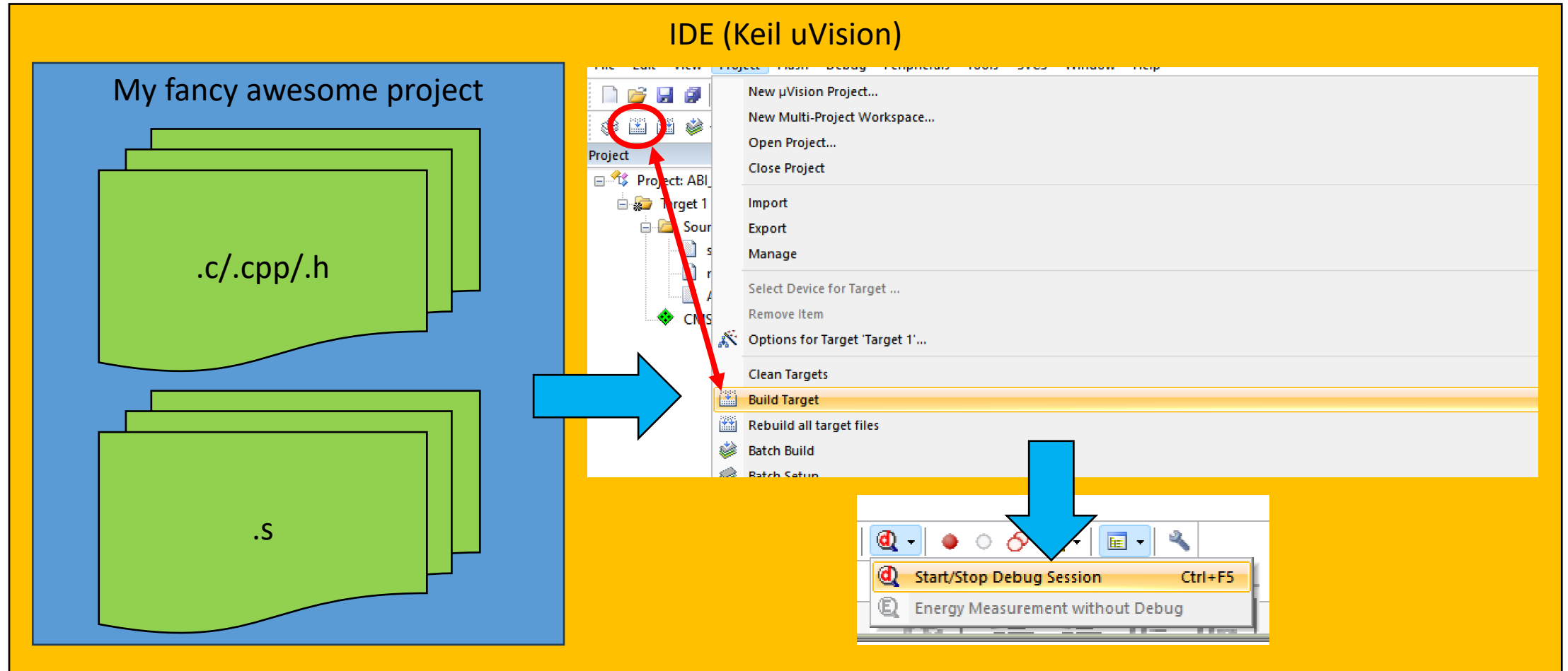


From Source code to Executable The Arm Toolchain for Embedded Systems

Francesco Angione, Paolo Bernardi

How is an executable produced from the source code?



Outline

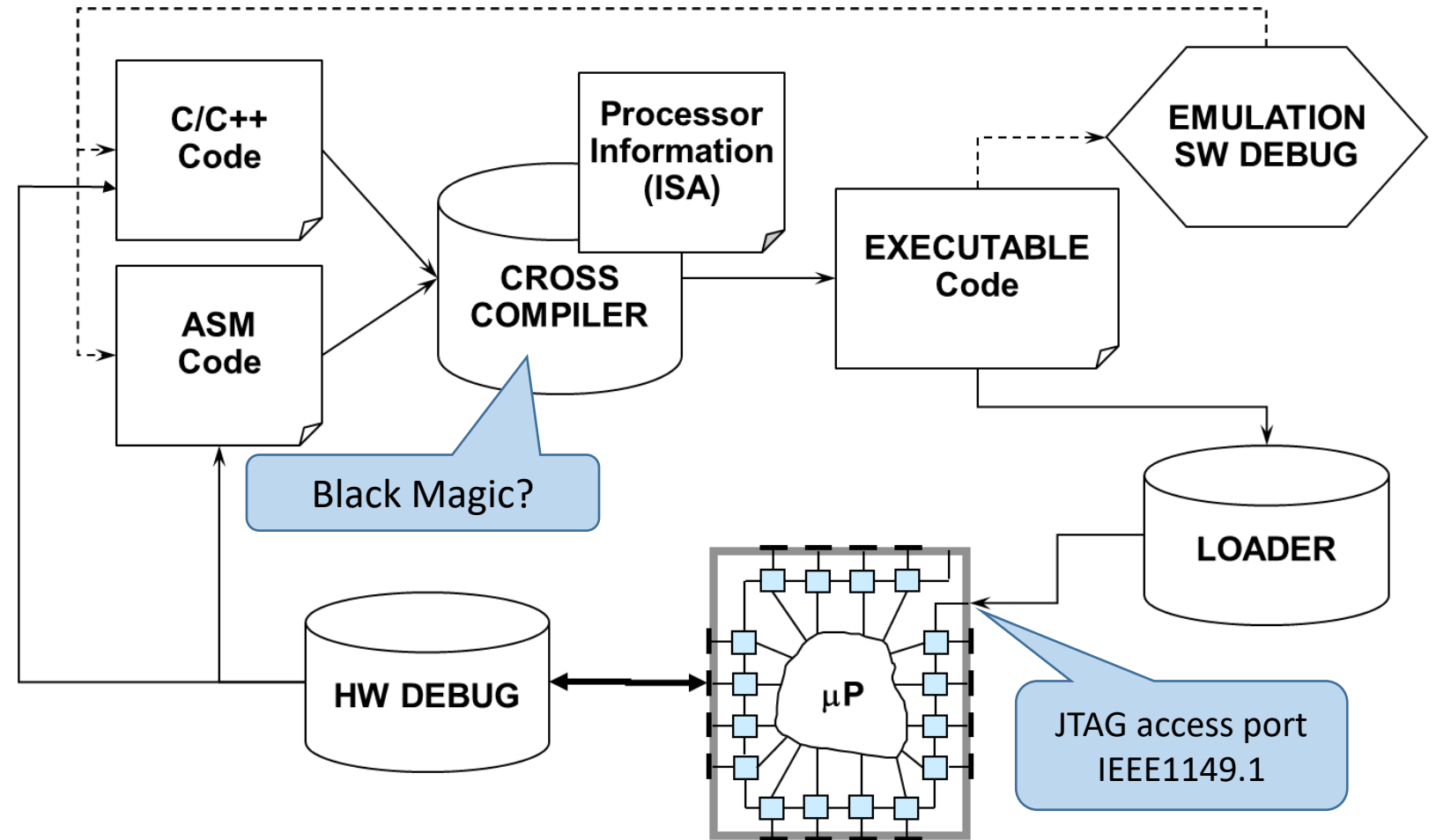
- What is a toolchain?
 - The Arm toolchain.
- Investigating the compilation output files.
- How does a System-on-Chip start the program?
- The Arm “Magic secret sauce”.
- HowTo and HowNotTo - Examples.

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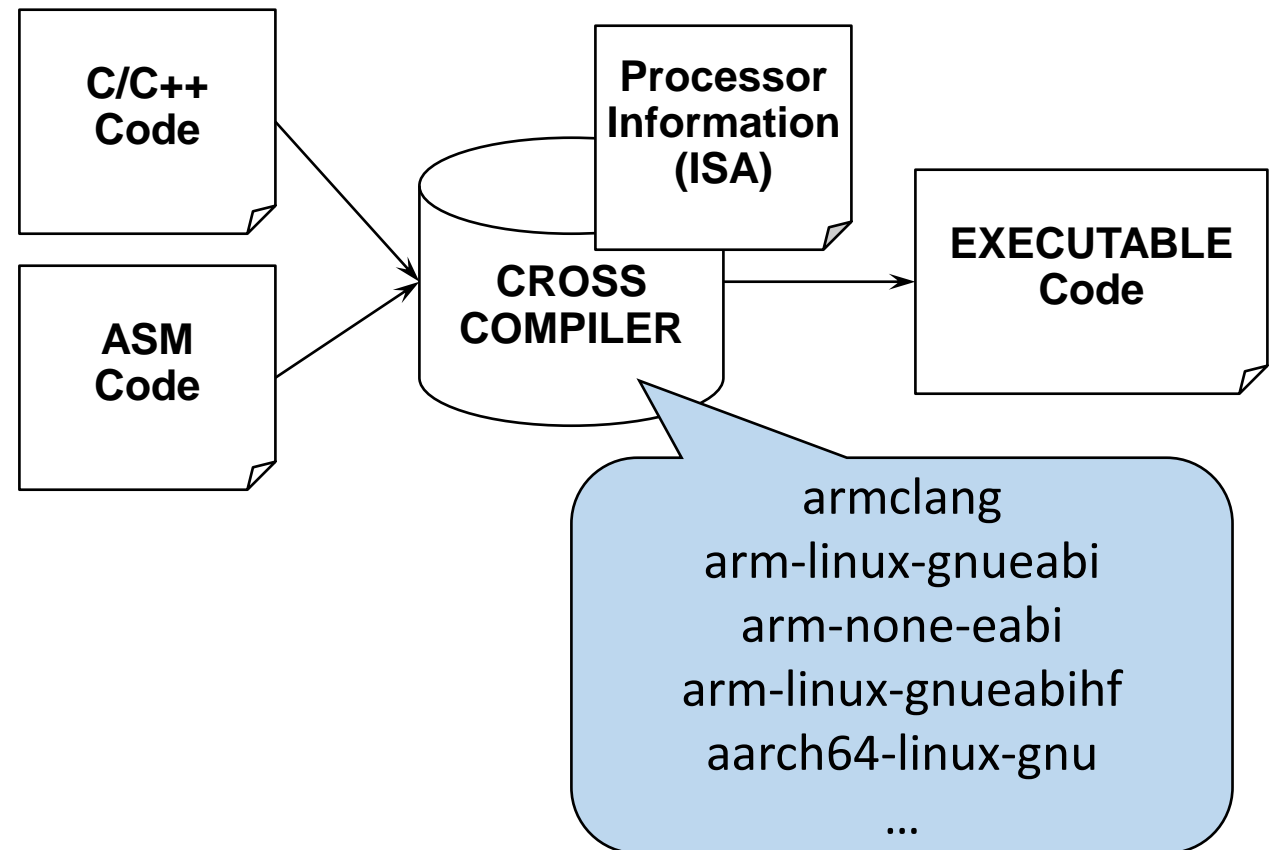
What is a toolchain?

- A set of **programming tools**.
- Used for complex development tasks or to create software products.



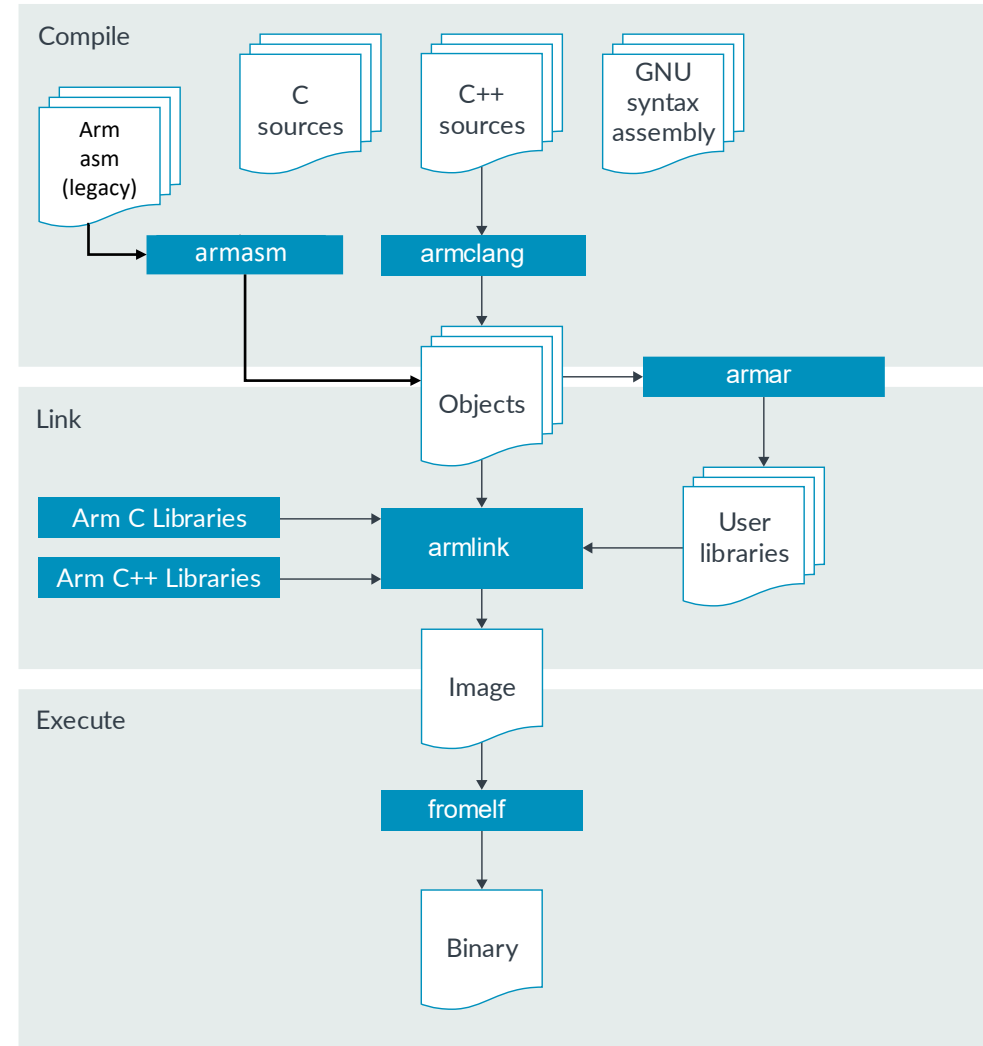
What does the cross-compiler?

- It is a compiler capable of creating executable code **for a platform other than the one on which the compiler is running.**
- It includes a set of programming tools (toolchain).
- Preprocess the source code.
- Translate high level code in machine code.
- Introduce already developed library.



