

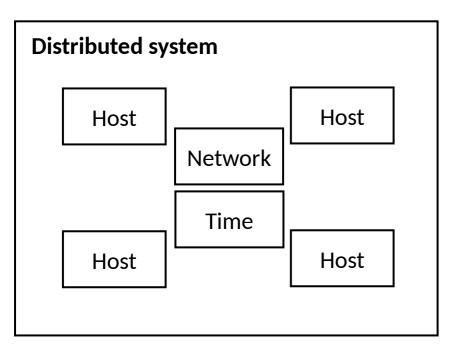
EECS498-008 Formal Verification of Systems Software

Material and slides created by

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A distributed system is composed of multiple hosts, a network and clocks



```
Distributed system: attempt #3
module DistributedSystem {
  datatype Variables =
    Variables(hosts:seq<Host.Variables>,
              network: Network. Variables,
              time: Time.Variables)
  predicate Next(v, v', hostid, msg0ps: Message0ps, clk:Time) {
       (&& HostAction(v, v', hostid, msg0ps)
        && Network.Next(v, v', msg0ps)
        && Time.Read(v.time, clk))
       (&& Time.Advance(v.time, v'.time)
        && v'.hosts == v.hosts
        && v'.network == v.network)
```



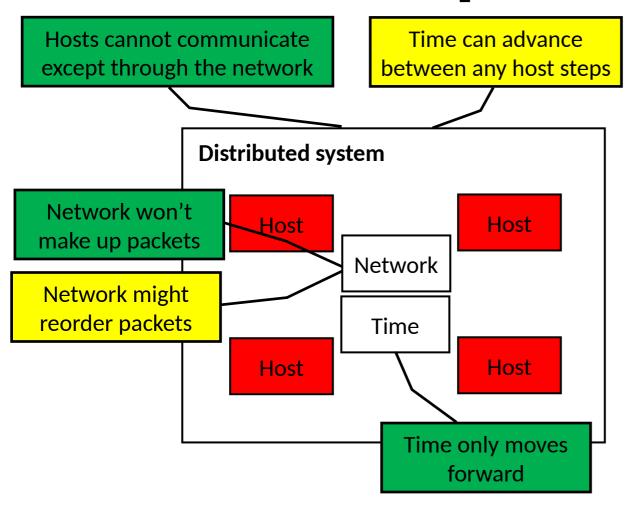
A "distributed" system

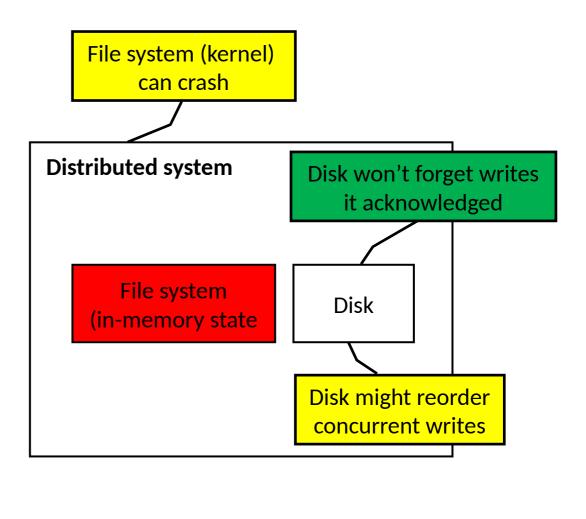
File system (in-memory state Disk

```
module DistributedSystem {
  datatype Variables =
    Variables(fs: FileSystem.Variables,
              disk: Disk.Variables)
  predicate Next(v, v') {
    || (exists io ::
        && FileSystem.Next(v.fs, v'.fs, boginding variable)
        && Disk.Next(v.disk, v'.disk, io)
    || ( // Crash!
        && FileSystem.Init(v'.fs)
        && v'.disk == v.disk
```



