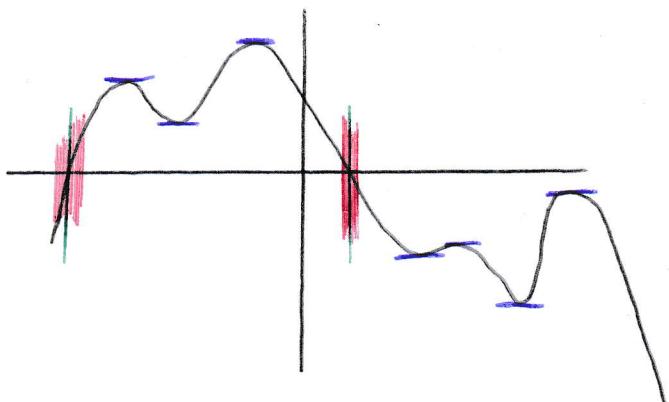


01/11/25

LRCS - VI

07/11/25

Locatis Roots of a polynomial
of any degree VIA change in
sign between the function's
Stationary points



† VAP. 4) 30% + 70%

† 100% 70% ΔΔ

Gavin Simpson - @gdsimpson3

① Input: $7n^5 + 9n^4 - 3n^2 + 4n + 2$ | input rules
 - - - - |
 split through ("+" or "-") | → Ascending powers
 $\rightarrow ["7n^5", \dots] \rightarrow [[A, Exponent]]$ | → missing powers allowed
 $\rightarrow A \text{ and } (\text{Exponent})$

$7n^5 \rightarrow 7, n^5 \rightarrow [7, 5]$ |
 If no $n \rightarrow [2, 0]$ |
 If no $A \rightarrow [7, 1]$ |
 $\rightarrow [[7, 5], [9, 4], [-3, 2], [4, 1], [2, 0]]$

missing powers → won't be needed

② (Creating $f(n)$)

$$f(n) = (7 * (n^{**} 5)) + (9 * (n^{**} 4)) \dots$$

Loop & sum to variable

③ Find min & maxes to check if there are roots

why?: if func has no roots, program will crash, find min & max & see if it changes sign via locatin's roots

③.1 ~~$\frac{d^n g}{d n^n} \rightarrow [7, 5] \rightarrow [7 * (5^{**}), 5 - 1]$~~ | loop & create according array
 not needed & $\int \rightarrow [7, 5] \rightarrow [7 / (5 + 1), 5 + 1]$ | → if 0, remove

③.2 technically only need stationary points, nature doesn't really matter?

Scenario 1 → 1 stationary point, nature: min → condition for root
 it must be Θ

Scenario 2 → 8 stationary points → to know root or not,
 find 4 of any two & see if crosses via sign change

Locating Roots via change of sign

(2)

(3) Find min/max. to check if there are roots $^{2/2}$

Scenario 3 \rightarrow No stationary points \Rightarrow what means: straight line

(3.2.1)

Programmatically analyse if function has roots via stationary point count:

cases: 0, 1, n

0 \rightarrow straight line

1 \rightarrow check position & two trailing edges

n \rightarrow check all roots & pass if diss sign's found

(3.3)

finding stationary points

roots of $f'(x)$, recursion? !potential!

(3.1.2)

