

EECS498-008 Formal Verification of Systems Software

Material and slides created by

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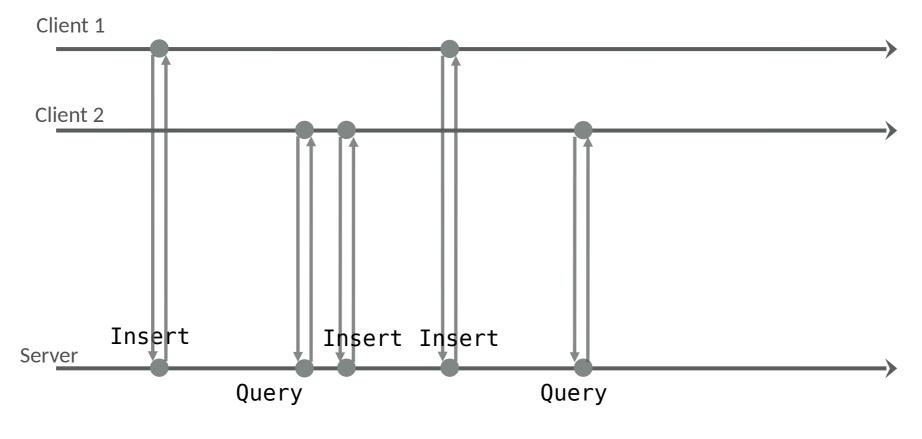


Synchronous specs

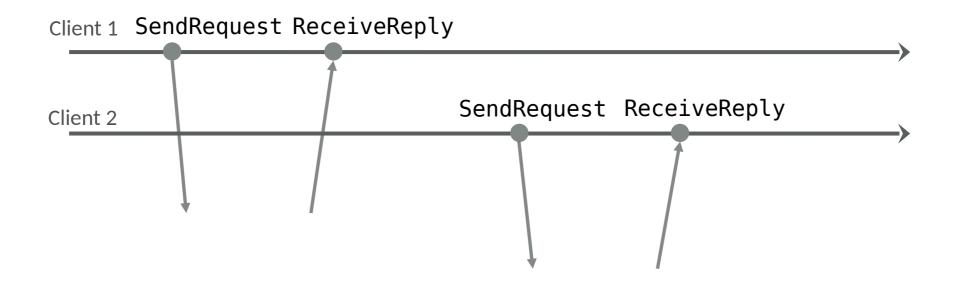
```
module MapSpec {
  datatype Variables = Variables(mapp:map<Key, Value>)
  predicate InsertOp(v:Variables, v':Variables, key:Key,
value:Value) {
  predicate QueryOp(v:Variables, v':Variables, key:Key,
output:Value) {
   Insert
                  Insert Insert
            Query
                                  Query
```