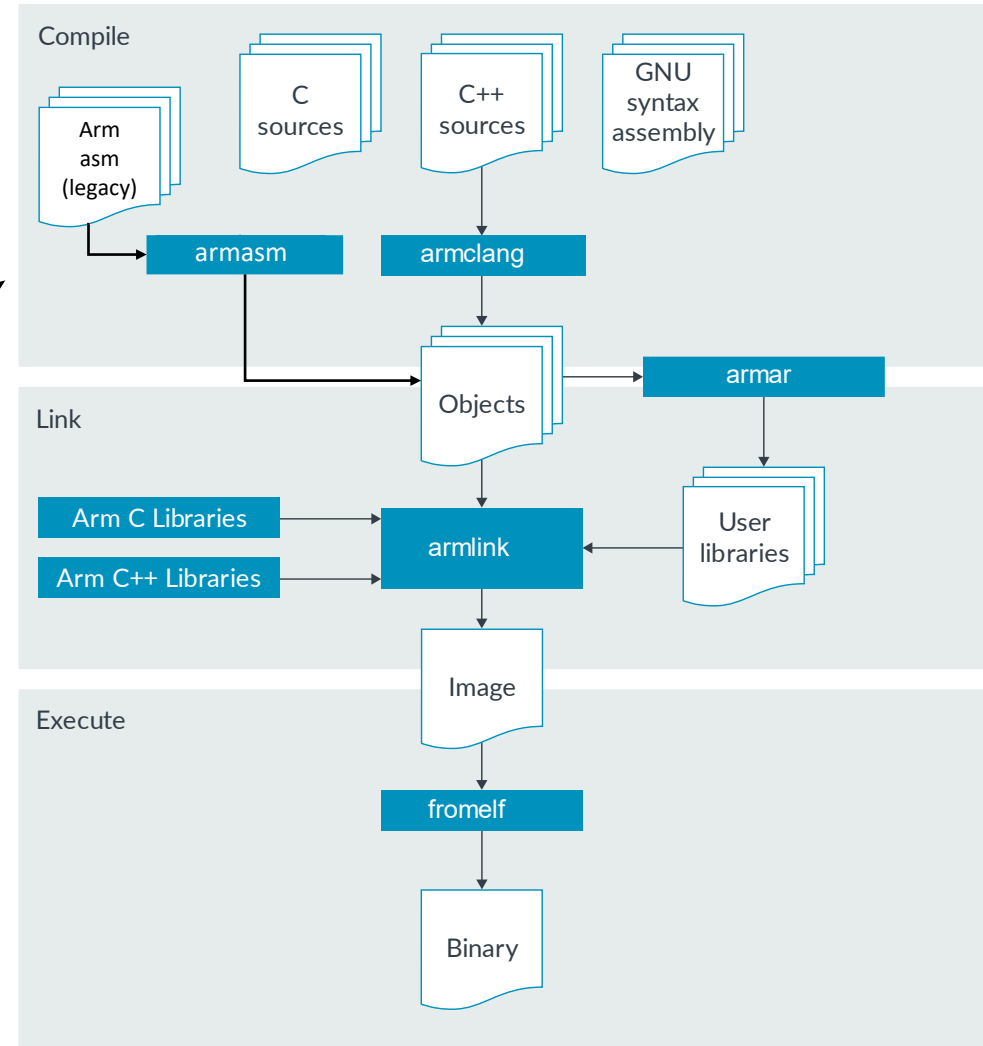
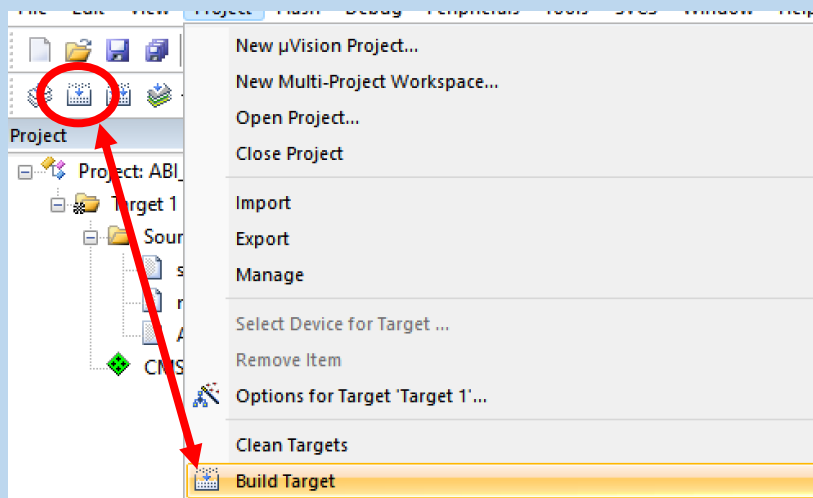
The Arm Toolchain

- Based on an enhanced version of *clang* (frontend) and *llvm project*, with proprietary customizations.
- The toolchain is composed of 3 different phases:
 - Preprocessing and compilation phase (armasm for legacy support).
 - Link phase.
 - Execute phase.



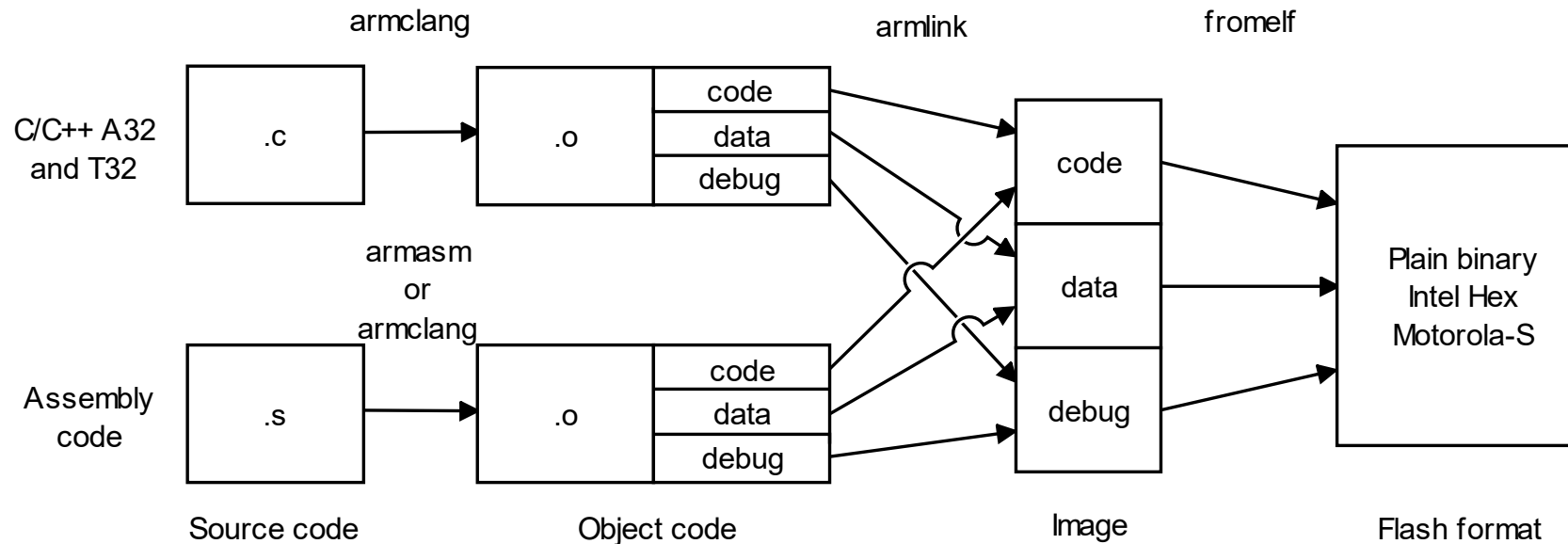
The Arm Toolchain

- Based on an enhanced version of *cl*
 - T
- All the three phases included in:

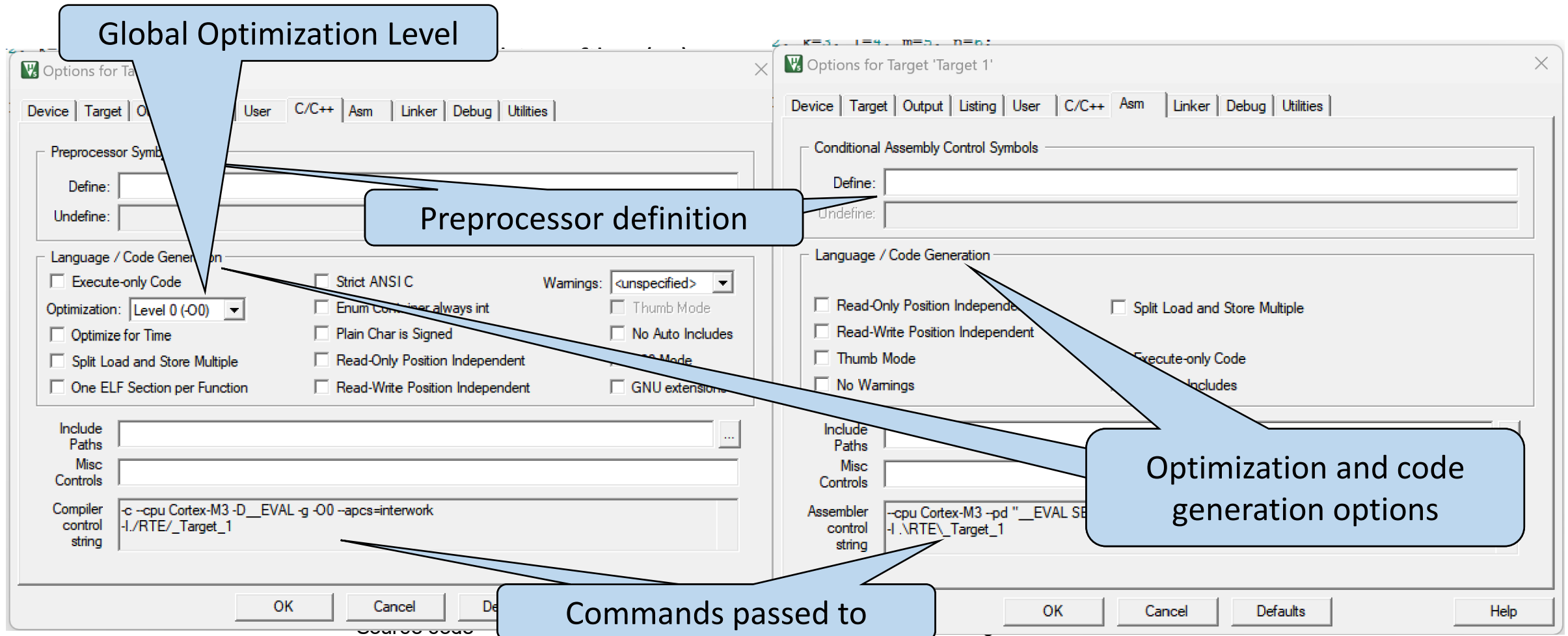


The compile phase – armclang and armasm

- The compile phase create object files (.o).
- Use armclang to compile high level code such as c or c++.
- Use armasm to assemble existing assembly code written in armasm syntax.
- Use armclang to assemble assembly language code, or inline assembly, written in GNU syntax.

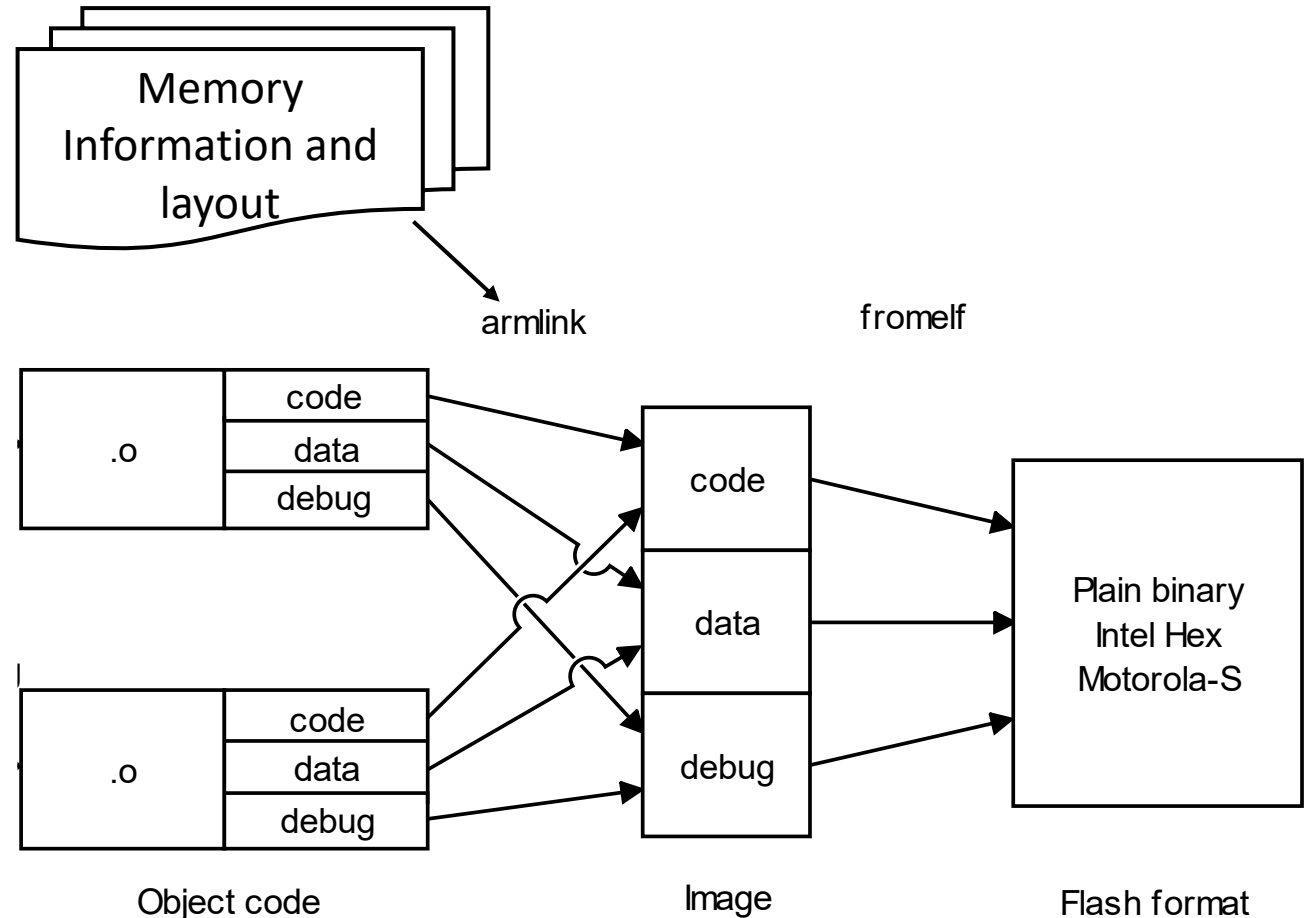


The compile phase – armclang and armasm



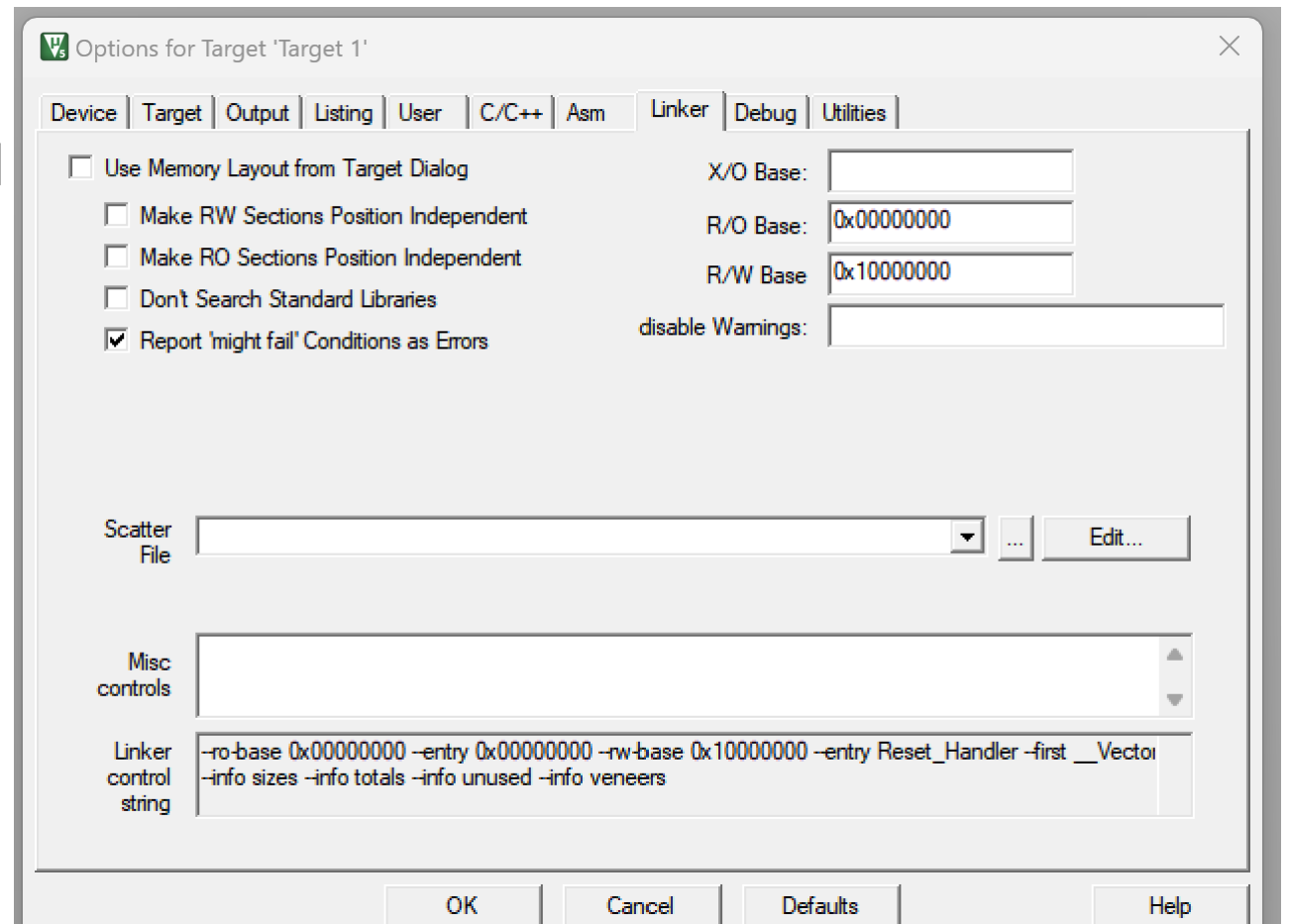
The link phase – armlink

- Link all object files into a single executable file (or another object file) by merging similar sections.
- It needs memory information to organize the image memory layout.
- It resolves:
 - Functions and variables (their symbols/label is substituted with an address).
 - Linker symbol (**different from functions and variables**).
- It eliminates unused sections **regardless of the optimization level**:
 - Removes unreachable code and data from the final image.



