

Shortcourse on

Next Generation Database Systems

11-12 February 1988

Portland State University

Organized by
Oregon Database Forum

Sponsored by
Oregon Center for Advanced Technology Education (OCATE)

In Cooperation with
Willamette Valley ACM



Schedule

Thursday, 11 February

9:00-10:00	Introduction: Maier
10:00-10:30	Break
10:30-12:00	Carey
12:00-1:00	Lunch
1:00-2:00	Carey
2:00-3:00	Tolbert
3:00-3:30	Break
3:30-5:00	Tolbert

Friday, 12 February

9:00-10:30	Stonebraker
10:30-11:00	Break
11:00-12:00	Stonebraker
12:00-1:00	Lunch
1:00-2:30	Zdonik
2:30-3:00	Break
3:00-4:00	Zdonik
4:00-4:30	Reception
4:30-5:30	Shootout at the OCATE Corral

<u>Speaker/Affiliation</u>	<u>Topic/Project</u>	<u>Contributing Technology</u>
Mike Carey U. Wisconsin	Extensible Database Systems Exodus	SE: Common Interfaces
Doug Tolbert Unisys	Semantic Database Systems SIM	KR: Semantic Data Models
Mike Stonebraker UC Berkeley	Enhanced Relational Systems Postgres	AI: Rules
Stan Zdonik Brown	Object-Oriented Databases Encore/Observer	PL: ADTs & Encaps.

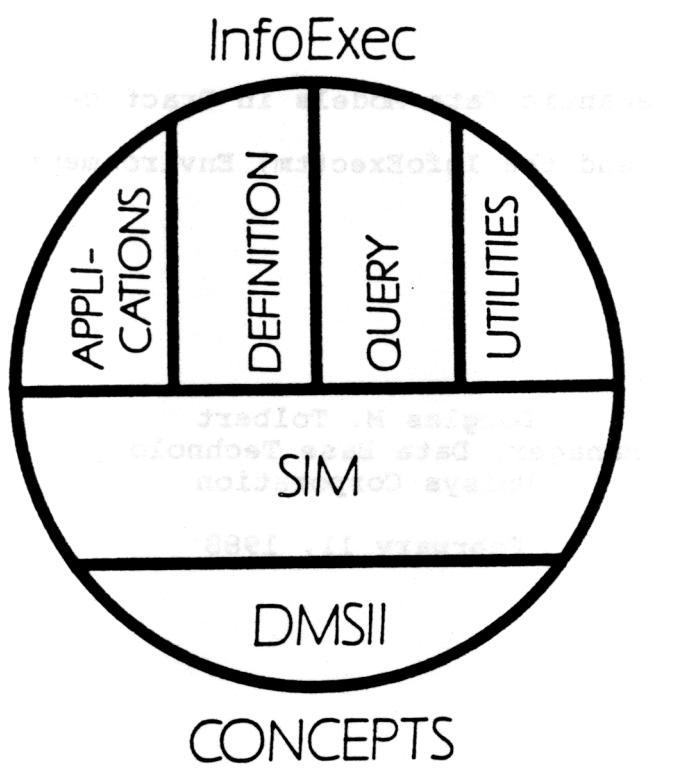
Semantic Data Models in Practice

SIM and the InfoExec(tm) Environment

Douglas M. Tolbert
Manager, Data Base Technology
Unisys Corporation

February 11, 1988

InfoExec is a trademark of Unisys Corporation



InfoExec Overview

Objectives

Data Models

File Systems

Database Systems

Semantic Systems

SIM Concepts

An Example

Application development effort

Query formation

InfoExec Product Environment

SIM Architecture and Implementation

Objectives

Create an environment that minimizes effort required to manage data by providing

Ease of Use

Productivity

Data Integrity

Coexistence

Performance

Objectives

Ease of Use

High level interface

Integrated interface

Flexibility

Data independence

Objectives

Productivity

Set-oriented, non-procedural interface

Naturally represent complex relationships

Easy to learn

Minimum programming effort

Data retrieval without programming

Objectives

Data Integrity

System enforced integrity

Shared data definitions

Referential integrity

Objectives

Coexistence

No impact on existing DMSII data bases

New capabilities for existing data bases

Performance

Production level

Data Model

A way of describing real-world application systems

Formalized into four components

Objects: Concepts to organize data

Operators: Instructions to manipulate data

Constraints: Rules to ensure correctness

Methodology: Instructions for designing "good"
data bases.

File Systems

Data Model

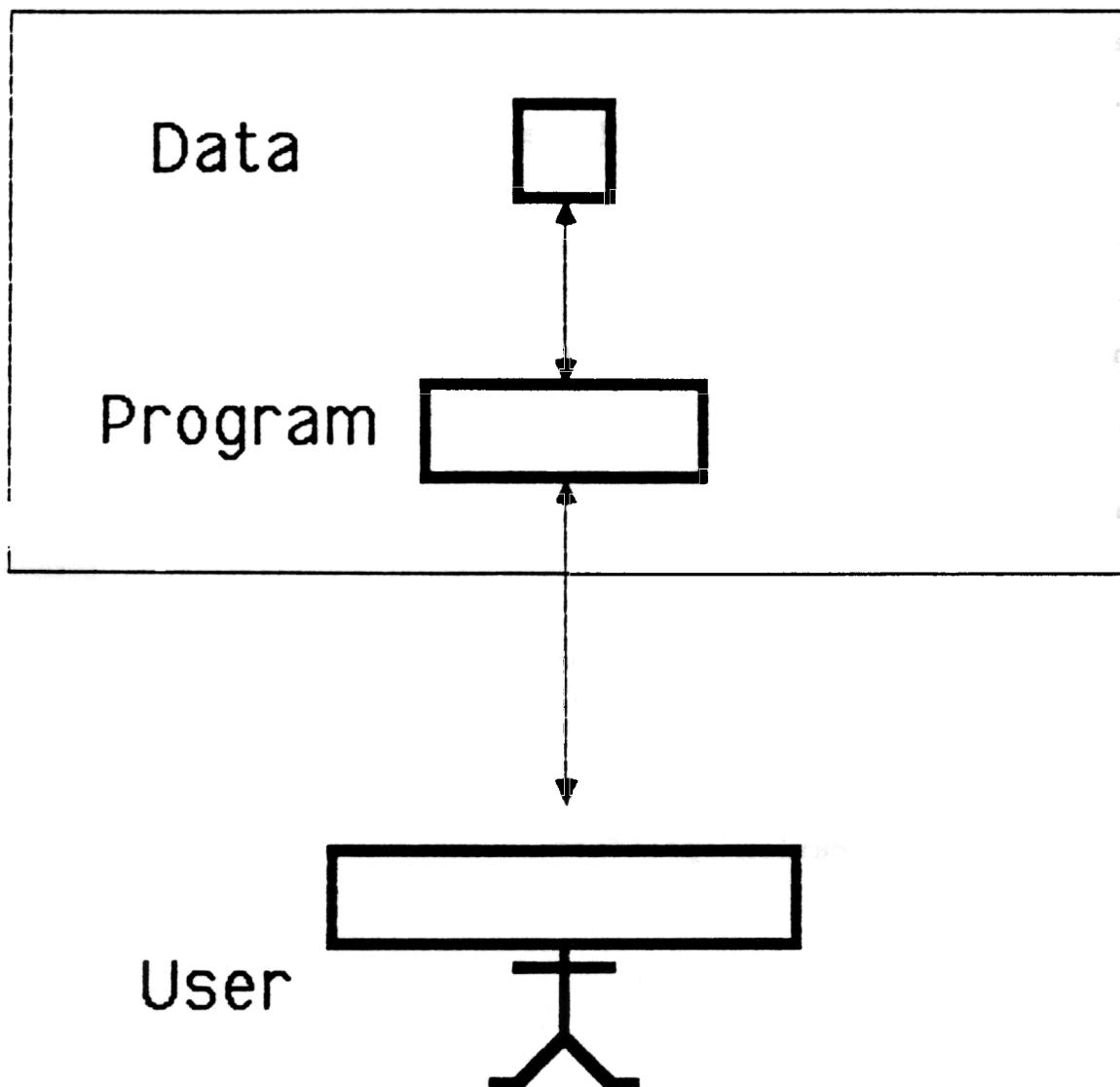
Objects: Records and files
Operators: Primitive (Read and Write)
Constraints: Almost none (parity?)
Methodology: Good programming practices

Knowledge Location

Data: I/O techniques
Program: Data item integrity
Simple relationships
Complex relationships
User: Complex relationships
Very complex relationships
Referential integrity
Resource sharing (concurrency)
Operational events (recovery)
Application specifics

File System

Computer



Database Systems

Data Model

Objects: Records (Network)
Tuples and Tables (Relational)

Operations: Navigational (Network)
Non-procedural (Relational)

Constraints: Limited

Methodology: Good programming practices (Network)
Normalization (Relational)

