

# ECS 122A B01-B03 FQ 2021 Homework 06

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TOTAL POINTS

**97 / 100**

QUESTION 1

1 Q1 25 / 25

✓ + 25 pts All Correct

+ 5 pts Incorrect Explanation in Part 2

+ 0 pts Invalid/No Submission

+ 0 pts Wrong proof in both the parts

QUESTION 2

2 Q2 30 / 30

✓ + 30 pts All Correct

+ 25 pts All Correct but no mention of  
distances/table

+ 25 pts All Correct but no mention of shortest  
paths

+ 20 pts Incomplete Solution

QUESTION 3

3 Q3 25 / 25

✓ + 25 pts Correct

+ 20 pts Incomplete algorithm

+ 15 pts No description of algorithm

+ 0 pts Invalid/No Submission

QUESTION 4

4 Q4 17 / 20

+ 20 pts Correct

✓ + 10 pts Part 1 Correct

+ 10 pts Part 2 Correct

✓ + 7 pts Part 2 Partially Correct

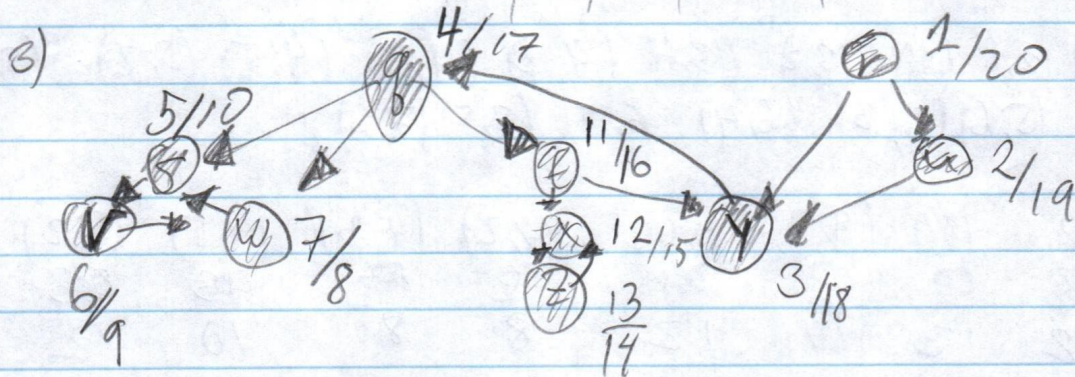
+ 5 pts Incorrect Explanation in Part 1

+ 7 pts Part 1 Partially Correct

1)  $\text{adj}[g] = [s, t, w]$      $\text{adj}[u] = [y]$      $\text{adj}[y] = [g]$   
 $\text{adj}[r] = [u, y]$      $\text{adj}[v] = [w]$      $\text{adj}[z] = [x]$   
 $\text{adj}[s] = [v]$      $\text{adj}[w] = [s]$   
 $\text{adj}[t] = [x, y]$      $\text{adj}[x] = [z]$

