

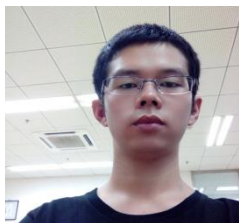
# CarrotSQL

<https://github.com/huangwentao0831/carrot>

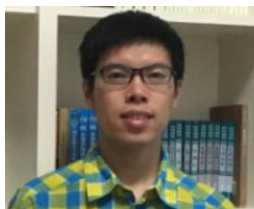
# Team



Han Han



Guan Hua



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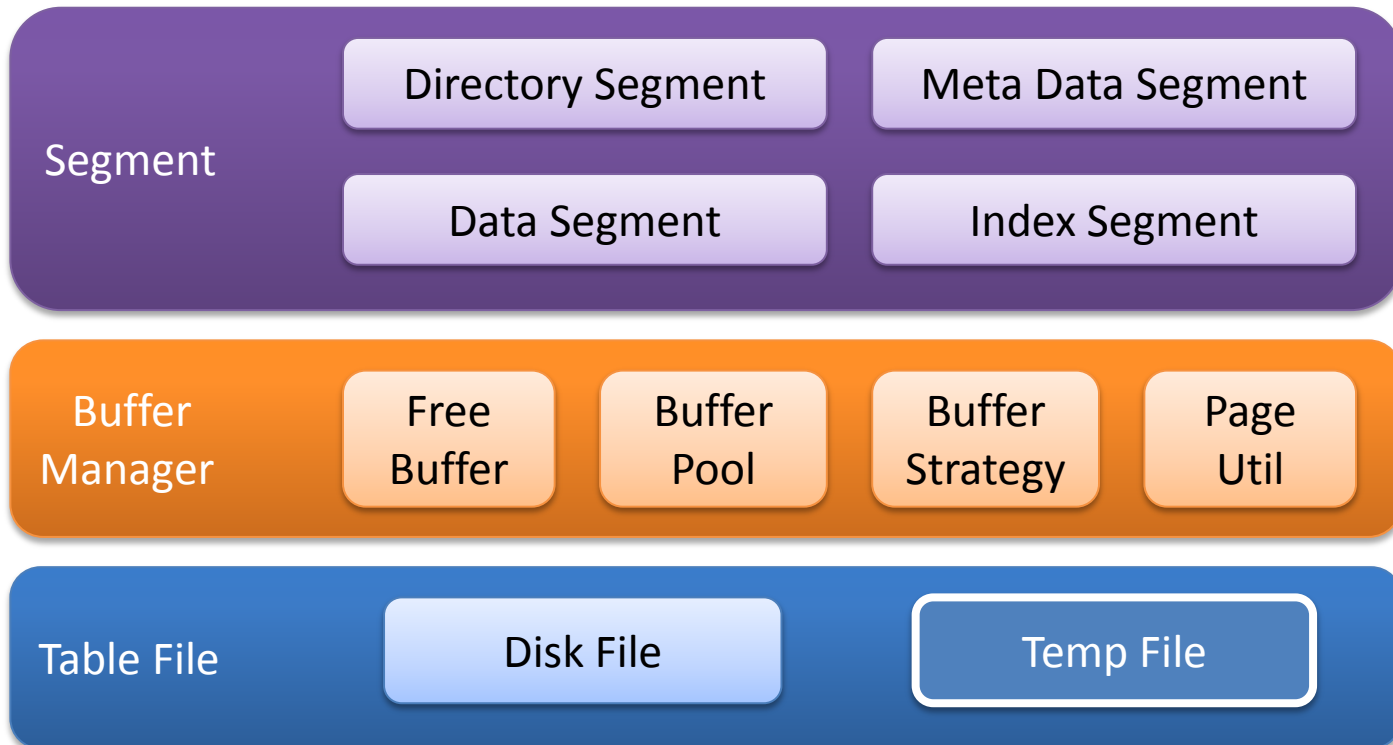


Shao Mingrui

# Content

- Storage Management
- Index
- SPJ
- The Execution Demo

# Storage Management



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# Index

0x00-0x03	magic	field ID	field Number	block Number
0x04-0x11	page Number			
0x12-0x19	next Addr			
.....	index Meta Item			

# B+ tree Index

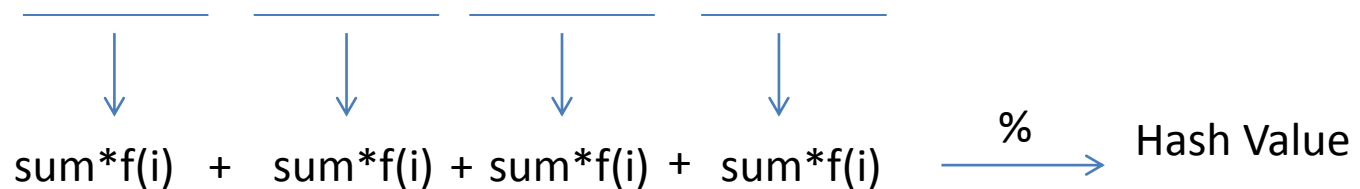
0x00-0x03	field ID	node Type	key Number
0x04-0x11	parent Addr		
0x12-0x19	previous		
0x20-0x27	Next		
.....	key array		
.....	children array		
.....	page number		

# Linear Hash Index

- Expand the range of hash function

We give the factor for each sum of bytes, so that the range of hash function can expand from hundreds to billions.

