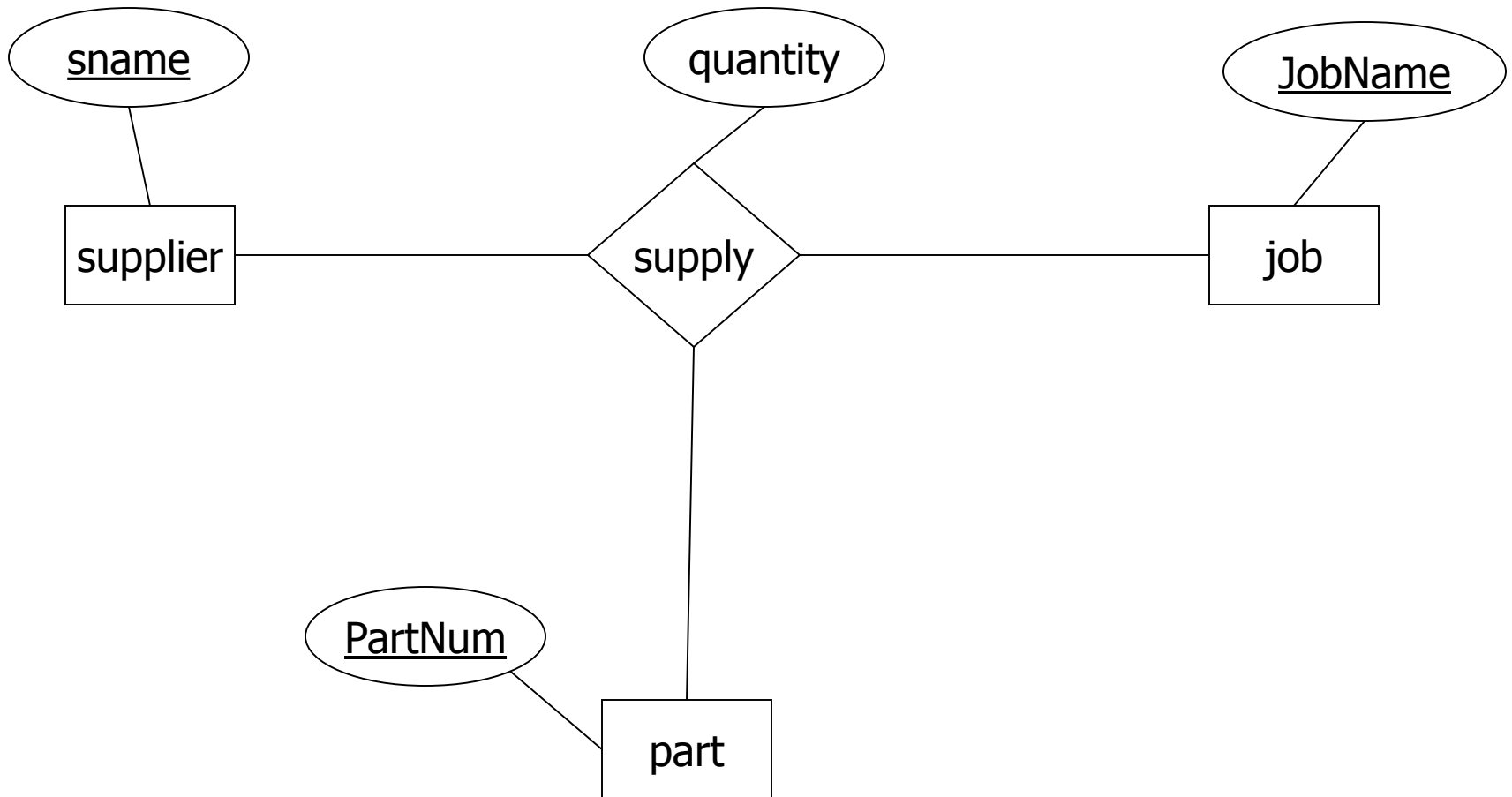
Step 6 – Multi-valued attributes

- For each multi-valued attribute A
 - create a relation R that contains the attribute A and the primary key of the relation S that corresponds to the entity type that contained A

Step 7: n-ary Relationship types

- For each n-ary relationship type L ($n > 2$)
 - create a new relation R to represent L
 - include as foreign keys in R the primary keys of all relations that correspond to the entity types participating in L
 - the primary key of R is the combination of all these foreign keys
 - include any simple attributes or simple components of L as attributes of R

N-ary example



Solution to n-ary example

Tables that already exist:

Supplier (sname, ...)

Job (jname, ...)

Part (partNum, ...)