Memory Information and layout

- You can specify the entry point at the startup (i.e., the function called at the system boot).
- You can specify the memory information:
 - By command line using armlink tool.
 - By passing a scatter file.
- You can specify additional custom code and data sections.



Memory Information and layout

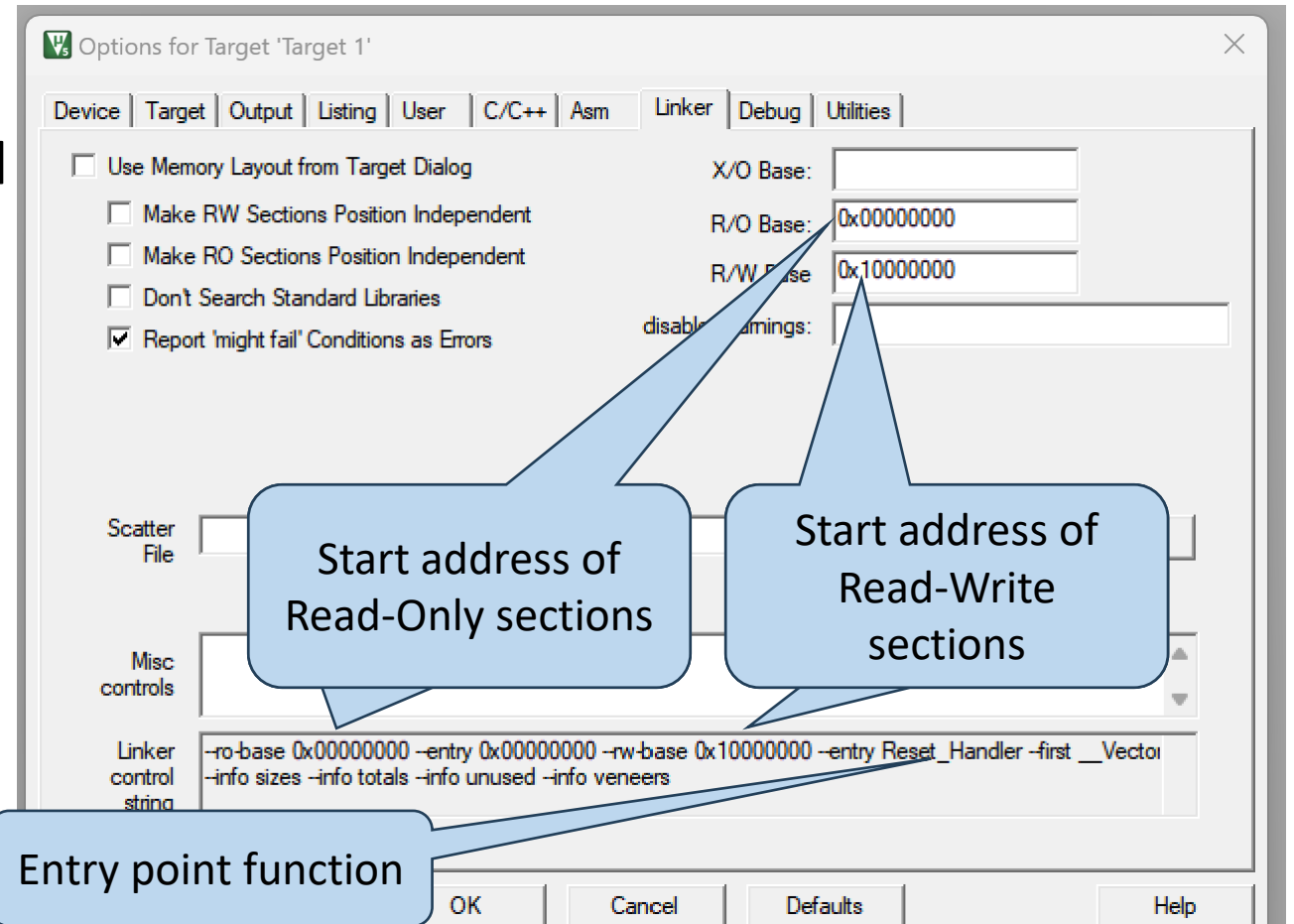
- You can specify the start address of the sections at the system level.

```
FLASH_LOAD 0x20000000
{
    RW 0x20000000 ; RW
    * (+RW-DATA)
}
```

Start address of
Zero Init sections

- By configuring the linker tool.
- By passing a scatter file.
- You can specify additional custom code and data sections.

```
ER_ZI 0x405000
{
    * (+ZI)
}
```



The execute phase - fromelf

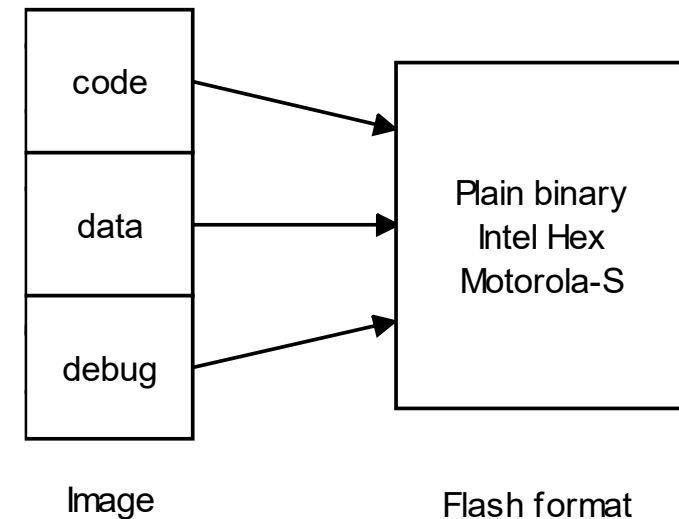
- Process object and image files.
- Convert ELF images into other formats for use by ROM tools or for direct loading into memory. The formats available are:
 - Plain binary.
 - Motorola 32-bit S-record.
 - Intel Hex-32.
 - Byte oriented hexadecimal.
- Display information about the input file, for example, disassembly output or symbol listings.

Binary

00000b60	88	00	00	00	10	02	00	00	08	00	00	00	0c	00	00	00	
00000b70	bc	00	00	00	18	02	00	00	08	00	00	00	0c	00	00	00	
00000b80	88	00	00	00	24	02	00	00	02	00	00	00	d8	00	00	00\$.....	
00000b90	03	00	04	05	00	00	00	04	01	43	54	4d	5f	66	75	6e	63ASM_func
00000ba0	74	2e	73	00	43	6f	6d	70	6f	6e	65	6e	74	3a	20	61	41t.s.Component: A
00000bb0	52	4d	20	43	6f	6d	70	69	6c	65	72	20	35	2e	30	36	36RM Compiler 5.06
00000bc0	20	75	70	64	71	74	65	20	36	20	28	62	75	69	6c	64	61update 6 (build
00000bd0	20	37	35	30	29	20	54	6f	6f	6c	3a	20	61	72	6d	61	61750) Tool: arma

Motorola

```
S31551807360000000000000000000000000000000000000000000000000000000046
S31551807370000000000000000000000000000000000000000000000000000000036
S31551807380000000000000000000000000000000000000000000000000000000026
S31551807390000000000000000000000000000000000000000000000000000000016
```



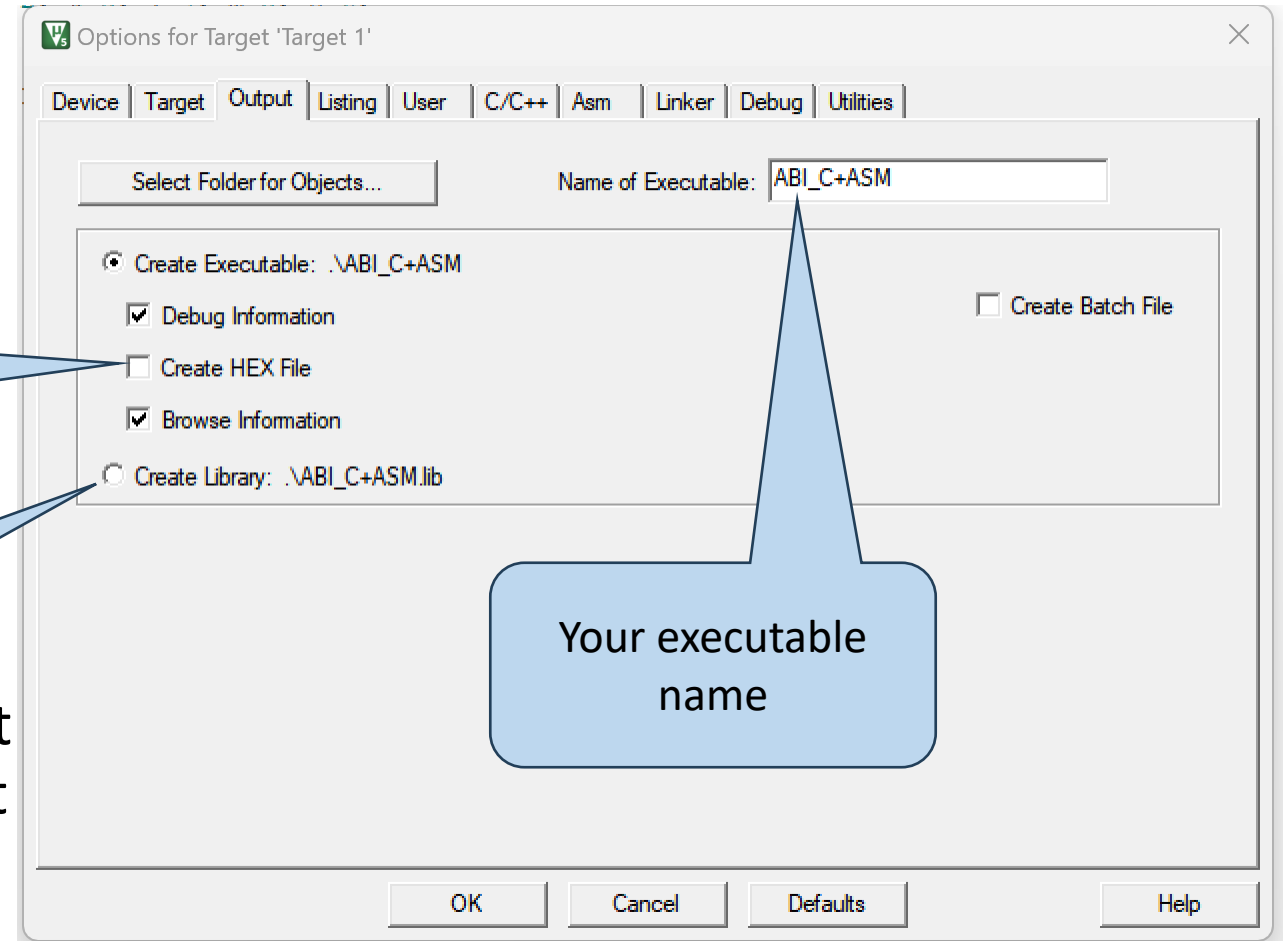
The execute phase - fromelf

- Process object and image files.
- Convert ELF images into other formats for use by other tools for direct loading in other formats available

If you want a hex file

- Plain binary.
- Motorola 32-bit S-record.
- Intel Hex-32.
- Byte oriented
- Display information in output file, for example, or symbol listing

If you want a library to be included into another project

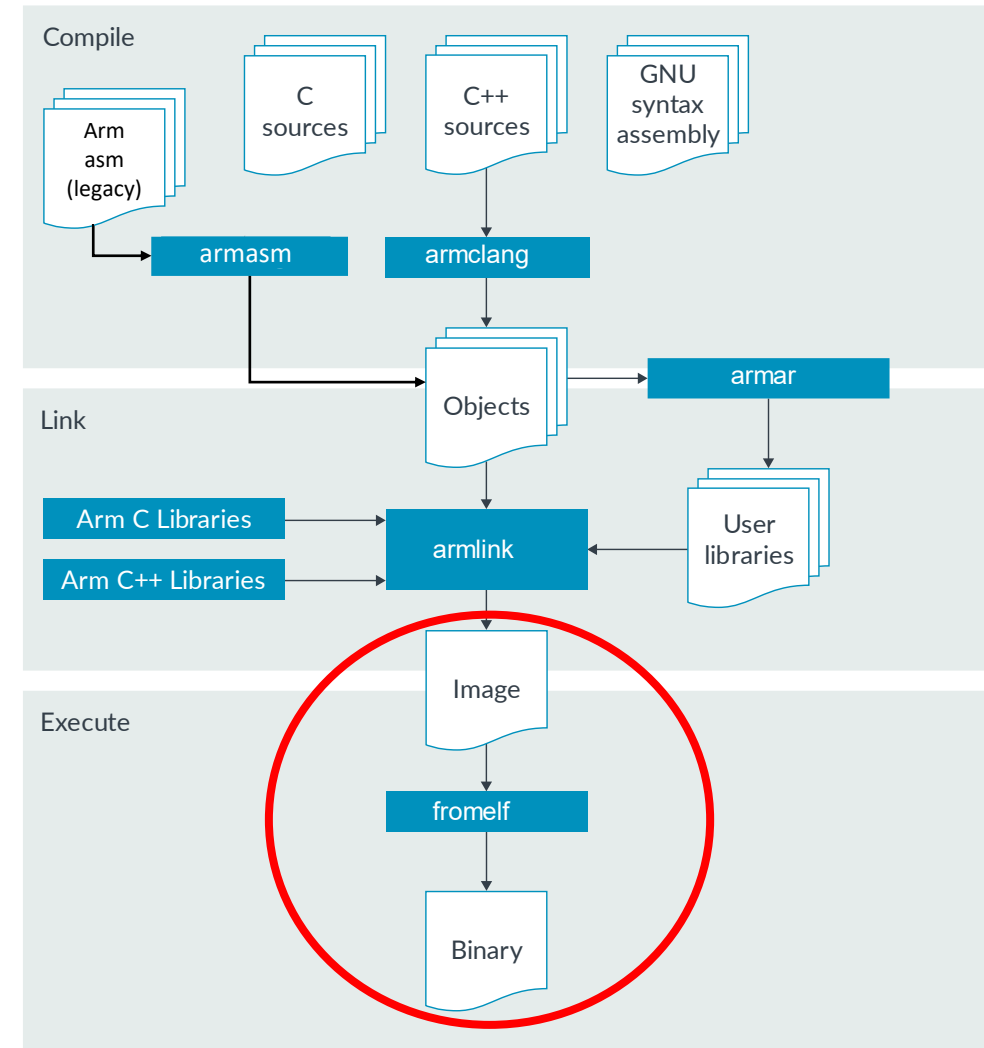


Outline

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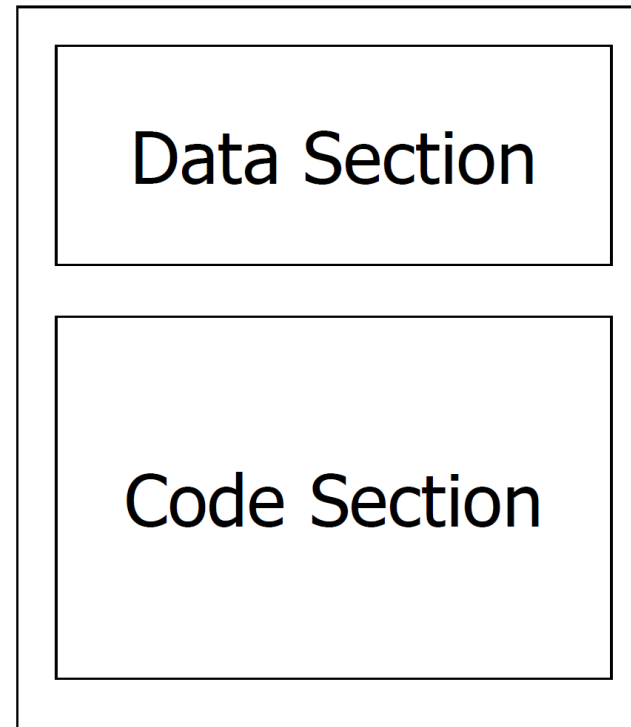
Investigating the compilation output files.