.... .... 0010      1001      0110      0001



$$\text{sum} \times a^i, a = 3, i$$

- Extendible and shrinkable
  - The linear hash table can extend and shrink dynamically.



# Linear Hash Index

- index structure

Bucket Size			Record Size
Split Rate	Key Size	Hash Size	Next Addr
Bucket Value ...			Bucket Addr ..

# Linear Hash Index

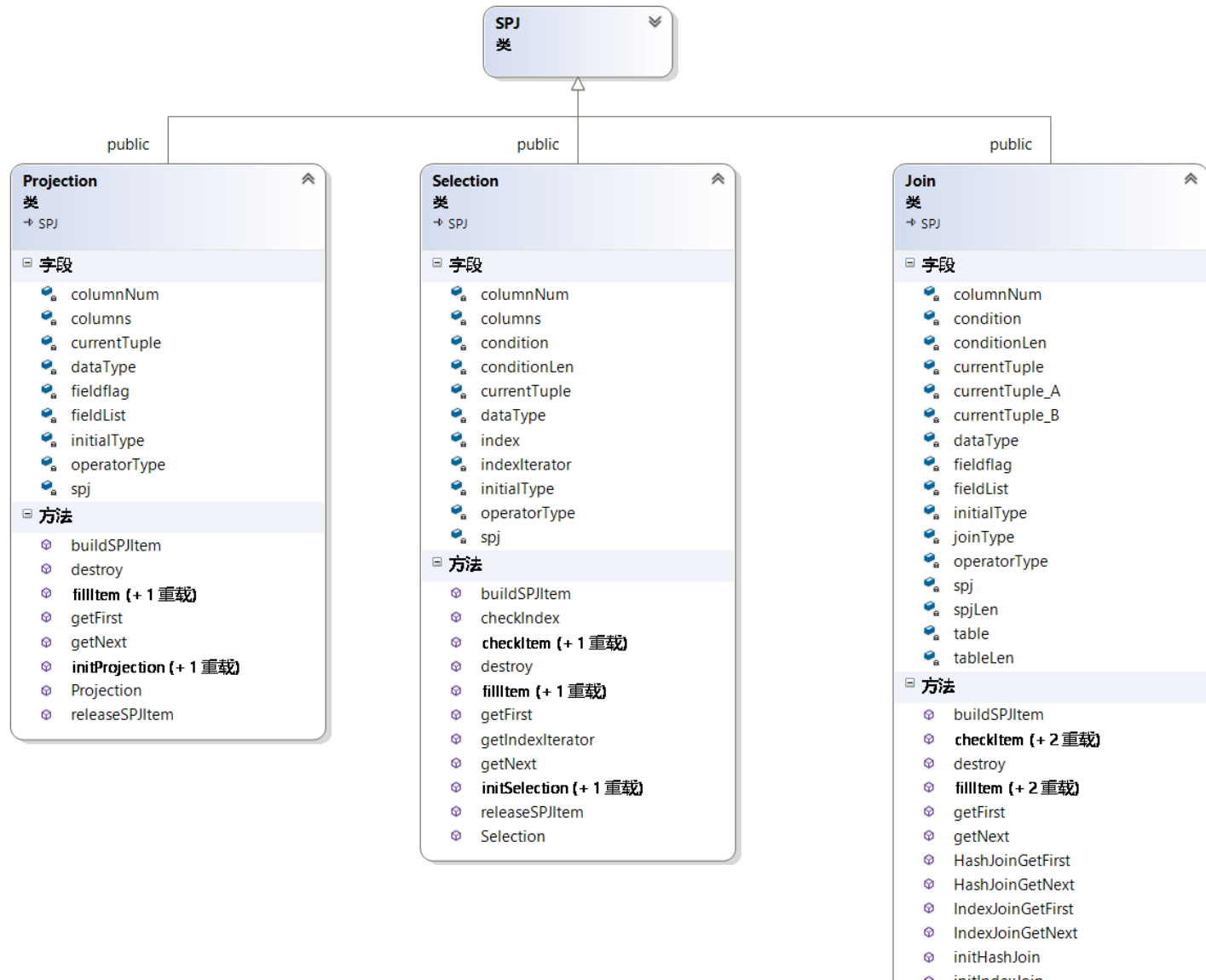
- bucket structure

Key num	Next Addr
Key value ...	Tuple Addr ..

