

Markov chains and their computation

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February 3, 2021

Abstract

The following describes what a Markov chain is and how the state of a Markov Chain system is computed at any time step in such a way as to reduce the computational time by multiple orders of magnitude and extrapolate to infinity iterations with just one computation. The Gambler's Ruin and Monopoly games are presented as examples of use cases, and follow expected patterns of behaviour. The PageRank algorithm is also described.

1 Markov Chain computation

A Markov chain is “a stochastic model describing a sequence of events in which the probability of an event depends only on the state attained in the previous event” (see Wikipedia). This is explained with an example such as the one below.

Say that there are three possible states you can hold: A , B and C . If you are in A , the probability of moving from A to B is $M_{B,A}$, if you are in C the probability of moving to A is $M_{A,C}$. For any state i , the probability of staying in that same position for the next measurement is $M_{i,i}$. Etc. An example is given in Fig. 1.

The probabilities can be represented by a *stochastic* matrix $M_{i,j}$, where (A, B, C) maps to $(1, 2, 3)$. For Fig. 1, we have the following matrix:

$$\begin{pmatrix} 1/3 & 1/2 & 1/10 \\ 1/3 & 1/4 & 2/10 \\ 1/3 & 1/4 & 7/10 \end{pmatrix} \quad (1)$$

Notice how, for each column, the elements total one.

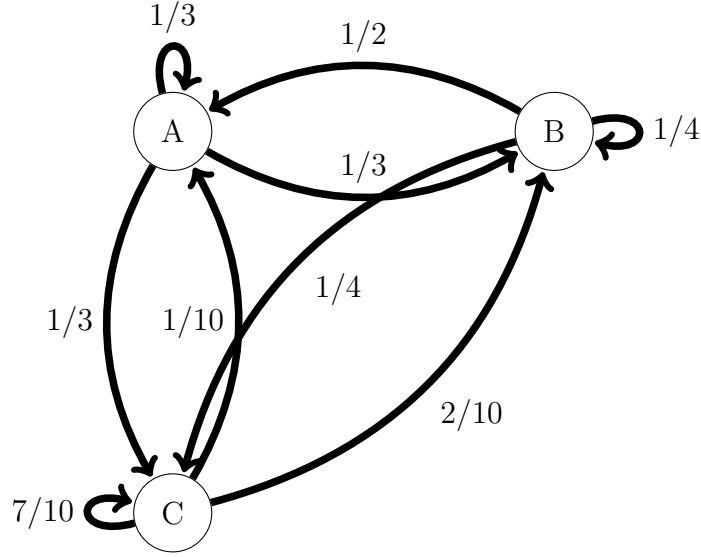


Figure 1: A display of a 3-state Markov-chain process.

We also need to define an initial state I , which will simply be the probability of being in each state at time step 0, for example $(1/3, 1/3, 1/3)$ or $(1, 0, 0)$. Again, this must sum to one. We label this P_0 .

The probability of being in any step at time t is

$$P_t = MP_{t-1} \quad (2)$$

In other words, by the definition that this is a Markov-chain process, the probability matrix is multiplied by the previous state in order to give the current state. Given P_0 and M , we can therefore calculate the probability of the state at any time t in one line:

$$P_t = M^t P_0 \quad (3)$$

Obviously this isn't ideal to do by hand, and even for a computer this could get expensive for a large number of states or large values of t . For this reason, we diagonalise the M matrix – or, in simpler terms, express P_0 in terms of the eigenvectors of M .

Define the eigenvalues of M to be m_i and the eigenvectors to be e_i . The initial state is written in terms of the eigenvectors to give:

$$\begin{aligned}
P_t &= M^t P_0 \\
&= \sum_{n=1}^i M^t c_n e_n \\
&= \sum_{n=1}^i c_n m_n^t e_n,
\end{aligned} \tag{4}$$

where c_i are scalar values, $P_0 = \sum_{n=1}^i c_n e_n$. This is far faster to compute.

Notes:

- It is given without proof that the eigenvectors of a stochastic matrix span the full space they inhabit.
- It is given, again without proof, that one of the eigenvalues is 1 and that the others are less than 1. As a result, clearly from the above equation, as $t \rightarrow \infty$, all of the terms in the sum vanish except for one. This is known as the *stationary state*.

2 Python code

The class for computing all of the above is given in MarkovClass.py. By running this script directly, the above example is computed.

```
python3 MarkovClass.py
```

```
Initial state: [1 0 0]
Probability matrix:
[[0.33333333 0.5      0.1      ]
 [0.33333333 0.25     0.2      ]
 [0.33333333 0.25     0.7      ]]
Constants: [ 0.49382716 -0.42910033  0.43247949]
Eigenvalues: [ 1.          0.39013419 -0.10680086]
Eigenvectors: [[ 0.525      0.5      1.          ]
 [-0.7183897 -0.2816103  1.          ]
 [ 1.          -0.85033548 -0.14966452]]
Equation with time:
P(t) = 0.4938271604938267*(1.0000000000000009^t)*[0.525 0.5  1.  ]
```

```

+-0.42910032543113297*(0.39013419193125953^t)*[-0.7183897 -0.2816103 1.          ]
+0.43247948804680914*(-0.10680085859792622^t)*[ 1.          -0.85033548 -0.14966452]
Stationary state: [0.25925926 0.24691358 0.49382716]

```

There are clearly issues with floating point errors, but these come directly from the numpy linear algebra computation and unfortunately cannot be helped.