- The Arm toolchain produces:
 - The executable file.
 - The listing, dependencies files.
 - The map file.
 - The build log and static call graph file.



The executable

- The overall image (from the source code) is converted into an executable (.exe, .elf, .axf for Arm).
- Data and Code sections are in the executable.

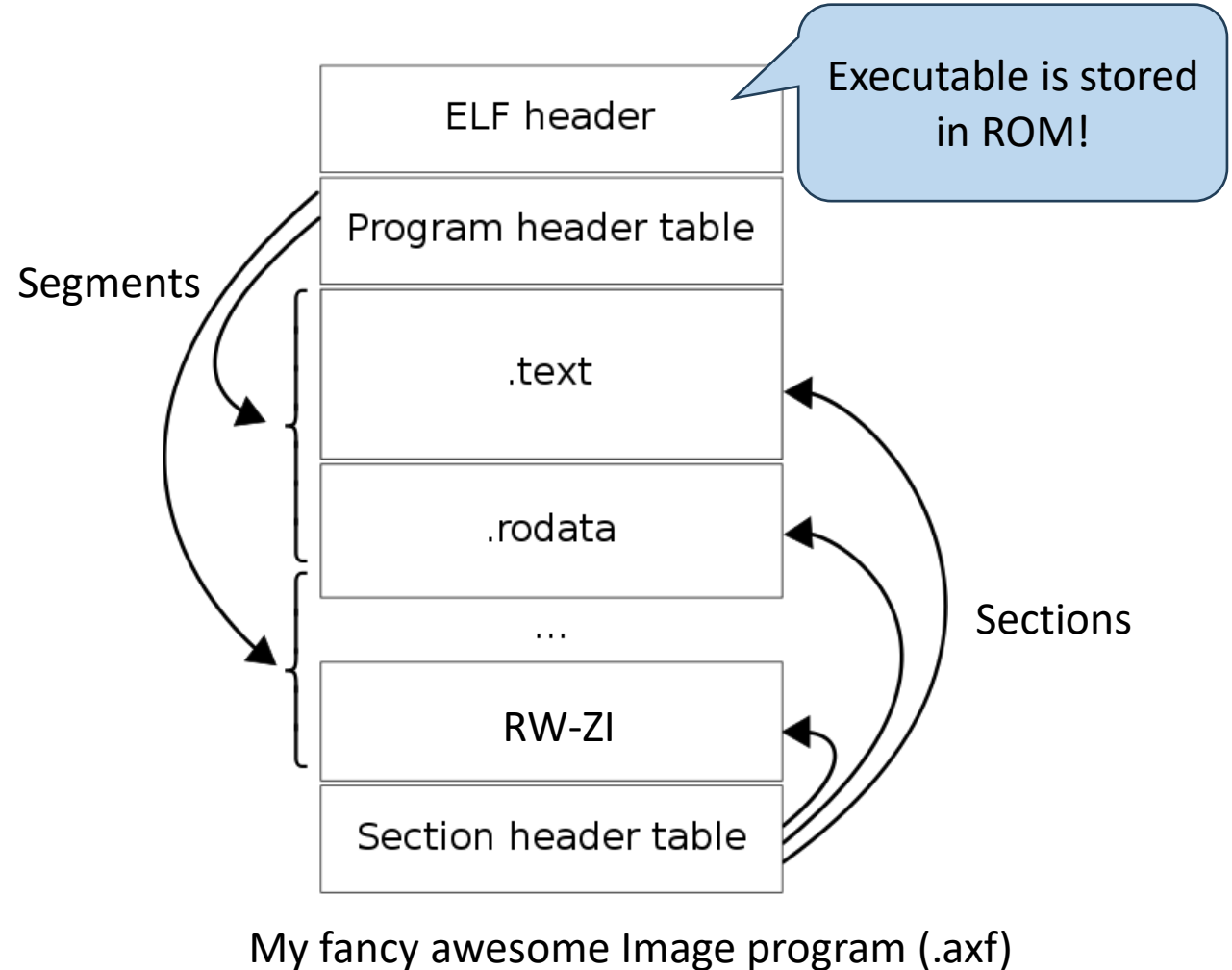


My fancy awesome Image program

- Data Section
 - Variables
 - Constants
- Code Section
 - Program
 - Routines
 - Subroutines

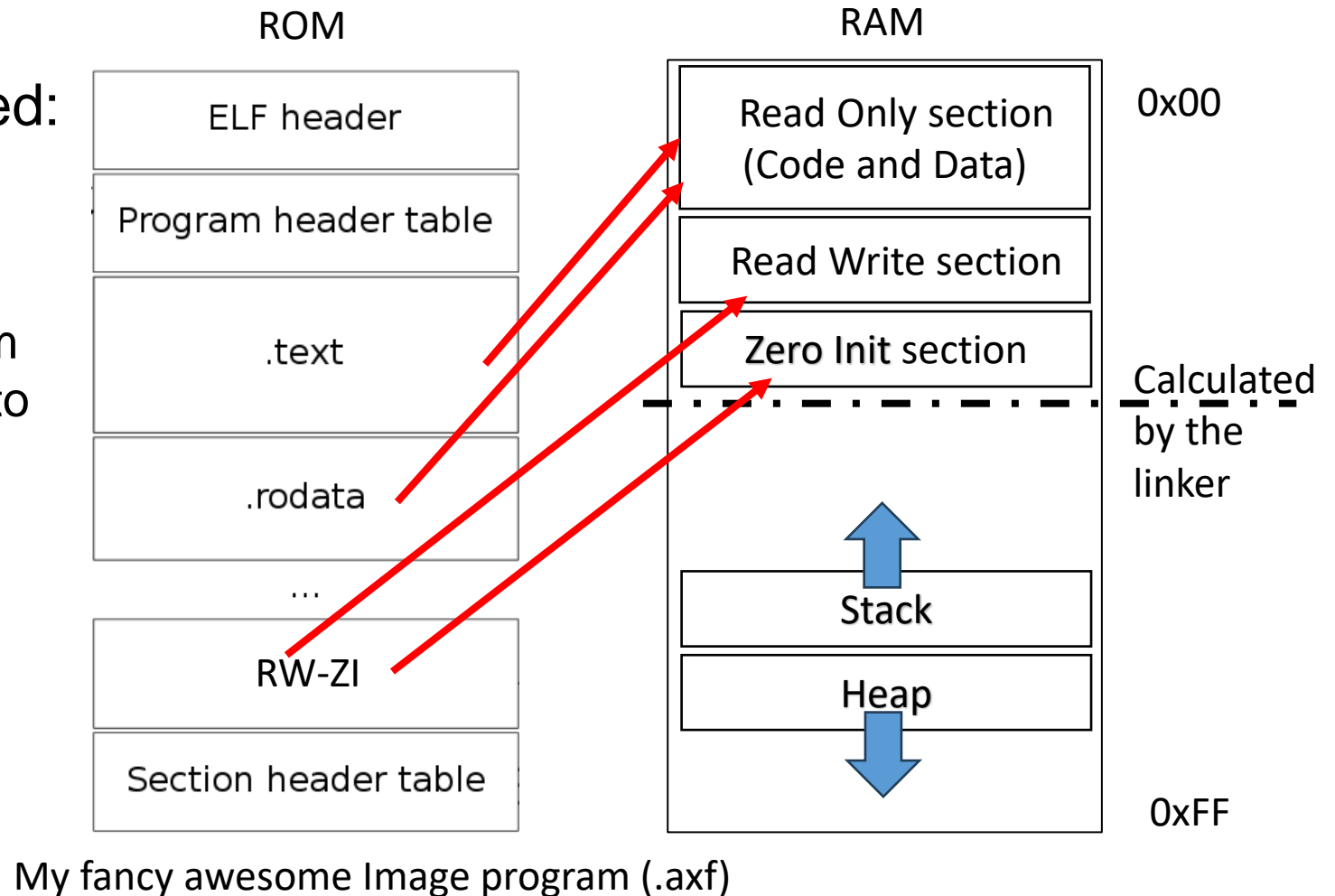
The executable – Load view

- The overall image (from the source code) is converted into an executable (.exe, .elf, .axf for Arm).
- Data and Code sections are in the executable.
- Composed of:
 - Entry address.
 - Stack and heap information.
 - Sections, used by the linker.
 - Segments, used by the loader (at runtime).



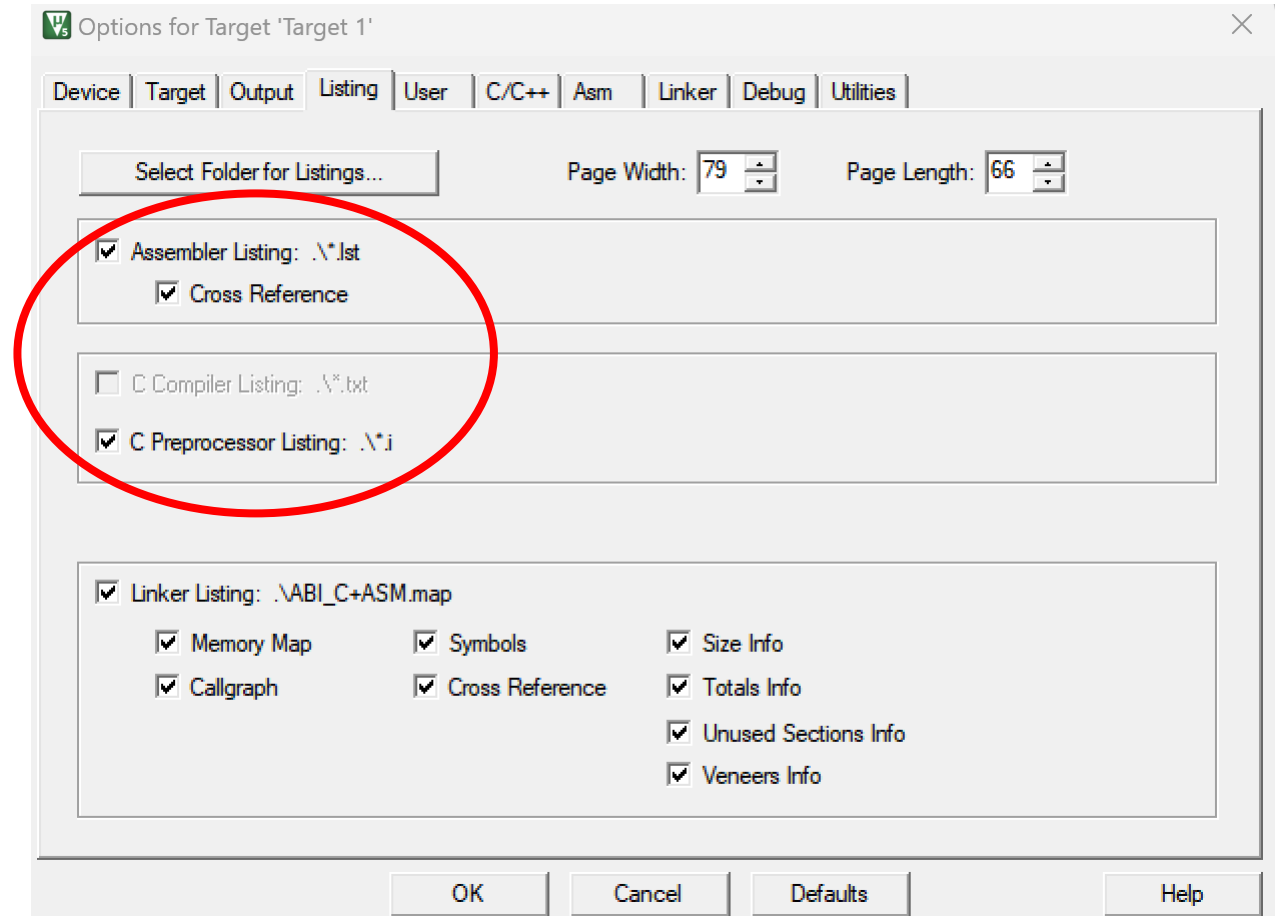
The executable – Execution view

- Before the image is executed:
 - Move executable segments from ROM to their execution addresses in RAM.
 - RW data must be copied from its load address in the ROM to its execution address in the RAM.
- Runtime memory layout information is calculated offline:
 - Stack and heap execution address and size.



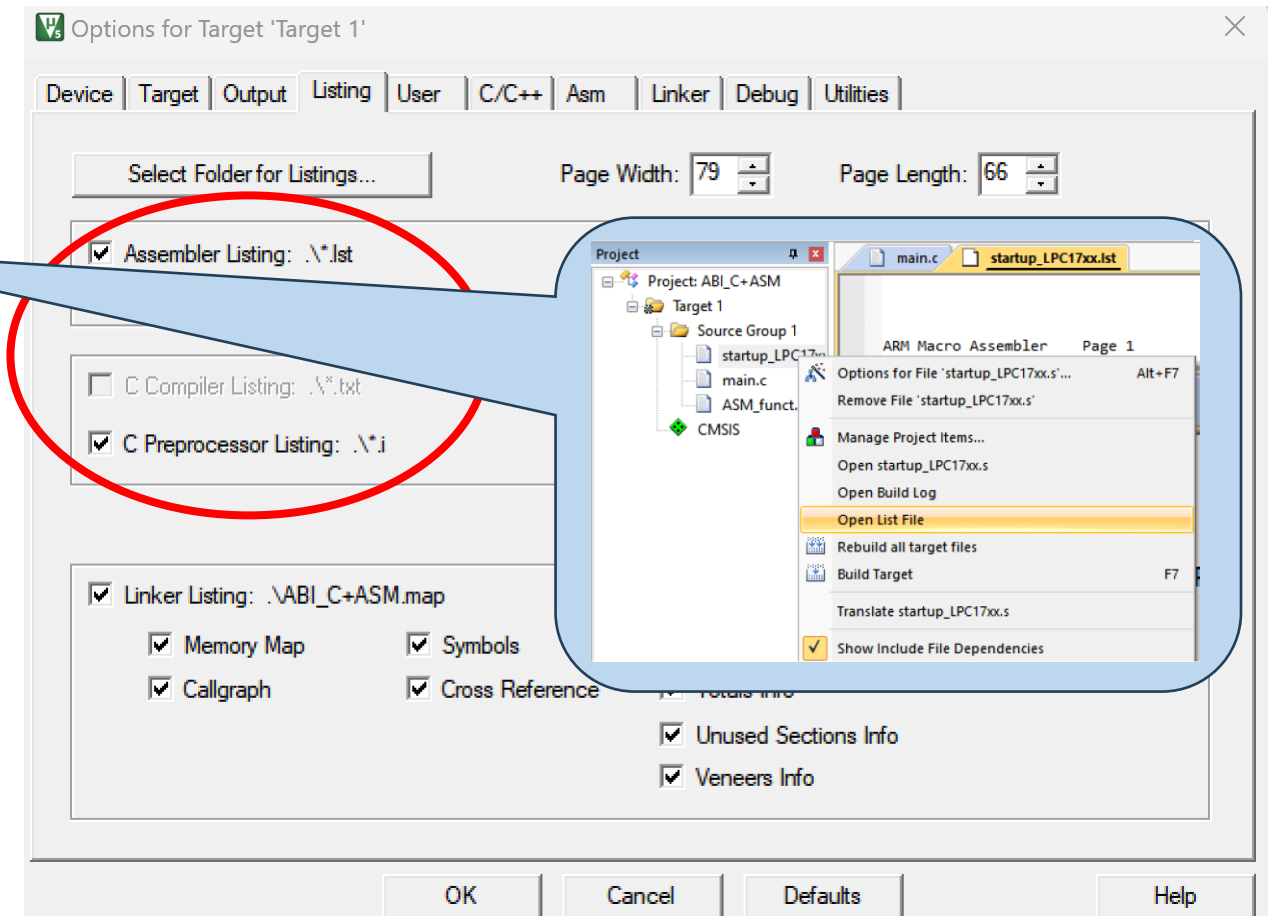
The Listing and dependencies files

- Dependencies files are generated (.d) and used by the toolchain (information needed during the link phase!).
- The project dependencies are in the .dep file.
- Listing files are debugging files showing how the code is translated in machine code.