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# SPJ



# SPJ

- Iterator

iterator->getFirst(item);

iterator->getNext(item);

- initial

only table

only iterator

hybrid

# SPJ

- Focus on Nest loop join and index join
  - Nest loop join based on tuple
  - Index join based on B+ tree and linear hash index
- Multiple tables join
  - Due to the join designed on iterators, so multiple tables join is very to realise and number of tables can be increased with no contrains.

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Thank you!

# Conclusion

- Creation

We focus on designing details, not perfect ,but better.

- Team Work

One people ,One experiment ?

- Future work

Work harder, work better

Coding&Debug(Segment Fault)

# Appendix: Page

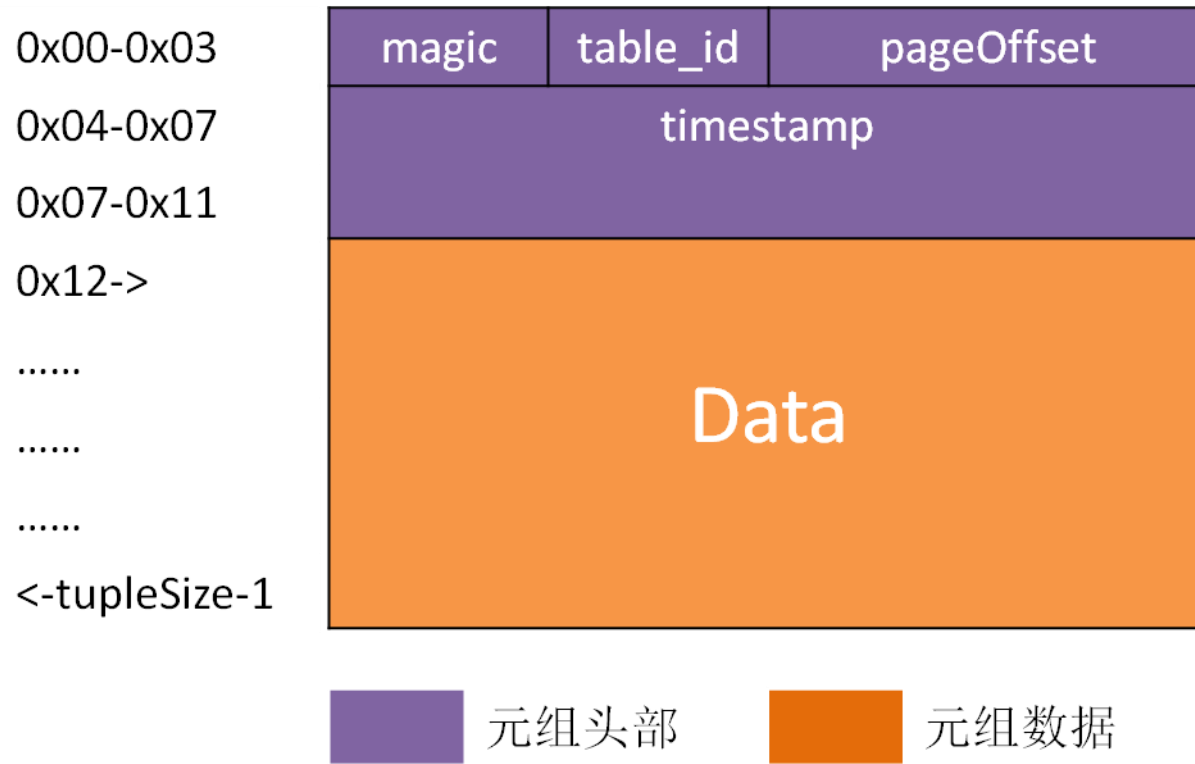
0x0000-0x0003	usedByte	flag
0x0004-0x0007	Control Data	
.....		
0x0390-0x0393	Page Data	
0x0394-0x0397		
.....		
.....		
.....		
0x01ffd-0x2000		