2) BFS,  $s$

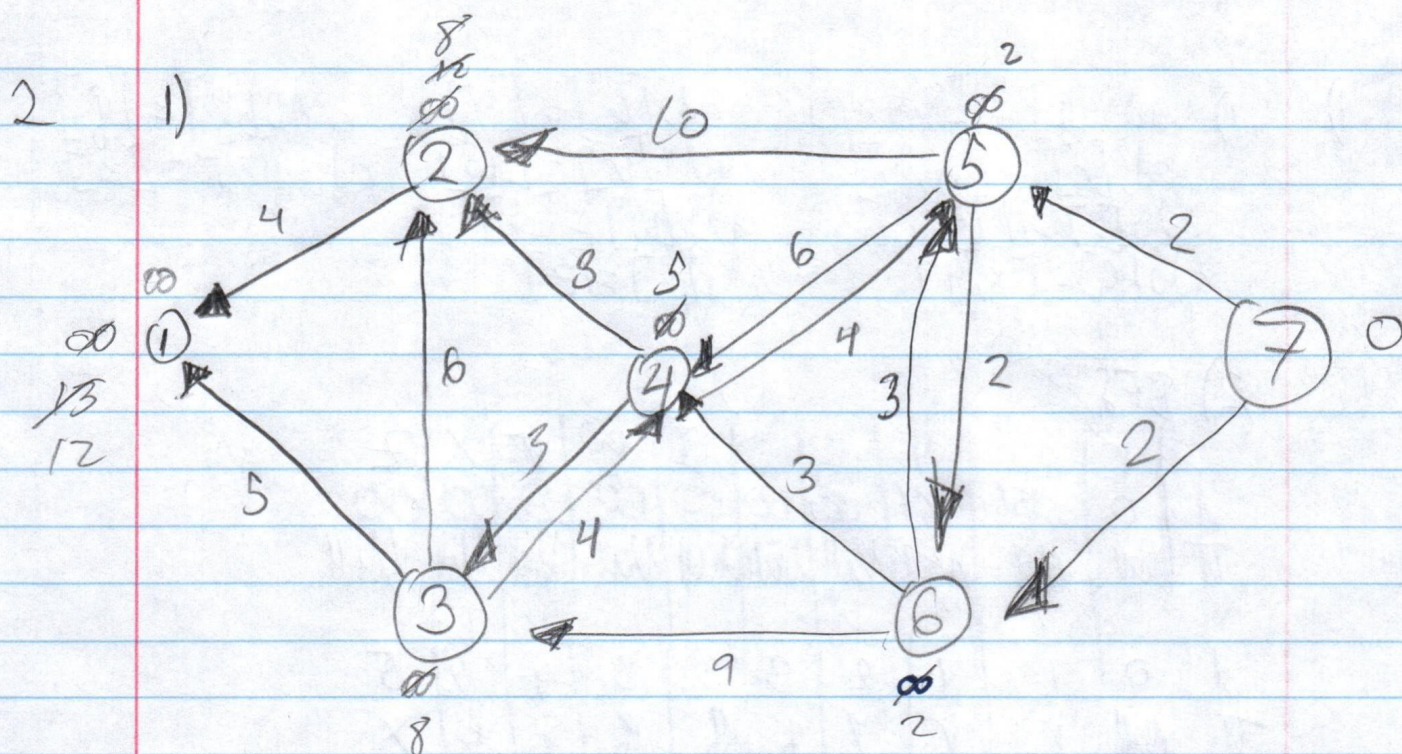
|   | r    | u            | y            | g            | s            | t            | w            | v            | x            | z            |
|---|------|--------------|--------------|--------------|--------------|--------------|--------------|--------------|--------------|--------------|
| d | 0    | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> |
| π | Null | Null         | Null         | Null         | Null         | Null         | Null         | Null         | Null         | Null         |
|   |      | r            |              |              |              |              |              |              |              |              |
| d | 0    | 1            | <del>1</del> | 2            | 3            | 3            | 3            | 4            | 4            | 5            |
| π | Null | r            | r            | y            | g            | g            | g            | s            | t            | x            |



1 Q1 25 / 25

✓ + 25 pts All Correct





2) (2,1) (3,1) (3,2) (3,4) (4,2) (4,3) (4,5) (5,2) (5,4)  
 (5,6) (6,3) (6,4) (6,5) (7,5) (7,6)

|              |              |              |              |              |              |              |              |
|--------------|--------------|--------------|--------------|--------------|--------------|--------------|--------------|
| (2,1)        | (3,1)        | (3,2)        | (3,4)        | (4,2)        | (4,3)        | (4,5)        | (5,2)        |
| <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> |
| 12           | 13           | 14           | 12           | 8            | 8            | 10           | 12           |

|              |              |              |              |              |              |              |
|--------------|--------------|--------------|--------------|--------------|--------------|--------------|
| (5,4)        | (5,6)        | (6,3)        | (6,4)        | (6,5)        | (7,5)        | (7,6)        |
| <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> | <del>∞</del> |
| 8            | 4            | 11           | 5            | 5            | 2            | 2            |

2 Q2 30 / 30

✓ + **30 pts** *All Correct*

+ **25 pts** All Correct but no mention of distances/table

+ **25 pts** All Correct but no mention of shortest paths

+ **20 pts** Incomplete Solution



3) Bellman Ford's algorithm as long as there are no negative cycles

$O(V \cdot E)$

$G(V, E)$

$d[s] = 0$

Every edge is negative weight

$S =$  Source vertex  
 $d =$  distance

for  $E$  in  $G[V, E]$

if  $d[v] < d[u]$

relax( $u, v, w$ )

if  $d[v] > d[u] + w(u, v)$

return false

Relax all edges  $K$  times

— Negative cycle

Relaxation

3 Q3 25 / 25

✓ + 25 pts *Correct*

+ 20 pts Incomplete algorithm

+ 15 pts No description of algoritihm

+ 0 pts Invalid/No Submission



4)  $G = (V, E, W)$   $S \in V$

$\pi[S] = 0$  to be visited again it must be a cyclic graph, & to be updated for  $\pi[S] = 0$  the edge must be negative  $\therefore$  it is a negative cycle through  $S$ .

2) Stop when find vertex  $S$  or the vertex is grayed & has been seen before. Stop & output if there is a negative cycle through  $S$  so that  $\pi[S]$  doesn't get updated



4 Q4 17 / 20

+ 20 pts Correct

✓ + 10 pts Part 1 Correct

+ 10 pts Part 2 Correct

✓ + 7 pts Part 2 Partially Correct

+ 5 pts Incorrect Explanation in Part 1

+ 7 pts Part 1 Partially Correct

+ 5 pts Incorrect Explanation in Part 2

+ 0 pts Invalid/No Submission

+ 0 pts Wrong proof in both the parts