Add new Table:

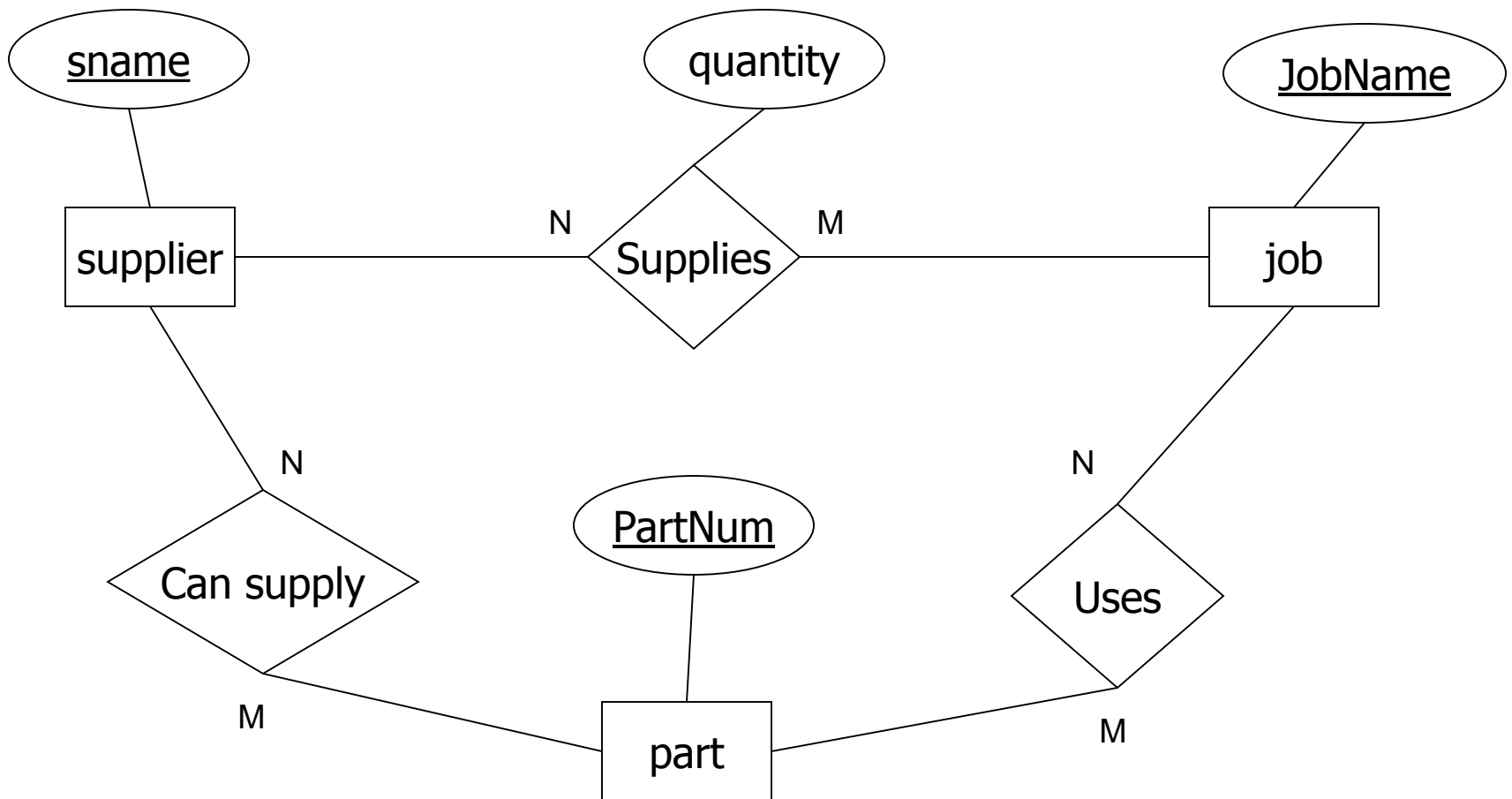
Supply (sname, jname, partNum, quantity, ...)

sname is FK -> Supplier (sname)

jname is FK -> Job (jname)

partNum is FK -> Part (partNum)

Binary version of n-ary example



Solution to binary version of n-ary

Supplier (sname, ...)

Job (jname, ...)

Part (partNum, ...)

Supplies (sname, jname)

Can_supply (sname, partNum)

Uses (partNum, jname)

Compare the two solutions...

N-ary Solution:

Supply (sname, jname, partNum, quantity, ...)