Knowledge Location

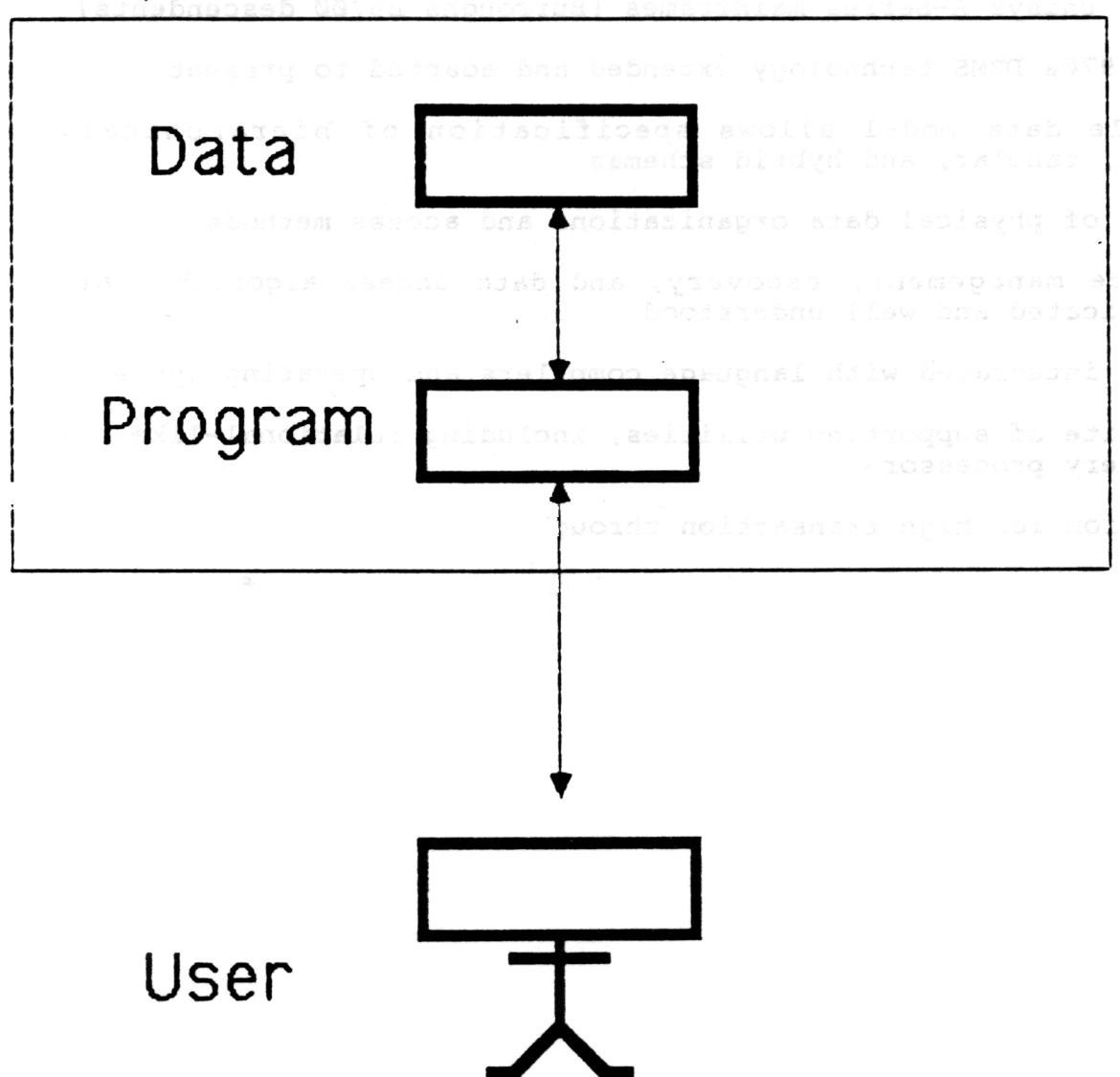
Data: Data item integrity
Simple relationships (limited in Relational)
Resource sharing
Operational events

Program: Complex relationships
Referential integrity
Application specifics

User: Very complex relationships
Referential integrity
Application specifics

Database System

Computer



DMS II

Runs on Unisys A-Series mainframes (Burroughs B6700 descendants)

Early 1970s DBMS technology extended and adapted to present

Flexible data model allows specification of hierarchical network, tabular, and hybrid schemas

Variety of physical data organizations and access methods

Resource management, recovery, and data access algorithms are sophisticated and well understood

Closely integrated with language compilers and operating system

Full suite of supporting utilities, including relational-like on-line query processor

Reputation for high transaction throughput

High customer base penetration: about 95 percent of A-Series customers license DMS II

DMS II

Provides efficient basic data management functions

But data model and interfaces do not provide benefits of newer,
developing technologies

Our challenge

To provide benefits of new technology without destroying
investment in existing database application systems

Options

Enhance DMS II

Ease of Use can be improved

Productivity and Data Integrity without sacrificing Coexistence

Performance already near architectural limits

Not likely to expand customer base

Build a Relational system

Ease of Use would improve

Productivity would be marginally improved

Data Integrity very limited

Coexistence

Extensive rework of basic data management software

No migration path for many existing DMS II data bases

"Two data base" strategy not desirable

Performance questionable

Would be one of many when released

Options

Explore new "semantic" data models

Examined

Semantic Networks (Quillian, Brachman)

Entity-Relationship (Chen)

RM/T (Codd)

SDM (Hammer and McLeod)

DAPLEX (Shipman)

GORDAS (El-Masri)

Meets Ease of Use, Productivity, and Data Integrity objectives

Many new features can be adapted to existing databases without migration

Performance based on proven DMS II algorithms

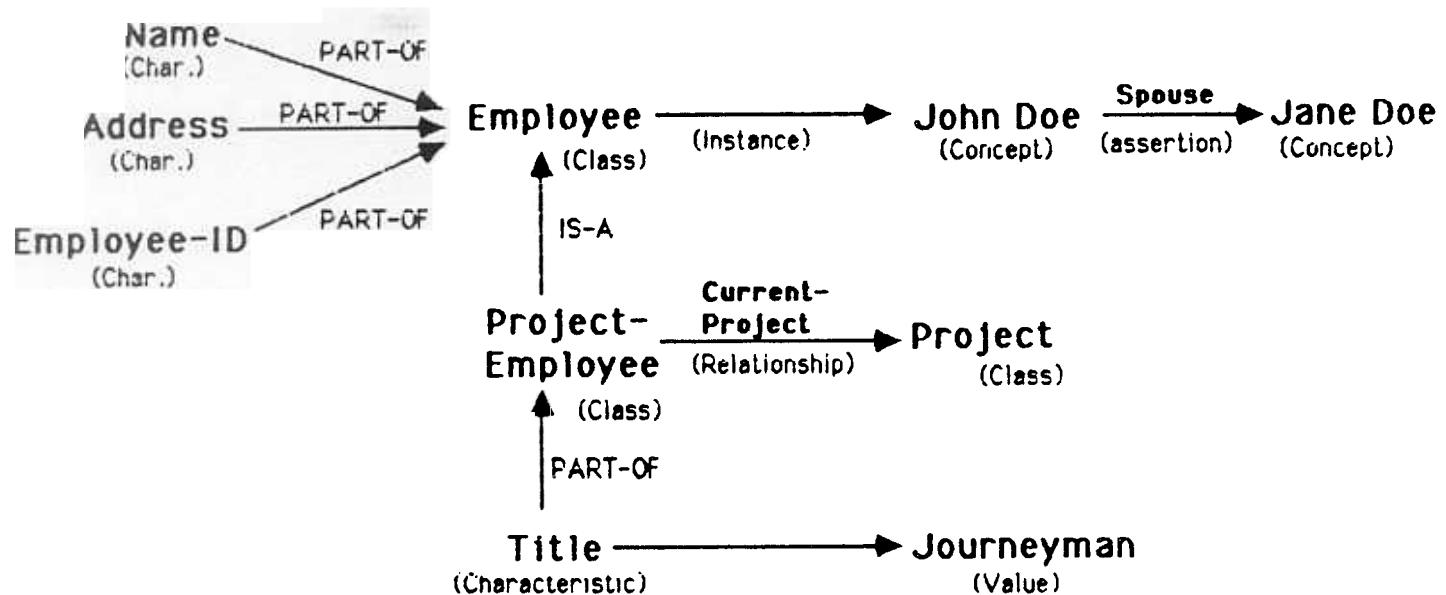
SDM provided

Clear conceptual basis

Best DMS II coexistence possibilities

Good vehicle for query language development

Semantic Net



Semantic Data Model References

Abiteboul, S., and R. Hull. 1987. IFO: A formal semantic database model. ACM Trans. Database Syst. 12, 525-565.

Brachman, R. J. 1979. "On the epistemological status of semantic networks," in Associative Networks (Findler, N., ed), pp 3-50. Academic Press, New York.

Chen, P. P. 1976. The entity-relationship model: Toward a unified view of data. ACM Trans. Database Syst. 1, 9-36.

Codd, E. F. 1979. Extending the database relational model to capture more meaning. ACM Trans. Database Syst. 4, 397-434.

El-Masri, R. A. 1981. GORDAS: A Data Definition, Query and Update Language for the Entity-Category-Relationship Model of Data. Honeywell Computer Science Technical Report HR-81-250.

Hammer, M., and D. McLeod. 1981. Database description with SDM: a semantic database model. ACM Trans. Database Syst. 6, 351-386.

Hull, R., and R. King. Semantic data modeling: Survey, applications, and research issues. ACM Comput. Surv. (to appear).

Quillian, M. R. 1968. "Semantic memory," in Semantic Information Processing (Minsky, M., ed), pp 227-270. MIT Press, Cambridge, MA.

Shipman, D. W. 1981. The functional data model and the data language DAPLEX. ACM Trans. Database Syst. 6, 140-173.

Smith, J. M., and D. C. P. Smith. 1977. Database Abstractions: Aggregation and Generalization. ACM Trans. Database Syst. 2, 105-133.

Stachowitz, R. A. 1985. A formal framework for describing and classifying semantic data models. Inform. Systems 10, 77-96.

Tsichritzis, D. C., and F. H. Lochovsky. 1982. Data Models. Prentice-Hall.

Tsur, S., and C. Zaniolo. 1984. An implementation of GEM-supporting a semantic data model on a relational back-end. Proc. ACM-SIGMOD Conference.