3 Example usages

Python scripts are available for both of the below examples, as well as gifs showing the state over time.

3.1 Gambler’s ruin

The Gambler’s ruin problem is well known. Given a starting value, x , a value at which the gambler would “cash out”, y , and a probability of increasing by 1 and decreasing by 1, what is the probability that the gambler would eventually cash out?

This can be solved using an iterative derivation. For example, for a fair coin the probability of ending penniless is:

$$P = \frac{x}{y}. \quad (5)$$

Using Markov chains, not only can this be calculated computationally, but you can also calculate the probability of the gambler having a certain number of coins at any time step. See `GamblersRuin.py` for more details.

The “`GamblerRuin.gif`” file shows evolution of a state with time, where the gambler starts on 4 coins and cashes out at 10. The patterns observed are as expected - the 0th iteration is a simple peak, the 1st iteration is two peaks either side that of the 0th iteration, and evolution follows that pattern as the probability slowly concentrates itself at 0 and 10.

3.2 PageRank

One of the first algorithms used by the Google search engine was PageRank, which in its most simple form made use of Markov Chain analysis. Essentially, given a list of links, and the hyperlinks to the other pages contained

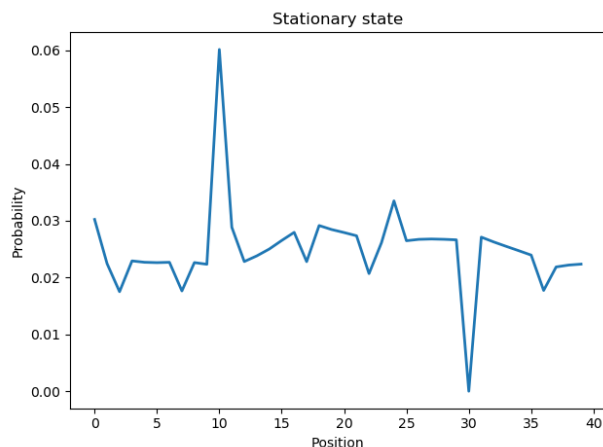


Figure 2: The stationary state of the Monopoly game as calculated using Markov-Chain analysis.

within, how do you rank them? This algorithm builds a Markov Chain system where the probability of moving to another page is split equally into the hyperlinks, and the stationary state determines the ranking used for the search page. The function to compute this is given in the PageRankExample.py script.

3.3 Monopoly

The monopoly game can also be simulated using Markov chain analysis. In the Monopoly.py script, a simple setup is created with dice rolls, chance and community chest “advance to” cards and “go to jail for three doubles in a row” accounted for. The stationary state gives the optimal position for properties. As has been observed elsewhere, the jail square is the most likely, followed by those around 7 squares afterward. The stationary state of the monopoly board is visualised in Fig. 2.

There are notable peaks and troughs. Peaks are at the positions of the “advance to” locations, and troughs are at the chance and community chest locations themselves. The biggest peak, however, is for the jail location, and a trough at the “go to jail” square as this is treated as a square that you cannot land on – it sends you straight to the jail cell. The shape of the distribution follows a wave, with upward parts following the peaks, particularly after the

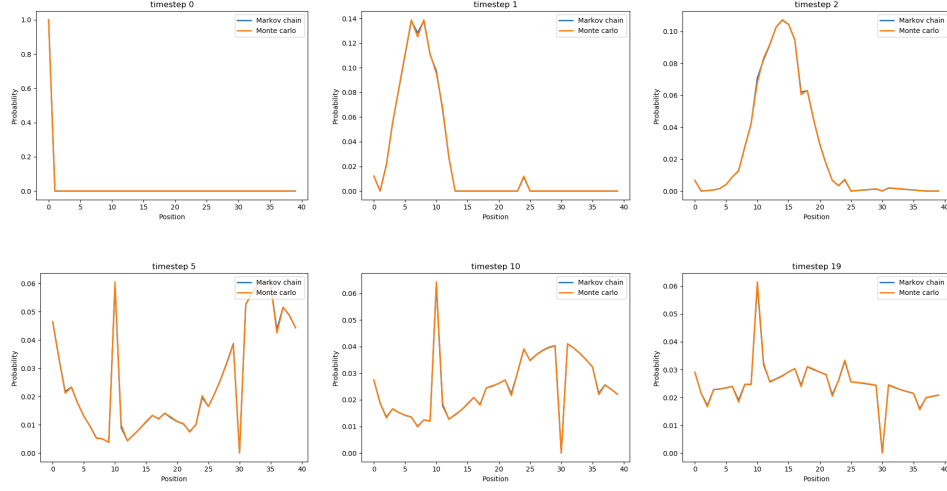


Figure 3: A comparison of the Monte Carlo and Markov Chain simulation techniques for the Monopoly game at a range of timesteps.

jail location.

To validate this evaluation of Monopoly, an additional more rigorous Monte-Carlo simulation is also written in the `Monopoly.py` script. This simulates the game in full for a number of players, and calculates the probability at any time based on the proportion of players occupying each state. In Fig. 3, a comparison is made between the two simulations – there is good agreement between the two. The `time_monopoly.py` script compares the processing time and CPU usage of each simulation. It was found that for 40 timesteps and 100,000 Monte-Carlo players, the Monte-Carlo simulation took 100 times as long and used 10,000 times as much memory in RAM.