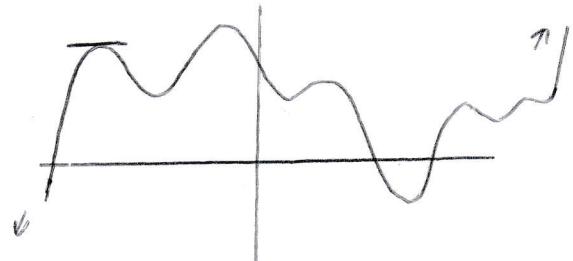
Derivative logic

$$[7, 5] \rightarrow [7*5, \cancel{5+1}]$$

$$[2, 0] \rightarrow [2*0, 0-1] \rightarrow \text{must let this pass}$$

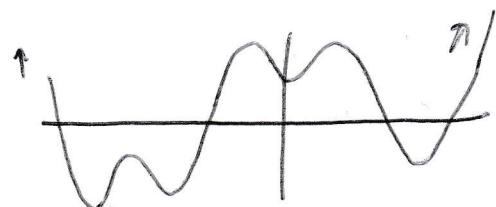
(4) finding Roots

1 \rightarrow find earliest stationary point and keep going backwards till theres a change in sign



2 \rightarrow then switch direction & go in smaller till next change

3 \rightarrow carry on till 0 or certain n.o. of d.p.



(4.1)

for n stationary points \rightarrow option 1 \rightarrow go left to right & look for sign changes

start with first stationary point

and go back, conduct Rootloc.

then go forwards till sign change and Rootloc

Locating Roots via Change of sign

(3)

4.2) finding 2 stationary points with root between

- Loop through all points & sign change, pair

- ① & ② \rightarrow no sign change

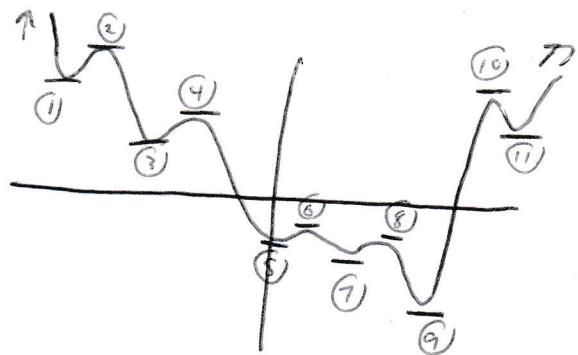
- ① & ③ \rightarrow no

- ① & ④ \rightarrow no

- ① & ⑤ \rightarrow yes

- ① & ② \rightarrow no - ③ & ④ \rightarrow no

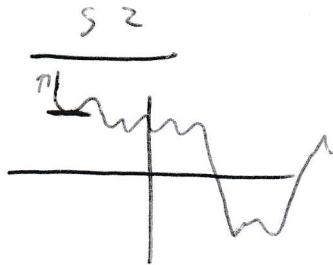
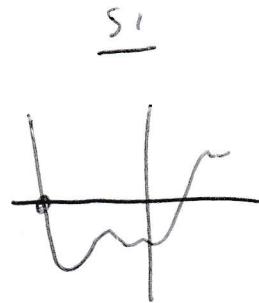
- ② & ③ \rightarrow no - ④ & ⑤ \rightarrow yes



4.3) first root

- set nature of point.

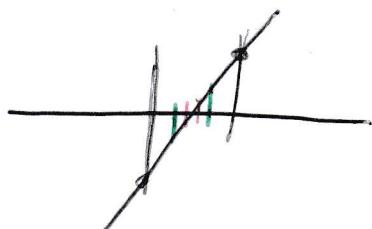
is min & above n axis,
then, search ~~to~~ right



is max & below n axis,
look right
similarly go left else

4) Root location

go from A to B in big intervals,
once sign changes, go back
in smaller intervals



Repeat till u find Root to contain D.P

④ Root location - programmatically

111

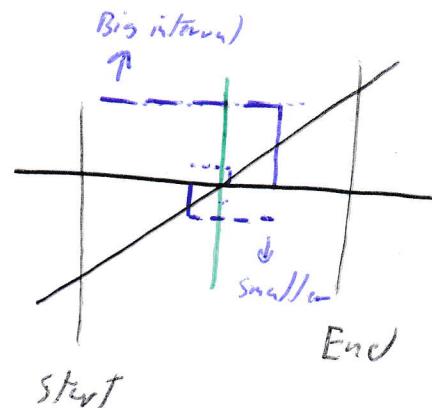
true = positive

False = negative

Interval = $2 * \text{(-iteration count)}$

Direction = ~~$+1 * (-1) * \text{(iteration count)}$~~

↳ swap directions every iteration
used in for loop range stops
(direction * interval)



concept : start \rightarrow sign change
New start
Sign change \leftarrow other direction
start \rightarrow change....