What makes a data model semantic?

"During the last few years numerous investigations have been aimed at capturing (in a reasonably formal way) more of the meaning of the data, while preserving independence of implementation. This activity is sometimes called semantic data modeling. Actually, the task of capturing the meaning of data is a never-ending one. So the label 'semantic' must not be interpreted in any absolute sense."

Codd, pp 397-398

That is, they attempt to capture the meaning of the data, not just physical data values.

So, the more "meaning" a data model captures, the more "semantic" it is.

How does a semantic data model capture meaning?

Abstract objects for modeling data, not physical storage containers

"... it is appropriate that the structure of a database mirror the structure of the system that it models. A database whose organization is based on naturally occurring structures will be easier for a database designer to construct and modify than one that forces him to translate the primitives of his problem domain into artificial specificaiton constructs."

Hammer and McLeod, pp 351-352

Abstraction techniques

Aggregation

Relationship between objects regarded as a higher level object ("PART_OF").

Generalization

Collection of individual objects viewed as single object ("IS_A").

How does a semantic data model capture meaning?

Database description allows expression of

Relationships between objects

one-to-one, one-to-many, many-to-many

generalization hierarchies

Limits on values that attributes may take

type mechanism (a la Pascal)

subrange enforcement

uniqueness

Cardinality constraints

required value

minimum and maximum instances

single- and multi-valued

Ad hoc constraints

non-structural constraints

Semantic Systems

Data Model

Objects: Entities (logical)

Operations: Non-procedural

Constraints: Attribute, Referential

Methodology: Application understand

Knowledge Location

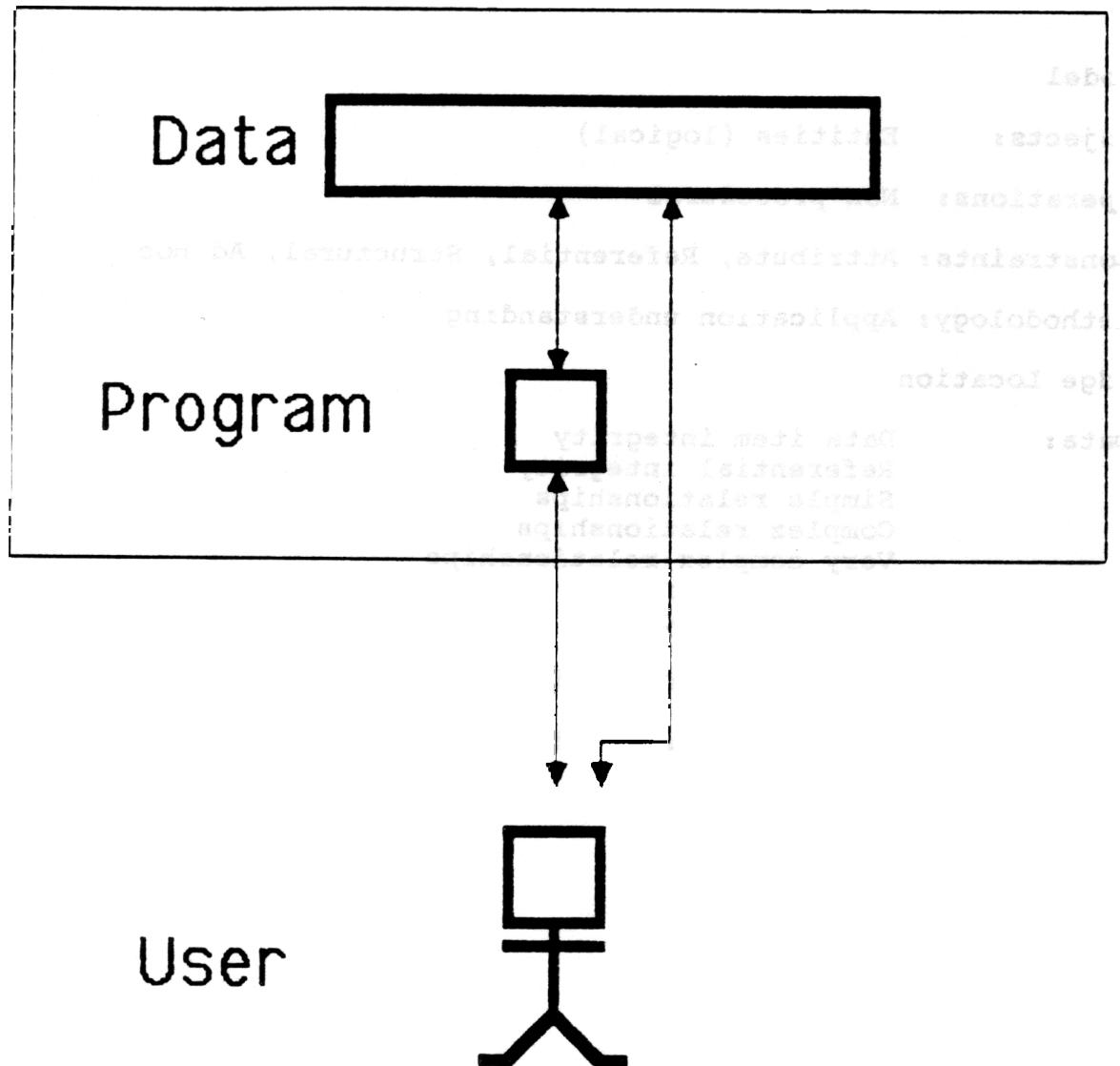
Data: Data item integrity
Referential integrity
Simple relationships
Complex relationships
Very complex relationships
Resource sharing
Operational events

Program: Application specifics

User: Application specifics

Semantic System

Computer



SIM Concepts

Entity

An object of interest in the application environment

For example: John Doe, an employee

Accounting, a department

Annual Report Preparation, a project

SIM Concepts

Attribute

A characteristic of an entity

For example: Employee John Doe has a name, address and employee ID.

SIM Concepts

Class

A collection of entities of the same type

Does not imply any specific physical implementation

For example: Employee, a collection of all employees
working for a company

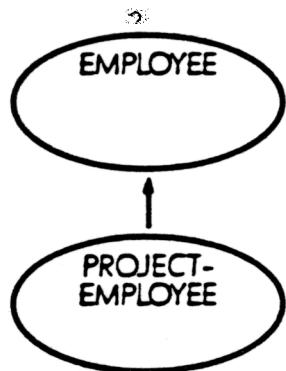


SIM Concepts

Subclass

A subset of entities in a class

For example: Project-Employee, all employees that work on projects



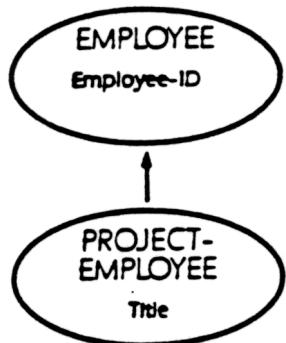
SIM Concepts

Attribute Inheritance

Subclasses inherit attributes from super classes

Additional attributes may be declared for subclasses

For example: Every employee has an employee ID,
but only project employees have titles

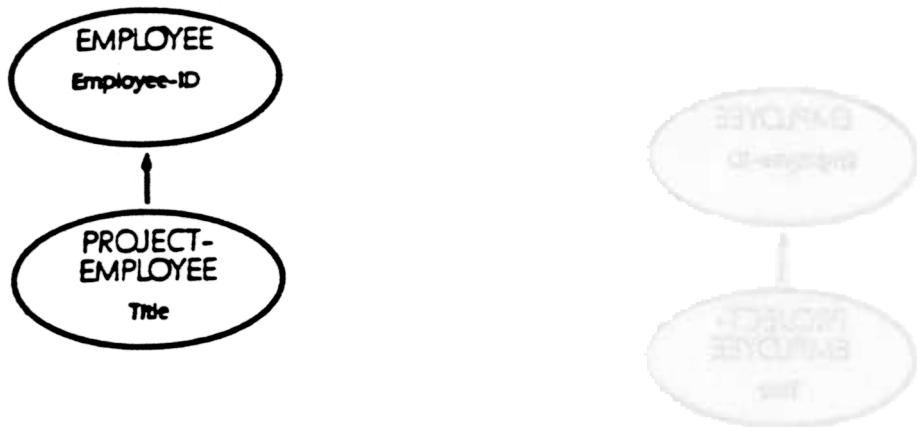


SIM Concepts

Data-Valued Attributes

An attribute whose values can be displayed

For example: An employee's ID number

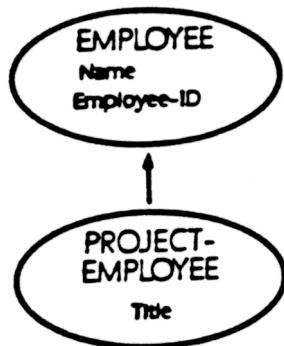


SIM Concepts

Compound Attributes

A collection of attributes that may be treated as a unit

For Example: An employee's name is a compound attribute made up of first name, middle initial, and last name.