Trusted vs proven







: the systems specification sandwich



trusted application spec

proof
proof
code

trusted environment assumptions

image: pixabay



Administrivia

- Midterm went great, congratulations!
- Problem set 3 due on Friday
- Project 1 will be released on Monday
 - Let me know if you can't find teammates
- Midterm evaluations due today!
 - They are really, really important
 - If we reach 66%, I'll design custom course stickers
- I'm giving another lecture right after this one, so I have to skip IOH



Chapter 6: Refinement

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State machines: a versatible ool

State machines can be used to

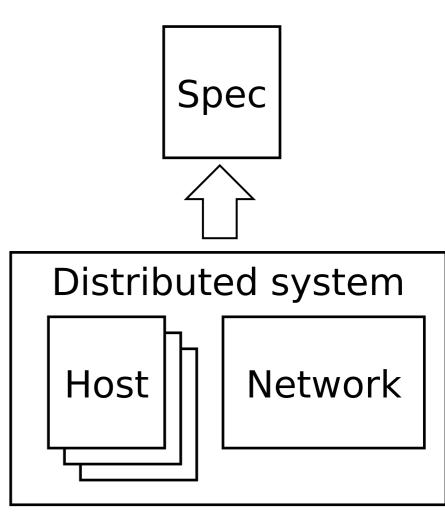
- Model the program
- Model environment components
- Model how the system (program+environment) fits together
- Specify the system behavior

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Different ways to specify behavior

- C-style assertions
- Postconditions
- Properties/invariants
- Refinement to a state machine



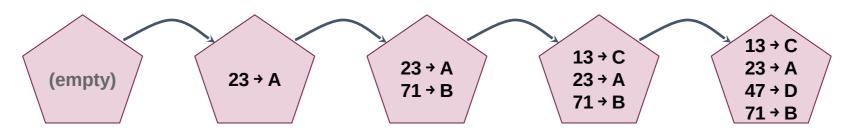


Example: hashtable

```
module HashTable {
  datatype Variables = Variables(tbl:seq<Pair<int, string>>)
  predicate Insert(v:Variables, v':Variables, key:int,
val:string) {
    var free := Probe(v.tbl, key);
    && free.Some?
    && v'.tbl == v.tbl[free.value := Pair(key, val)]
                                               13 → C
                                                            13 → C
                    23 → A
                                 23 → A
                                               23 → A
                                                            23 → A
                                                            47 → D
                                 71 → B
                                               71 → B
                                                            71 → B
```



The spec: a simple map

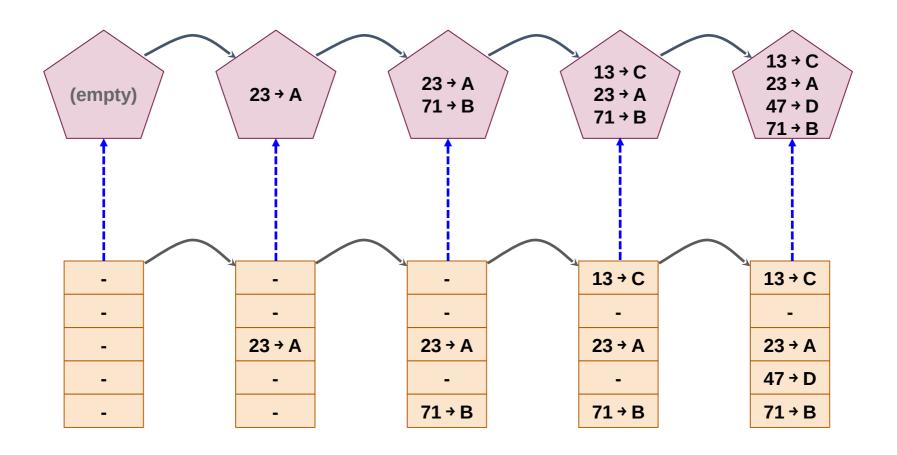


```
module MapSpec {
   datatype Variables = Variables(mapp:map<Key, Value>)

   predicate InsertOp(v:Variables, v':Variables, key:Key,
   value:Value) {
     && v'.mapp == v.mapp[key := value]
   }
}
```