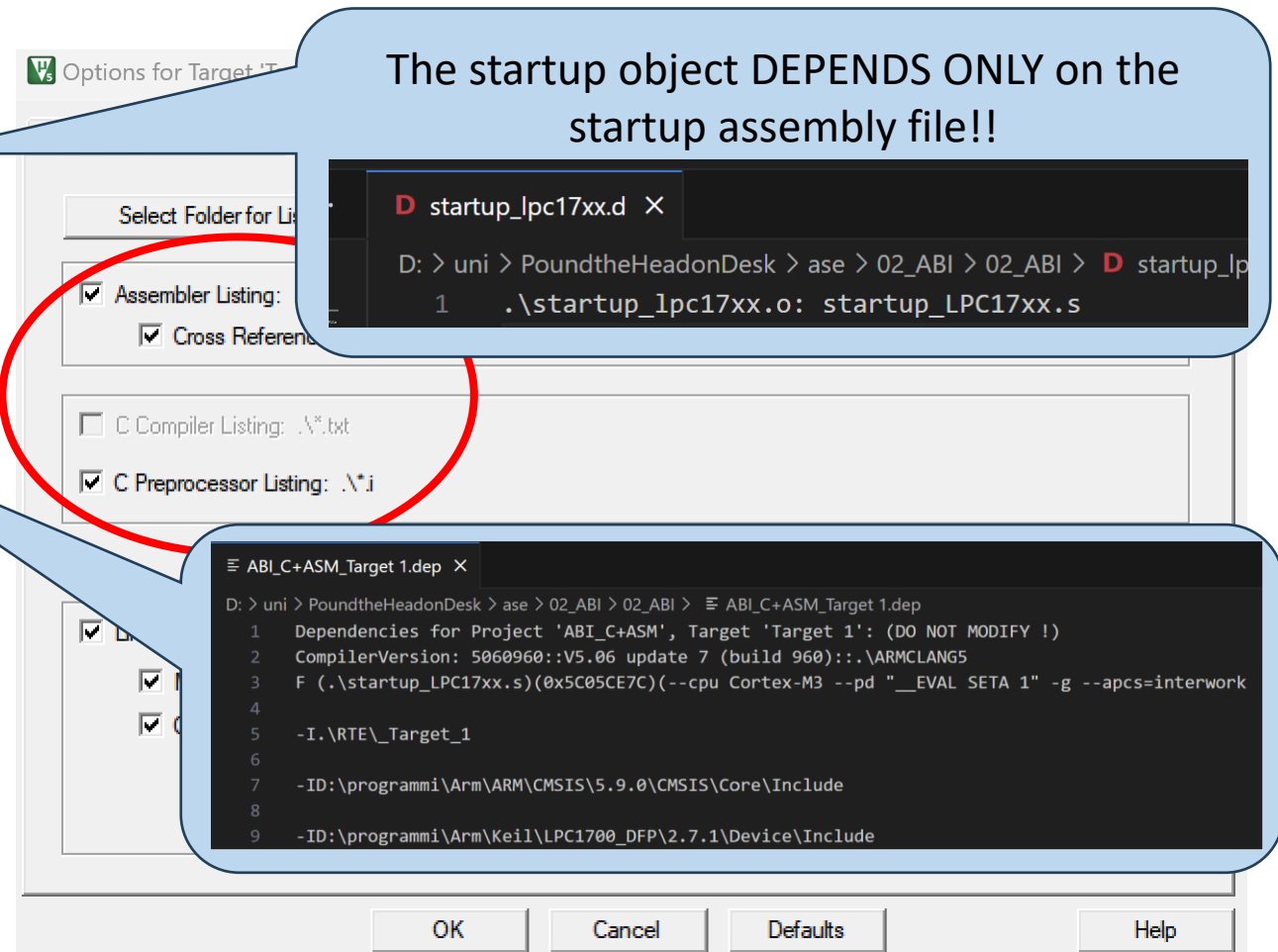
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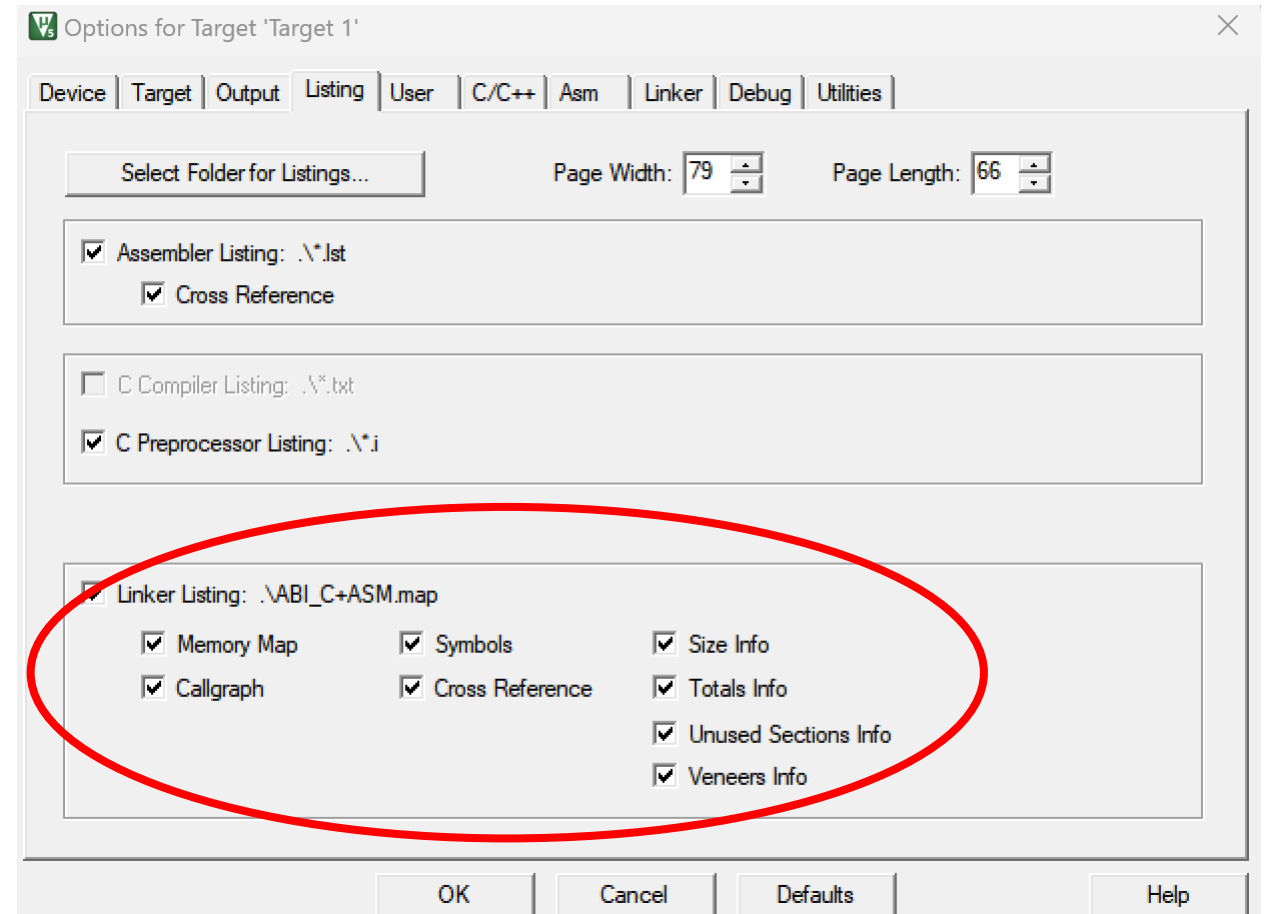
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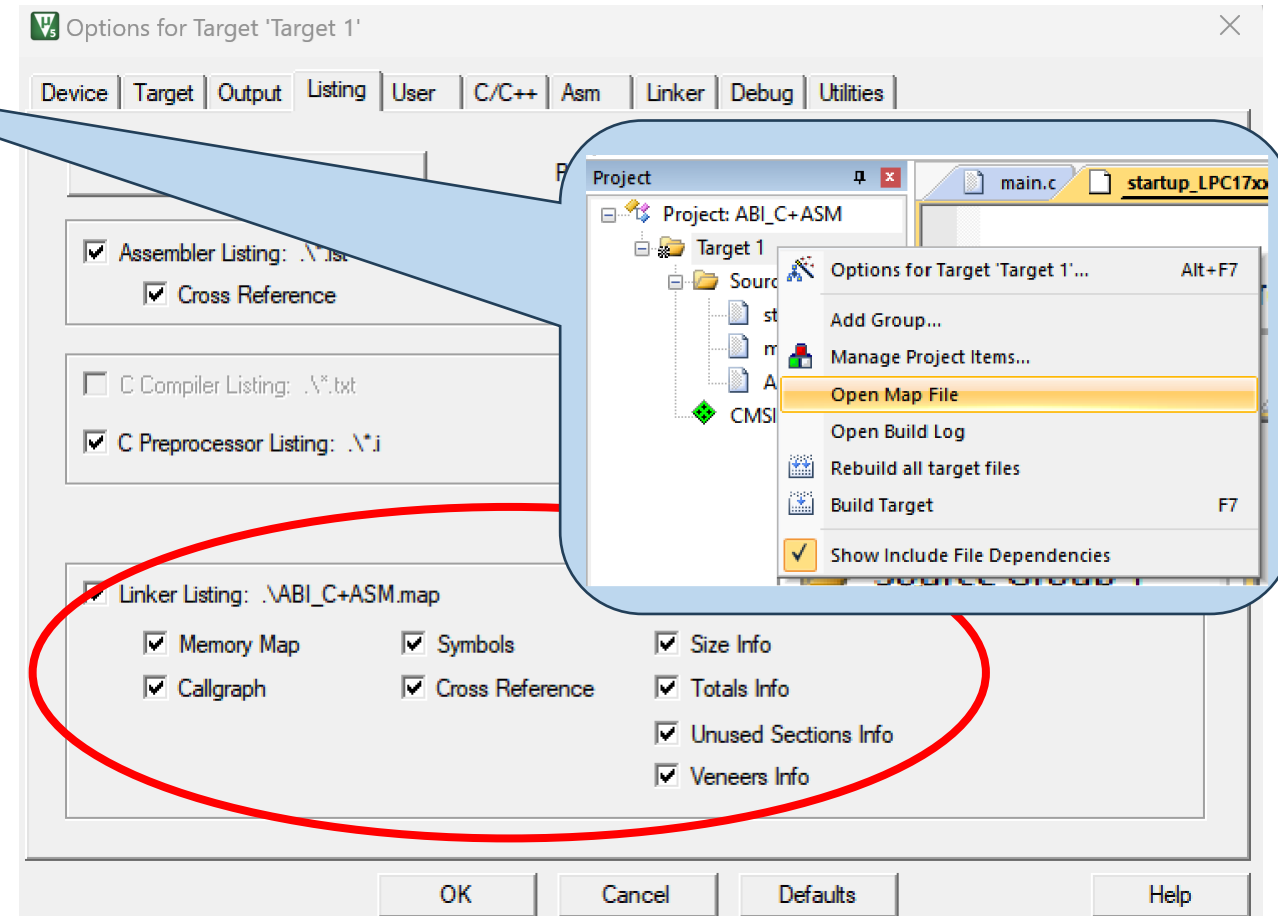
The map

- It is the output “log” of the link phase.
- It includes the memory map, symbols table, cross references and sizes.
- It may be used by debugging tools.



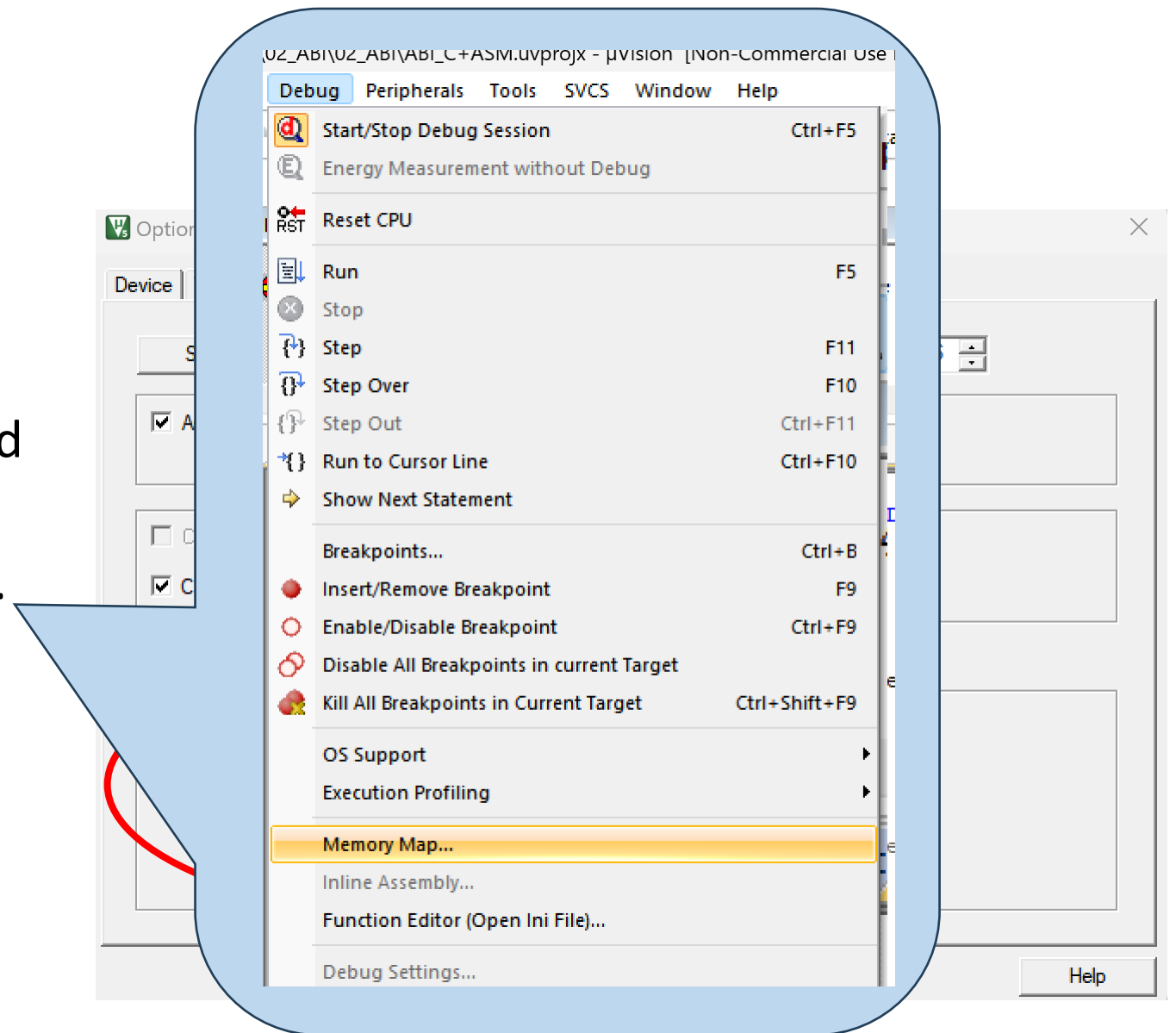
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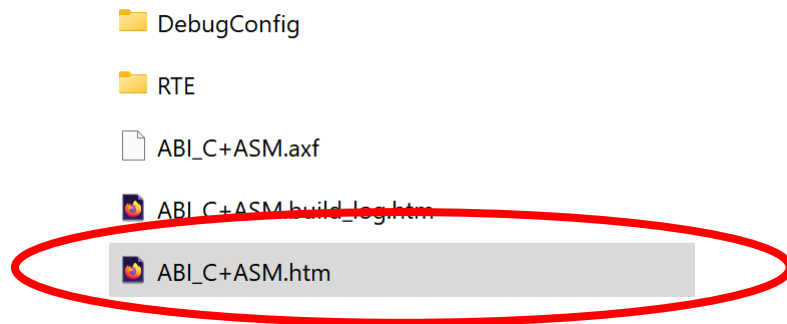
The map

- It is the output “log” of the link phase.
- It includes the memory map, symbols table, cross references and sizes.
- It may be used by debugging tools.



The static call graph file

- It is another debugging output “log” of the link phase.
- It is a control-flow graph.
- It represents the calling relationships between functions in the executable.