Binary Solution:

Supplies (sname, jname) (supplier can supply to a job)

Can_supply (sname, partNum) (supplier can supply a part)

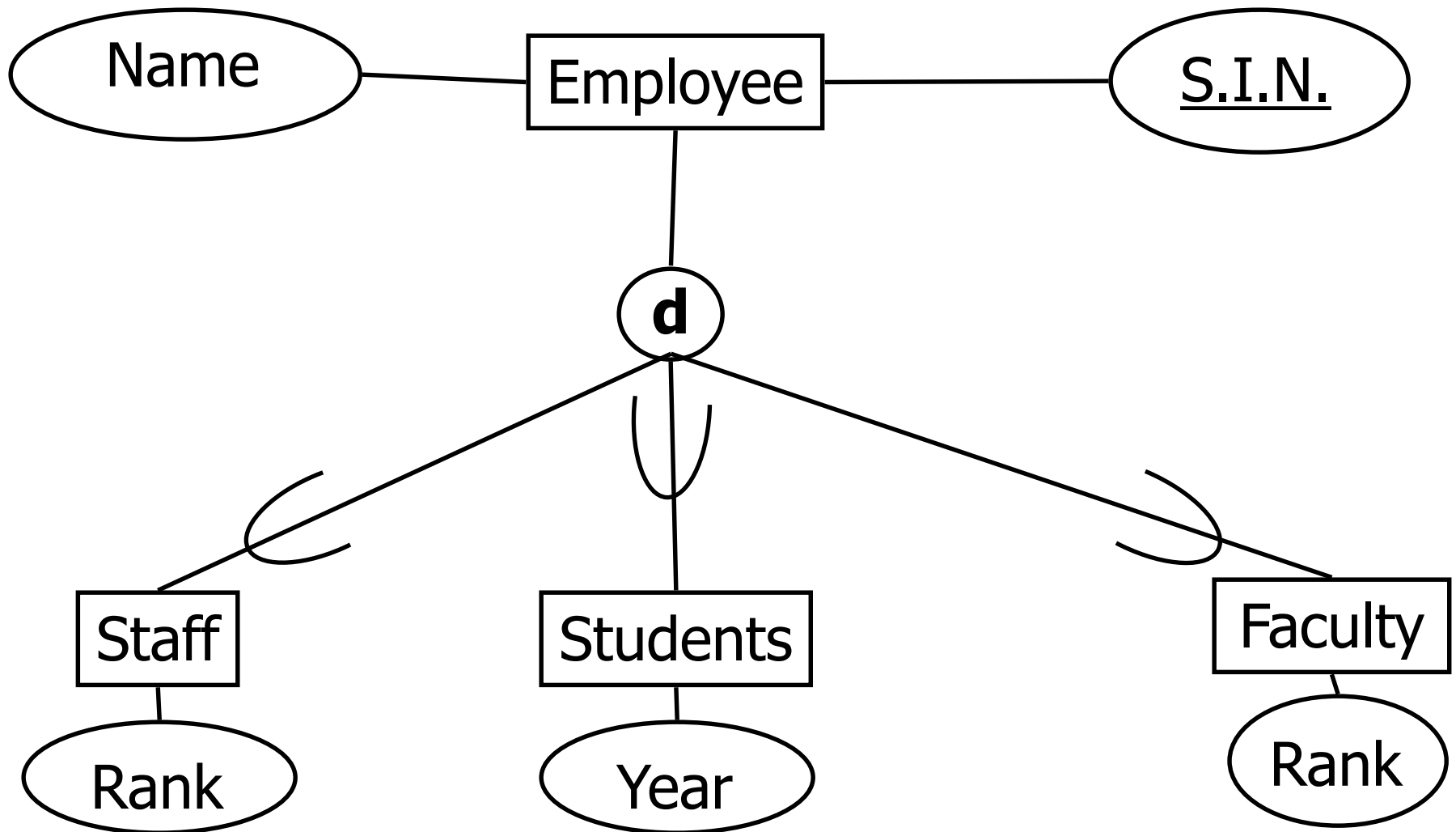
Uses (partNum, jname) (job uses a part)

(s,j), (s,p), (p,j) != (s, p, j)

Step 8 - Specializations

- For each specialization with subclasses $\{S_1, S_2, \dots, S_m\}$ and superclass **C** where attributes of **C** are $\{k, a_1, a_2, \dots, a_n\}$ and **k** is the primary key
- Convert using one of 4 options:

Example for 8a-c



Step 8 – Option 8a

- create relation **L** for **C** where
Attr(L) = {k, a₁, a₂, ..., a_n} and
PK(L) = k
- create relation **L_i** for each **S_i** where **Attr(L_i) = {k} U {attributes of S_i}** and **PK(L_i) = k**