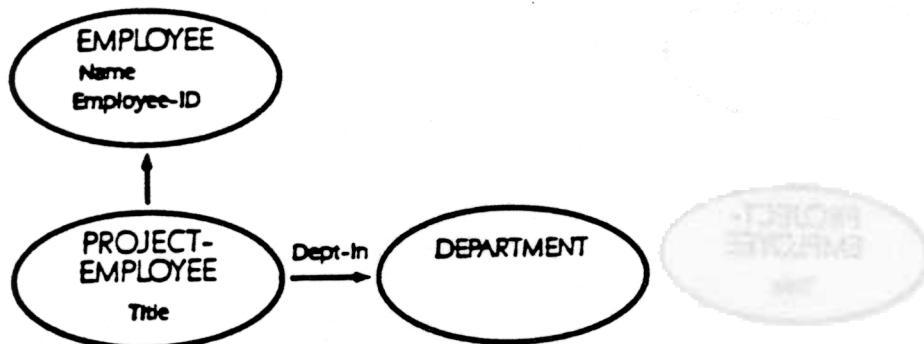
SIM Concepts

Entity-Valued Attributes

An attribute whose values are entities in a class

Indicate relationships between entities

For example: A project employee's department

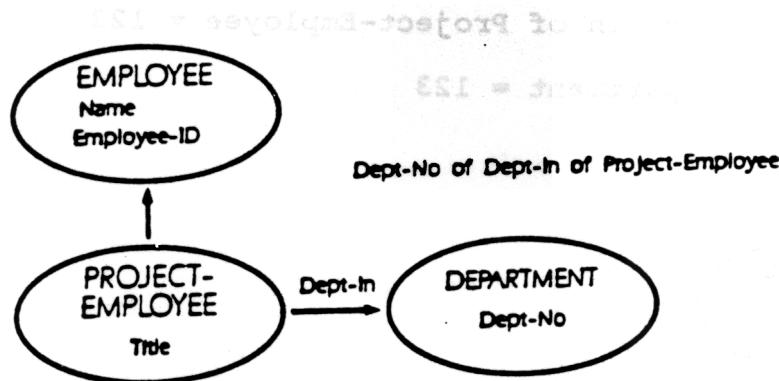


SIM Concepts

Extended Attributes

Attributes of related entities may be considered extended attributes of an entity

For example: The department number of a project employee



SIM Concepts

Perspective

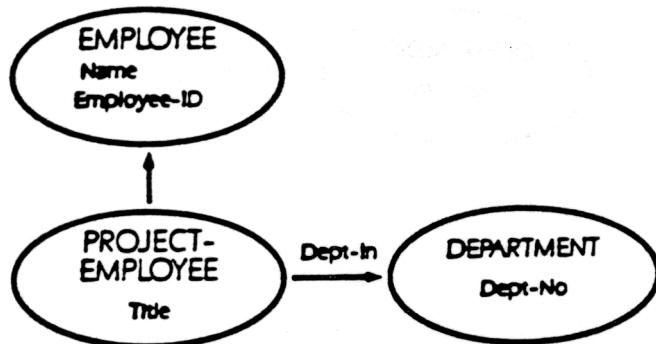
A point of view chosen for a query

Extended attributes used to view information elsewhere in the database

Same data viewed from different perspectives may have different meanings

Dept-No of Dept-In of Project-Employee = 123

Dept-No of Department = 123



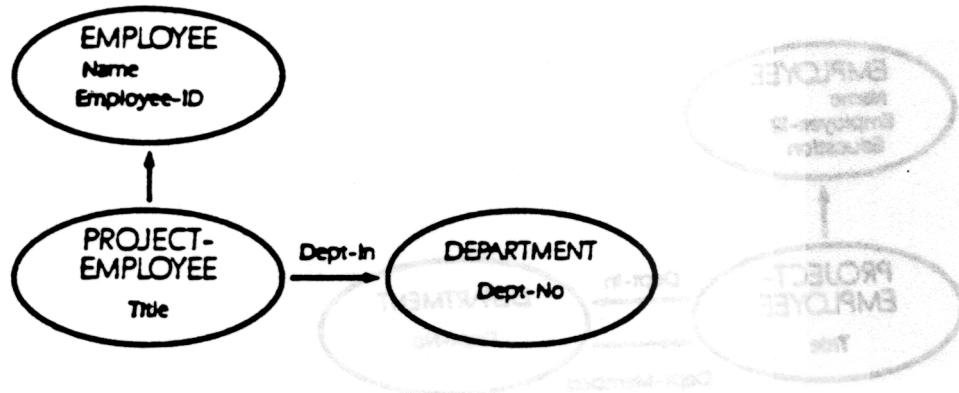
SIM Concepts

Single-Valued Attributes

An attribute that can have only one value for an entity

For example: An employee's ID number

A project employee's department



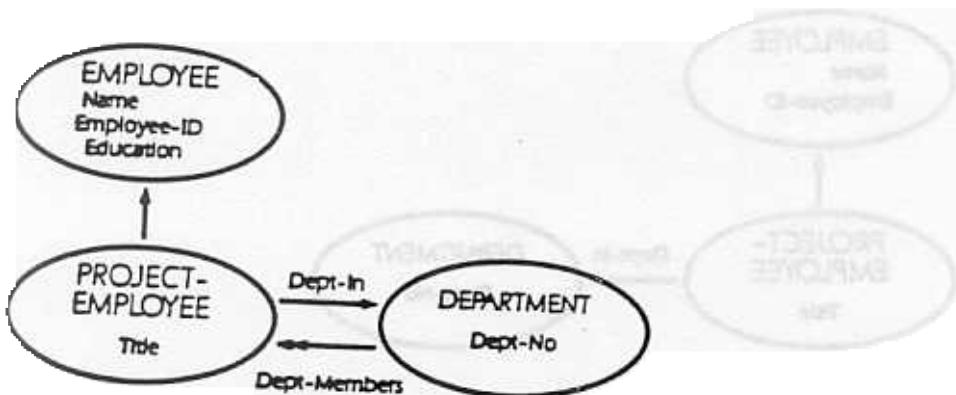
SIM Concepts

Multi-Valued Attributes

An attribute that may have more than one value for an entity

For example: An employee's education degrees

The project employees in a department

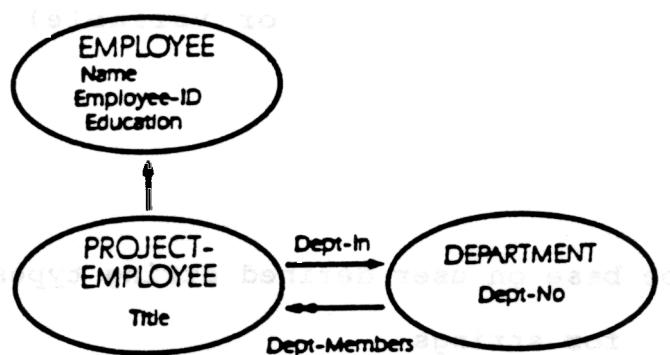


SIM Concepts

Bidirectional Relationships

Inverse relationship between two entity-valued attributes

For example: The project employees in a department, and
The members of a department



SIM Concepts

Types

Similar to Pascal

String type checking relaxed

More widely understood than SDM class-based domains

System Defined

Integer, Real, Boolean, Date, Time, Character, Number
(packed decimal), String (fixed or variable), Kanji,
Symbolic, Ordered Symbolic

User Defined

Based on system defined types or other user defined types

Strings may be base on user defined string types

Set membership for strings

Enforced subranges

Employee-Age : Integer (18..70)

SIM Concepts

Class Attributes

Belong to class as a whole, not any particular entity

One value per class

Data-valued only

SIM Concepts

Generalized Verify

General constraints not related to schema structure

Allows user specified error messages

For example, the spouse of an employee hired after December 31, 1987 may not work for the company

```
VERIFY NoCouples ON Employee  
ASSERT NOT (Spouse ISA Employee)  
WHERE Employee-Hire-Date > 12/31/87  
ELSE "Spouse may not work for company"
```

available on first release

SIM Concepts

Security

Access limits

Visibility of attributes

Retrieval and update operations

Permission associates Access with user or program

Example, the Accounting department may see but not
change the salary of journeyman and higher
employees

```
ACCESS LookButDontTouch ON Employee
      (Employee-Salary) RETRIEVE
WHERE Employee-Status >= Journeyman
```

PERMISSION

```
USERCODE = Accounting,
ACCESS = LookButDontTouch
```

SIM Concepts

Index

Visible only to database administrator and SIM Optimizer

Not visible to programmers or query users

Multiple data-valued attributes, ascending, and descending

Used for performance improvement

For Example,

```
INDEX (Employee-ID) ON Employee ASCENDING
```

ORGANIZATION -- An Example

A projects and employees data base.

Data base description

DMSII A record oriented system

DB2 A relational system

SIM A semantic system

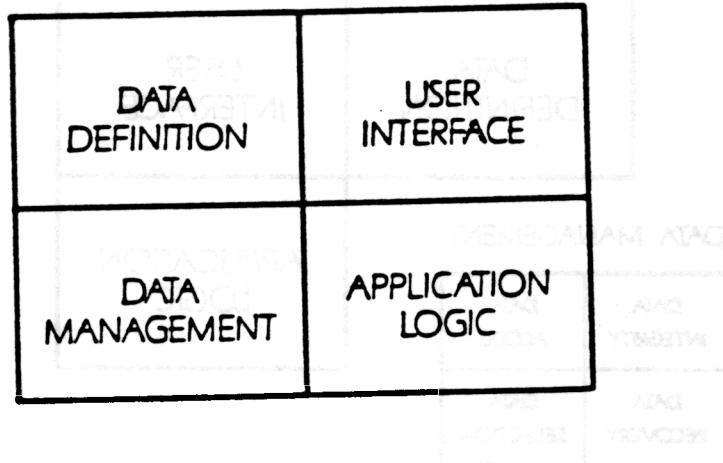
Application Responsibilities

Query Formation

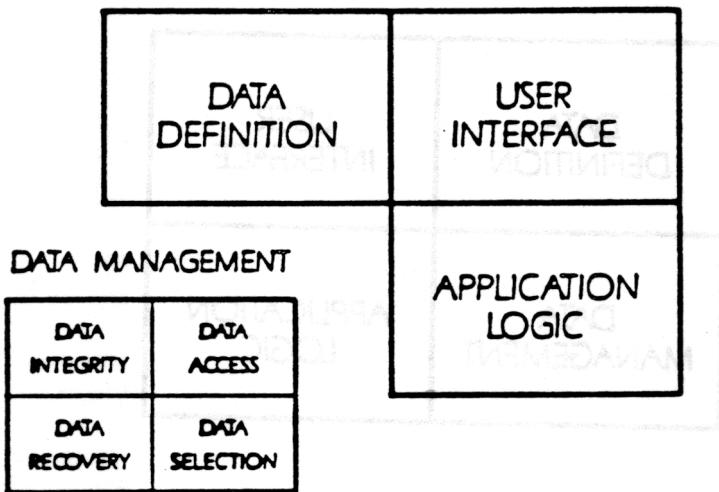
ORGANIZATION Entities

People	Projects
Employees	Assignments
Project Employees	Departments
Managers	Previous Employees
Interim-Managers	