Static Call Graph for image .\ABI_C+ASM.axf

#<CALLGRAPH># ARM Linker, 5060960: Last Updated: Sun Dec 03 13:41:14 2023

Maximum Stack Usage = 16 bytes + Unknown(Functions without stacksize, Cycles, Untraceable Function Pointers)

Call chain for Maximum Stack Depth:

__rt_entry_main ⇒ main

Functions with no stack information

- [_user_initial_stackheap](#)
- [ASM_func](#)

Mutually Recursive functions

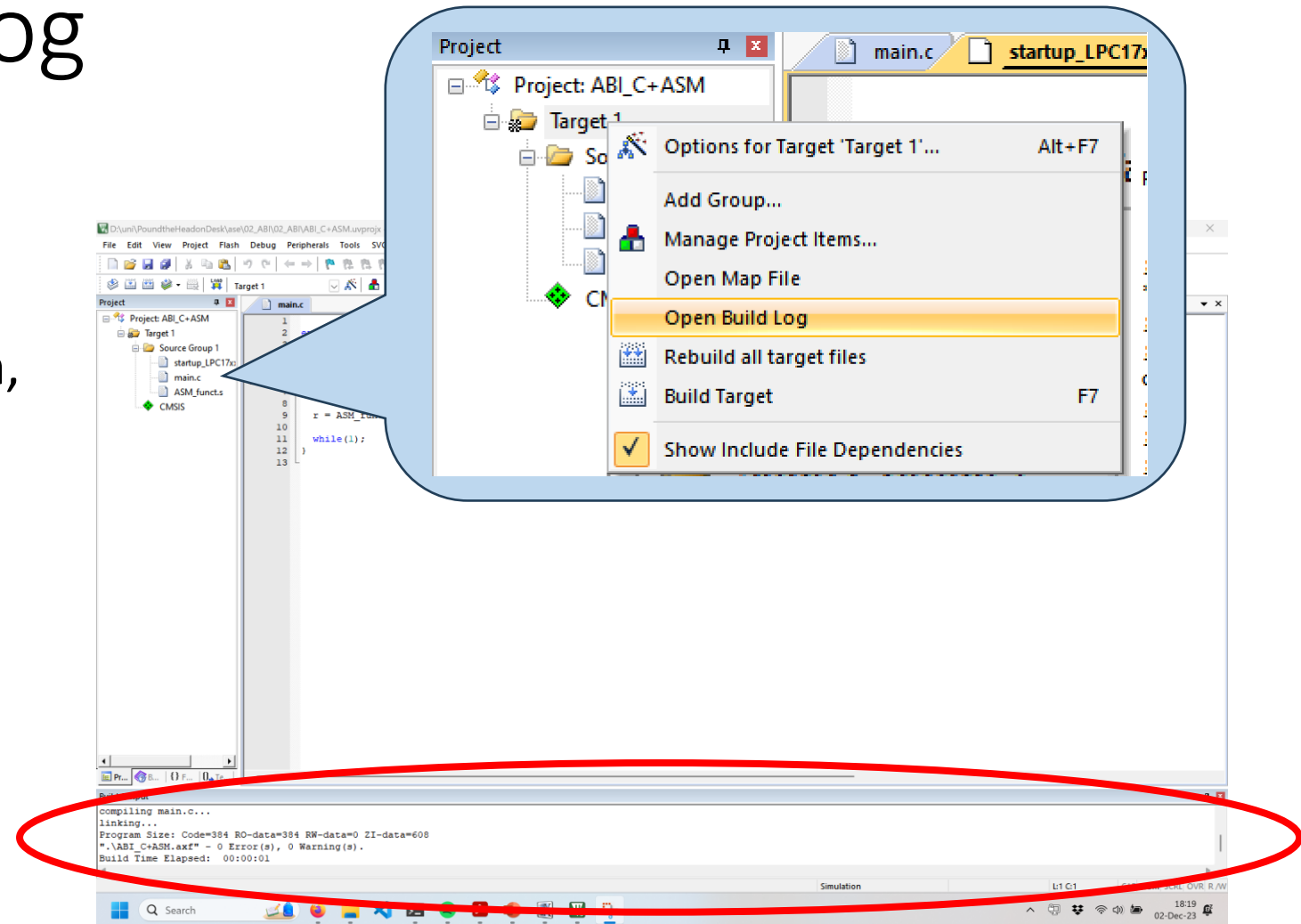
- [NMI_Handler](#) ⇒ [NMI_Handler](#)
- [HardFault_Handler](#) ⇒ [HardFault_Handler](#)
- [MemManage_Handler](#) ⇒ [MemManage_Handler](#)
- [BusFault_Handler](#) ⇒ [BusFault_Handler](#)
- [UsageFault_Handler](#) ⇒ [UsageFault_Handler](#)
- [SVC_Handler](#) ⇒ [SVC_Handler](#)
- [DebugMon_Handler](#) ⇒ [DebugMon_Handler](#)
- [PendSV_Handler](#) ⇒ [PendSV_Handler](#)
- [SysTick_Handler](#) ⇒ [SysTick_Handler](#)
- [ADC_IRQHandler](#) ⇒ [ADC_IRQHandler](#)

Function Pointers

- [ADC_IRQHandler](#) from startup_lpc17xx.o(text) referenced from startup_lpc17xx.o(RESET)
- [ASM_func](#) from asm_func.o(asm_functions) referenced from asm_func.o(asm_functions)
- [BOD_IRQHandler](#) from startup_lpc17xx.o(text) referenced from startup_lpc17xx.o(RESET)

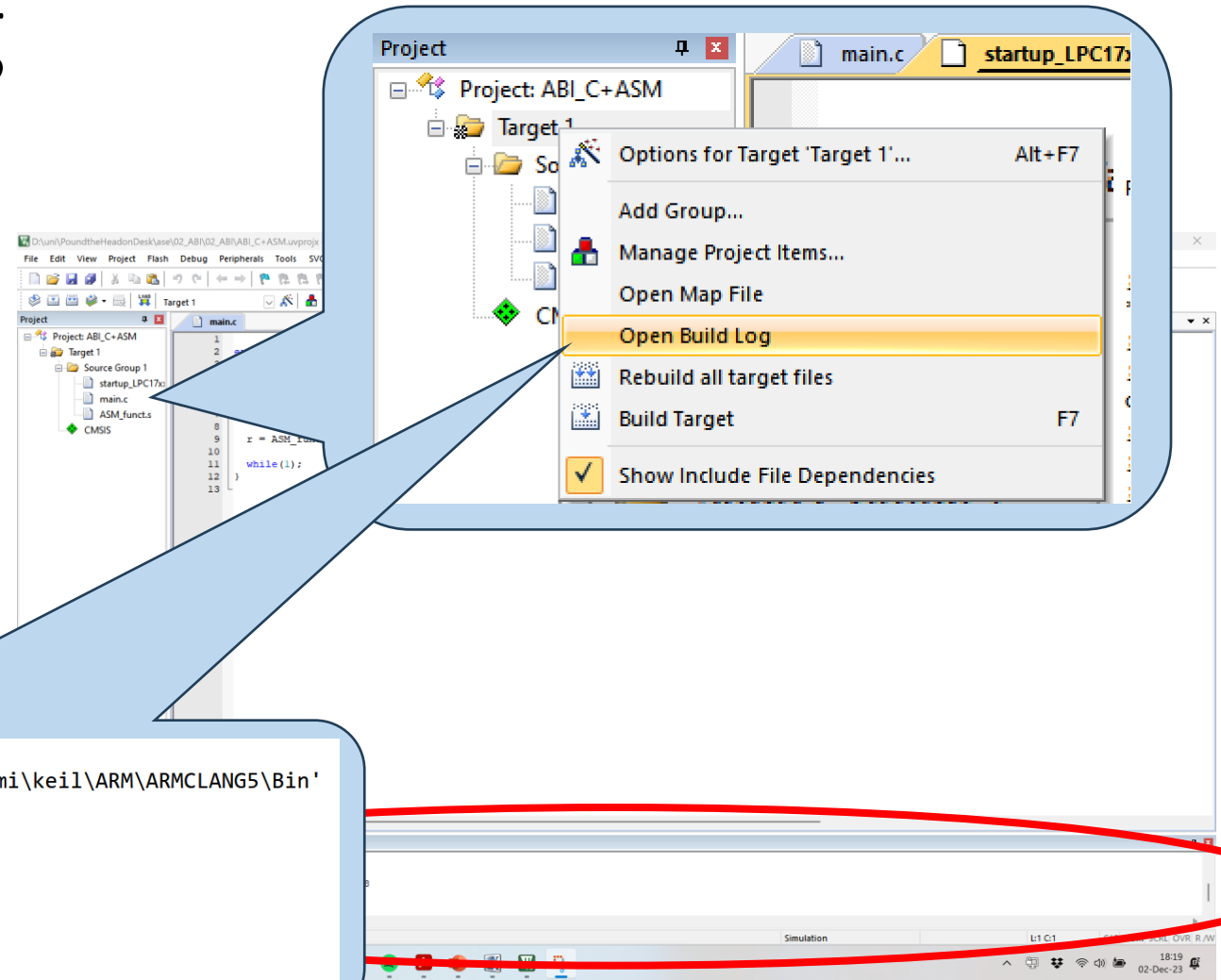
The build output log

- It is the log of the entire build process for a given project.
- It includes a log of tools version, software packages and components used.



The build output log

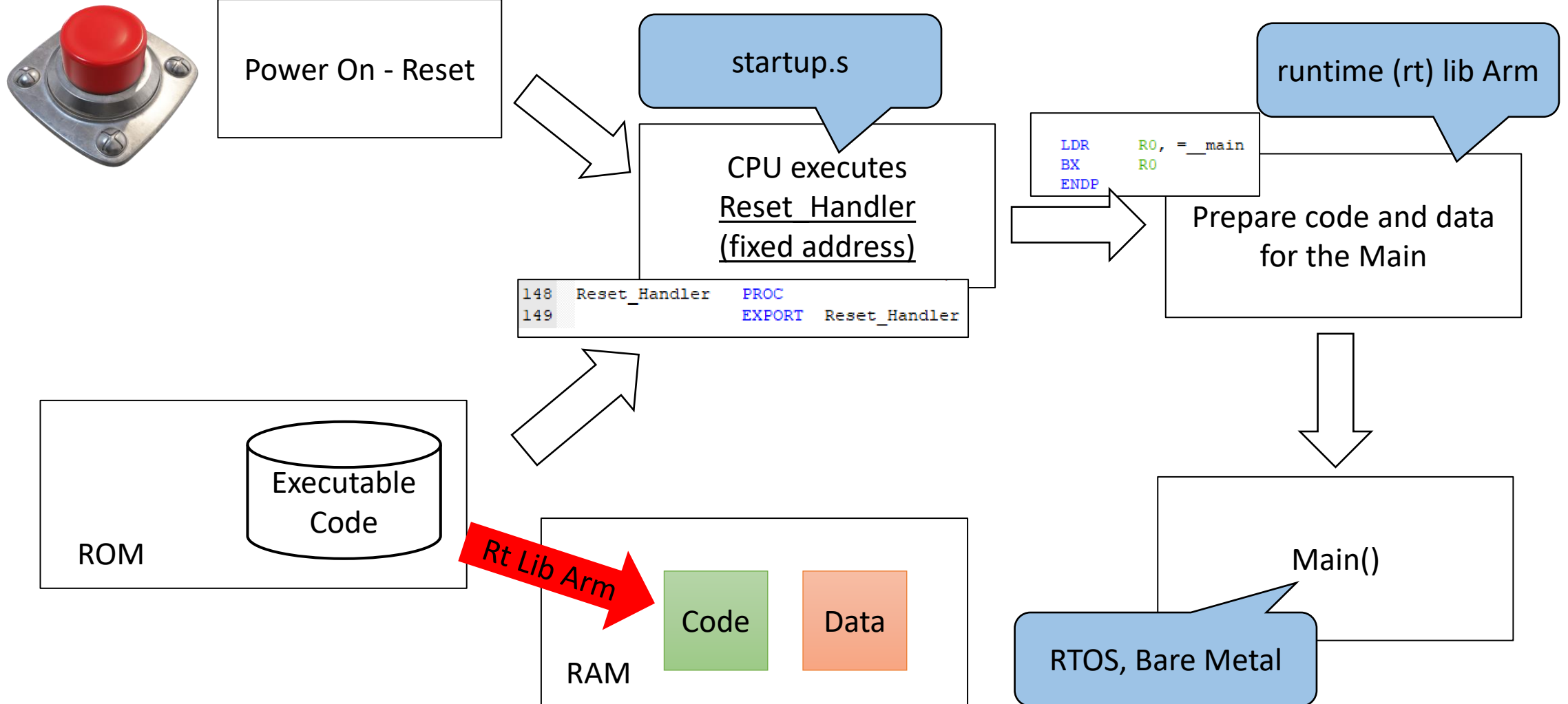
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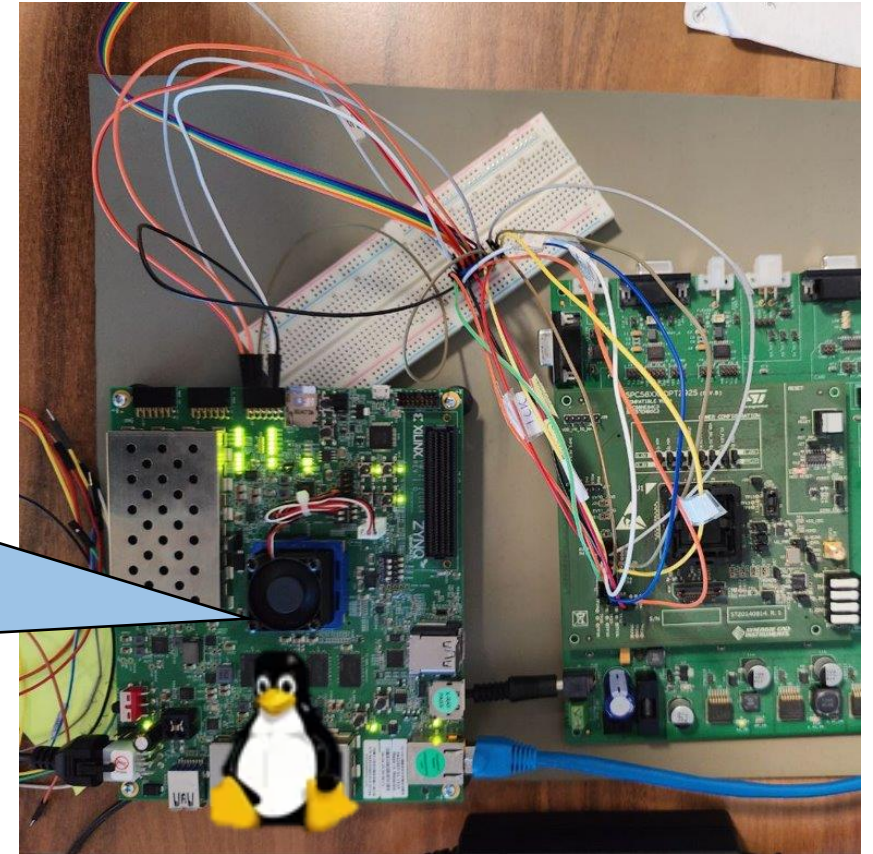
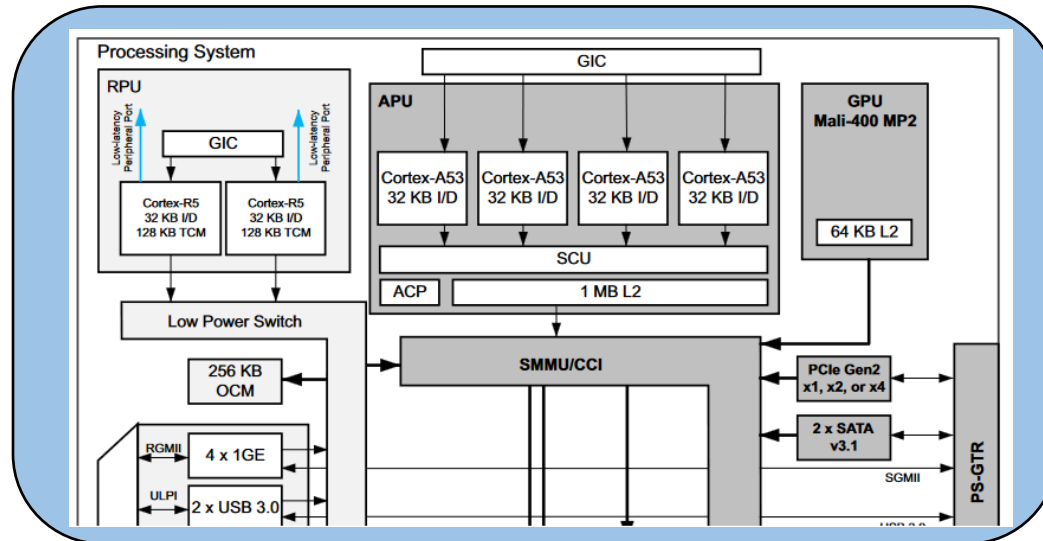
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How does a System-on-Chip start the program?