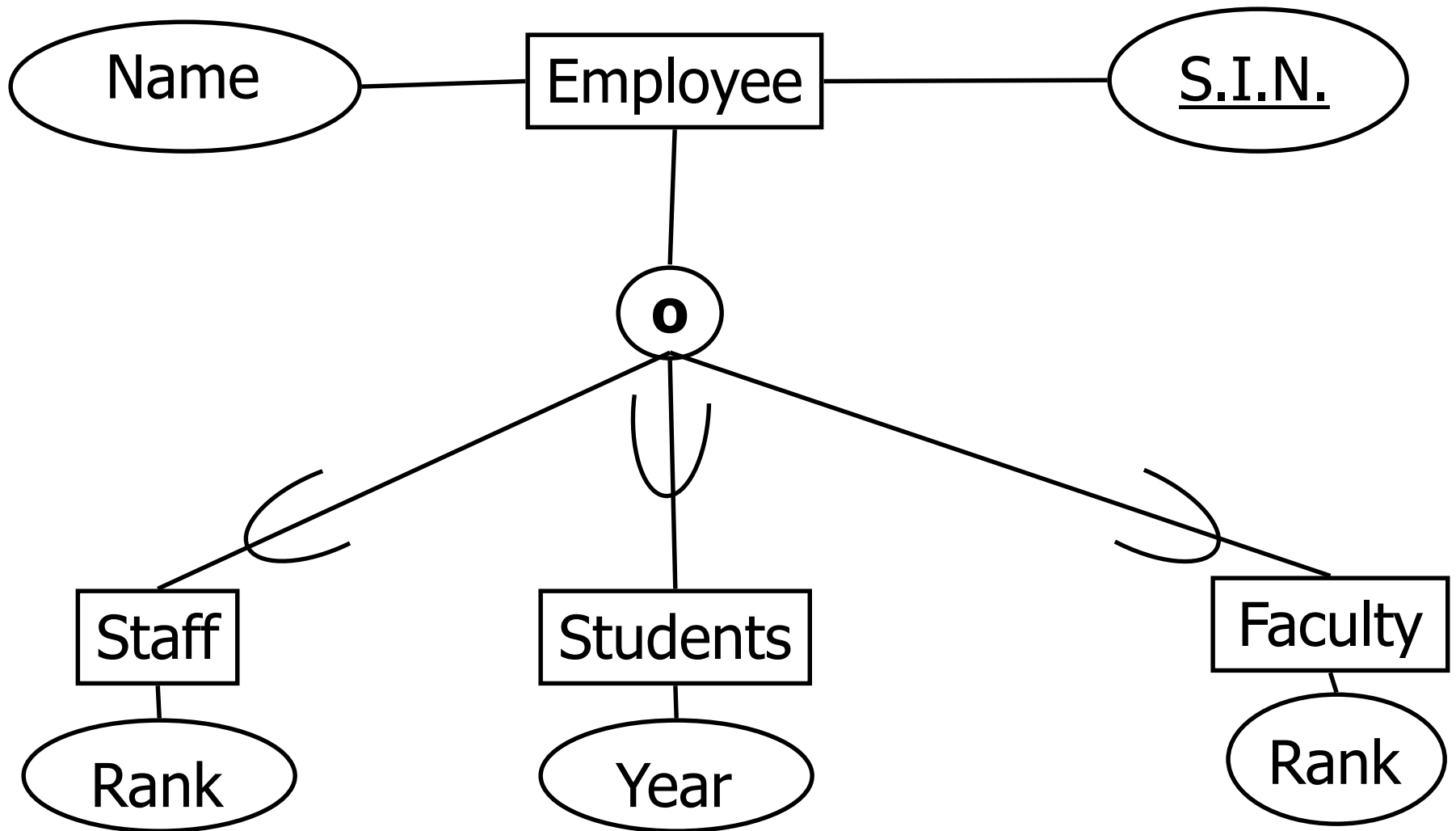
Step 8 – Option 8b

- create relation L_i for each S_i where $\text{Attr}(L_i) = \{k, a_1, a_2, \dots, a_n\} \cup \{\text{attributes of } S_i\}$ and $\text{PK}(L_i) = k$

Step 8 – Option 8c

- for cases where subclasses are disjoint
- create relation **L** where **Attrs(L) = {k, a₁, a₂, ..., a_n} ∪ {attributes of S₁} ∪ {attributes of S₂} ∪ ... ∪ {attributes of S_m} ∪ {t}** and
PK(L) = k, where **t** is a type (or discriminating) attribute

Example for 8d



Step 8 – Option 8d

- for cases where subclasses are overlapping
- create relation **L** where **Attrs(L) = {k, a₁, a₂, ..., a_n} ∪ {attributes of S₁} ∪ {attributes of S₂} ∪ ... ∪ {attributes of S_m} ∪ {t₁, t₂, ..., t_m}** and **PK(L) = k**, where t_i is a boolean attribute indicating whether a tuple belongs to S_i