↓ interval narrows
& directions swaps
+ 1 iteration count

Pseudocode

Start = ...

End =

StartSign = true

EndSign = false

if (RootSign != startSign) or (RootSign == EndSign)
either start works

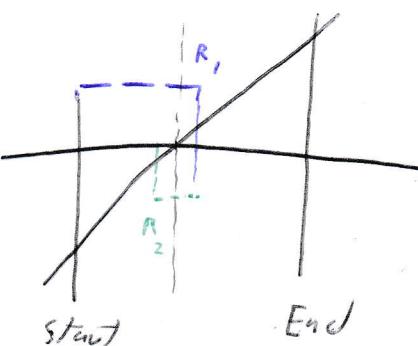
Root is new start

start = Root

iteration count += 1 \rightarrow changes interval & direction

next loop must go back from Root to start

loop after must go from next to root



Start \rightarrow end
~~R₁ \rightarrow start~~
R₂ \rightarrow R₁

Swap start & end

R₂ is new start

④ Root location - Pseudo code

Start = ...

End = ...

Start Sign = ...

End Sign = ...

Interval = $2 * * (-1 * \text{IterationCount})$ Direction = $(-1) * * (\text{IterationCount})$

IterationCount = 0

D.P Accuracy = ... # when to stop

while D.P Accuracy not met:

[for Root in Range (start, end, interval * direction):

if (Root Sign != Start Sign):

End = Start

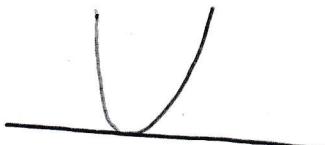
Start = Root

IterationCount = +1

steps when $g(n) = 0$

← or

Limitation: this mechanism would work in most cases except Roots that touch



workaround: A touching Root is a stationary Point... Hence all stationary points could just be checked

Python limitation: Range(), increment steps cannot be type int

Custom loop \Rightarrow

Condition = ...

if direction > 0 :

Condition = ...

else ...

Root = Start

while Root < end :

Root += interval

doesn't show condition

Change condition

C1 = Start < End
C2 = Start > end

Locating Roots VIA change in sign

⑥

④ custom loop

flow : ① Start \rightarrow end : Pass Root @ RI
② Start \leftarrow RI = RI \rightarrow start
③ R2 \rightarrow RI

conditions : ① Root = start
while Root < End Root ++

*) Ver Reassig:

① Root = start
End = End

② Root = RI
End = start
direction = \rightarrow

③ Root = R2
End = RI

② Root = RI
~~start = End~~ End = start
while Root > End Root --

③ Root = R2
End = RI
while Root < End Root ++

we can have toggle for directions, true = \rightarrow , false = \leftarrow

while True

start = start Param

End = End Param

interval = 2 * * (-1 * iterationCount)

direction = (-1) * * (iterationCount)

iterationCount = 0

① start = start
Root = start
End = End
direction = \rightarrow
② Root = RI
End = start
start = Root
direction = \leftarrow
③ Root = R2
End = start (RI)
start = Root