Synchronous specs



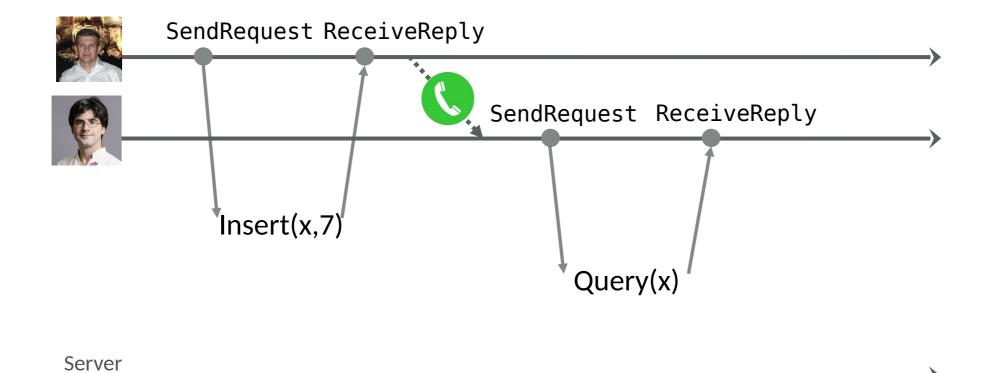
Asynchrony in real life



Server

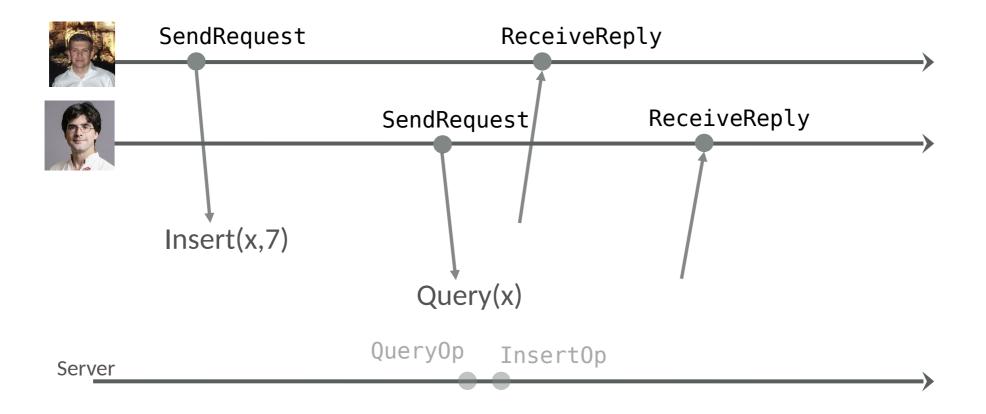


Linearizability



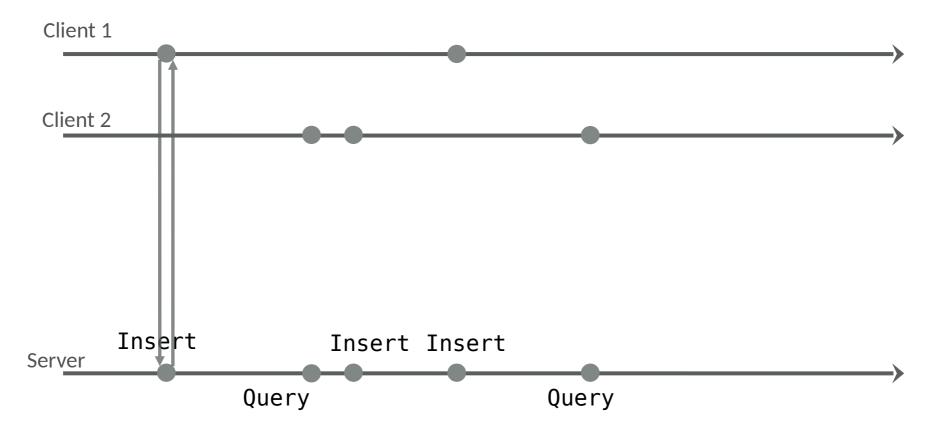


Linearizability





The limitation of Synchronous specs





Defining an asynchronous interface

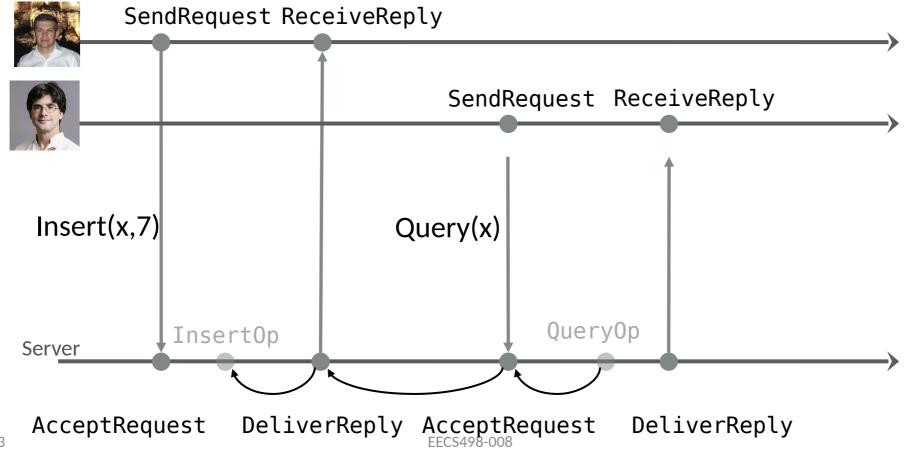
```
module MapSpec {
  datatype Variables = Variables(mapp:map<Key, Value>,
                                    requests:set<Input>, replies:set<Output>)
  predicate InsertOp(v:Variables, v':Variables, request: Input) {...}
  predicate QueryOp(v:Variables, v':Variables, request: Input, output:Value)
{...}
  predicate AcceptRequest(v:Variables, v':Variables, request: Input) {
   // add request to requests, if it's not there already
  predicate DeliverReply(v:Variables, v':Variables, reply: Output) {
   // remove reply from replies
```

AcceptRequest ProcessRequest DeliverReply Server

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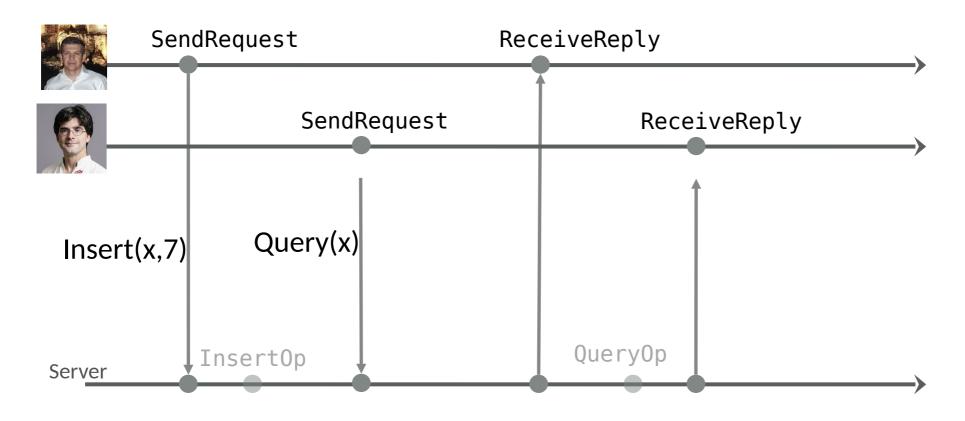