# Appendix: Meta Data

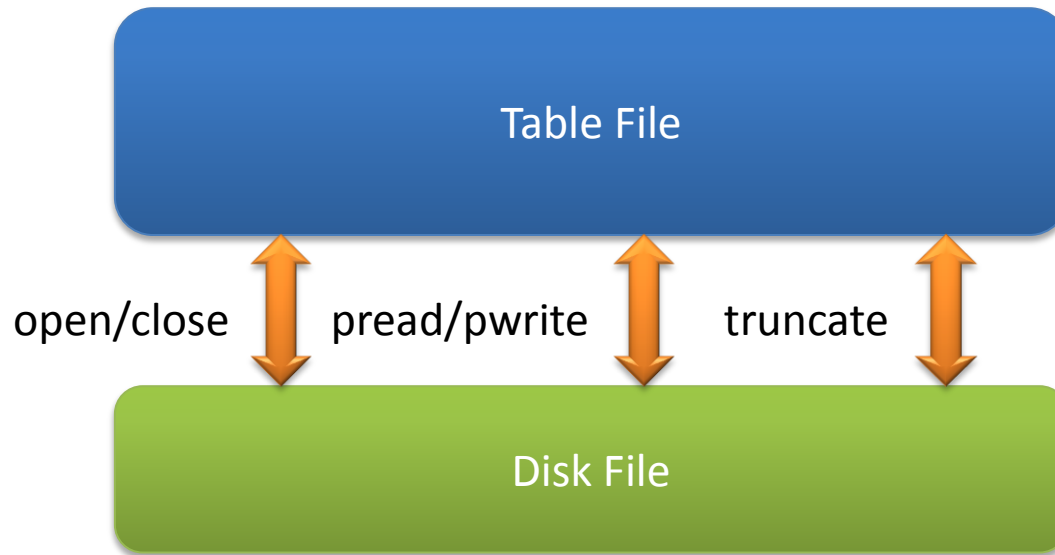
0x00-0x03	magic	Table Id	Field Num	Block Num
0x04-0x07	Table Name			
.....				
0x40-0x43				
0x44-0x47	Field Info			
.....				
.....				
.....				