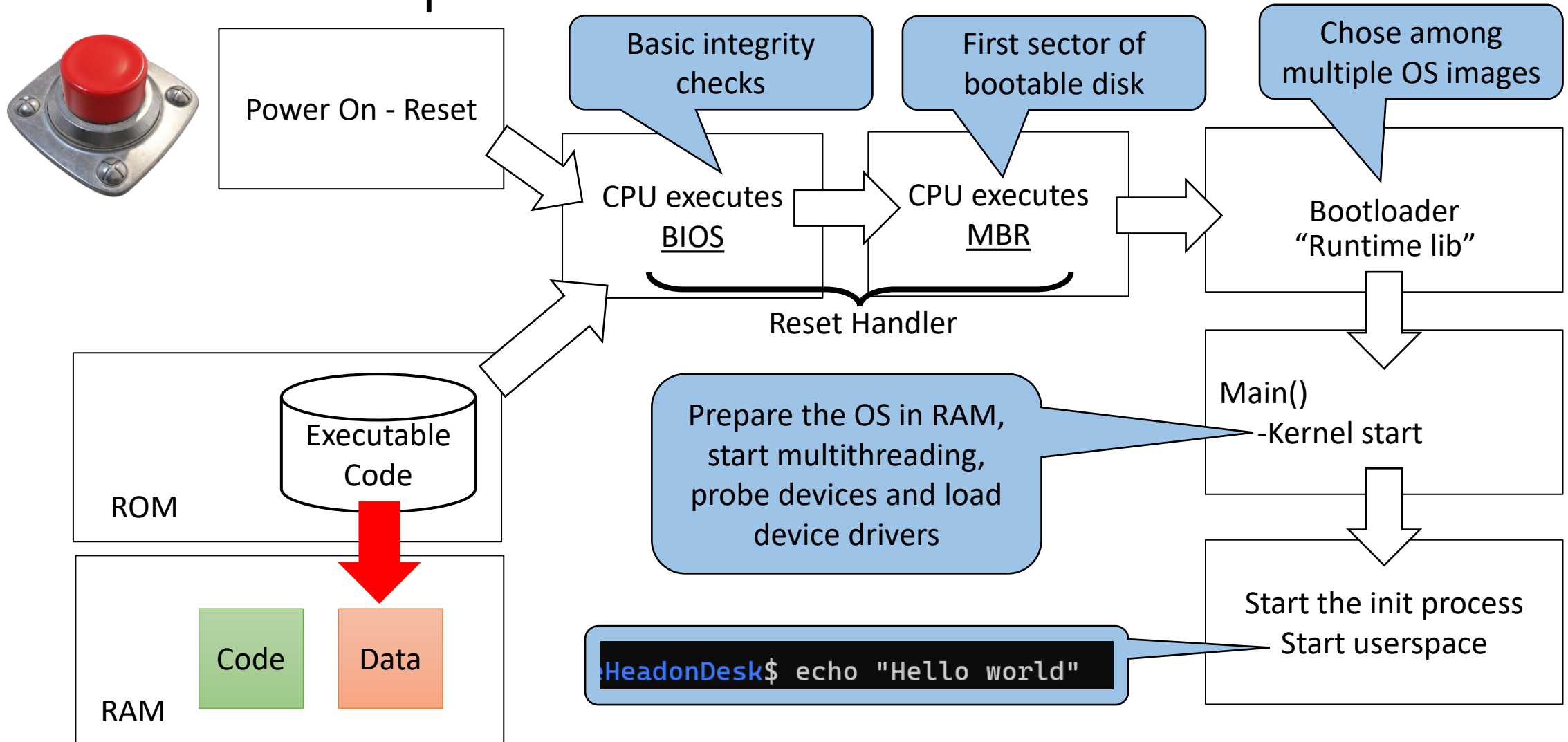


OS Bootstrap – The linux case

- Xilinx MPSoC zcu 104 Evaluation board.
- Four Cortex-A53 Arm Core.
- Running a custom linux-based operating system.



OS Bootstrap – The linux case



The BIOS

- BIOS stands for Basic Input/Output System.
- Performs some system integrity checks.
- Searches, loads, and executes the boot loader program.
- It looks for boot loader in floppy, cd-rom, or hard drive. You can press a key (typically F12 or F2, but it depends on your system) during the BIOS startup to change the boot sequence.
- Once the boot loader program is detected and loaded into the memory, BIOS gives the control to it.
- The BIOS loads and executes the MBR boot loader.

The MBR

- MBR stands for Master Boot Record.
- It is located in the 1st sector of the bootable disk. Typically `/dev/hda`, or `/dev/sda`.
- MBR is less than 512 bytes in size. This has three components:
 - Primary boot loader info in 1st 446 bytes
 - Partition table info in next 64 bytes
 - MBR validation check in last 2 bytes.
- It contains information about the bootloader.
- MBR loads and executes the boot loader.

The Bootloader

- If you have multiple kernel images installed on your system, you can choose which one to be executed.
- A common bootloader is GRUB (Grand Unified Bootloader).
- GRUB displays a splash screen, waits for few seconds, if you don't enter anything, it loads the default kernel image as specified in the grub configuration file.
- GRUB has the knowledge of the filesystem.
- Grub configuration file is `/boot/grub/grub.conf` (`/etc/grub.conf` is a link to this).
- GRUB just loads and executes Kernel and initrd images.

The Kernel

- Mounts the root file system as specified in the “root=” in grub.conf.
- Kernel executes the /sbin/init program.
- Since init was the 1st program to be executed by Linux Kernel, it has the process id (PID) of 1.
- initrd stands for Initial RAM Disk.
- initrd is used by kernel as temporary root file system until kernel is booted and the real root file system is mounted. It also contains necessary drivers compiled inside, which helps it to access the hard drive partitions, and other hardware.
- It starts other cores, and probe devices (loading their device drivers).

Start the user space – the init process

- Looks at the `/etc/inittab` file to decide the Linux run level.
- Following are the available run levels
 - 0 – halt, 1 – Single user mode, 2 – Multiuser, without NFS, 3 – Full multiuser mode, 4 – unused, 5 – X11, 6 – reboot.
- Init identifies the default initlevel from `/etc/inittab` and uses that to load all appropriate program.
- Execute `'grep initdefault /etc/inittab'` on your system to identify the default run level.
- If you want to get into trouble, you can set the default run level to 0 or 6. Since you know what 0 and 6 means, probably you might not do that.
- Typically you would set the default run level to either 3 or 5.

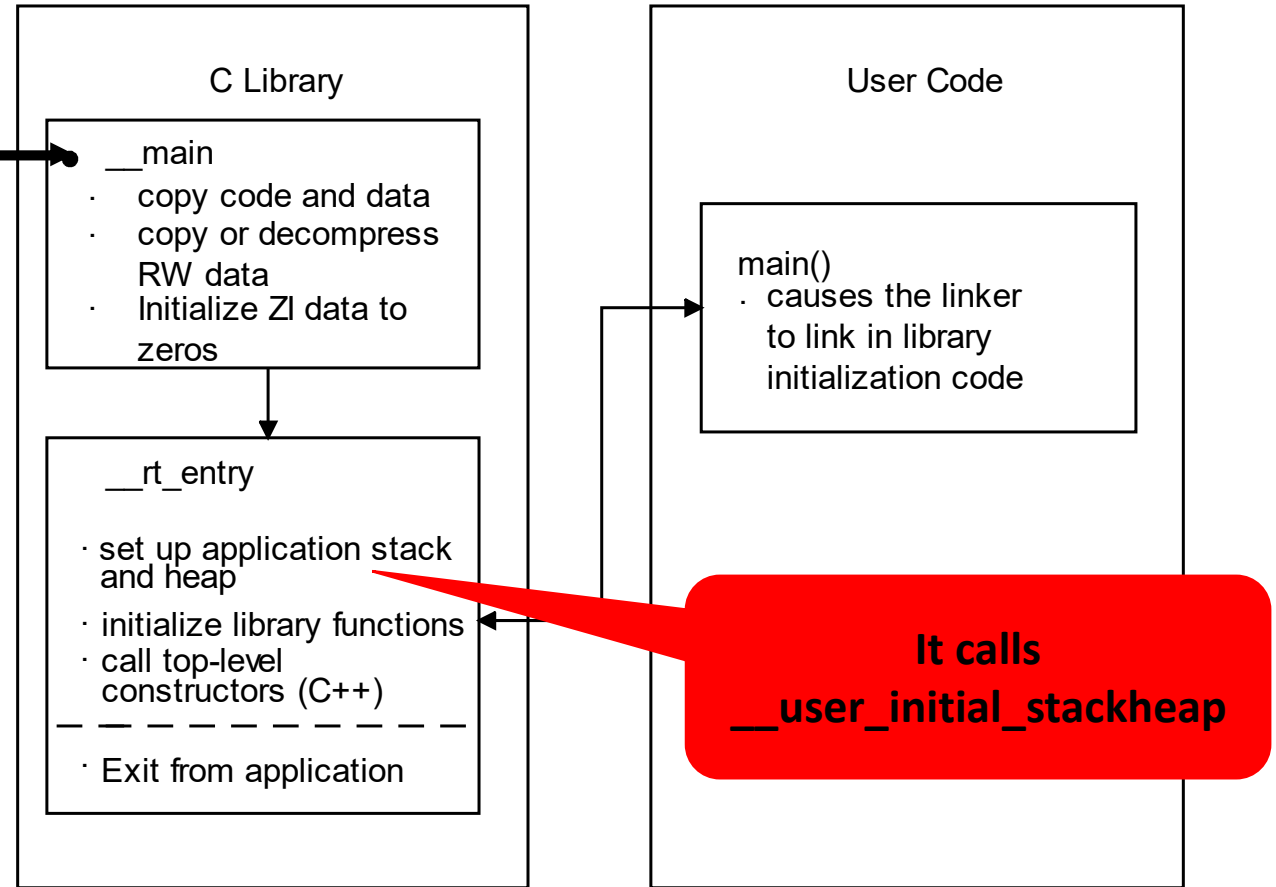
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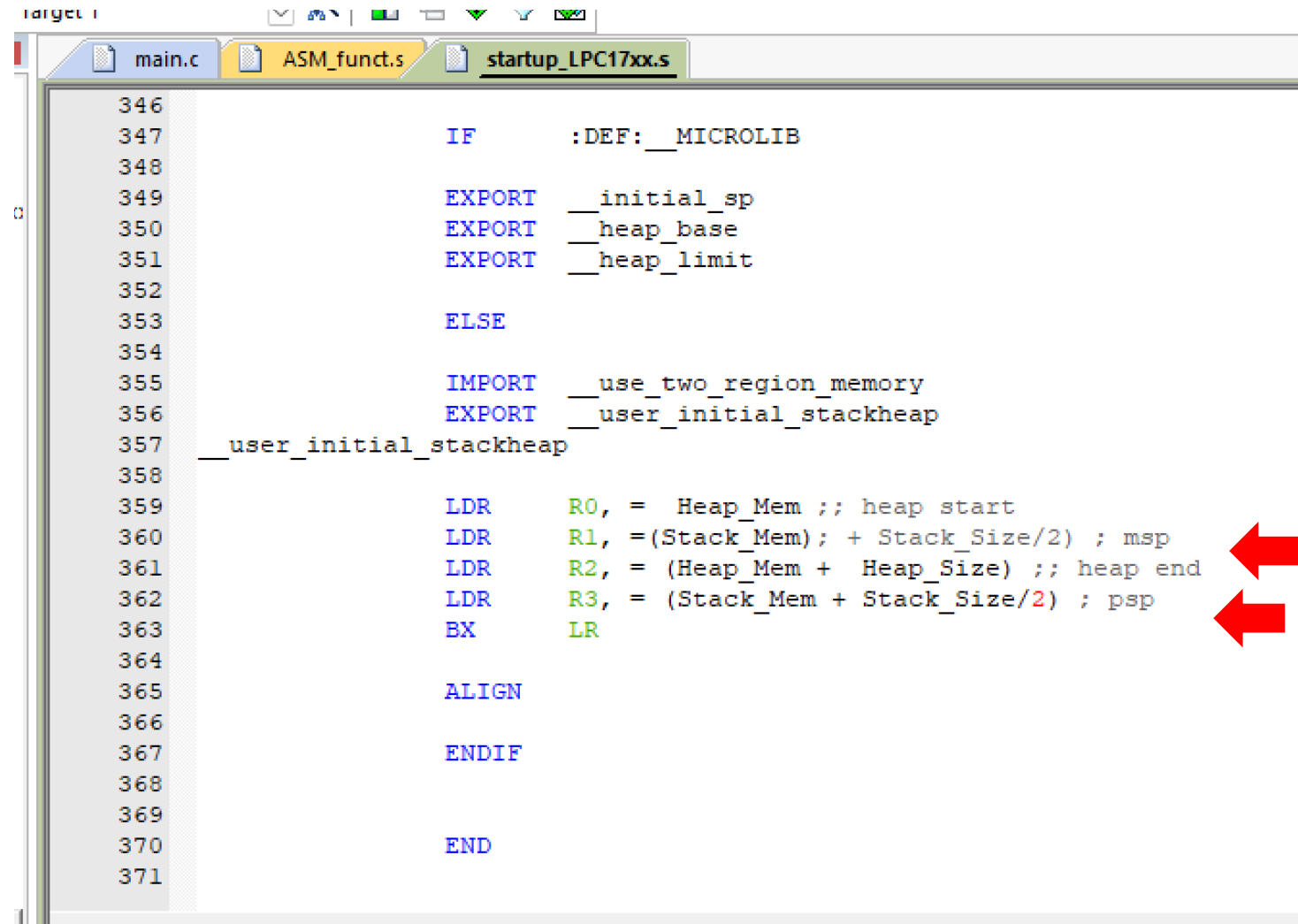
The Arm “Magic secret sauce”

```
Reset_Handler  PROC
                 EXPORT Reset_Handler            [WEAK]
                 IMPORT __main
                 LDR     R0, =__main
                 BX      R0
                 ENDP
```

- `__main` responsible for:
 - Setting up the memory code and data.
- `__rt_entry` responsible for:
 - Setting up **stack(s)** and **heap**.
 - Initializing the lib functions and static data.
 - Calling any top level constructors.



Setting up stack(s)



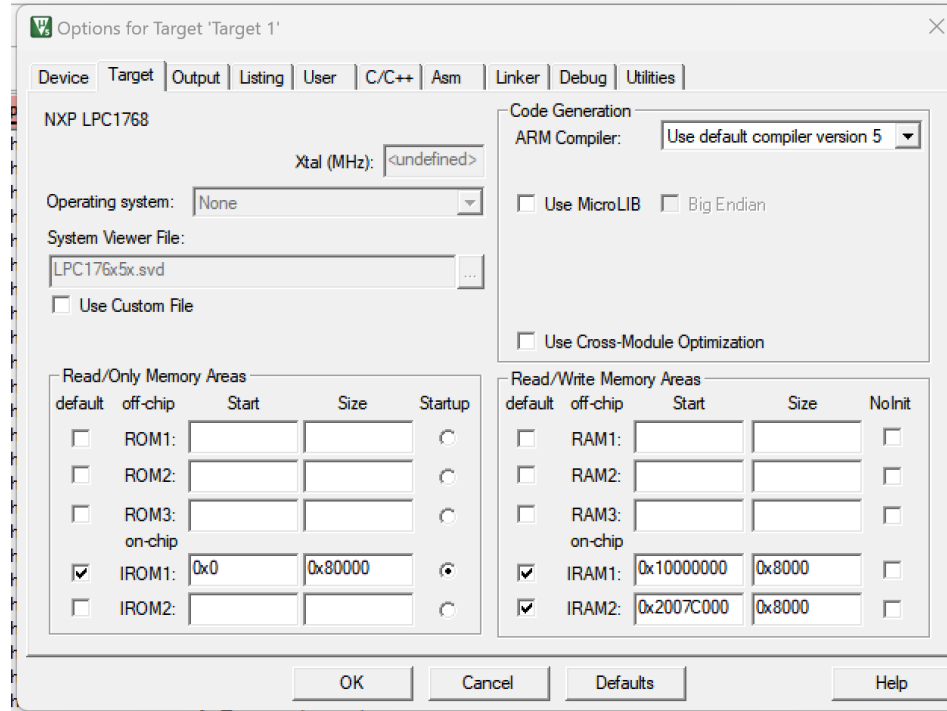
```
target 1
main.c ASM_funct.s startup_LPC17xx.s
346
347         IF      :DEF:__MICROLIB
348
349         EXPORT  __initial_sp
350         EXPORT  __heap_base
351         EXPORT  __heap_limit
352
353         ELSE
354
355         IMPORT  __use_two_region_memory
356         EXPORT  __user_initial_stackheap
357 __user_initial_stackheap
358
359         LDR      R0, = Heap_Mem ;; heap start
360         LDR      R1, =(Stack_Mem); + Stack_Size/2) ; msp
361         LDR      R2, = (Heap_Mem + Heap_Size) ;; heap end
362         LDR      R3, = (Stack_Mem + Stack_Size/2) ; psp
363         BX      LR
364
365         ALIGN
366
367         ENDIF
368
369
370         END
371
```

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Example – Where do they belong?