Locating Roots VIA change in sign

②

$f(u)$



$$\frac{dy}{du}$$



Stationary points



Check if any point is a root



Look for sign changes between Roots & create Pairs



Find ~~convergence~~^{root} between pairs via convergence



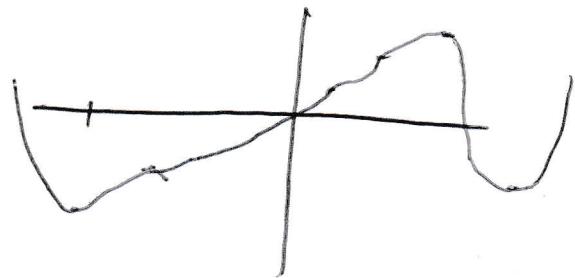
Then look at Edge roots

Finding stationary point pairs~~[9, 8, 4]~~

$$[9, 7, 5, 3, -2, -4]$$

$\downarrow s(7) = +$ $\downarrow s(3) = +$ $\downarrow s(-4) = -$
 $s(9) = -$ $s(5) = +$ $s(-2) = -$

Best to take from left



$$[-4, -2, 3, 5, 7, 9]$$

$\ominus \ominus \oplus \oplus \oplus \ominus$

For Point in Range ()

If not $f_N(\text{point})$ sign == $f_N(\text{Next point})$ sign:
Root between them

make sure to exclude stationary roots from pair finding

Finding stationary points

Scenario 1: 4 des polynomial

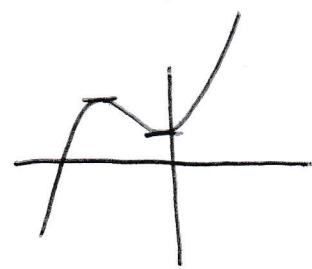
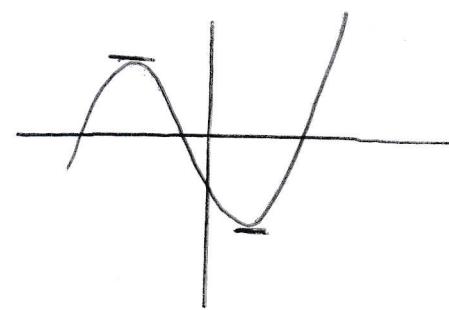
$$f(n) = n^4 \dots$$

$$f'(n) = n^3$$

Root Finder ($f'(n)$)↳ $f''(n)$ Root Finder ($f''(n)$)

↳ ...

cubic



Theoretically if Root Finder returns true Roots, we don't need to worry about its long recursion

if $f(n)$ is 2nd degree, we can use the quadratic formula for the Roots which will give us the 2 stationary points of a cubic. This can be the end of recursion

Input: 3rd degree

Start point will jump @ des 2 & change it to quadratic

Locating Roots VIA change in sign

10

Root finder function

Param: $f(u)$

Derivative $\Rightarrow f'(u)$

stationary point finder $\Rightarrow [[\alpha, \beta], [y, z]]$

\downarrow
Pairs

2 edge Roots

Signchange func $(\alpha, \beta) \rightarrow \text{Roots}[]$

edge Root (α, β) $\rightarrow \text{Roots}[]$

Return Roots

Edge Roots

Stationary Point $[0]$

Stationary Point $[-1]$

Pairs:

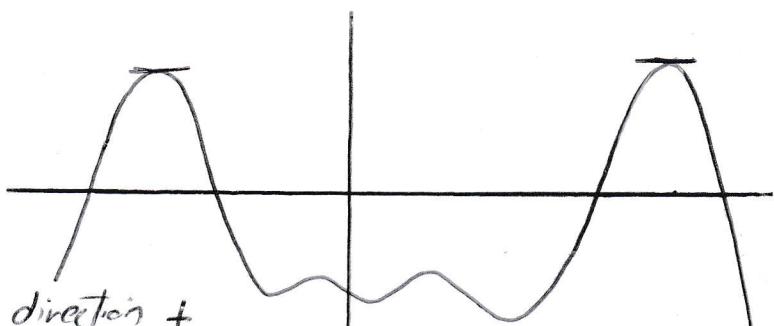
direction -



direction +



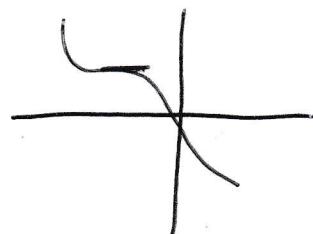
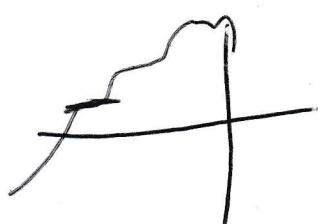
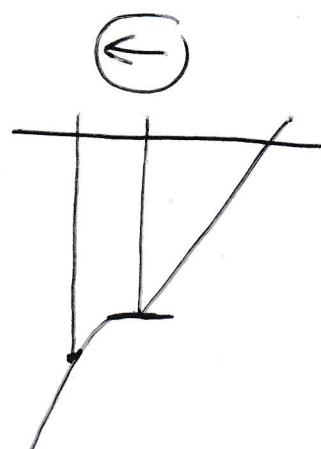
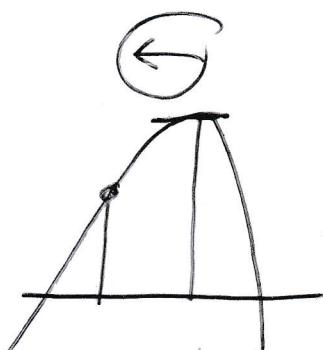
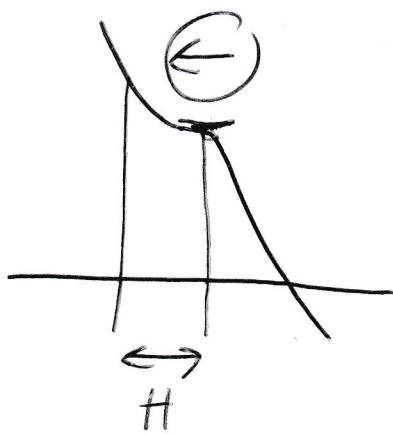
$[[SP[0], -1ns], [SP[-1], +1ns]]$



checks: if ~~SP[0]~~ max $f > 0$, \Rightarrow yes root

min $f < 0 \Rightarrow$ yes root

deals with inflection \rightarrow look either side



Locating Roots VIA change in sign

(12)

Edge Roots 2

Scenarios

(L) $u \geq m$



(R) $u \geq m$



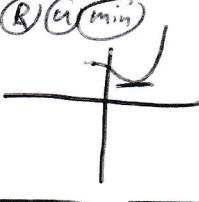
(L) $u \leq m$



(L) $d \leq m$



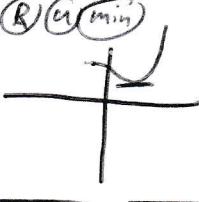
(R) $d \leq m$



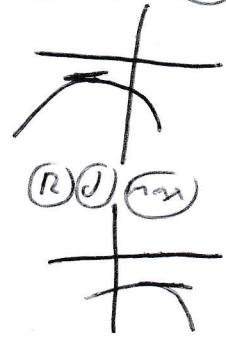
(L) $d \geq m$



(R) $d \geq m$



(L) $d \geq m$



with inflection



conditioning \Rightarrow (L) $u \geq m \Rightarrow$ look left

(L) $u \leq m \Rightarrow$ ignore

(L) $d \geq m \Rightarrow$ ignore

(L) $d \leq m \Rightarrow$ look left

(R) $u \geq m \Rightarrow$ look right

(R) $u \leq m \Rightarrow$ ignore

(R) $d \geq m \Rightarrow$ ignore

(R) $d \leq m \Rightarrow$ look right

with inflection points \Rightarrow we want know if its a max or min

(L) u ↗ look from point
↗ look left, is
its going down
then traces root

(R) ↗
↗ same, since its (d) is
its closer to x-axis

(R) u ↗

Classification of nature could take inflection & call max/min
dependent on direction its go from - add too in direction
& see if going up or down