0x00-0x03	magic	field id	flag	data type
0x04-0x43	fname			
0x44 - 0x47	len			

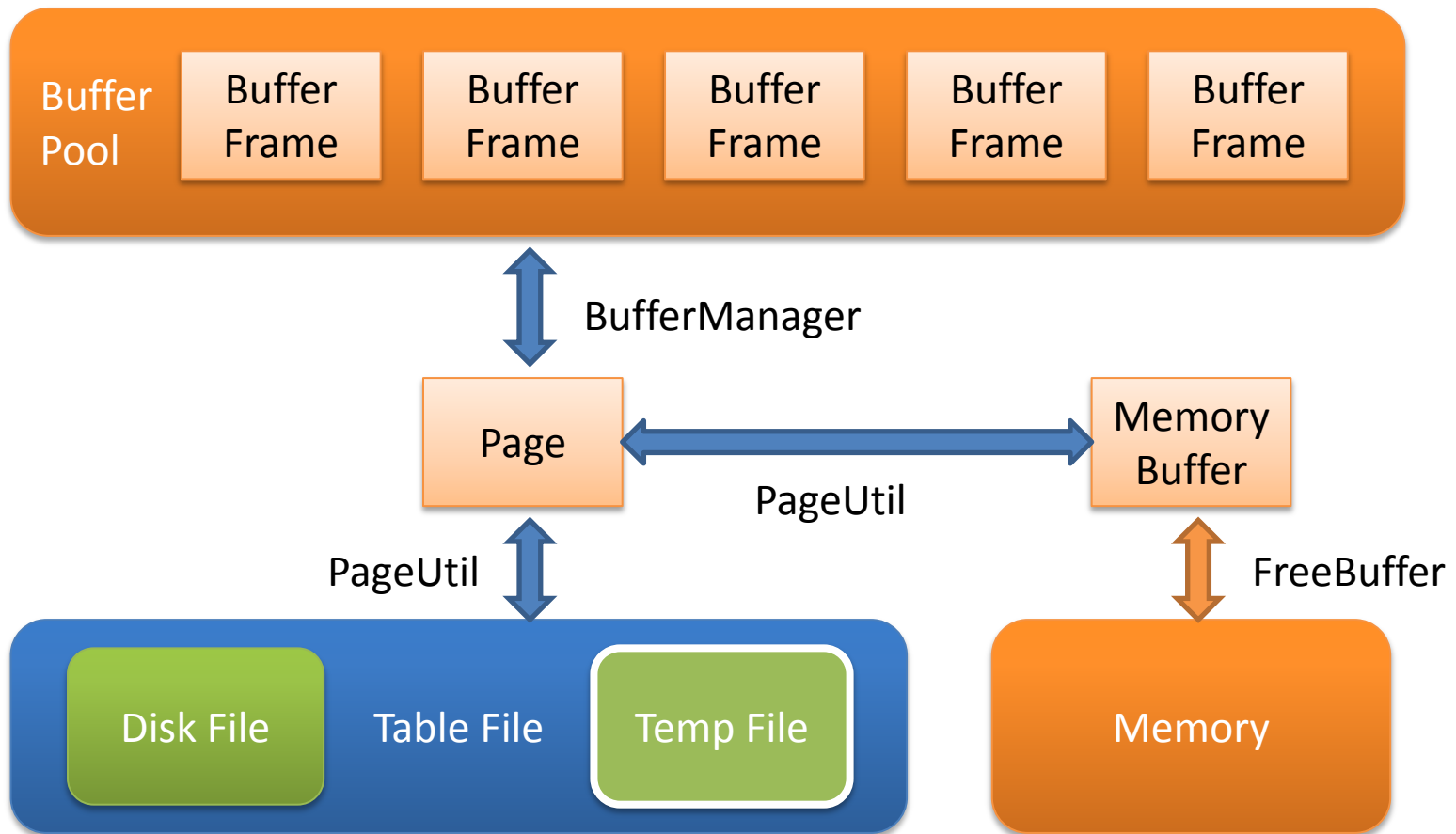
# Appendix: Tuple