Refinement





The benefits of refinement

Refinement allows for good specs

- Abstract: elide implementation details
- Concise: simple state machine
- Complete: better than a "bag of properties"
 - But if you want, you can prove properties about the spec

Refinement is very powerful

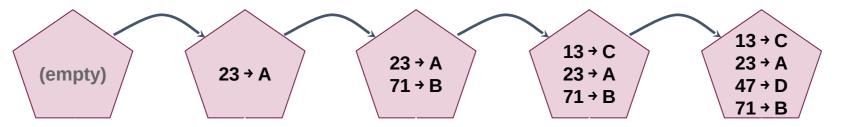
- Can specify systems that are hard to specify otherwise
 - E.g. linearizability

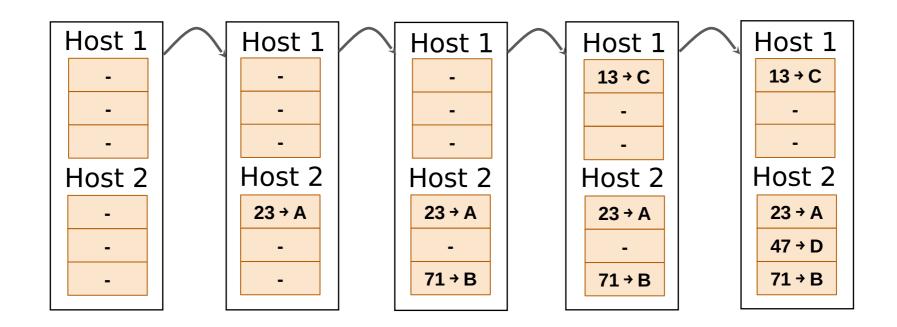
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A sharded key-value store

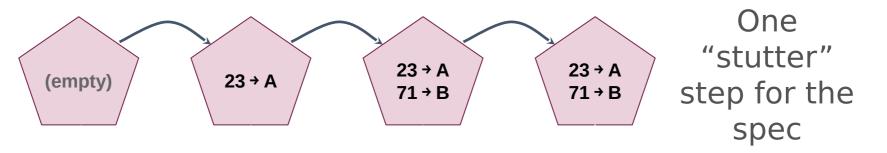
Logically centralized, physically distributed

