

# Merge Sort

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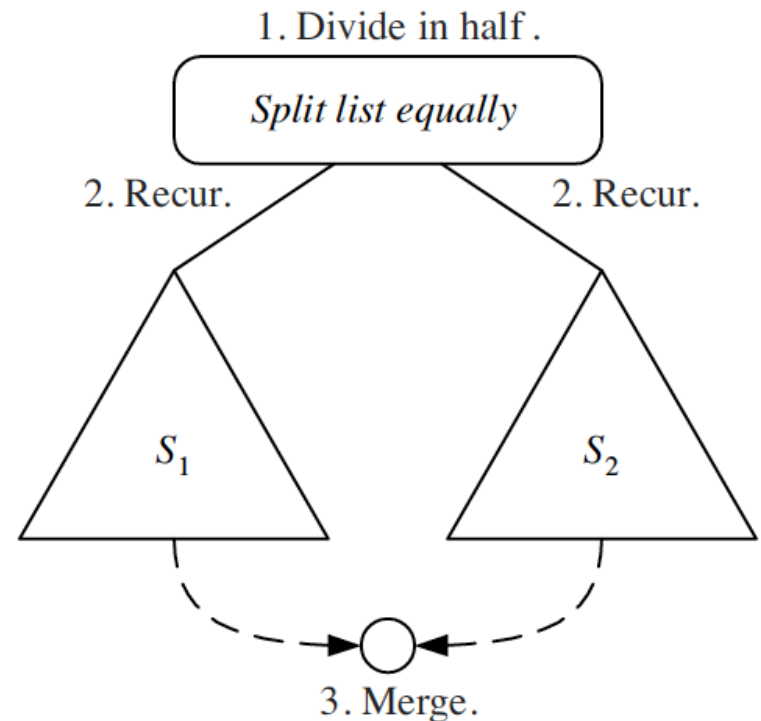
Information on slides originating from *Algorithm Design & Applications* by Michael T. Goodrich and Roberto Tamassia, © 2015 John Wiley & Sons, Inc. Goodrich and Tamassia, ISBN: 978-1118335918.

# Reading Material

- *Algorithm Design & Applications* by Michael T. Goodrich and Roberto Tamassia
  - Chapter 8 Section 8.1 (**Note, the book has a typo for the MergeSort algorithm. I fixed it in these slides.**)

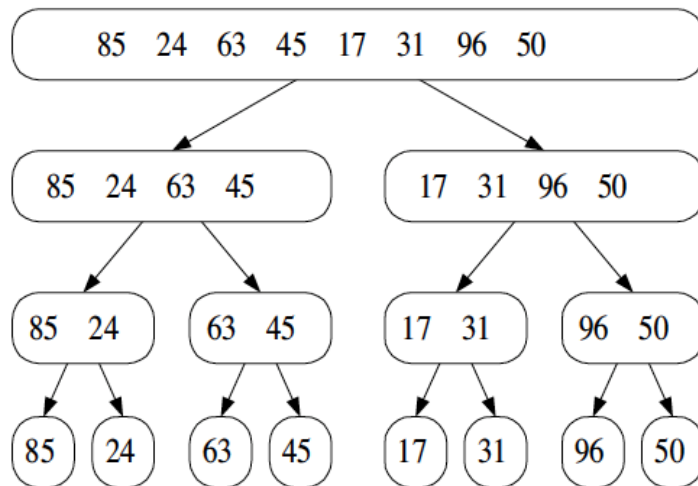
# Divide-and-Conquer

- Divide-and conquer is a general algorithm design paradigm:
  - Divide: divide the input data  $S$  in two disjoint subsets  $S_1$  and  $S_2$
  - Recur: solve the subproblems associated with  $S_1$  and  $S_2$
  - Conquer: combine the solutions for  $S_1$  and  $S_2$  into a solution for  $S$
- The base case for the recursion are subproblems of size 0 or 1