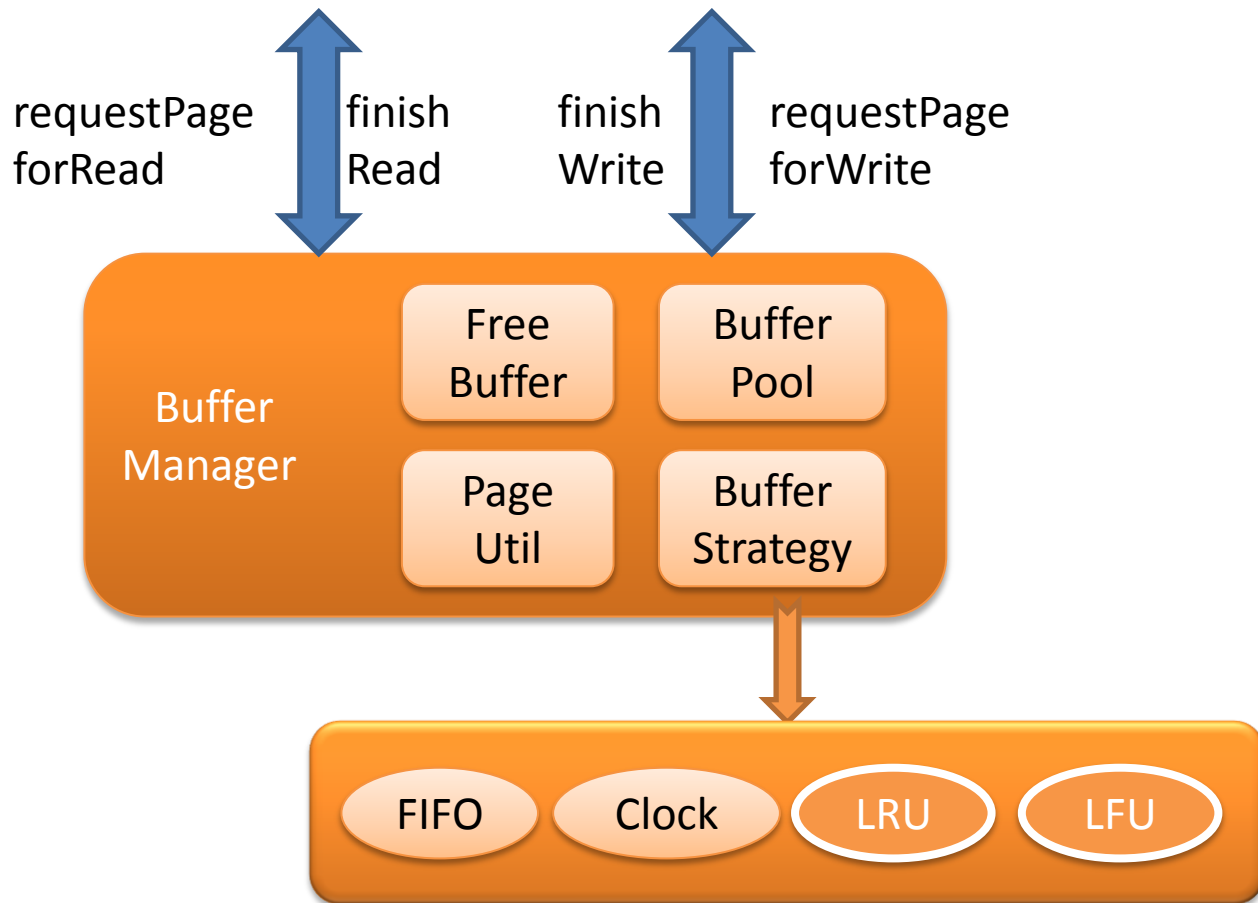
# Storage: IO



# Storage: Buffer



# Storage: Buffer

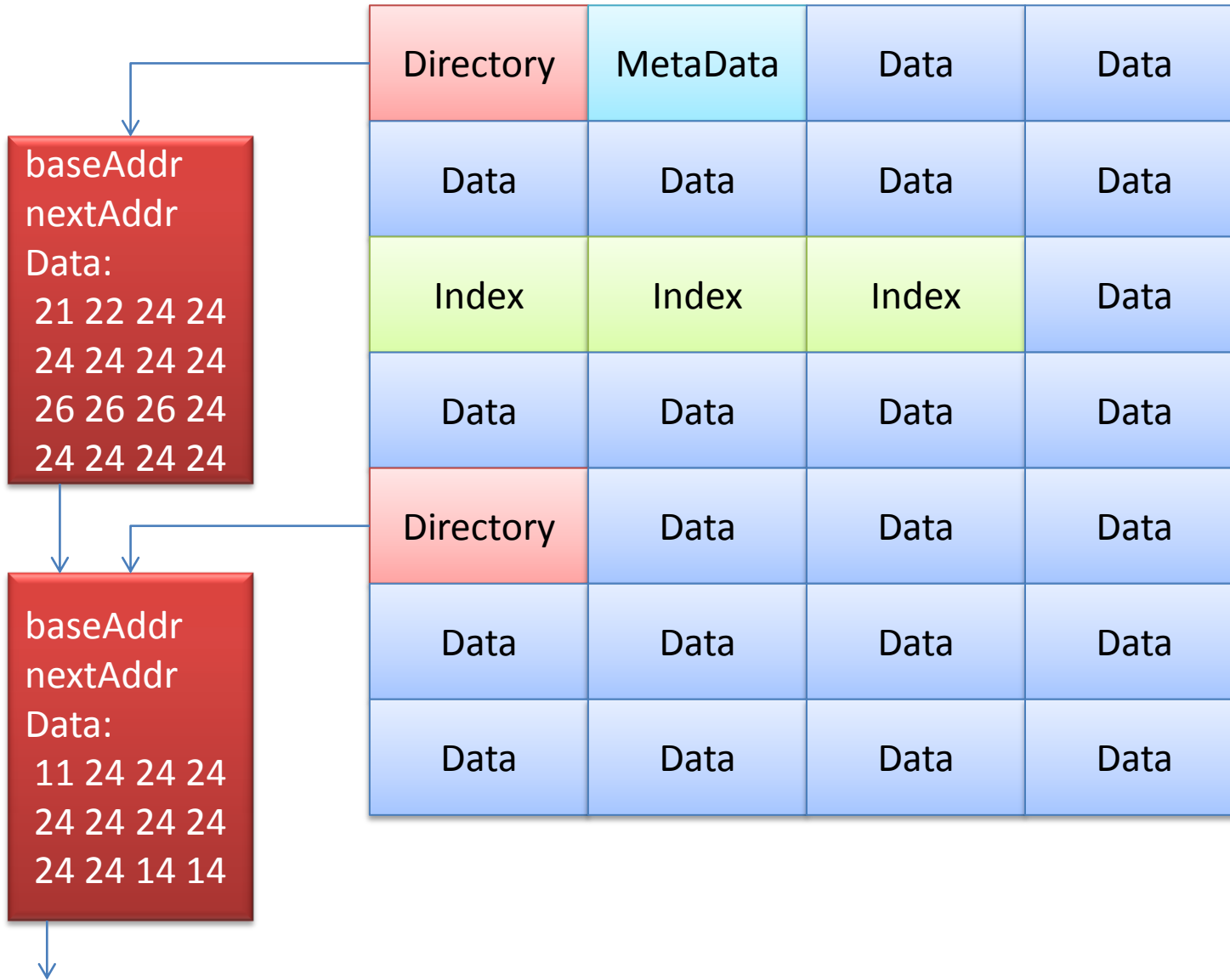




# Storage: Segment