Example run



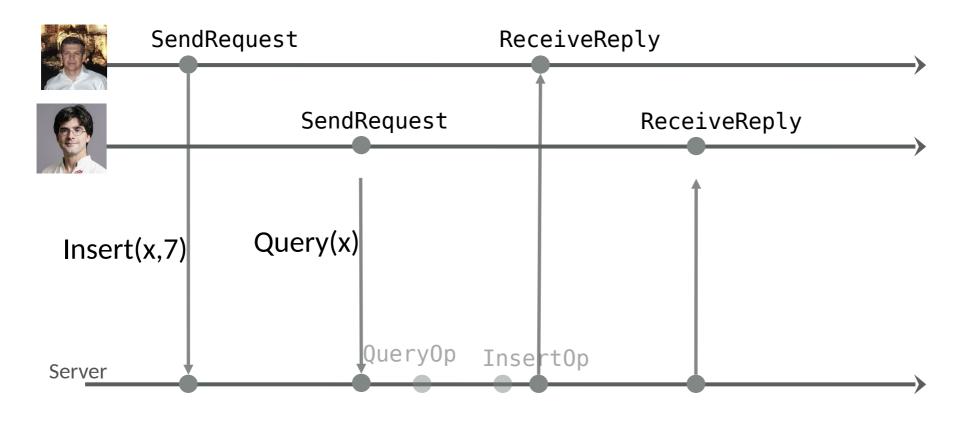


Example run #2



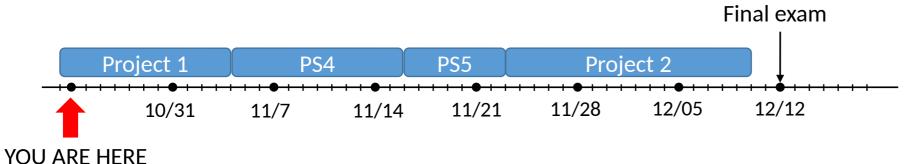


Example run #2



Administrivia

- Project 1 has been released
 - Deadline: Nov 4
 - Groups of (up to 2)
- We will have Jon live with us next Monday!
- No lecture on Nov 2 and Nov 14
- Assignment timeline



01/20/2023



Dafny: finite set heuristics

```
predicate IsEven(x:int) {
  x/2*2==x
predicate IsModest(x:int) {
 0 <= x < 10
lemma IsThisSetFinite() {
  var modestEvens := set x | IsModest(x) &&
IsEven(x);
  assert modestEvens == \{0,2,4,6,8\};
                 Error: the result of a set comprehension must be finite, but Dafny's
                 heuristics can't figure out how to produce a bounded set of values for 'x'
```



Dafny: finite set heuristics

```
predicate IsEven(x:int) {
 x/2*2==x
predicate IsModest(x:int) {
 0 <= x < 10
function ModestNumbers() : set<int> {
 set x \mid 0 \le x < 10
lemma IsThisSetFinite() {
 var modestEvens := set x | x in ModestNumbers() &&
IsEven(x);
 assert modestEvens == \{0,2,4,6,8\};
```



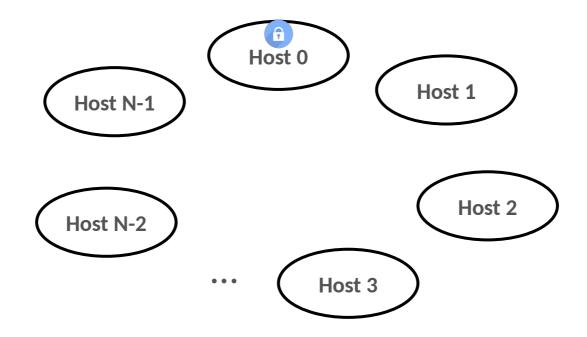
Distributed lock service

Differences from centralized lock server

- No centralized server that coordinates who holds the lock
 - The hosts pass the lock amongst themselves
- The hosts communicate via asynchronous messages
 - A single state machine transition cannot read/update the state of two hosts



Distributed lock server



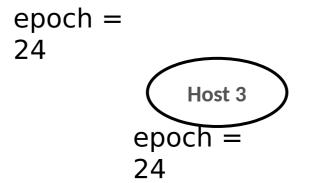
- N = numHosts, defined in network.t.dfy
- Messages are asynchronous (i.e. sending and receiving are two separate steps)



Distributed lock server

The lock is associated with a monotonically increasing epoch number





Accept an incoming message only if it has a higher epoch number than your current epoch



Distributed lock server

Safety property:

The desirable property is the same as the centralized lock server: at most one node holds the lock at any given time



Project files

Framework files

(trusted/immutable)

network.t.dfy

distributed_system.t. dfy

Host and proof files

(for you to complete)

host.v.dfy

exercise01.dfy