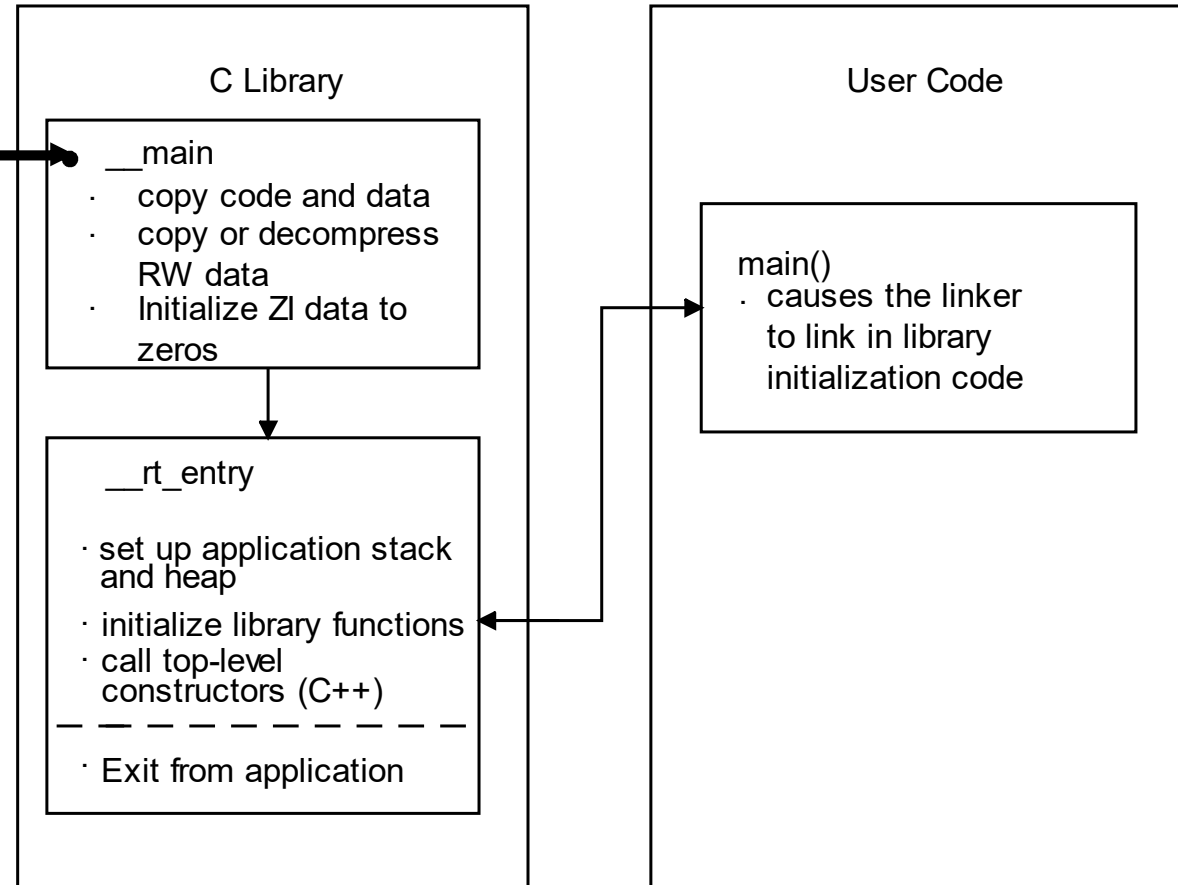
- See the map file.



```
4
5
6  const int pippo[]={21,32};
7  int this_is_zero;
8  int caught=234;
9  volatile int array[N] = {1,2,3,4,5,6,7,8,9,10};
10
11 void my_fancy_function(int * a ){
12     if (a==0) {
13         return ;
14     }
15     (*a)++;
16 }
17
18 int main(void) {
19
20     int i=0;
21     volatile int value=0;
22     volatile int my_var=pippo[0];
23
24     my_fancy_function(&my_var);
25
26     for (i=0; i<N; i++) {
27         value=array[i];
28     }
29
30
31     while(1);
32 }
33
34
35
```

Example – Skipping the “Magic secret sauce”

```
Reset_Handler  PROC
EXPORT  Reset_Handler      [WEAK]
IMPORT  __main
LDR     R0, =__main
BX      R0
ENDP
```



Example – Skipping the “Magic secret sauce”

```
Reset_Handler PROC
EXPORT Reset_Handler [WEAK]
IMPORT main
LDR R0, =main
BX R0
ENDP
```

```
Reset_Handler PROC
EXPORT Reset_Handler [WEAK]
IMPORT main
LDR R0, =main
BX R0
ENDP
```

- Compilation without errors.
- See the map file.
- Is it correct?

```
4
5
6 const int pippo[]={21,32};
7 int this_is_zero;
8 int caught=234;
9 volatile int array[N] = {1,2,3,4,5,6,7,8,9,10};
10
11 void my_fancy_function(int * a ){
12     if (a==0) {
13         return ;
14     }
15     (*a)++;
16 }
17
18 int main(void){
19
20
21     int i=0;
22     volatile int value=0;
23     volatile int my_var=pippo[0];
24
25     my_fancy_function(&my_var);
26
27     for (i=0; i<N; i++) {
28         value=array[i];
29     }
30
31
32
33     while(1);
34 }
35
```

Example – On the fly variable declaration

- See build log.

```
*** Using Compiler 'V5.06 update 7 (build 960)', folder: 'D:\programmi\keil\ARM\ARMCLANG5\Bin'
Build target 'Target 1'
compiling main.c...
main.c(25): warning: #167-D: argument of type "volatile int *" is incompatible with parameter of type
    my_fancy_function(&my_var);
main.c(27): error: #29: expected an expression
    for (int i=0; iSoftware Packages used:

Package Vendor: ARM
    http://www.keil.com/pack/ARM.CMSIS.5.9.0.pack
    ARM.CMSIS.5.9.0
    CMSIS (Common Microcontroller Software Interface Standard)
* Component: CORE Version: 5.6.0

Package Vendor: Keil
    http://www.keil.com/pack/Keil.LPC1700_DFP.2.7.1.pack
    Keil.LPC1700_DFP.2.7.1
    NXP LPC1700 Series Device Support, Drivers and Examples for MCB1700 and LPC1788-32
```

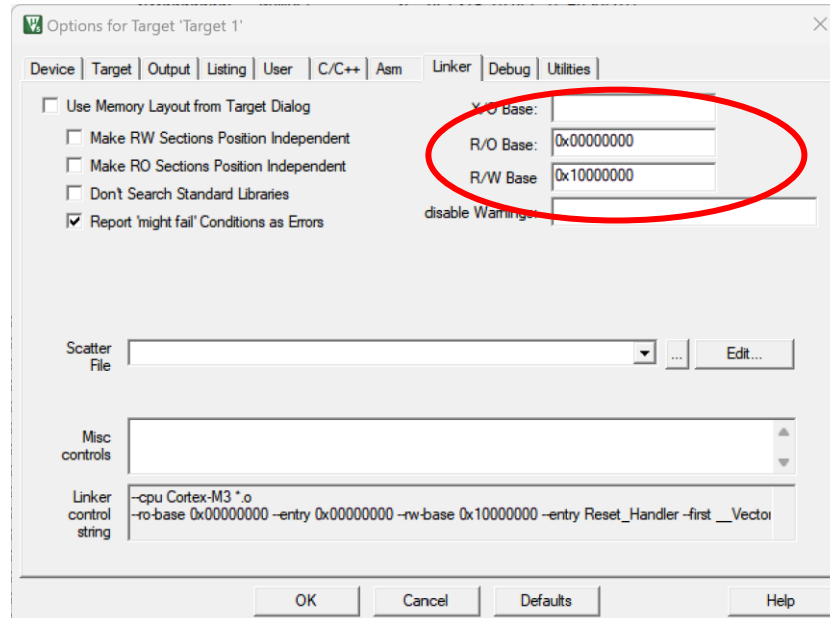
```

4
5
6 const int pippo[]={21,32};
7 int this_is_zero;
8 int cached=234;
9 volatile int array[N] = {1,2,3,4,5,6,7,8,9,10};
10
11 void my_fancy_function(int * a ){
12     if (a==0) {
13         return ;
14     }
15     (*a)++;
16 }
17
18 int main(void){
19
20
21     int i=0;
22     volatile int value=0;
23     volatile int my_var=pippo[0];
24
25     my_fancy_function(&my_var);
26
27     for (i=0; i<N; i++) {
28         value=array[i];
29     }
30
31
32
33     while(1);
34 }
35

```

Example - Changing the Section start address

- See the map file.
- Let's put the IROM1 start address at 0x00100.
- Same for RW base address (see array in map file).



```
4
5
6 const int pippo[]={21,32};
7 int this_is_zero;
8 int caught=234;
9 volatile int array[N] = {1,2,3,4,5,6,7,8,9,10};
10
11 void my_fancy_function(int * a ){
12     if (a==0) {
13         return ;
14     }
15     (*a)++;
16 }
17
18 int main(void){
19
20
21     int i=0;
22     volatile int value=0;
23     volatile int my_var=pippo[0];
24
25     my_fancy_function(&my_var);
26
27     for (i=0; i<N; i++) {
28         value=array[i];
29     }
30
31
32     while(1);
33 }
34
35
```

Example – Adding a custom section

- See the map file.

```
Removing Unused input sections from the image.
```

```
Removing asm_funct.o(my_asm_functions), (32 bytes).
```

```
1 unused section(s) (total 32 bytes) removed from the image.
```

- Let's call the function from the main.

```
2
3      AREA my_asm_functions, CODE, READONLY
4      EXPORT ASM_funct
5
6      ASM_funct
7
8      ; save current SP for a faster access
9      ; to parameters in the stack
10     MOV    r12, sp
11     ; save volatile registers
12     STMFD  sp!, {r4-r8, r10-r11, lr}
13     ; extract argument 4 and 5 into R4 and R5
14     LDR    r4, [r12]
15     LDR    r5, [r12, #4]
16
17     LDR    r12, =ASM_funct
18     STR    r5, [r12]
19
20     ; setup a value for R0 to return
21     MOV    r0, r5
22     ; restore volatile registers
23     LDMFD  sp!, {r4-r8, r10-r11, pc}
24
25     END
```


Example – Adding a custom section

- See the map file.

Removing Unused input sections from the image.

Removing asm_funct.o(my_asm_functions), (32 bytes).

1 unused section(s) (total 32 bytes) removed from the image.

- Let's call the function from the main.

main.c	startup_LPC17xx.s	ASM_funct.s	ABI_C+ASM.map
.text	0x00000364	Section	
.text	0x00000378	Section	
my_asm_functions	0x00000384	Section	
.constdata	0x000003a4	Section	
.data	0x10000000	Section	
.bss	0x10000030	Section	
HEAP	0x10000090	Section	

```
2
3          AREA my_asm_functions, CODE, READONLY
4          EXPORT ASM_funct
5 ASM_funct
6          ; save current SP for a faster access
7          ; to parameters in the stack
```

```
22 int main(void) {
23
24     int azeroth=0xFFFFFFFF, kalimdor=2,
25         outlands=3, northrend=4,
26         pandaria=5, shadowlands=6;
27     volatile int r=0;
28     int i=0;
29     volatile int value=0;
30     volatile int my_var=pippo[0];
31
32     //my_fancy_function(&my_var);
33
34     for (i=0; i<N; i++) {
35         value=array[i];
36     }
37
38
39     r = ASM_funct(azeroth, kalimdor,
40                 outlands, northrend,
41                 pandaria, shadowlands);
42
43     while(1);
44 }
45
```

Example – Adding a custom section

- You can also use `__attribute__((section(".ARM.__at_<address>")))` to specify the absolute address of a variable or a section name.
- The linker help us.

This is unsafe!!!

```
#if TEST_UNSAFE
__attribute__((section(".ARM.__at_0x10000008"))) int fancy_value[]={
    0xcafecafe, 0xc1a0c1a0};
#endif
```

```
volatile int array[N] = {1,2,3,4,5,6,7,8,9,10};
```

- It directly handles all the data placement during the link phase.

```
31
32 int *p=(int *)0x10000008;
33 *p=101010;
34
35
36 my_fancy_function(&my_var);
37
38 for (i=0; i<N; i++) {
39     value=array[i];
40 }
41
```

startup_LPC17xx.s	ASM_funct.s	main.c	ABI_C+ASM.map
fancy_value	0x10000008	Data	8 main.o(.ARM.__at_0x10000008)
this_is_zero	0x10000010	Data	4 main.o(.data)
cached	0x10000014	Data	4 main.o(.data)
array	0x10000018	Data	40 main.o(.data)
__libspace_start	0x10000040	Data	96 libspace.o(.bss)
__temporary_stack_top\$libspace	0x100000a0	Data	0 libspace.o(.bss)

Example – Exporting data from asm to C

- Data can be exported from asm to C (like functions).
- Why in my_amazing_asm_vector do we have the first data?

```
AREA my_amazing_data, DATA, READWRITE
EXPORT _my_asm_vector
_my_asm_vector DCD 0x1234, 0x3214
```

```
extern unsigned int * _my_asm_vector;
```

```
volatile unsigned int * my_amazing_asm_vector=_my_asm_vector;
```

Example – Exporting data from asm to C

- Data can be exported from asm to C (like functions).
- Why in my_amazing_asm_vector do we have the first data?
- Let's have a look at the asm_func.lst file.

```
76 ARM Macro Assembler      Page 1 Alphabetic symbol ordering
77 Relocatable symbols
78
79 _my_asm_vector 00000000
80
81 Symbol: _my_asm_vector
82   Definitions
83     At line 28      file ASM_func.s
84   Uses
85     At line 27      file ASM_func.s
86   Comment: _my_asm_ used once
87   my_amazing_data
88
```

```
25 0000001C
26 0000001C 00000000      AREA      my_amazing_data, DATA, READWRIT
E
27 00000000      EXPORT      _my_asm_vector
28 00000000 00001234
   | | | | 00003214 _my_asm_vector
   | | | | DCD      0x1234,0x3214
29 00000008
30 00000008
```

It is a symbol (label) to the first value!!

Example – Exporting data from asm to C

- The linker is in charge of substituting the symbols with addresses during the link phase.
- A linker symbol is not equivalent to a variable declaration in high level language, it is instead a symbol that does not have a value.
- The label `_my_asm_vector` is a linker symbol (it is directly defined in assembly and used in C), when the linker resolves the symbols it retrieves the first value of the vector (independently from the data type).
- The correct way to export data defined in assembly is to force the use of the address.

```
AREA my_amazing_data, DATA, READWRITE
EXPORT _my_asm_vector
_my_asm_vector DCD 0x1234, 0x3214
```

```
extern unsigned int * _my_asm_vector;
```

```
extern unsigned int _my_asm_vector;
```

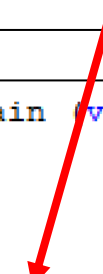
```
volatile unsigned int * my_amazing_asm_vector=&_my_asm_vector;
```

Example – ASM SVC vs C-ASM SVC calling

- In assembly, you have full control over the stacks.
- What happens if we call the SVC from C?

```
111 Reset_Handler PROC
112     EXPORT Reset_Handler
113     import __main
114     ; your code here
115
116     MOV     R0, #3
117     MSR     CONTROL, R0
118     LDR     SP, =Stack_Mem
119
120     nop
121
122     SVC     0x10      ;0x000000DA
123
```

```
3 int main (void) {
4
5
6
7
8     __asm volatile("svc 0x10");
9
10    while(1);
11 }
12
```



Remainder!!

- Debugging is a very, very, very long painful process.
- Tools, especially the compiler, are your best (and worst) friends!
- Knowledge of toolchains easily allows you to debug your code.
- The more information you provide to the toolchain, the fewer chances to have different results than the one you have in mind!



References

- [Clang and the LLVM project](#)
- [Arm Compiler](#)
- [Arm Embedded Software Development](#)
- [Arm Image Structure and Generation](#)
- [Debugging With Arbitrary Record Formats \(DWARF\)](#) and [Executable and Linkable Format \(ELF\)](#)
- [Linux Boot Process](#)