APPLICATION PROGRAM



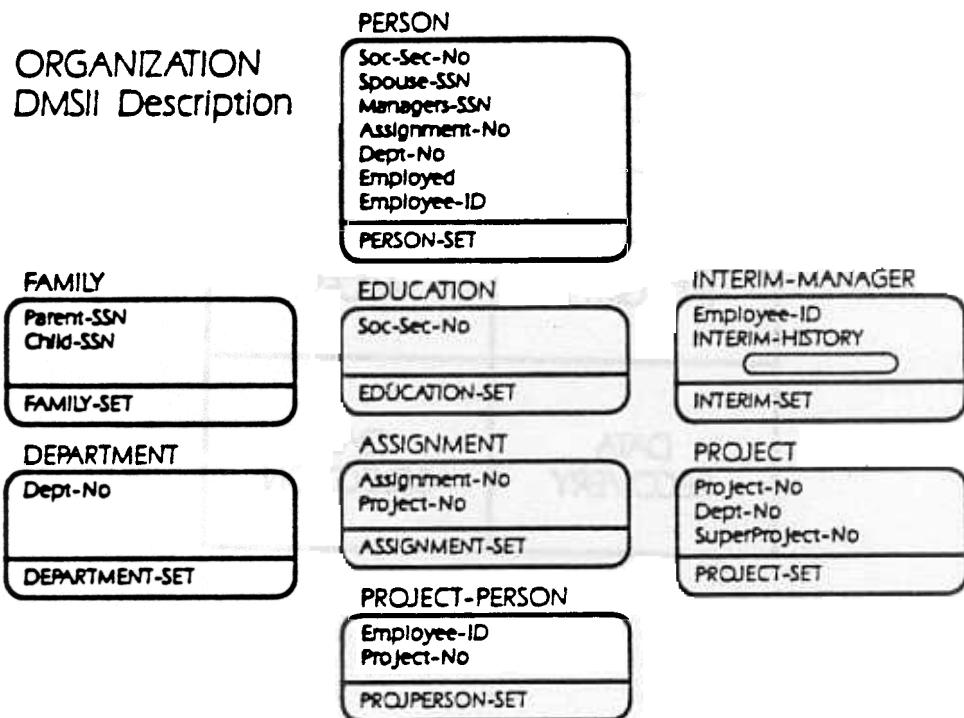
APPLICATION PROGRAM



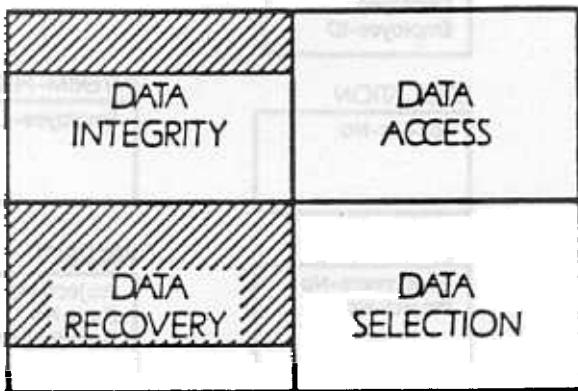
DATA MANAGEMENT

DATA INTEGRITY	DATA ACCESS
DATA RECOVERY	DATA SELECTION

ORGANIZATION
DMSII Description



DMSII APPLICATION



ORGANIZATION
DB2 DESCRIPTION

PERSON

Soc-Sec-No
Spouse-SSN
Managers-SSN
Assignment-No
Dept-No
Employed
Employee-ID

FAMILY

Parent-SSN
Child-SSN

DEPARTMENT

Dept-No

EDUCATION

Soc-Sec-No

INTERIM-HISTORY

Employee-ID

ASSIGNMENT

Assignment-No
Project-No

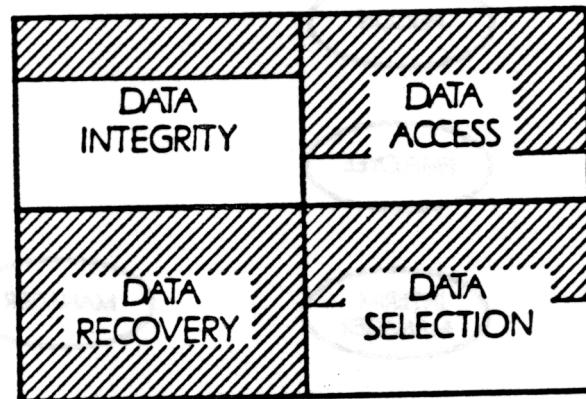
PROJECT

Project-No
Dept-No
SuperProject-No

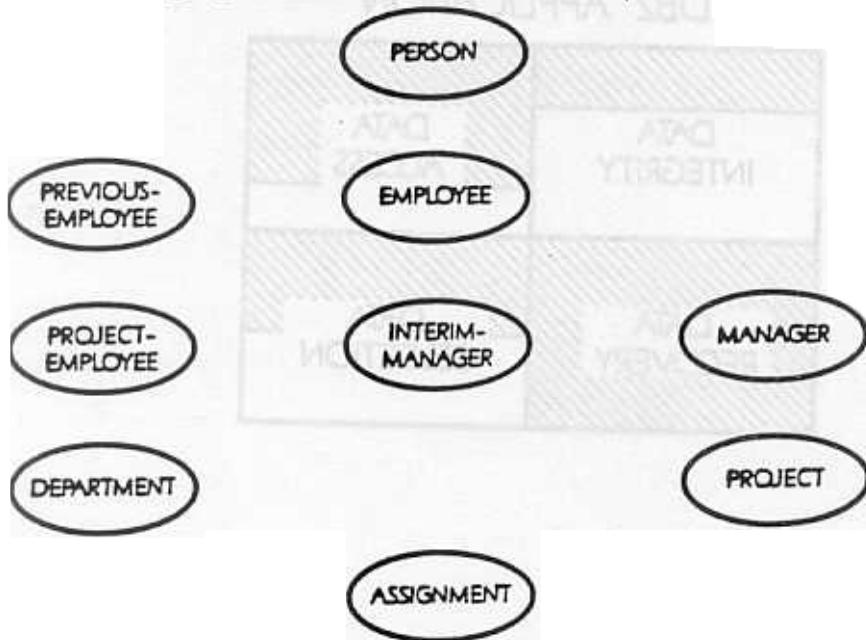
PROJECT-PERSON

Employee-ID
Project-No

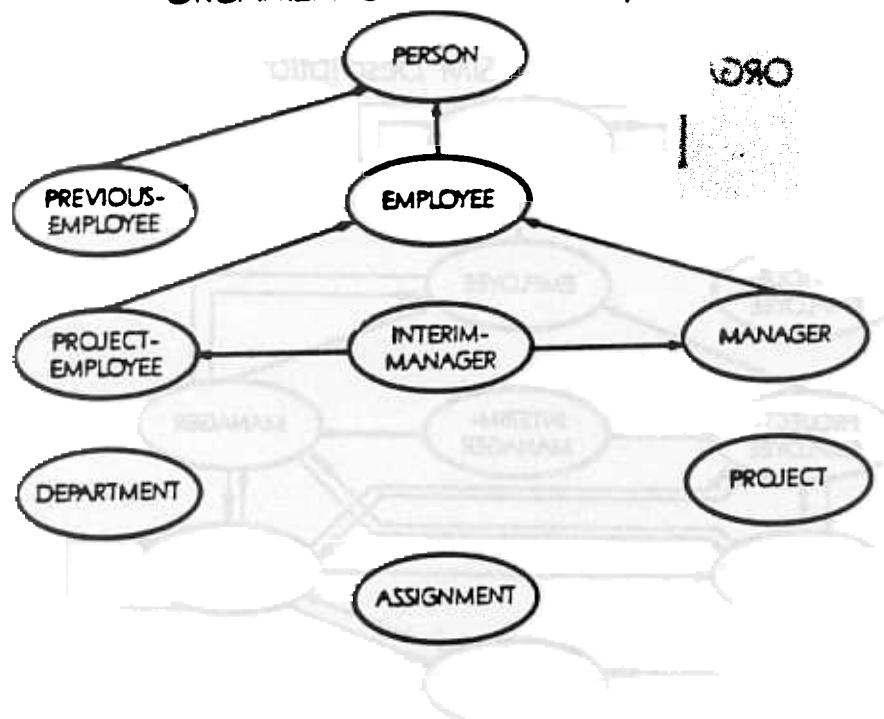
DB2 APPLICATION



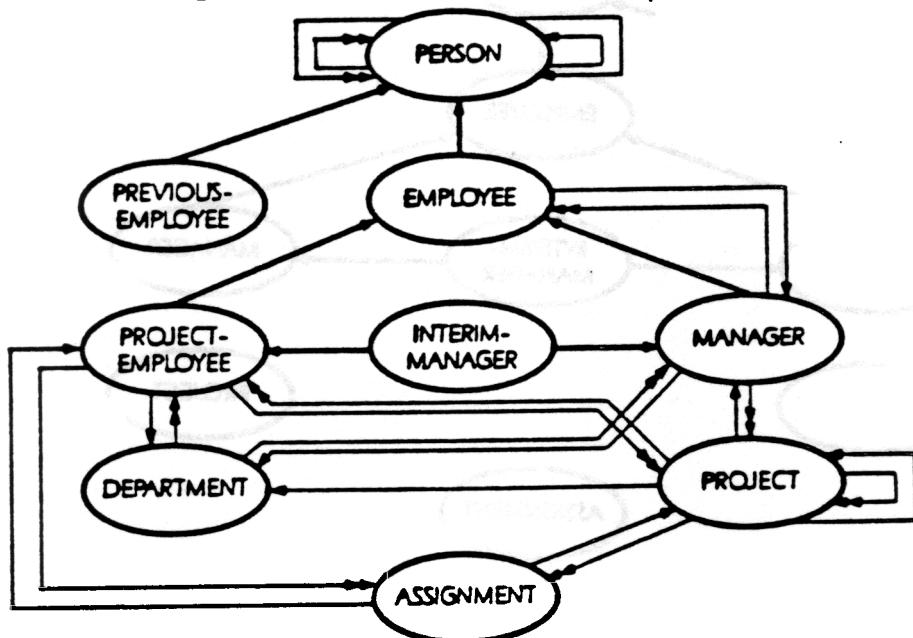
ORGANIZATION - SIM Description



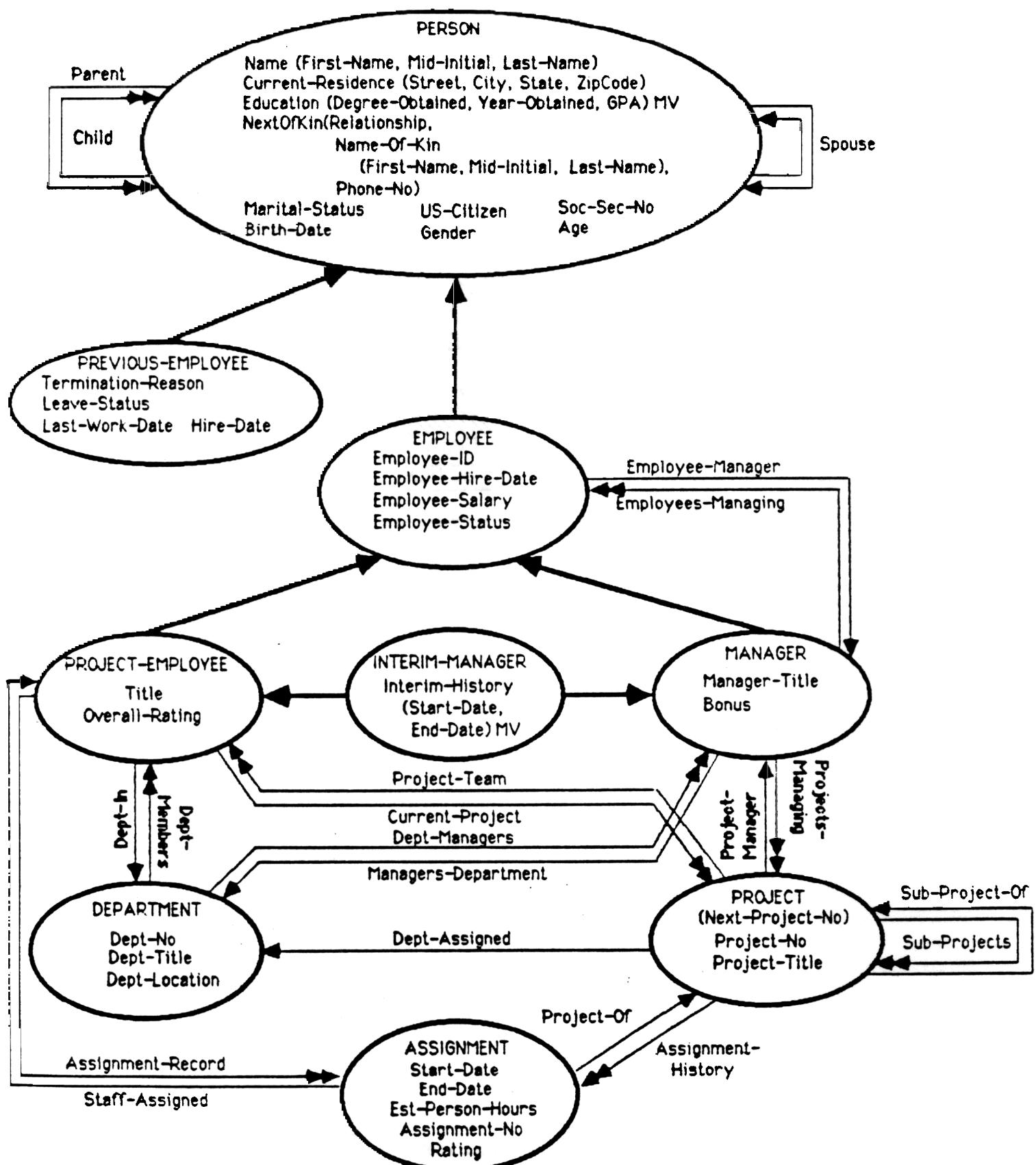
ORGANIZATION - SIM Description



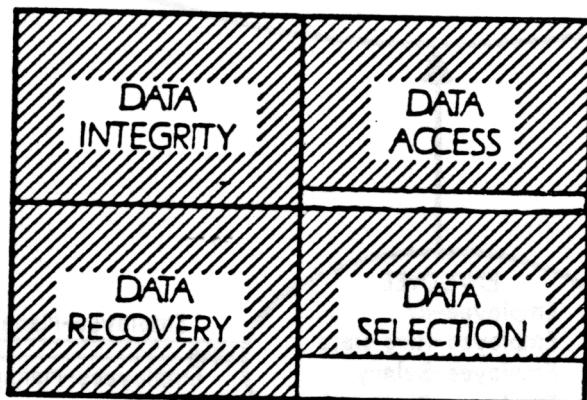
ORGANIZATION - SIM Description



Organization SIM Description



SIM APPLICATION



ORGANIZATION Data Integrity Enforcement

System	Data Item Validation		Referential Integrity	
	DBMS	Application	DBMS	Application
DMSII	77	25	4	86
	66	36	0	90
	102	0	90	0

ORGANIZATION COBOL Source Lines

Operation	DMS II	SIM	SIM / DMS II
Population DB	482	227	0.47
Retrieve 1	49	44	0.90
Retrieve 2	79	42	0.53
Modify 1	43	12	0.28
Modify 2	76	12	0.16
Delete 1	122	12	0.10
Delete 2	90	8	0.09

Query Examples

Simple Retrieval

Print the names of all non-managers and the title of the department in which they work

Complex Retrieval

Print the names of employees and the titles of all their projects if they work on any project assigned to the Accounting Department

More Complex Retrieval

For the Annual Report Preparation project, print the titles of its subprojects and the names of employees currently assigned

Simple Query

Print the names of all non-managers and the title of the department in which they work

DMSII

RELATE Person TO Department BY MATCHING Dept-No WITH Dept-No
AS Person-Dept;

TAB First-Name, Mid-Initial, Last-Name, Dept-Title
