

Design and Analysis of Algorithms

Practice-sheet : **Amortized Analysis and Fibonacci heaps**

**1. Deleting half elements**

Design a data structure to support the following two operations for a dynamic multiset  $S$  of integers, which allows the duplicate values:

- *Insert*( $S, x$ ): inserts  $x$  into  $S$ .
- *Delete-Larger-Half*( $S$ ): delete the largest  $\lceil |S|/2 \rceil$  elements from  $S$ .
- *Report-max*: report the largest element from the set.

Explain how to implement this data structure so that any sequence of  $m$  operations mentioned above run in  $O(m)$  time. Your implementation should also include a way to output the elements of  $S$  in  $O(|S|)$  time.

**2. Simulating a queue using stacks**

Show how to implement a queue with two ordinary stacks so that amortized cost of each *Enqueue* and each *Dequeue* operation is  $O(1)$ .

**3. Computing  $\pi$  function of a pattern** In the KMP algorithm for pattern matching, we assumed that the  $\pi$  function associated with the pattern  $P[1..m]$  was available. You have to design  $O(m)$  time algorithm for computing  $\pi$ .

Execute your algorithm to compute the  $\pi$  function for the pattern “ababbabbabbababb”.

**Hint:** There is a great amount of similarities between the  $\pi$  function and **Computing-F**( $i$ ) function discussed in the class.

**4. Multi-stack**

A multistack consists of an infinite series of stacks  $S_0, S_1, S_2, \dots$ , where the  $i$ th stack  $S_i$  can hold up to  $3^i$  elements. Whenever a user attempts to push an element onto any full stack  $S_i$ , we first pop all the elements off  $S_i$  and push them onto stack  $S_{i+1}$  to make room. (Thus, if  $S_{i+1}$  is already full, we first recursively move all its members to  $S_{i+2}$ .) Moving a single element from one stack to the next takes  $O(1)$  time.

- (a) In the worst case, how long does it take to push one more element onto a multi-stack containing  $n$  elements?
- (b) Prove that the amortized cost of a push operation is  $O(\log n)$ , where  $n$  is the maximum number of elements in the multistack.

**5. cyclic pattern matching**

Give a linear time algorithm to determine whether a text  $T$  is a cyclic rotation of another string  $T'$ . For example, *arc* and *car* are cyclic rotations of each other.

**6. A short and clean code for Decrease-key in Fibonacci Heap**

Write a neat pseudocode for the Decrease-key( $H, x$ ) in a Fibonacci Heap ?

**7. Delete-key in a Fibonacci heap**

Design an efficient algorithm for deleting an element from a Fibonacci Heap. The amortized cost must be  $O(\log n)$ .

**8. A surprising property for Fibonacci Heap**

Let  $v$  be any node in a Fibonacci heap. We showed that if the size of the subtree rooted at  $v$  is  $m$ , then the degree of  $v$  is  $O(\log m)$ . Can we say the same thing about the height as well ? That is, will the height of  $v$  be bounded by  $O(\log m)$  ? Note that all operations, including merging of Fibonacci heaps is allowed.

**Hint:** There exists a sequence of operations that may result in a Fibonacci heap which will be a single tree that is just a vertical chain of  $m$  elements. Invent one such sequence.