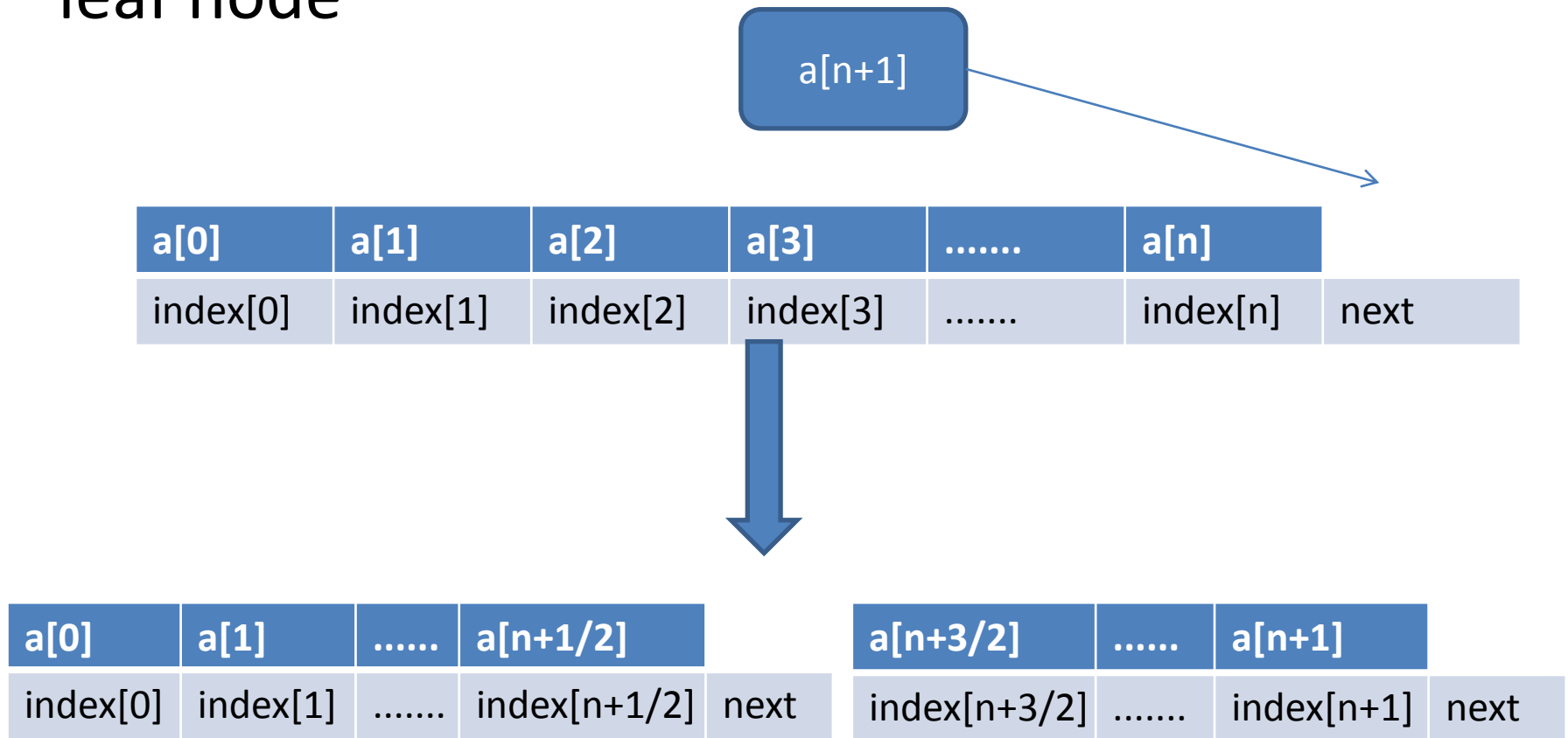
- How to manage pages that store different data?
- Can one page store both a tuple and a index structure(eg. a tree node)?