WHERE Employed NEQ "Manager"
FROM Person TO Department;

Simple Query

Print the names of all non-managers and the title of the department in which they work

DB2

```
SELECT First-Name, Mid-Initial, Last-Name, Dept-Title  
      FROM Person, Department  
     WHERE Person.Dept-No = Department.Dept-No AND  
           Person.Employed NEQ "Manager";
```

Simple Query

Print the names of all non-managers and the title of the department in which they work

SIM

**RETRIEVE Name of Project-Employee,
Dept-Title of Dept-In**

Complex Query

Print the names of employees and the titles of all their projects if they work on any project assigned to the Accounting Department

DMSII

```
RELATE Department TO Project BY MATCHING Dept-No WITH Dept-No  
AS Dept-Proj;
```

```
RELATE Project TO Project-Person BY MATCHING Project-No WITH  
Project-No AS Proj-Person;
```

```
RELATE Project-Person TO Person BY MATCHING Soc-Sec-No WITH  
Soc-Sec-No AS Proj-Emp;
```

```
EXTRACT Soc-Sec-No, COUNT AS SSN-Cnt WHERE Dept-Title =  
"Accounting" FROM Person TO Project-Person TO Project TO  
Department : EXTRACTFILE = Extfile;
```

```
OPEN FILE Extfile;
```

```
RELATE Extfile TO Project-Person BY MATCHING Soc-Sec-No WITH  
Soc-Sec-No AS Extrel;
```

```
TAB First-Name, Mid-Initial, Last-Name, Project-Title  
FROM Extfile TO Project-Person TO Project  
WHERE SSN-Cnt > 0;
```

Complex Query

Print the names of employees and the titles of all their projects if they work on any project assigned to the Accounting Department

DB2

```
SELECT First-Name, Mid-Initial, Last-Name, Project-Title
      FROM Person, Project-Person, Project
 WHERE Person.Soc-Sec-No = Project-Person.Soc-Sec-No
   AND Project.Project-No = Project-Person.Project-No
   AND EXISTS
   ( SELECT *
       FROM Project-Person, Project, Department
     WHERE Project-Person.Soc-Sec-No = Person.Soc-Sec-No
       AND Project-Person.Project-No =
             Project.Project-No
     AND Department.Dept-No = Project.Dept-No
     AND Department.Dept-Title = "Accounting")
```