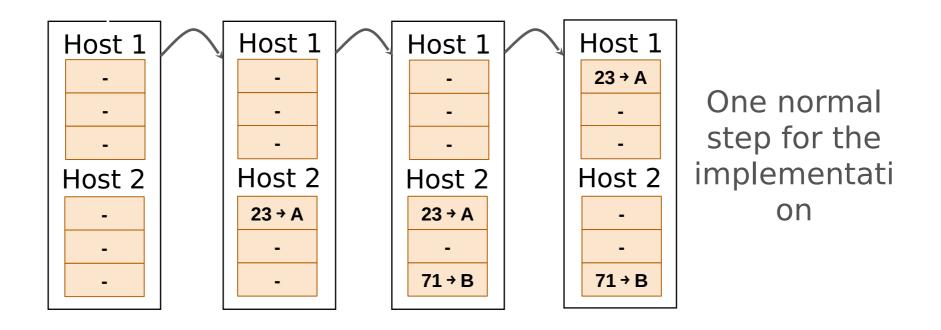






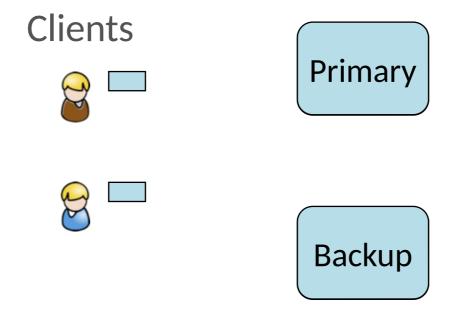
Stutter steps





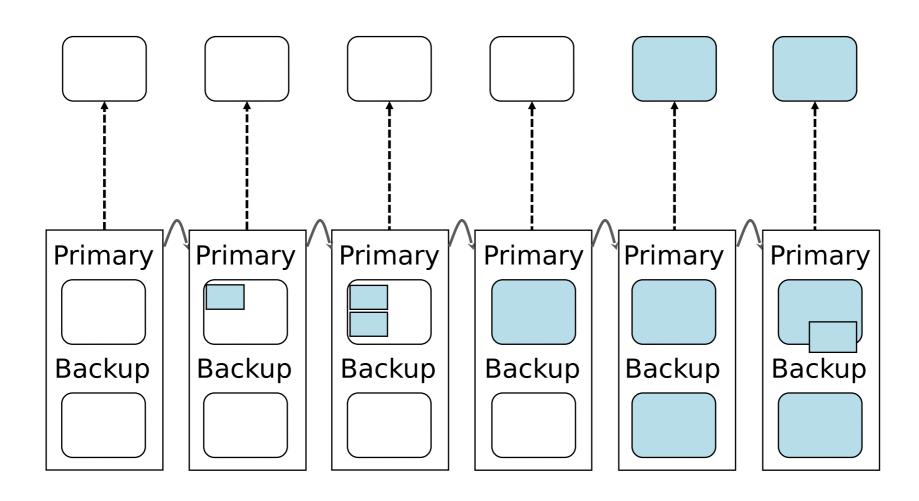


A primary-backup protocol



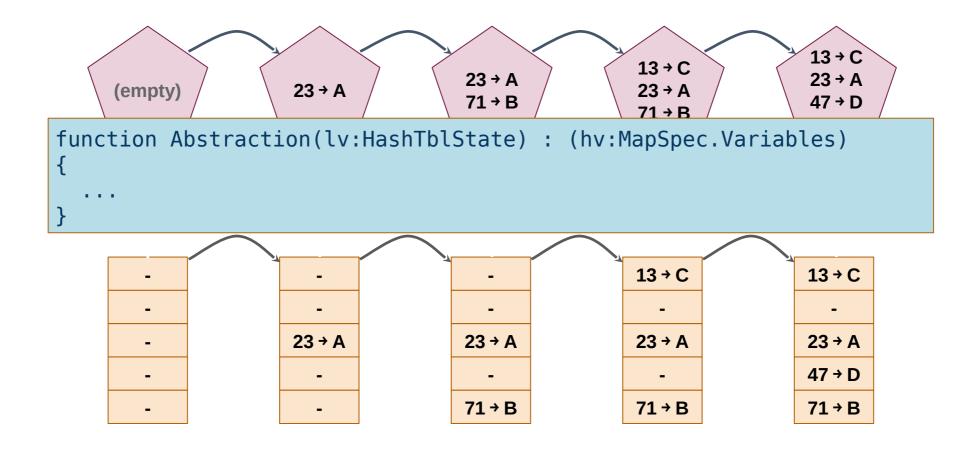


A primary-backup protocol





The interpretation (Abstraction) function





A refinement proof

```
function Abstraction(v:Variables) : MapSpec.Variables
predicate Inv(v:Variables)
lemma RefinementInit(v:Variables)
requires Init(v)
ensures Inv(v) // Inv base case
ensures MapSpec.Init(Abstraction(v)) // Refinement base case
lemma RefinementNext(v:Variables, v':Variables)
    requires Next(v, v')
    requires Inv(v)
    ensures Inv(v') // Inv inductive step
    ensures MapSpec.Next(Abstraction(v), Abstraction(v')) // Refinement inductive
step
          Abstraction(v) == Abstraction(v')
                                                                       OR stutter
step
```