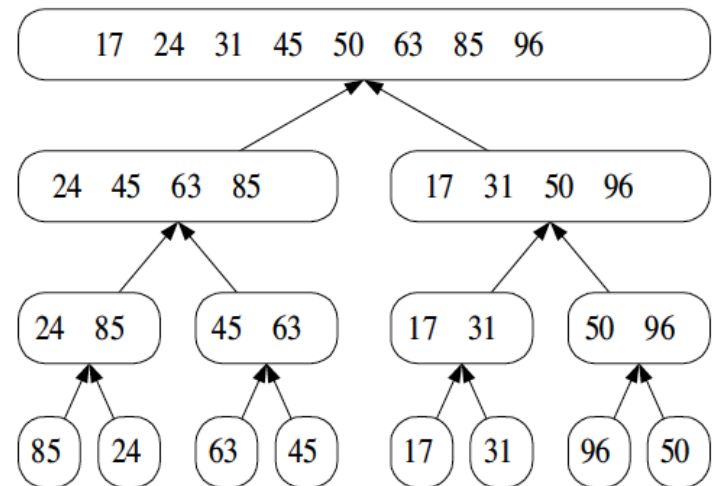


# Merge-Sort

- Merge-sort is a sorting algorithm based on the divide-and-conquer paradigm
- Like heap-sort
  - It has  $O(n \log n)$  running time
- Unlike heap-sort
  - It does not use an auxiliary priority queue
  - It accesses data in a sequential manner (suitable to sort data on a disk)



Merge Sort



# The Merge-Sort Algorithm

- Merge-sort on an input sequence  $S$  with  $n$  elements consists of three steps:
  - Divide: partition  $S$  into two sequences  $S_1$  and  $S_2$  of about  $n/2$  elements each
  - Recur: recursively sort  $S_1$  and  $S_2$
  - Conquer: merge  $S_1$  and  $S_2$  into a unique sorted sequence

**Algorithm** *mergeSort*( $S$ )

**Input** sequence  $S$  with  $n$  elements

**Output** sequence  $S$  sorted according to  $C$

**if**  $S.size() > 1$

$(S_1, S_2) \leftarrow partition(S, n/2)$

*mergeSort*( $S_1$ )

*mergeSort*( $S_2$ )

$S \leftarrow merge(S_1, S_2)$

# Merging Two Sorted Sequences

- The conquer step of merge-sort consists of merging two sorted sequences  $A$  and  $B$  into a sorted sequence  $S$  containing the union of the elements of  $A$  and  $B$
- Merging two sorted sequences, each with  $n/2$  elements and implemented by means of a doubly linked list, takes  $O(n)$  time

**Algorithm** merge( $S_1, S_2, S$ ):

**Input:** Two arrays,  $S_1$  and  $S_2$ , of size  $n_1$  and  $n_2$ , respectively, sorted in non-decreasing order, and an empty array,  $S$ , of size at least  $n_1 + n_2$

**Output:**  $S$ , containing the elements from  $S_1$  and  $S_2$  in sorted order

```
 $i \leftarrow 1$   
 $j \leftarrow 1$   
while  $i \leq n_1$  and  $j \leq n_2$  do  
    if  $S_1[i] \leq S_2[j]$  then  
         $S[i + j - 1] \leftarrow S_1[i]$   
         $i \leftarrow i + 1$   
    else  
         $S[i + j - 1] \leftarrow S_2[j]$   
         $j \leftarrow j + 1$   
while  $i \leq n_1$  do  
     $S[i + j - 1] \leftarrow S_1[i]$   
     $i \leftarrow i + 1$   
while  $j \leq n_2$  do  
     $S[i + j - 1] \leftarrow S_2[j]$   
     $j \leftarrow j + 1$ 
```

Type-o in book fixed on these slides. While conditions are not correct in the book.

# Thank You !



# Questions ?