Complex Query

Print the names of employees and the titles of all their projects if they work on any project assigned to the Accounting Department

SIM

```
RETRIEVE Name of Project-Employee,  
          Project-Title of Current-Project  
WHERE Dept-Title of SOME (Dept-Assigned of Current-Project)  
      = "Accounting"
```

More Complex Query

For the Annual Report Preparation project, print the titles of its subprojects and the names of employees currently assigned

DMSII

```
OPEN DMSII Organization;  
OPEN DMSII OrgCopy (DMI = DMINTERPRETER/ORGANIZATION);  
RELATE Project of Organization TO Project of OrgCopy BY  
MATCHING Project-No WITH SuperProject-No AS SubProj;  
RELATE Project of Organization TO Person of OrgCopy  
MATCHING Project-No WITH Project-No AS Project-Emp;  
RELATE Project-Person of OrgCopy TO Person of OrgCopy BY  
MATCHING Employee-ID WITH Employee-ID AS Proj-Person;  
TAB Project-Title of Project of OrgCopy, First-Name of  
Person of OrgCopy, Mid-Initial of Person of  
OrgCopy, Last-Name of Person of OrgCopy  
FROM Project of Organization TO Project of OrgCopy TO  
Project-Person of OrgCopy TO Person of OrgCopy  
WHERE Project-Title of Project of Organization =  
"Annual Report Preparation";
```

More Complex Query

For the Annual Report Preparation project, print the titles of its subprojects and the names of employees currently assigned

DB2

```
SELECT Project-Title, First-Name, Mid-Initial, Last-Name  
      FROM Project, Project SubProj, Person, Project-Person  
     WHERE Project.Project-No = SubProj.SuperProject-No  
       AND SubProj.Project-No =  
             Project-Person.Project_No  
       AND Project-Person.Soc-Sec-No = Person.Soc-Sec-No  
       AND Project.Project-Title = "Annual Report Preparation"
```

More Complex Query

For the Annual Report Preparation project, print the titles of its subprojects and the names of employees currently assigned

SIM

```
RETRIEVE Project-Title of Sub-Projects of Project,  
          Name of Project-Team of Sub-Projects  
WHERE Project-Title of Project = "Annual Report Preparation"
```

Host Language Interface

High-level interface to full SIM functionality

Extended language grammar

Transaction orientation

Full data independence

Structured, Tabular, Hybrid retrievals

Language statements may be interleaved with SIM constructs

Language variables allowed in queries

Host Language Interface

Hybrid Retrieval, COBOL

```
QD DEPARTMENT-QUERY.  
 01 DEPARTMENT-RECORD.  
    02 DEPT-TITLE PIC X(20).  
  
QD MANAGER-QUERY.  
 01 MANAGER-RECORD.  
    02 LAST-NAME PIC X(20).  
    02 PROJECT-TITLE PIC X(30).
```

Host Language Interface

Hybrid Retrieval, COBOL

```
SELECT DEPARTMENT-QUERY FROM DEPARTMENT
  (DEPT-TITLE,
   SELECT MANAGER-QUERY FROM DEPT-MANAGERS
     (LAST-NAME = LAST-NAME OF NAME,
      PROJECT-TITLE = PROJECT-TITLE
        OF PROJECTS-MANAGING)).
```

RETRIEVE DEPARTMENT-QUERY.

RETRIEVE MANAGER-QUERY.

Host Language Interface

Hybrid Retrieval, Sample Data

Departments	Managers	Projects
-----	-----	-----
Accounting		Payroll
	Adams	Year-end Statements
	Adams	Accounts Payable
Engineering	Burns	
	Duncan	Maintenance
	Eaton	New Products
Purchasing	Eaton	
	Carr	Inventory

Update: Insert

Create a new employee named John Doe and assign him to the manager named Smith

```
INSERT Employee
  (Name := (First-Name := "John", Last-Name := "Doe"),
   Gender := Male,
   Soc-Sec-No := 123-45-6789,
   US-Citizen := True,
   Employee-ID := 46060,
   Spouse := Person WITH
     (First-Name of Name = "Mary" AND
      Last-Name of Name = "Doe"),
   Child := INCLUDE Person WITH
     (First-Name of Name = "Junior" AND
      Last-Name of Name = "Doe"),
   Employee-Manager := Manager WITH
     (Last-Name of Name = "Smith")
  )
```

Promote John Doe to a department manager with a bonus of \$5000

```
INSERT Manager FROM Employee
  WHERE Last-Name of Name = "Doe"
  (Manager-Title := Department-Manager,
   Bonus := 5000
  )
```

Update: Modify

Reassign all journeyman Project-Employees in the Construction Department to the Maintenance Department

```
MODIFY Project-Employee
  (Dept-In := Department WITH
   (Dept-Title = "Maintenance"))
WHERE Title = Journeyman AND
  Dept-Title of Dept-In = "Construction"
```

Update: Delete

Remove all projects assigned to managers that manage departments located in Los Angeles

```
DELETE Project
WHERE Dept-Location of Managers-Department
      of Project-Manager = "Los Angeles"
```

Expressions

Allowed in target list and where expression
Automatic null handling (tri-state logic)

Operators

- Arithmetic (+, -, *, /, DIV, MOD, **)
- Boolean (NOT, AND, OR)
- Relational (<, >, =, <=, >=, <>)
- String (&)
- Existence (EXISTS)

Functions

- Arithmetic (ABS, ROUND, TRUNC, SQRT)
- String (LENGTH, EXT, POS, RPT)
- Symbolic (PRED, SUCC)
- Date (YEAR, MONTH, DAY, ELAPSED_DAYS, ADD_DAYS, DAY_OF_WEEK,
MONTH_NAME, CURRENT_DATE)
- Time (HOUR, MINUTE, SECOND, ELAPSED_TIME, ADD_TIME,
CURRENT_TIME)

INVERSE

TRANSITIVE

Pattern matching

Multi-valued Expressions

Operators

INCLUDE

EXCLUDE

Functions

Aggregate (AVG, SUM, COUNT, MIN, MAX)

Quantifiers (SOME, ALL, NO)

~~Advanced Query Topics~~

Subrole attribute

Read-only enumeration of subclasses of a class

May be used in queries

Retrieve the names of employees that are managers

```
RETRIEVE Name of Employee  
WHERE Profession = Manager
```

Role testing

Retrieve the names of employees whose spouses are managers

```
RETRIEVE Name of Employee  
WHERE Spouse ISA Manager
```

Role qualification

Retrieve the names of US citizens and the employee ID of their spouses

```
RETRIEVE Name of Person, Employee-ID of Spouse AS Employee  
WHERE US-Citizen
```

Advanced Query Topics

Multiple perspective queries: Value-based joins

Retrieve the names of project employees and managers that are the same age

```
RETRIEVE Name of Project-Employee,  
          Name of Manager  
WHERE Age of Project-Employee = Age of Manager
```

Reference variables

Created implicitly by SIM for aggregate functions and quantifiers

Created explicitly in query

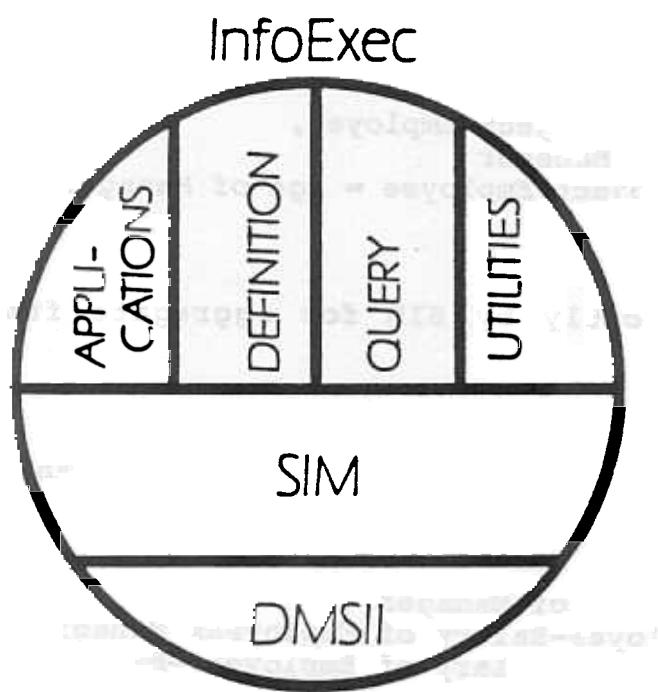
Retrieve the names of all managers who manage employees making more than \$40,000 and employees making less than \$20,000

```
RETRIEVE Name of Manager  
WHERE Employee-Salary of Employees-Managing > 40000  
      AND Employee-Salary of Employees-Managing  
           CALLED OtherEmp < 20000
```

Local Selection

Retrieve the names of managers of all departments and the salaries of only the division managers

```
RETRIEVE Name of Dept-Managers of Department  
          Employee-Salary of Dept-Managers  
          WITH (Manager-Title of Dept-Managers  
                 Division-Manager)
```



PRODUCT OVERVIEW

InfoExec Environment

General

Single, screen-based environment

Seamless, function-oriented screen flow

Screen flows can be record and played back later

Multi-lingual support

All new documentation -- user oriented

Classes available from Joseph & Cogan Associates

objeto

InfoExec Environment

Applications

Host Language Interfaces

COBOL74, Pascal, ALGOL

Definition

Advanced Data Dictionary System (ADDS)

