

LAIS, Lecture #22

"Spectral Clustering"

$$G = (V, E, W)$$

$$G_i = (V_i, E_i) \quad i \in [c]$$

connected components:

$w_{ab} = 0$ whenever $a \in V_i$,
 $b \in V_j$ and $i \neq j$.

The G_i 's are connected:

$G_i: \forall a, b \in V_i$
 $\exists \alpha_1, \dots, \alpha_k \in V_i$ s.t.

$(\alpha, \alpha_1) \in E_i, (\alpha_1, \alpha_2) \in E_i$
 $\dots (\alpha_{k-1}, \alpha_k) \in E_i, (\alpha_k, b) \in E_i$

Thm 1 e_{G_1}, \dots, e_{G_c} are
a basis of $V(L)$.

$$\text{Prf } \alpha^T L \alpha = \sum_{i < j} w_{ij} (\alpha(i) - \alpha(j))^2$$

$$= \sum_{\substack{i < j \\ i, j \in V_k \\ \text{for some } k}} w_{ij} (\alpha(i) - \alpha(j))^2$$

$$e_{G_k}^T L e_{G_k} = \sum_{\substack{i < j \\ i, j \in V_k \\ \text{for some } k}} w_{ij} (e_{G_k}(i) - e_{G_k}(j))^2$$

$$= \sum_{\substack{i < j \\ i, j \in V_k}} w_{ij} (e_{G_k}(i) - e_{G_k}(j))^2 = 0$$

$$z \in \mathcal{V}(L)$$

$$\sum_{\substack{i < j \\ i, j \in V_K \\ \text{for some } K}} w_{ij} (z(i) - z(j))^2 = 0$$

Take $i, j \in V_s$. Since G_s is connected \exists

$i_1, \dots, i_\ell \in V_s$ s.t.

$(i, i_1), (i_1, i_2), \dots, (i_{\ell-1}, i_\ell), (i_\ell, j) \in E_s$

$$\Rightarrow w_{ii_1}, w_{i_1 i_2}, \dots, w_{i_{\ell-1} i_\ell}, w_{i_\ell j} > 0$$

$$z(i) = z(i_1), z(i_1) = z(i_2), \dots, z(i_\ell) = z(j)$$

$$\Rightarrow z(i) = z(j)$$

So for any $i, j \in V_K$

$$z(i) = z(j) = c_K$$

$$\Rightarrow z = \sum_{K \in \mathcal{C}} c_K \mathbf{1}_{V_K} \quad \square$$

"Subspace Clustering"

- * special case of clustering
- * generalized principal component analysis

Lemma 2 [motivation]. In

many cases the data

$$X \in \mathbb{R}^{D \times N}$$

are heterogeneous (come from different sources) and each source is modeled by a linear subspace. \square

Def 3 $X = [X_1 \dots X_c]$ unknown permutation

$X_i = [x_{i1} \dots x_{iN_i}] \in \mathbb{R}^{D \times N_i}$

$x_{ij} \in \mathcal{S}_i \quad \forall j \in [N_i]$
 linear subspace of \mathbb{R}^D \square

Lemma 4 $\mathcal{S} = \sum_{i \in [c]} \mathcal{S}_i$

$x_{ij} \in \mathcal{S} \quad \forall i \in [c], \forall j \in [N_i]$
 and so we are back to the case of PCA. \square

Ex 5 $D=10, c=6$

$d_i = \dim \mathcal{S}_i = 2$. In this case it may happen that $\mathcal{S} = \mathbb{R}^{10}$. \square

* dimensionality reduction via PCA might not be possible

* need to cluster

"Sparse Subspace Clustering"

Def 6 [sparse representation, compressed sensing]

$b \in \mathbb{R}^D, A \in \mathbb{R}^{D \times N}$
 find the sparsest $x \in \mathbb{R}^N$ s.t.
 $Ax = b$ (usually $N \gg D$)