# Segment Design



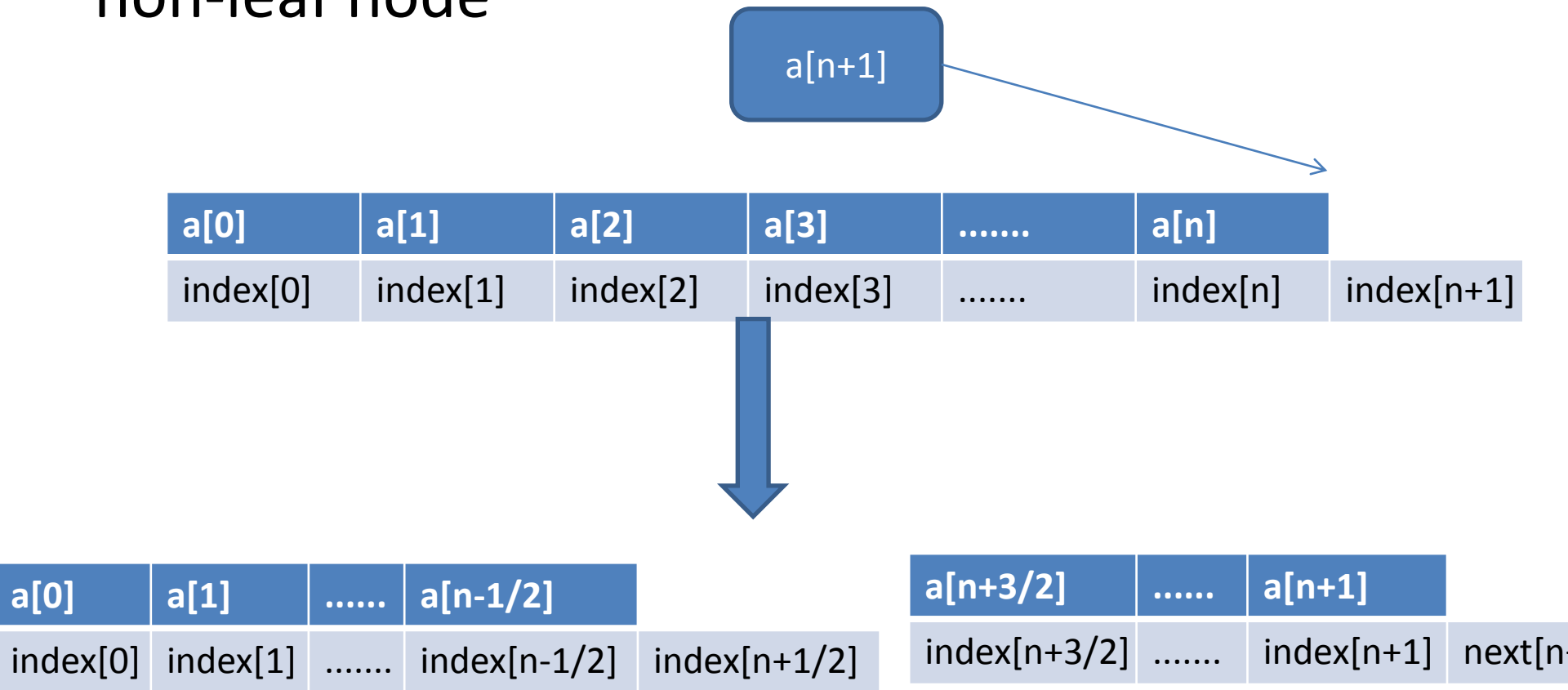
# split of node

leaf node



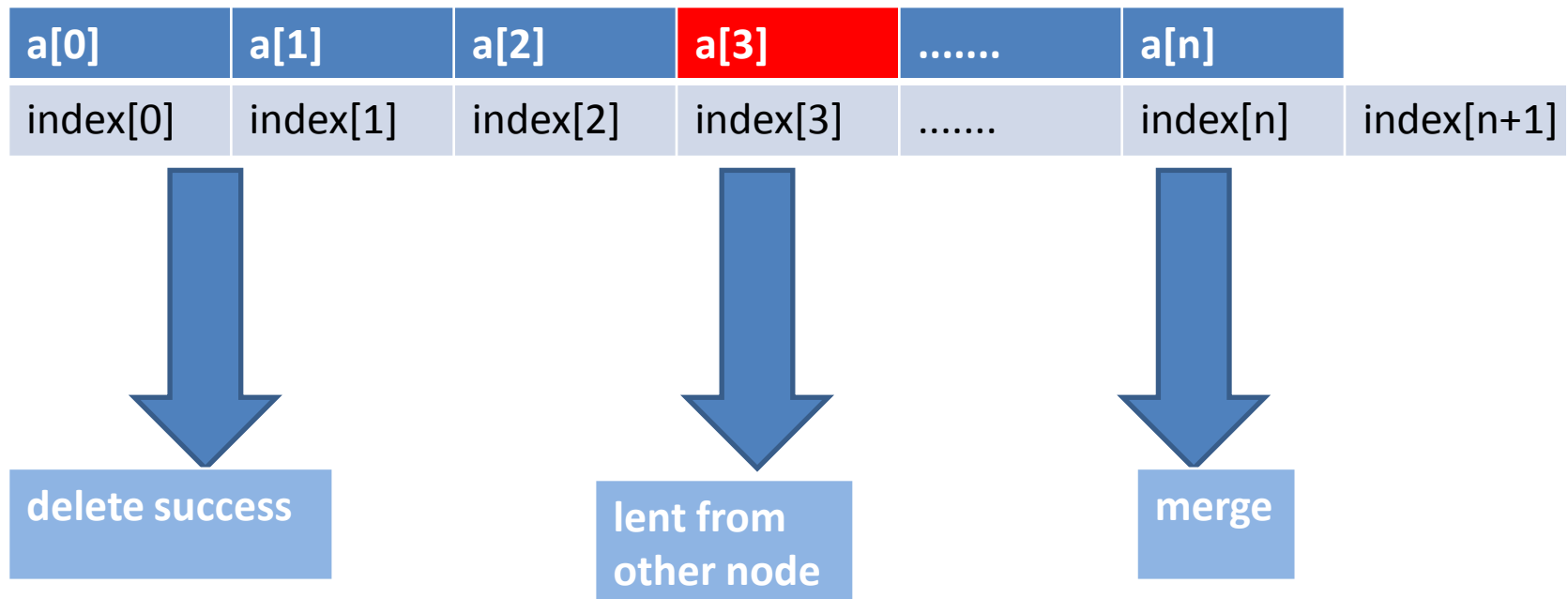
# split of node

non-leaf node



# merge of node

leaf node



# search for node