SIM, DMS II, COBOL74, Screen Formats, Saved Queries

Logical entities: Program, Process, Keyword

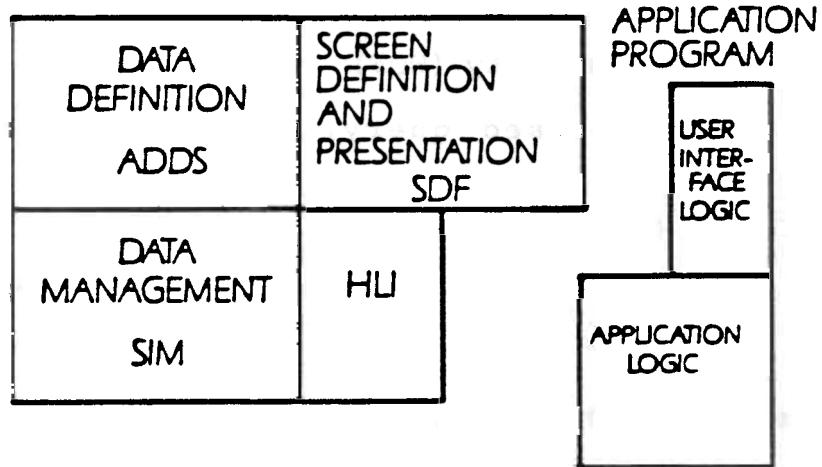
Usage tracking and reporting

A SIM database application

Screen Design Facility (SDF)

Screen definition and manipulation outside user applicaiton program

InfoExec ENVIRONMENT



InfoExec Environment

Query

Interactive Query Facility (IQF)

Host-based query, update, browsing, and reporting

Screen and menu based

Workstation Query Facility (WQF)

Workstation-based query, update, browsing, and reporting

Mouse and window based

Unisys B25 and PC families, IBM PC compatibles

Utilities

Operations Control Manager (OCM)

Screen-based operations for SIM and DMS II

Dictionary Utilities

Database schema management and generation

Dictionary management and security

COBOL74 program loader

DMS.View and LINC.View

Inquiry-only SIM access to existing DMS II databases and LINC systems

SIM Architecture

Goals

Performance

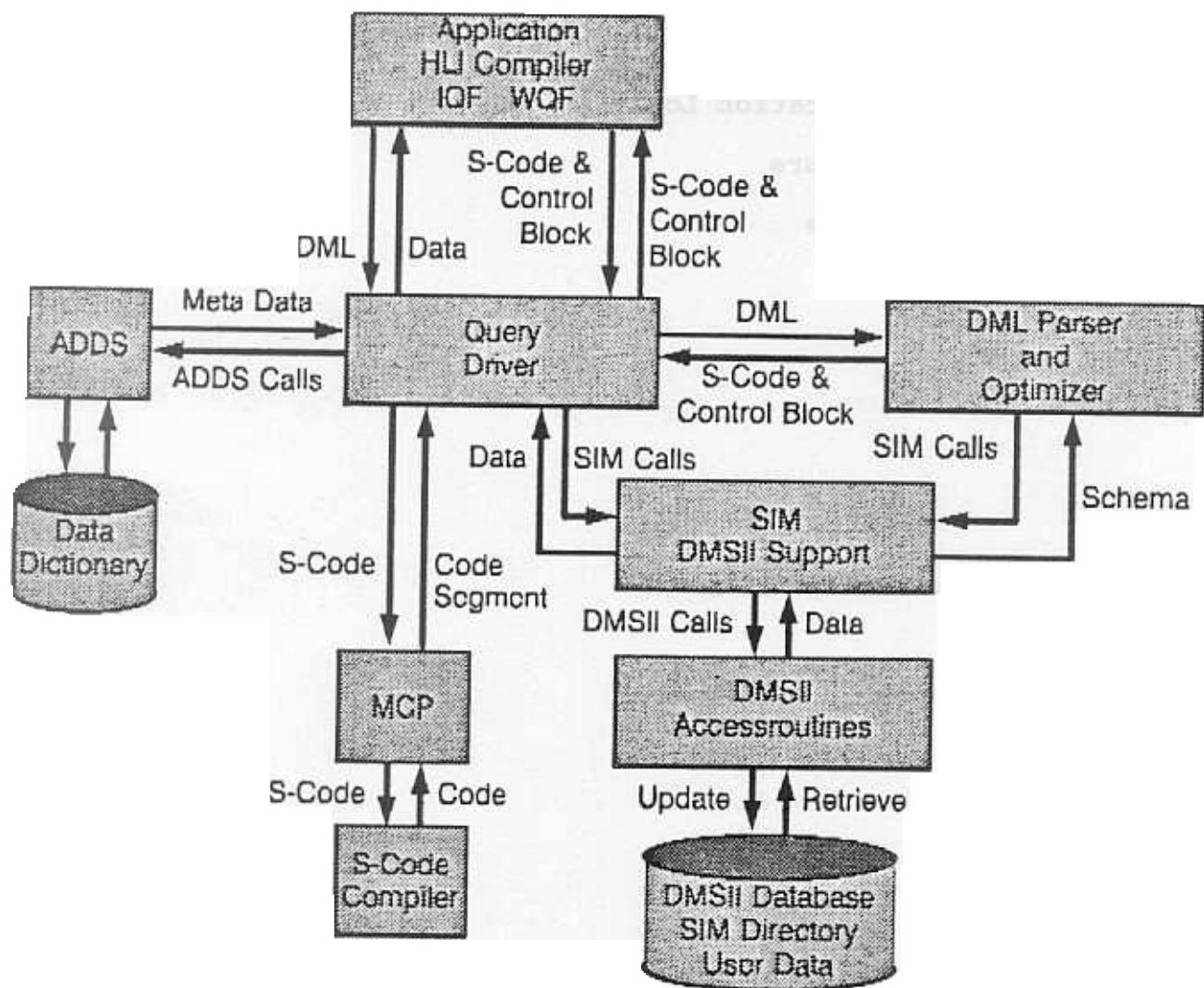
Date Independence

Simplified Application Logic

Compilation Architecture

Execution Architecture

SIM Architecture



Interesting Implementation Problems

Bootstrapping

ADDS required to describe SIM schemas, but ADDS is itself a SIM database

DMS II-mapped SIM databases contain an internal, self-describing SIM database

Solution: Special bootstrapping programs

Automatic Reparsing

Some queries may need reparsing and reoptimization between compilation and execution because of conceptual, interface, or physical schema changes

Solution: Queries are timestamped and recompiled as needed

Postponed Updates

Some integrity constraints are enforced during query parsing while others are delayed until execution. Some of the latter require postponing updates until integrity checking is complete.

Solution: Implement update tanking scheme

Interesting Implementation Problems

Locking Protocols

Many interesting locking problems surfaced due to predefined relationships and full referential integrity.

Solution: Implement two-phase locking protocol employing shared and exclusive locks in DMS II.

Performance Tuning

Many mapping options available for each conceptual component, but choice of defaults that perform well in a wide range of applications was not obvious.

Solution: Performance of alternatives studied for surrogates, relationships, and generalization hierarchies.

Host Language Interface

Stylistically difficult to blend query language with host language without destroying data independence.

Solution: Hybrid retrieval.

On-going Issues

Views

Challenging for strongly typed data models

May not be as important as in relational model

Derived attributes

System-maintained ordering of classes and EVAs

Dependent classes

Nested queries

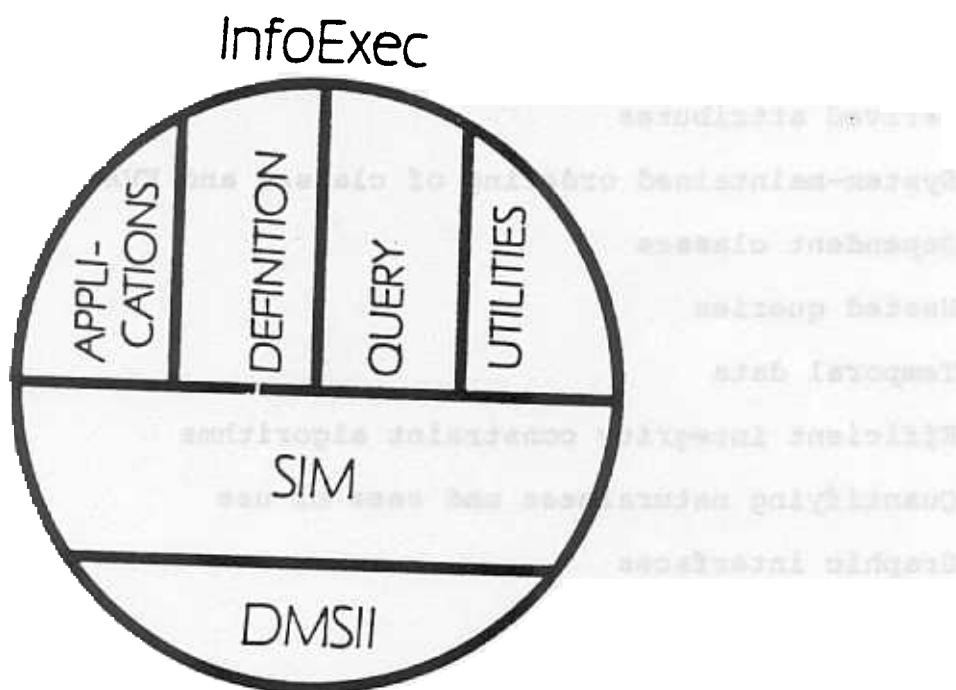
Temporal data

Efficient integrity constraint algorithms

Quantifying naturalness and ease of use

Graphic interfaces

)



Three Layered Architecture

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