$\min \|x\|_0 \quad \text{s.t. } Ax = b$
 $x \in \mathbb{R}^N$ counts # of non-zero entries

if noise $\min_{x \in \mathbb{R}^D} \|x\|_0 \quad \text{s.t. } \|Ax - b\|_2 \leq \epsilon$
 IP hard

$\min \|x\|_1 \quad \text{s.t. } \|Ax - b\|_2 \leq \epsilon$
 $x \in \mathbb{R}^D$
 convex, Lasso

$\min_{x \in \mathbb{R}^D}$
 basic

Lemma 7
 i) $\dim \mathcal{S}$

Idea:
 point
 linear
 rest
 $\min_{C \in \mathbb{R}}$

unknown
connection

* dimensionality reduction
via PCA might not
be possible

* need to cluster

"Sparse Subspace Clustering"

Def 6 [sparse representation,
compressed sensing]

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NP hard

$$\min \|x\|_1 \text{ s.t. } \|Ax - b\|_2 \leq \epsilon$$

convex, Lasso

$$\min \|x\|_1 \text{ s.t. } Ax = b$$

$x \in \mathbb{R}^D$

basics pursuit

Lemma 7 hypothesis:

$$i) \dim S_i = d_i \ll D \quad ii) \forall i, \text{rank}(X_i | \{x_j\}) = d_i$$

Idea: express every
point x_j as a sparse
linear combination of the
rest of the points

$$\min \|c\|_0 \text{ s.t. } x_j = Xc, c_{(j)} = 0$$

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Def 8 The solution c is called subspace preserving if $c(kl) = 0 \ \forall \ k \neq i$. \square

Thm 9 Suppose the S_i 's are independent. Then the solution c to (*) will be subspace preserving.

Prf $Xc = [X_i \ \underbrace{X_{-i}}_{\substack{\text{all points} \\ \text{outside } X_i}}] \begin{bmatrix} c_i \\ c_{-i} \end{bmatrix}$
 We will show $c_i = 0$. Suppose $c_i \neq 0$.

$$X_{ij} = Xc = X_i c_i + X_{-i} c_{-i}$$

$$\Rightarrow \underbrace{X_{ij} - X_i c_i}_{\in S_i} = \underbrace{X_{-i} c_{-i}}_{\in \sum_{j \neq i} S_j}$$

By independence $X_{ij} = X_i c_i$
 and $X_{-i} c_{-i} = 0$.
 So $\begin{bmatrix} c_i \\ 0 \end{bmatrix}$ is a feasible

solution to (*). But
 $\| \begin{bmatrix} c_i \\ 0 \end{bmatrix} \| < \| \begin{bmatrix} c_i \\ c_{-i} \end{bmatrix} \| \Rightarrow \leftarrow \square$

Def 10 [disjoint subspaces]

$$S_i \cap S_j = 0 \quad \forall i \neq j \quad \square$$

Thm 11 For each $i \in [c]$
let Φ_i be the set of all

$\tilde{X}_i \in \mathbb{R}^{D \times d_i}$ submatrices of X_i of rank k d_i .

θ_{ij} : first principal angle
between S_i, S_j

$$\cos \theta_{ij} = \max_{\substack{u \in S_i \\ \|u\|_2=1 \\ v \in S_j \\ \|v\|_2=1}} |u^T v|$$

Then the solution
to (*) is subspace
preserving if

$$\max_{\tilde{X}_i \in \Phi_i} \sigma_{d_i}(\tilde{X}_i) > \sqrt{d_i} \max_{\substack{k \neq i \\ l \in [m_k]}} \|X_k\|_2 \cdot \max_{k \neq i} \cos \theta_{ik}$$

Thm 12 Auxiliary problems:

$$\min_{\alpha} \|\alpha\|_1 \quad \text{s.t.} \quad y = X_i \alpha \quad (1)$$

$$\min_b \|\mathbf{b}\|_1 \quad \text{s.t.} \quad y = X_i \mathbf{b} \quad (2)$$

$$\text{for } y \in S_i \cap \left(\sum_{j \neq i} S_j \right)$$